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# C++

## IN A NUTSHELL

*A Language & Library Reference*

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*Ray Lischner*

# **C++ In a Nutshell**

**Ray Li schner**

Editor

**Jonat han Genni ck**

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## Preface

*C++ in a Nutshell* is a reference to the C++ language and library. Being a Nutshell guide, it is not a comprehensive manual, but it is complete enough to cover everything a working professional needs to know. Nonetheless, C++ is such a large and complex language that even this Nutshell guide is a large book.

This book covers the C++ standard, the international standard published as ISO/IEC 14882:1998(E), Programming Languages–C++, plus Technical Corrigendum 1. Many implementations of C++ extend the language and standard library. Except for brief mentions of language and library extensions in the appendixes, this book covers only the standard. The standard library is large—it includes strings, containers, common algorithms, and much more—but it omits much that is commonplace in computing today: concurrency, network protocols, database access, graphics, windows, and so on. See [Appendix B](#) for information about nonstandard libraries that provide additional functionality.

This book is a reference. It is not a tutorial. Newcomers to C++ might find portions of this book difficult to understand. Although each section contains some advice on idioms and the proper use of certain language constructs, the main focus is on the reference material. Visit <http://www.tempest-sw.com/cpp/> for links to sites and lists of books that are better suited for beginners.

## Structure of This Book

This book is divided into two interleaved sections that cover the language and the library, and a section of appendixes. Roughly speaking, the language is the part of C++ that does not require any additional `#include` headers

or files. The library is the part of C++ that is declared in the standard headers.

[Chapter 1-Chapter 7](#), [Chapter 11](#), and [Chapter 12](#) cover the language. The first seven chapters form the main language reference, organized by topic. It is customary for a programming reference to contain a formal grammar, and this book does so in [Chapter 12](#), which is organized alphabetically by keyword (with some additional entries for major syntactic categories, such as expressions). [Chapter 11](#) is a reference for the preprocessor.

[Chapter 13](#) is the library reference, organized alphabetically by header. [Chapter 8-Chapter 10](#) present an overview of the library and introduce the topics that span individual headers.

Sometimes, information is duplicated, especially in [Chapter 12](#). My goal has been to present information when you need it, where you need it. I tried to balance the need for a single, clear, complete description of each language feature with the desire to reduce the number of cross references you must chase before you can understand that language feature.

Here are more detailed descriptions of each chapter.

[Chapter 1](#) describes the basic rules for the C++ language: character sets, tokens, literals, and so on.

[Chapter 2](#) describes how objects, types, and namespaces are declared and how names are looked up.

[Chapter 3](#) describes operators, precedence, and type casts.

[Chapter 4](#) describes all the C++ statements.

[Chapter 5](#) describes function declarations and definitions, overload resolution, argument passing, and related topics.



[Chapter 6](#) describes classes (and unions and structures), members, virtual functions, inheritance, accessibility, and multiple inheritance.

[Chapter 7](#) describes class and function template declarations, definitions, instantiations, specializations, and how templates are used.

[Chapter 8](#) introduces the standard library and discusses some overarching topics, such as traits and allocators.

[Chapter 9](#) introduces the I/O portion of the standard library. Topics include formatted and unformatted I/O, stream buffers, and manipulators.

[Chapter 10](#) introduces the suite of container class templates, their iterators, and generic algorithms. This is the portion of the library that has traditionally been called the Standard Template Library (STL).

[Chapter 11](#) is an alphabetical reference for the preprocessor, which is part of the language, but with a distinct set of syntactic and semantic rules.

[Chapter 12](#) is an alphabetical reference for the language and grammar. Backus-Naur Form (BNF) syntax descriptions are given for each keyword and other language elements, with pointers to the first seven chapters for the main reference material.

[Chapter 13](#) is a reference for the entire standard library, organized alphabetically by header, and alphabetically by name within each header section.

[Appendix A](#) describes ways that some compilers extend the language: to satisfy customer need, to meet platform-specific requirements, and so on.

[Appendix B](#) describes a few interesting, open source C++ projects. You can find information about additional projects on this book's web site (<http://www.tempest-sw.com/cpp/>).

The [Glossary](#) defines some words and phrases used throughout this book and in the C++ community.

## About the Examples

Whenever possible, the examples in this book are complete, compilable programs. You can tell which examples fall into this category because they start with `#include` directives and contain a `main( )` function. You can download these examples as text files from the book's web site at <http://www.tempest-sw.com/cpp/> or from O'Reilly's catalog page for this book: <http://www.oreilly.com/catalog/cplsian/>.

Most examples are shortened to eliminate excess code that might interfere with the clarity of the example. In particular, these examples are fragments that lack a `main` function. Sometimes, an ellipsis indicates missing code, such as a function body. In other cases, the omissions are clear from the context. Most abbreviated examples have complete and compilable versions available for download.

All of the examples have been checked with several different compilers, including Comeau Computing's compiler with the Dinkumware standard library (widely acknowledged as the most complete and correct implementations of the C++ standard). Not all compilers can compile all the examples due to limitations and bugs in the compilers and libraries. For best results, try to work with the latest version of your compiler. Recent releases of several major compilers have made dramatic progress toward conformance with the standard. When possible, I have tried to alter the example files to work around the bugs without interfering with the intent of the example.

I have checked all the examples with the following compilers:

### Linux

- Borland Kylix 3.0

- Comeau 4.3.0.1
- GNU 3.2
- Intel 7.0

#### Microsoft Windows

- Borland C++ Builder 6.4
- Metrowerks CodeWarrior 8.3
- Microsoft Visual Studio.NET 7.0

# Conventions Used in This Book

This book uses the following conventions:

## Constant Width

Used for identifiers and symbols, including all keywords. In the language reference, constant width shows syntax elements that must be used exactly as shown. For example, the `if` keyword, parentheses, and `else` keyword must be used exactly as follows:

```
if ( condition ) statement else statement
```

A function name that is followed by parentheses refers to a function call, typically to obtain the function result. The function name without the parentheses refers to the function in more general terms. For example:

The empty function returns `true` if the container is empty, e.g., `size( ) == 0`.

## Constant Width Italic

Used in the language reference chapters for syntax elements that must be replaced by your code. In the previous example, you must supply the *condition* and the two *statements*.

## Constant Width Bold

Used in examples to highlight key lines, and in complex declarations to highlight the name being declared. In some C++ declarations, especially for templates, the name gets buried in the middle of the declaration and can be hard to spot.

## Italic

Used in the language reference for nonterminal syntax elements. Italic is also used for filenames, URLs,

emphasis, and for the first use of a technical term.

. . .

Indicates statements and declarations that have been removed for the sake of brevity and clarity. An ellipsis is also a symbol in C++, but context and comments make it clear when an ellipsis is a language element and when it represents omitted code.

[first, last)

Indicates a range of values from first to last, including first and excluding last.

### **Ti p**

This icon indicates a tip, suggestion, or general note.

### **War ni ng**

This icon indicates a warning or caution.



This icon indicates an issue or feature that might affect the portability of your code. In particular, some aspects of C++ are implementation-defined, such as the size of an integer, which allows the compiler or library author to decide what the best implementation should be.



## For More Information

Visit the *C++ in a Nutshell* web site at <http://www.tempest-sw.com/cpp/> to find links to newsgroups, frequently asked questions, tool and library web sites, free compilers, open source projects, other C++ books, and more. The web site also has information about the ongoing activities of the C++ Standardization Committee.

If you are a glutton for punishment, or if you need more details than are provided in this book, you might want to read the actual standard: ISO/IEC 14882:1998(E), *Programming Language—C++*. The standard is not easy to read, and even its authors sometimes disagree on its interpretation. Nonetheless, it is the one specification for the C++ language, and all other books, including this one, are derivatives, subject to error and misinterpretation. The C++ standard library includes the entire C standard library, which is documented in ISO/IEC 9899:1990, *Programming Language—C*, plus Amendment 1:1995(E), *C Integrity*.

The C and C++ standards are evolving documents; the committees meet regularly to review defect reports and proposals for language extensions. As I write this, the C++ standard committee has approved a technical corrigendum (TC1), which is an update to the C++ standard that corrects defects and removes ambiguities in the original standard. TC1 is winding its way through the ISO bureaucracy. By the time you read this, TC1 will have probably completed its journey and been added to the official standard for the C++ programming language. The book's web site has up-to-date information about the status of the C++ and C standards.

## Comments and Questions

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# Chapter 1. Language Basics

C++ is a case-sensitive, free-form programming language that supports procedural, object-oriented, and generic programming. This chapter presents the basic rules for the language, such as lexical rules and basic syntax elements.

## Compilation Steps

A C++ source file undergoes many transformations on its way to becoming an executable program. The initial steps involve processing all the `#include` and conditional preprocessing directives to produce what the standard calls a *translation unit*. Translation units are important because they have no dependencies on other files. Nonetheless, programmers still speak in terms of source files, even if they actually mean translation units, so this book uses the phrase *source file* because it is familiar to most readers. The term "translation" encompasses compilation and interpretation, although most C++ translators are compilers. This section discusses how C++ reads and compiles (translates) source files (translation units).

A C++ program can be made from many source files, and each file can be compiled separately. Conceptually, the compilation process has several steps (although a compiler can merge or otherwise modify steps if it can do so without affecting the observable results):

1. Read physical characters from the source file and translate the characters to the source character set (described in [Section 1.4](#) later in this chapter). The source "file" is not necessarily a physical file; an implementation might, for example, retrieve the source from a database. Trigraph sequences are reduced to their equivalent characters (see [Section 1.6](#) later in this

chapter). Each native end-of-line character or character sequence is replaced by a newline character.

2. If a backslash character is followed immediately by a newline character, delete the backslash and the newline. The backslash/newline combination must not fall in the middle of a universal character (e.g., `\u1234`) and must not be at the end of a file. It can be used in a character or string literal, or to continue a preprocessor directive or one-line comment on multiple lines. A non-empty file must end with a newline.
3. Partition the source into preprocessor tokens separated by whitespace and comments. A preprocessor token is slightly different from a compiler token (see the next section, [Section 1.2](#)). A preprocessor token can be a header name, identifier, number, character literal, string literal, symbol, or miscellaneous character. Each preprocessor token is the longest sequence of characters that can make up a legal token, regardless of what comes after the token.
4. Perform preprocessing and expand macros. All `#include` files are processed in the manner described in steps 1-4. For more information about preprocessing, see [Chapter 11](#).
5. Convert character and string literals to the execution character set.
6. Concatenate adjacent string literals. Narrow string literals are concatenated with narrow string literals. Wide string literals are concatenated with wide string literals. Mixing narrow and wide string literals results in an error.
7. Perform the main compilation.
8. Combine compiled files. For each file, all required template instantiations (see [Chapter 7](#)) are identified,

and the necessary template definitions are located and compiled.

9. Resolve external references. The compiled files are linked to produce an executable image.



## Tokens

All source code is divided into a stream of *tokens*. The compiler tries to collect as many contiguous characters as it can to build a valid token. (This is sometimes called the "max munch" rule.) It stops when the next character it would read cannot possibly be part of the token it is reading.

A token can be an identifier, a reserved keyword, a literal, or an operator or punctuation symbol. Each kind of token is described later in this section.

Step 3 of the compilation process reads *preprocessor tokens*. These tokens are converted automatically to ordinary compiler tokens as part of the main compilation in Step 7. The differences between a preprocessor token and a compiler token are small:

- The preprocessor and the compiler might use different encodings for character and string literals.
- The compiler treats integer and floating-point literals differently; the preprocessor does not.
- The preprocessor recognizes `< header >` as a single token (for `#include` directives); the compiler does not.

## Identifiers

An identifier is a name that you define or that is defined in a library. An identifier begins with a nondigit character and is followed by any number of digits and nondigits. A nondigit character is a letter, an underscore, or one of a set of universal characters. The exact set of nondigit universal characters is defined in the C++ standard and in ISO/IEC PDTR 10176. Basically, this set contains the universal characters that represent letters. Most programmers

restrict themselves to the characters a...z, A...Z, and underscore, but the standard permits letters in other languages.



Not all compilers support universal characters in identifiers.

Certain identifiers are reserved for use by the standard library:

- Any identifier that contains two consecutive underscores (like `__this`) is reserved, that is, you cannot use such an identifier for macros, class members, global objects, or anything else.
- Any identifier that starts with an underscore, followed by a capital letter (A-Z) is reserved.
- Any identifier that starts with an underscore is reserved in the global namespace. You can use such names in other contexts (i.e., class members and local names).
- The C standard reserves some identifiers for future use. These identifiers fall into two categories: function names and macro names. Function names are reserved and should not be used as global function or object names; you should also avoid using them as "C" linkage names in any namespace. Note that the C standard reserves these names regardless of which headers you `#include`. The reserved function names are:
  - `is` followed by a lowercase letter, such as `isblank`
  - `mem` followed by a lowercase letter, such as `memxyz`
  - `str` followed by a lowercase letter, such as `strtof`
  - `to` followed by a lowercase letter, such as `toxyz`

- `wcs` followed by a lowercase letter, such as `wcstof`
- In `<cmath>` with `f` or `l` appended, such as `cosf` and `sinl`
- Macro names are reserved in all contexts. Do not use any of the following reserved macro names:
  - Identifiers that start with `E` followed by a digit or an uppercase letter
  - Identifiers that start with `LC_` followed by an uppercase letter
  - Identifiers that start with `SIG` or `SIG_` followed by an uppercase letter

## Keywords



A *keyword* is an identifier that is reserved in all contexts for special use by the language. The following is a list of all the reserved keywords. (Note that some compilers do not implement all of the reserved keywords; these compilers allow you to use certain keywords as identifiers. See [Section 1.5](#) later in this chapter for more information.)

<code>and</code>	<code>continue</code>	<code>goto</code>	<code>public</code>	<code>try</code>
<code>and_eq</code>	<code>default</code>	<code>if</code>	<code>register</code>	<code>typedef</code>
<code>asm</code>	<code>delete</code>	<code>inline</code>	<code>reinterpret_cast</code>	<code>typeid</code>
<code>auto</code>	<code>do</code>	<code>int</code>	<code>return</code>	<code>typename</code>
<code>bitand</code>	<code>double</code>	<code>long</code>	<code>short</code>	<code>union</code>
<code>bitor</code>	<code>dynamic_cast</code>	<code>mutable</code>	<code>signed</code>	<code>unsigned</code>
<code>bool</code>	<code>else</code>	<code>namespace</code>	<code>sizeof</code>	<code>using</code>

break	enum	new	static	virtual
case	explicit	not	static_cast	void
catch	export	not_eq	struct	volatile
char	extern	operator	switch	wchar_t
class	false	or	template	while
compl	float	or_eq	this	xor
const	for	private	throw	xor_eq
const_cast	friend	protected	true	

## Literals

A *literal* is an integer, floating-point, Boolean, character, or string constant.

### Integer Literals

An integer literal can be a decimal, octal, or hexadecimal constant. A prefix specifies the base or radix: `0x` or `0X` for hexadecimal, `0` for octal, and nothing for decimal. An integer literal can also have a suffix that is a combination of `U` and `L`, for unsigned and long, respectively. The suffix can be uppercase or lowercase and can be in any order. The suffix and prefix are interpreted as follows:

- If the suffix is `UL` (or `ul`, `LU`, etc.), the literal's type is unsigned long.
- If the suffix is `L`, the literal's type is long or unsigned long, whichever fits first. (That is, if the value fits in a long, the type is long; otherwise, the type is unsigned long. An error results if the value does not fit in an unsigned long.)

- If the suffix is U, the type is unsigned or unsigned long, whichever fits first.
- Without a suffix, a decimal integer has type int or long, whichever fits first.
- An octal or hexadecimal literal has type int, unsigned, long, or unsigned long, whichever fits first.



Some compilers offer other suffixes as extensions to the standard. See [Appendix A](#) for examples.

Here are some examples of integer literals:

```

314           // Legal
314u         // Legal
314LU        // Legal
0xFeeL      // Legal
0ul         // Legal
078         // Illegal: 8 is not an octal digit
032UU       // Illegal: cannot repeat a suffix

```

## Fl o a t i n g - p o i n t L i t e r a l s

A floating-point literal has an integer part, a decimal point, a fractional part, and an exponent part. You must include the decimal point, the exponent, or both. You must include the integer part, the fractional part, or both. The signed exponent is introduced by e or E. The literal's type is double unless there is a suffix: F for type float and L for long double. The suffix can be uppercase or lowercase.

Here are some examples of floating-point literals:

```

3.14159      // Legal
.314159F     // Legal
314159E-5L   // Legal
314.         // Legal
314E         // Illegal: incomplete exponent
314f         // Illegal: no decimal or exponent
.e24        // Illegal: missing integer or fraction

```

## Boo l e a n L i t e r a l s

There are two Boolean literals, both keywords: `true` and `false`.

## Character Literals

Character literals are enclosed in single quotes. If the literal begins with `L` (uppercase only), it is a wide character literal (e.g., `L'x'`). Otherwise, it is a narrow character literal (e.g., `'x'`). Narrow characters are used more frequently than wide characters, so the "narrow" adjective is usually dropped.



The value of a narrow or wide character literal is the value of the character's encoding in the execution character set. If the literal contains more than one character, the literal value is implementation-defined. Note that a character might have different encodings in different locales. Consult your compiler's documentation to learn which encoding it uses for character literals.

A narrow character literal with a single character has type `char`. With more than one character, the type is `int` (e.g., `'abc'`). The type of a wide character literal is always `wchar_t`.

### Tip

In C, a character literal always has type `int`. C++ changed the type of character literals to support overloading, especially for I/O (e.g., `cout << '\n'` starts a new line and does not print the integer value of the newline character).





A character literal can be a plain character (e.g., 'x'), an escape sequence (e.g., '\b'), or a universal character (e.g., '\u03C0'). [Table 1-1](#) lists the possible escape sequences. Note that you must use an escape sequence for a backslash or single-quote character literal. Using an escape for a double quote or question mark is optional. Only the characters shown in [Table 1-1](#) are allowed in an escape sequence. (Some compilers extend the standard and recognize other escape sequences.)

*Table 1-1. Character escape sequences*

Escape sequence	Meaning
\\	\ character
\'	' character
\"	" character
\?	? character (used to avoid creating a trigraph, e.g., \?\?-)
\a	Alert or bell
\b	Backspace
\f	Form feed
\n	Newline
\r	Carriage return
\t	Horizontal tab
\v	Vertical tab
\ooo	Octal number of one to three digits

Escape sequence	Meaning
<code>\x hh . . .</code>	Hexadecimal number of one or more digits

## String literals

String literals are enclosed in double quotes. A string contains characters that are similar to character literals: plain characters, escape sequences, and universal characters. A string cannot cross a line boundary in the source file, but it can contain escaped line endings (backslash followed by newline).



A wide string literal is prefaced with `L` (always uppercase). In a wide string literal, a single universal character always maps to a single wide character. In a narrow string literal, the implementation determines whether a universal character maps to one or multiple characters (called a multibyte character). See [Chapter 8](#) for more information on multibyte characters.

Two adjacent string literals (possibly separated by whitespace, including new lines) are concatenated at compile time into a single string. This is often a convenient way to break a long string across multiple lines. Do not try to combine a narrow string with a wide string in this way.

After concatenating adjacent strings, the null character (`'\0'` or `L'\0'`) is automatically appended after the last character in the string literal.

Here are some examples of string literals. Note that the first three form identical strings.

```
"hello, reader"
"hello, \
reader"
"hello, " "rea" "der"
```

```
"Alert: \a; ASCII tab: \010; portable tab: \t"
"illegal: unterminated string
L"string with \"quotes\""
```

A string literal's type is an array of `const char`. For example, `"string"`'s type is `const char[7]`. Wide string literals are arrays of `const wchar_t`. All string literals have static lifetimes (see [Chapter 2](#) for more information about lifetimes).

As with an array of `const` anything, the compiler can automatically convert the array to a pointer to the array's first element. You can, for example, assign a string literal to a suitable pointer object:

```
const char* ptr;
ptr = "string";
```

As a special case, you can also convert a string literal to a non-`const` pointer. Attempting to modify the string results in undefined behavior. This conversion is deprecated, and well-written code does not rely on it.

## Symbols

Nonalphabetic symbols are used as operators and as punctuation (e.g., statement terminators). Some symbols are made of multiple adjacent characters. The following are all the symbols used for operators and punctuation:

{	(	%,	.	^	.	=	!=	--	&=
}	)	%,%,	+	&	.*	==	<<	+=	=
[	<:	;	-		->	<	>>	*=	^=
]	:>	:	*	?	->*	>	<<=	/=	++
#	<%	...	/	:	~	<=	>>=	%=	--

##	%>	,	%	::	!	>=				
----	----	---	---	----	---	----	--	--	--	--

You cannot insert whitespace between characters that make up a symbol, and C++ always collects as many characters as it can to form a symbol before trying to interpret the symbol. Thus, an expression such as `x+++y` is read as `x ++ + y`. A common error when first using templates is to omit a space between closing angle brackets in a nested template instantiation. The following is an example with that space:

```
std::list<std::vector<int> > list;↑
                        Note the space here.
```

The example is incorrect without the space character because the adjacent greater than signs would be interpreted as a single right-shift operator, not as two separate closing angle brackets. Another, slightly less common, error is instantiating a template with a template argument that uses the global scope operators:

```
::std::list< ↑ ::std::list<int> > list;↑
                ↑
                Space here                and here
```

Again, a space is needed, this time between the angle-bracket (<) and the scope operator (::), to prevent the compiler from seeing the first token as <: rather than <. The <: token is an alternative token, as described in [Section 1.5](#) later in this chapter.

# Comments

Comments start with `/*` and end with `*/`. These comments do not nest. For example:

```
/* this is a comment /* still a comment */  
int not_in_a_comment;
```

A comment can also start with `//`, extending to the end of the line. For example:

```
const int max_widget = 42; // Largest size of a widget
```

Within a `/*` and `*/` comment, `//` characters have no special meaning. Within a `//` comment, `/*` and `*/` have no special meaning. Thus, you can "nest" one kind of comment within the other kind. For example:

```
/* Comment out a block of code:  
const int max_widget = 42; // Largest size of a widget  
*/  
  
/** Inhibit the start of a block comment  
const int max_widget = 10; // Testing smaller widget limit  
**/
```

A comment is treated as whitespace. For example, `str/*comment*/ing` describes two separate tokens, `str` and `ing`.

# Character Sets



The character sets that C++ uses at compile time and runtime are implementation-defined. A source file is read as a sequence of characters in the *physical character set*. When a source file is read, the physical characters are mapped to the compile-time character set, which is called the *source character set*. The mapping is implementation-defined, but many implementations use the same character set.

At the very least, the source character set always includes the characters listed below. The numeric values of these characters are implementation-defined.

Space

Horizontal tab

Vertical tab

Form feed

Newline

a ... z

A ... Z

0 ... 9

\_ { } [ ] # ( ) < > % : ; . ? \* + - / ^ & | ~ !

= , \ " '

The runtime character set, called the *execution character set*, might be different from the source character set (though it is often the same). If the character sets are different, the compiler automatically converts all character and string literals from the source character set to the execution character set. The *basic execution character set* includes all the characters in the source character set, plus the characters listed below. The execution character set is



a superset of the basic execution character set; additional characters are implemented-defined and might vary depending on locale.

Alert

Backspace

Carriage return

Null

Conceptually, source characters are mapped to Unicode (ISO/IEC 10646) and from Unicode to the execution character set. You can specify any Unicode character in the source file as a *universal character* in the form `\u XXXX` (lowercase u) or `\U XXXXXXXX` (uppercase U), in which `0000 XXXX` or `XXXXXXXX` is the hexadecimal value for the character. Note that you must use exactly four or eight hexadecimal digits. You cannot use a universal character to specify any character that is in the source character set or in the range `0-0x20` or `0x7F-0x9F` (inclusive).



How universal characters map to the execution character set is implementation-defined. Some compilers don't recognize universal characters at all, or support them only in limited contexts.

Typically, you would not write a universal character manually. Instead, you might use a source editor that lets you edit source code in any language, and the editor would store source files in a manner that is appropriate for a particular compiler. When necessary, the editor would write universal character names for characters that fall outside the compiler's source character set. That way, you might write the following in the editor:

```
const long double pi = 3.1415926535897932385L;
```

and the editor might write the following in the source file:

```
const long double \u03c0 = 3.1415926535897932385L;
```

The numerical values for characters in all character sets are implementation-defined, with the following restrictions:

- The null character always has a value that contains all zero bits.
- The digit characters have sequential values, starting with 0.

The space, horizontal tab, vertical tab, form feed, and newline characters are called *whitespace characters*. In most cases, whitespace characters only separate tokens and are otherwise ignored. (Comments are like whitespace; see [Section 1.3](#) earlier in this chapter.)

# Alternative Tokens

Some symbols have multiple representations, as shown in [Table 1-2](#). These *alternative tokens* have no special meaning in a character or string literal. They are merely alternative representations of common symbols. Most programmers do not use alternative tokens, especially the nonalphabetic ones. Some programmers find `and`, `or`, and `not` to be easier to read and understand than `&&`, `||`, and `!`.

*Table 1-2. Alternative tokens*

Alternative token	Primary token
<code>&lt;%</code>	<code>{</code>
<code>%&gt;</code>	<code>}</code>
<code>&lt;:</code>	<code>[</code>
<code>;&gt;</code>	<code>]</code>
<code>%:</code>	<code>#</code>
<code>%.%:</code>	<code>##</code>
<code>and</code>	<code>&amp;&amp;</code>
<code>and_eq</code>	<code>&amp;=</code>
<code>bitand</code>	<code>&amp;</code>
<code>bitor</code>	<code> </code>
<code>compl</code>	<code>~</code>
<code>not</code>	<code>!</code>

Alternative token	Primary token
not_eq	!=
or	
or_eq	=
xor	^
xor_eq	^=



Many compilers do not support some or even all of the alternative tokens. In particular, some compilers do not treat alternative keywords (and, or, etc.) as reserved keywords, but allow you to use them as identifiers. Fortunately, this problem is becoming less common as more vendors hew closer to the standard. Compilers that do not support the alternative keywords might provide them as macros in the `<ciso646>` header. See [Chapter 13](#) for more information.

# Trigraphs

A few characters have an alternative representation, called a trigraph sequence. A *trigraph* is a three-character sequence that represents a single character. The sequence always starts with two question marks. The third character determines which character the sequence represents. All the trigraph sequences are shown in [Table 1-3](#). If the third character is not one of those in the table, the sequence is not a trigraph and is left alone. For example, the characters ???- represent the two characters ?~. Note that trigraphs are expanded anywhere they appear, including within string literals and character literals, in comments, and in preprocessor directives.

*Table 1-3. Trigraph sequences*

Tr i gr ap h	Repl acement
??=	#
??/	\
??'	^
??(	[
??)	]
??!	
??<	{
??>	}
??-	~



Not all compilers support trigraphs. Some compilers require an extra switch or option. Others use a separate program to convert all trigraphs to their equivalent characters.

Do not use trigraphs. They are confusing to read and use. If you ever write multiple, adjacent question marks, make sure you are not accidentally writing a trigraph. For example:

```
std::cout << "What do you mean??!\n";
```

To avoid the trigraph interpretation, separate the string literal:

```
std::cout << "What do you mean?" "?!!\n";
```

or escape the question marks:

```
std::cout << "What do you mean\\?\\?!!\n";
```

## Chapter 2. Declarations

A C++ source file contains a series of zero or more *declarations*. A declaration can be a function, type, object (constant or variable), namespace, template, or a related entity. The first part of this chapter covers general topics that pertain to all declarations. The second part discusses types, objects, and namespaces specifically. Function, class, and template declarations each get their own chapters: [Chapter 5](#) covers functions; [Chapter 6](#) covers classes and friends; and [Chapter 7](#) covers template declarations, specializations, and instantiations.

The syntax descriptions in this chapter are informal. See [Chapter 12](#) for a precise BNF grammar.

### Declarations and Definitions

A declaration is the all-encompassing term for anything that tells the compiler about an identifier. In order to use an identifier, the compiler must know what it means: is it a type name, a variable name, a function name, or something else? Therefore, a source file must contain a declaration (directly or in an `#include` file) for every name it uses.

### Definitions

A *definition* defines the storage, value, body, or contents of a declaration. The difference between a declaration and a definition is that a declaration tells you an entity's name and the external view of the entity, such as an object's type or a function's parameters, and a definition provides the internal workings of the entity: the storage and initial value of an object, a function body, and so on.

In a single source file, there can be at most one definition of an entity. In an entire program, there must be exactly one definition of each function or object used in the program, except for inline functions; an inline function must be defined in every source file that uses the function, and the definitions must all be identical.

A program can have more than one definition of a given class, enumeration, inline function, or template, provided the definitions are in separate source files, and each source file has the same definition.

These rules are known as the One Definition Rules, or ODR.

Before you can use an entity (e.g., calling a function or referring to an object), the compiler needs the entity's declaration, but not necessarily its definition. You can use a class that has an incomplete declaration in some contexts, but usually you need a complete definition. (See [Chapter 6](#) for details about incomplete classes.) The complete program needs definitions for all the declared entities, but those definitions can often reside in separate source files. The convention is to place the declarations for classes, functions, and global objects in a header file (whose name typically ends with `.h` or `.hpp`), and their definitions in a source file (whose name typically ends with `.cpp`, `.c`, or `.C`). Any source file that needs to use those entities must `#include` the header file. Templates have additional complications concerning declarations and definitions. (See [Chapter 7](#) for details.)

In this and subsequent chapters, the description of each entity states whether the entity (type, variable, class, function, etc.) has separate definitions and declarations, states when definitions are required, and outlines any other rules pertaining to declarations and definitions.

## **Ambiguity**



Some language constructs can look like a declaration or an expression. Such ambiguities are always resolved in favor of declarations. A related rule is that a declaration that is a type specifier followed by a name and empty parentheses is a declaration of a function that takes no arguments, not a declaration of an object with an empty initializer. (See [Section 2.6.3](#) later in this chapter for more information about empty initializers.) [Example 2-1](#) shows some examples of how declarations are interpreted.

### *Example 2-1. Disambiguating declarations*

```
#include <iostream>
#include <ostream>

class T
{
public:
    T( )    { std::cout << "T( )\n"; }
    T(int) { std::cout << "T(int)\n"; }
};

int a, x;

int main( )
{
    T(a);          // Variable named a of type T, not an invocation
of the T(int)    // constructor
    T b( );       // Function named b of no arguments, not a
variable named b of // type T
    T c(T(x));   // Declaration of a function named c, with one
argument of       // type T
}
```

The last item in [Example 2-1](#) deserves further explanation. The function parameter `T(x)` could be interpreted as an expression: constructing an instance of `T` with the argument `x`. Or it could be interpreted as a declaration of a function parameter of type `T` named `x`, with a redundant set of parentheses around the parameter name. According to

the disambiguation rule, it must be a declaration, not an expression. This means that the entire declaration cannot be the declaration of an object named `c`, whose initializer is the expression `T(x)`. Instead, it must be the declaration of a function named `c`, whose parameter is of type `T`, named `x`.

If your intention is to declare an object, not a function, the simplest way to do this is not to use the function-call style of type `cast`. Instead, use a keyword cast expression, such as `static_cast<>`. (See [Chapter 3](#) for more information about type casts.) For example:

```
T c(static_cast<T>(x)); // Declares an object named c whose
initial value is
                        // x, cast to type T
```

This problem can crop up when you least expect it. For example, suppose you want to construct a vector of integers by reading a series of numbers from the standard input. Your first attempt might be to use an `istream_iterator`:

```
using namespace std;
vector<int> data(istream_iterator<int>(cin),
istream_iterator<int>( ));
```

This declaration actually declares a function named `data`, which takes two parameters of type `istream_iterator<int>`. The first parameter is named `cin`, and the second is nameless. You can force the compiler to interpret the declaration as an object definition by enclosing one or more arguments in parentheses:

```
using namespace std;
vector<int> data((istream_iterator<int>(cin)),
(istream_iterator<int>( )));
```

or by using additional objects for the iterators:

```
std::istream_iterator<int> start(std::cin), end;
std::vector<int> data(start, end);
```

## Scope

A *scope* is a region of source code that contains declarations. Every declaration adds a name to a scope, and every use of a name requires the compiler to identify which scope contains that name's declaration. Sometimes you tell the compiler exactly which scope contains the name, and at other times the compiler determines the scope. Once the compiler knows the scope, it can look up the name to learn what the name is (object, function, class, etc.) and how the name can be used. Thus, you can think of a scope as a dictionary of names mapped to declarations.

A scope can be named or unnamed. Classes and namespaces (see [Section 2.7](#) later in this chapter) define named scopes. Statement blocks, function bodies, and unnamed namespaces define unnamed scopes. You can *qualify* a name with a scope name to tell the compiler where to look up the qualified name, but you cannot qualify names from unnamed scopes. In a typical program, most names are unqualified, so the compiler must determine which scope declares the name. (See [Section 2.3](#) later in this chapter.)

Scopes can be nested, and names in inner scopes can hide names that are declared in outer scopes. [Example 2-2](#) illustrates abuses of the simple scope rules: the body of the for loop is a scope, in which the variable `x` is declared as an `int`; the `if` statement creates a nested scope, in which another declaration of `x` hides the outer `x`. The reference to `x` at the end of `main` is invalid: no `x` is in scope at that point.

*Example 2-2. Names in inner scopes can hide names in outer scopes*

```
#include <iostream>
#include <ostream>
```

```

int main( )
{
    for (int i = 0; i < 100; ++i)
    {
        int x = 42;
        if (x < i)
        {
            double x = 3.14;
            std::cout << x; // Prints 3.14
        }
        std::cout << x;    // Prints 42
    }
    std::cout << x;      // Error: no x declared in this scope
}

```

At the same scope level, you cannot have multiple declarations for the same name, unless every declaration of that name is for an overloaded function or function template, or if the declarations are identical typedef declarations.

[Example 2-2](#) shows that a name can be hidden by a different declaration in an inner scope level. Also, a class name or enumeration type name can be hidden by an object, function, or enumerator at the same scope level. For example:

```

enum x { a, b, c };
const int x = 42;    // OK: hides enum x
const int a = 10;   // Error: int cannot hide enumerator a
{
    const int a = 10; // OK: inner scope can hide outer a
}

```

Different entities (functions, statements, classes, namespaces, etc.) establish scope in different ways. The description of each entity (in this and subsequent chapters) includes scoping rules. The general rule of thumb is that curly braces delimit a scope region. The outermost scope region is outside of all the curly braces and is called the *global scope*. (See [Section 2.7](#) later in this chapter.)

## Name Lookup

When the compiler reads an identifier, it must look up the identifier to determine which declaration it comes from. In most cases, you can readily tell which identifier is which, but it is not always so simple. A small mistake can sometimes lead to code that compiles successfully but runs incorrectly because an identifier refers to a different object from the one you intended. To understand name lookup fully, you must first understand namespaces (covered later in this chapter), functions ([Chapter 5](#)), classes ([Chapter 6](#)), and templates ([Chapter 7](#)).

Name lookup takes place before overloaded functions are resolved and before the access level of class members is checked. If a name is found in an inner scope, the compiler uses that declaration, even if a better declaration would be found in an outer scope. [Example 2-3](#) shows how problems can arise when an overloaded function is declared in more than one namespace. The function `func(int)` is global, and `func(double)` is defined in namespace `N`. Inside `call_func`, the compiler looks up the name `func` by searching first in the local scope (that is, the function body), then in namespace `N`, where it finds `func(double)`. Name lookup stops at that point because the compiler found a match. Therefore, `func(3)` converts `3` to type `double` and calls `func(double)`. The main function brings all the overloaded `func` functions into its scope (with `using` declarations, which are described at the end of this chapter), so name lookup can find the best match, which is `func(int)`.

*Example 2-3. Name lookup trumps overload resolution*

```
void func(int i)
{
    std::cout << "int: " << i << '\n';
}
```

```

namespace N {
    void func(double d)
    {
        std::cout << "double: " << std::showpoint << d << '\n';
    }

    void call_func( )
    {
        // Even though func(int) is a better match, the compiler
        // finds
        // N::func(double) first.
        func(3);
    }
}

int main( )
{
    N::call_func( );          // Prints "double: 3.000000"
    using N::func;
    using ::func;
    // Now all overloaded func( )s are at the same scope level.
    func(4);                 // Prints "int: 4"
}

```

Refer to [Chapter 5](#) for more information about overloaded functions and to [Chapter 6](#) for information about access levels in a class declaration.

## Qualified Name Lookup

To specify a particular namespace for looking up a name, qualify the name using the scope operator (::). A name that follows the global scope operator (the unary ::) is looked up in the global scope (outside of all namespaces). The name must be declared in the global scope, not in a nested scope, or the name must have been introduced into the global scope by a using directive or using declaration (see [Section 2.7](#) later in this chapter).

Use the global scope operator to access names that have been hidden by an inner scope. [Example 2-4](#) shows this use of the :: operator to access the global x from within main after an inner x has been declared.

### *Example 2-4. The global scope operator*

```
#include <iostream>
#include <ostream>

int x = 42;

int main( )
{
    double x = 3.1415;           // Hides the global x
    std::cout << x << '\n';     // Prints 3.1415
    std::cout << ::x << '\n';  // Prints 42
}
```

The binary scope resolution operator (also ::) requires a namespace or class name as its lefthand operand, and an identifier as its righthand operand. The identifier is looked up in the scope of the lefthand namespace or class.

[Example 2-5](#) shows the scope resolution operator untangling a mess made by using the same names for different kinds of entities. Notice how the inner counter hides the outer counter, so the simple name counter refers to the int variable. The lefthand operand to ::, however, must be a class or namespace, so in the expression counter::c, the inner counter does not hide the outer counter.

### *Example 2-5. The scope resolution operator*

```
#include <iostream>
#include <ostream>

namespace n {
    struct counter {
        static int n;
    };
    double n = 2.71828;
}

int n::counter::n = 42; // Defines static data member

int main( )
{
    int counter = 0;    // Unrelated to n::counter
    int n = 10;        // Hides namespace n
}
```

```
::n::counter x;          // Refers to namespace n

std::cout << n::counter::n; // Prints 42
std::cout << n::n;          // Prints 2.71828
std::cout << x.n;          // Prints 42
std::cout << n;            // Prints 10
std::cout << counter;     // Prints 0
}
```

## Unqualified Name Lookup

The compiler looks up an unqualified name, that is, a bare identifier or operator symbol, in a series of scopes in order to find its declaration. The simple rule is that the innermost scope is searched first, and succeeding outer scopes are searched until a matching declaration is found. Additional named scopes are searched, depending on the context of the usage, as explained in this section and the next ([Section 2.3.3](#)). An associated simple rule is that a name must be declared before it is used, reading the source file from top to bottom.

In a class definition, the class is searched first; the declaration must appear before it is used, reading the class definition from top to bottom. The immediate base classes are searched next (in declaration order), and their base classes are searched. If the class is nested in another class, the containing class is searched, and its base classes are searched. If the class is local to a function, the block that contains the class is searched, then enclosing blocks are searched. Finally, the namespaces that contain the class declaration are searched. That is, the namespace that immediately contains the class is searched first; if that namespace is nested within another namespace, the outer namespace is searched next, and so on. If the class is nested, the namespaces that contain the outer class or classes are searched.



In the body of a member function, names are looked up first in the local scope and nested scopes of the function body. Then the class is searched; the name can be declared anywhere in the class definition, even if that declaration appears after the member function definition. (This is an exception to the "declare-before-use" rule.) Such names can be used in a parameter type, a default argument value, or the function body, but not in the function's return type. The name is then looked up in the manner described earlier for other names used in a class definition.

## Warning

The name lookup rules permit a member function to use a name that is declared later in the class definition, but the name lookup rules do not trump the syntax and parsing rules. In particular, the parser must be able to distinguish between names used as types from other names. In the following example, the name `big` is used as a type, and is declared as a type in the class definition, but when the compiler first sees the name `big` in the member function declaration, the only `big` in scope is the global object, so the compiler parses the function declaration incorrectly:

```
int big;
// typedef big float;
struct demo {
    void func(big); // Error
    typedef big long;
};
```

If the declaration `int big` were a typedef instead, the declaration of `func` would be parsed as intended, and name lookup would find the nested type `demo::big` for the parameter type of `func`. A simpler solution is to move the `typedef big long;` to the start of the class definition, which is the style used throughout this book.

If a class or namespace contains a `using` directive (described later in this chapter), the used class or namespace is also searched.

A friend declaration does not add its name to the class scope. Thus, the rules for name lookup are slightly different than they are for a member function. If a friend function is defined within the body of the class granting friendship, the name is looked up in the class scope, just as it would be for a member function. If the function is defined outside the

class, the class is not searched, so the rules are the same as for an ordinary function. [Example 2-6](#) shows how the two different lookup rules can cause confusion.

### *Example 2-6. Looking up names in a friend function*

```
class foo {
public:
    friend void bar1(foo& f) {
        ++y;           // OK: refers to foo::y
    }
    friend void bar2(foo& f);
private:
    static int y;
    int x;
};

void bar2(foo& f) {
    ++y;           // Error: y not in scope
}
```

If the friend is a member function, the function and other names in its declaration are looked up first in the class granting friendship, and, if not found, in the class that contains the friend declaration. (See [Example 2-7](#).)

### *Example 2-7. Declaring friend member functions*

```
class test {
public:
    typedef int T1;
    typedef float T2;
    void f(T1);
    void f(T2);
};
class other {
    typedef char T2;
    friend void test::f(T1); // Look up f and T1 in test.
    friend void test::f(T2); // Look up f and T2 in test before
                             // looking it up in
                             // other.
};
```

In the definition of a class member (function or static data) outside of the class declaration, the lookup searches the class and ancestor classes, but only after the class name

appears in the declarator. Thus, the type specifier (a function's return type or the type of a static data member) is not looked up in the class declaration unless the type name is explicitly qualified. (Declarators and type specifiers are covered later in this chapter.) [Example 2-8](#) shows the consequences of this rule.

*Example 2-8. Defining members outside of a class declaration*

```
class node {
public:
    enum color { red, black };
    node(color x);
    color toggle(color c);
private:
    color c;
    static color root_color;
};

// Must qualify node::color and node::root_color, but initializer
// is in the scope
// of node, so it doesn't need to be qualified

node::color node::root_color = red;

// Similarly, return type must be qualified, but parameter type
// does not need to
// be.
node::color node::toggle(color c)
{
    return static_cast<color>(1 - c);
}
```

In a template declaration, the lookup rules for unqualified names have an additional wrinkle that depends on the template parameters and arguments. See [Chapter 7](#) for details.

## Argument-Dependent Name Lookup

Argument-dependent name lookup is an additional rule for looking up unqualified function names. The rule specifies

additional classes and namespaces to search based on the types of the function arguments. Argument-dependent name lookup is also known as *Koenig lookup*, named after Andrew Koenig, the creator of this lookup rule. The short version of the rule is that the compiler looks up a function name in the usual places, and in the namespaces that contain the user-defined types (classes and enumerations) of the function arguments.

The slightly longer version of the Koenig lookup rule is that the compiler first searches all the usual places, as described earlier in this chapter. If it does not find a declaration, it then searches an additional list of classes and namespaces. The additional list depends on the types used for all of the function's argument types:

- For a class type, the compiler searches the class and its namespaces, plus all ancestor classes and their namespaces.
- For a pointer to a data member, the compiler searches its class and its namespaces, plus all ancestor classes and their namespaces.
- For a function pointer or reference, the compiler searches the classes and namespaces associated with the return type and all parameter types. For a pointer or reference to a member function, the compiler also searches its class and all ancestor classes, and their namespaces.
- For a union or enumerated type, the namespace that contains the declaration is searched. If the type is a class member, its class is searched.

[Example 2-9](#) shows a typical case in which argument-dependent name lookup is needed. The operator<< function is declared in the std namespace in the <string> header. It is not a member function of ostream, and the

only way the compiler can find the operator is to search the `std` namespace. The fact that its arguments are in the `std` namespace tells the compiler to look there.

*Example 2-9. Argument-dependent name lookup*

```
#include <string>
#include <iostream>
#include <ostream>
int main( )
{
    std::string message("Howdy!\n");
    std::cout << message;
}
```

Another way to look at argument-dependent lookup is to consider a simple class declaration. For example, with rational numbers, to support the customary and usual arithmetic operators, you might choose to declare them as member functions:

```
namespace numeric {
    class rational {
        ...
        rational operator+(const rational& other);
        rational operator+(int i);
    };
}
```

Expressions such as `r1 + 42` compile successfully because they are equivalent to member function calls such as `r1.operator+(42)`, so the `+` operator is found as a member of `rational` (assuming `r1` is an instance of `rational`). But `42 + r1` does not work because an integer cannot have a member function—e.g., `42.operator+(r1)`. Instead, you must declare an operator at the namespace level:

```
namespace numeric {
    class rational { ... };
    rational operator+(int i, const rational& r);
    ...
}
```

In order to compile the expression `42 + r1`, the compiler needs to find `operator+` in the `numeric` namespace.

Without a `using` directive, the compiler has no way of knowing it needs to search the `numeric` namespace, but the compiler knows that the type of the second argument is declared in `numeric`. Thus, argument-dependent lookup allows for the everyday, expected use of overloaded operators.

# Linkage

Every name has *linkage* , which determines how the compiler and linker can use the name. Linkage has two aspects: scope and language. *Scope linkage* dictates which scopes have access to an entity. *Language linkage* dictates an entity's properties that depend on programming language.

## Scope Linkage

Scope linkage can be one of the following:

### Internal linkage

A name with internal linkage can be referred to from a different scope within the same source file. At namespace scope (that is, outside of functions and classes), static declarations have internal linkage, as do const declarations that are not also extern. Data members of anonymous unions have internal linkage. Names in an unnamed namespace have internal linkage.

### External linkage

A name with external linkage can be referred to from a different scope, possibly in a different source file. Functions and objects declared with the extern specifier have external linkage, as do entities declared at namespace scope that do not have internal linkage.

### No linkage

A name with no linkage can be referred to only from within the scope where it is declared. Local declarations that are not extern have no linkage.

## Language Linkage





Every function, function type, and object has a language linkage, which is specified as a simple character string. By default, the linkage is "C++". The only other standard language linkage is "C". All other language linkages and the properties associated with different language linkages are implementation-defined.

You can specify the language linkage for a single declaration (not a definition) or for a series of declarations and definitions. When you specify linkage for a series of declarations and definitions, you must enclose the series in curly braces. A language linkage declaration does not define a scope within the curly braces. For example:

```
extern "C" void cfunction(int);
extern "C++" {
    void cppfunc(int);
    void cppfunc(double);
}
```

Language linkage is part of a function's type, so typedef declarations keep track of the language linkage. When assigning a function to a function pointer, the function and pointer must have the same linkage. In the following example, `funcptr` is a function pointer with "C" linkage (note the need for curly braces because it is a definition, not a declaration). You can assign a "C" function to `funcptr`, but not a "C++" function, even though the rest of the function type matches.

```
extern "C" { void (*funcptr)(int); }
funcptr = cfunction; // OK
funcptr = cppfunc;  // Error
```

C does not support function overloading, so there can be at most one function with "C" linkage of a given name. Even if you declare a C function in two different namespaces, both declarations refer to the same function, for which there must be a single definition.

Typically, "C" linkage is used for external functions that are written in C (such as those in the C standard library), but that you want to call from a C++ program. "C++" linkage is used for native C++ code. Sometimes, though, you want to write a function in C++ that can be called from C; in that case, you should declare the C++ function with "C" linkage.



An implementation might support other language linkages. It is up to the implementation to define the properties of each language: how parameters are passed to functions, how values are returned from functions, whether and how function names are altered, and so on. In many C++ implementations, a function with "C++" linkage has a "mangled" name, that is, the external name encodes the function name and the types of all its arguments. So the function `strlen(const char*)` might have an external name of `strlen__FCcP`, making it hard to call the function from a C program, which does not know about C++ name-mangling rules. Using "C" linkage, the compiler might not mangle the name, exporting the function under the plain name of `strlen`, which can be called easily from C.

## Type Declarations

One of the hallmarks of C++ is that you can define a type that resembles any built-in type. Thus, if you need to define a type that supports arbitrary-sized integers—call it `bigint`—you can do so, and programmers will be able to use `bigint` objects the same way they use `int` objects.

You can define a brand new type by defining a class (see [Chapter 6](#)) or an enumeration (see [Section 2.5.2](#) later in this chapter). In addition to declaring and defining new types, you can declare a `typedef`, which is a synonym for an existing type. Note that while the name `typedef` seems to be a shorthand for "type definition," it is actually a type declaration. (See [Section 2.5.4](#) later in this chapter.)

## Fundamental Types

This section lists the fundamental type specifiers that are built into the C++ language. For types that require multiple keywords (e.g., `unsigned long int`), you can mix the keywords in any order, but the order shown in the following list is the conventional order. If a type specifier requires multiple words, one of which is `int`, the `int` can be omitted. If a type is signed, the signed keyword can be omitted (except in the case of `signed char`).

`bool`

Represents a Boolean or logical value. It has two possible values: `true` and `false`.



`char`

Represents a narrow character. Narrow character literals usually have type `char`. (If a narrow character literal contains multiple characters, the type is `int`.)

Unlike the other integral types, a plain char is not necessarily equivalent to signed char. Instead, char can be equivalent to signed char or unsigned char, depending on the implementation. Regardless of the equivalence, the plain char type is distinct and separate from signed char and unsigned char.

All the narrow character types (char, signed char, and unsigned char) share a common size and representation. By definition, char is the smallest fundamental type—that is, `sizeof(char)` is 1.

double

Represents a double-precision, floating-point number. The range and precision are at least as much as those of float. A floating-point literal has type double unless you use the F or L suffix.

float

Represents a single-precision, floating-point number.

long double

Represents an extended-precision, floating-point number. The range and precision are at least as much as those of double.

signed char

Represents a signed byte.

signed int

Represents an integer in a size and format that is natural for the host environment.

signed long int

Represents an integer whose range is at least as large as that of int.

signed short int

Represents an integer such that the range of an `int` is at least as large as the range of a `short`.

`unsigned char`

Represents an unsigned byte. Some functions, especially in the C library, require characters and character strings to be cast to `unsigned char` instead of plain `char`. (See [Chapter 13](#) for details.)

`unsigned long int`

Represents an unsigned long integer.

`unsigned short int`

Represents an unsigned short integer.

`void`

Represents the absence of any values. You cannot declare an object of type `void`, but you can declare a function that "returns" `void` (that is, does not return a value), or declare a pointer to `void`, which can be used as a generic pointer. (See `static_cast<>` under [Section 3.5.2](#) for information about casting to and from a pointer to `void`.)

`wchar_t`

Represents a wide character. Its representation must match one of the fundamental integer types. Wide character literals have type `wchar_t`.



The representations of the fundamental types are implementation-defined. The integral types (`bool`, `char`, `wchar_t`, `int`, etc.) each require a binary representation: signed-magnitude, ones' complement, or two's complement. Some types have alignment restrictions, which are also

implementation-defined. (Note that new expressions always return pointers that are aligned for any type.)

The unsigned types always use arithmetic modulo  $2^n$ , in which  $n$  is the number of bits in the type. Unsigned types take up the same amount of space and have the same alignment requirements as their signed companion types. Nonnegative signed values are always a subset of the values supported by the equivalent unsigned type and must have the same bit representations as their corresponding unsigned values.

### Tip

Although the size and range of the fundamental types is implementation-defined, the C++ standard mandates minimum requirements for these types. These requirements are specified in `<climits>` (for the integral types) and `<cfloat>` (for the floating-point types). See also `<limits>`, which declares templates for obtaining the numerical properties of each fundamental type.

## Enumerated Types

An *enumerated type* declares an optional type name (the *enumeration*) and a set of zero or more identifiers (*enumerators*) that can be used as values of the type. Each enumerator is a constant whose type is the enumeration. For example:

```
enum logical { no, maybe, yes };
logical is_permitted = maybe;

enum color { red=1, green, blue=4 };
const color yellow = static_cast<color>(red | green);

enum zeroes { a, b = 0, c = 0 };
```

You can optionally specify the integral value of an enumerator after an equal sign (=). The value must be a constant of integral or enumerated type. The default value of the first enumerator is 0. The default value for any other enumerator is one more than the value of its predecessor (regardless of whether that value was explicitly specified). In a single enumeration declaration, you can have more than one enumerator with the same value.

The name of an enumeration is optional, but without a name you cannot declare use the enumeration in other declarations. Such a declaration is sometimes used to declare integer constants such as the following:

```
enum { init_size = 100 };  
std::vector<int> data(init_size);
```



An enumerated type is a unique type. Each enumerated value has a corresponding integer value, and the enumerated value can be promoted automatically to its integer equivalent, but integers cannot be implicitly converted to an enumerated type. Instead, you can use `static_cast<>` to cast an integer to an enumeration or cast a value from one enumeration to a different enumeration. (See [Chapter 3](#) for details.)

The range of values for an enumeration is defined by the smallest and largest bitfields that can hold all of its enumerators. In more precise terms, let the largest and smallest values of the enumerated type be  $v_{min}$  and  $v_{max}$ . The largest enumerator is  $e_{max}$  and the smallest is  $e_{min}$ . Using two's complement representation (the most common integer format),  $v_{max}$  is the smallest  $2^n - 1$ , such that  $v_{max} \geq \max(\text{abs}(e_{min}) - 1, \text{abs}(e_{max}))$ . If  $e_{min}$  is not negative,  $v_{min} = 0$ ; otherwise,  $v_{min} = -(v_{max} + 1)$ .



In other words, the range of values for an enumerated type can be larger than the range of enumerator values, but the exact range depends on the representation of integers on the host platform, so it is implementation-defined. All values between the largest and smallest enumerators are always valid, even if they do not have corresponding enumerators.

In the following example, the enumeration `sign` has the range (in two's complement) -2 to 1. Your program might not assign any meaning to `static_cast<sign>(-2)`, but it is semantically valid in a program:

```
enum sign { neg=-1, zero=0, pos=1 };
```



Each enumeration has an underlying integral type that can store all of the enumerator values. The actual underlying type is implementation-defined, so the size of an enumerated type is likewise implementation-defined.

The standard library has a type called `iostate`, which might be implemented as an enumeration. (Other implementations are also possible; see `<ios>` in [Chapter 13](#) for details.) The enumeration has four enumerators, which can be used as bits in a bitmask and combined using bitwise operators:

```
enum iostate { goodbit=0, failbit=1, eofbit=2, badbit=4 };
```

The `iostate` enumeration can clearly fit in a `char` because the range of values is 0 to 7, but the compiler is free to use `int`, `short`, `char`, or the unsigned flavors of these types as the underlying type.

Because enumerations can be promoted to integers, any arithmetic operation can be performed on enumerated



values, but the result is always an integer. If you want to permit certain operations that produce enumeration results, you must overload the operators for your enumerated type. For example, you might want to overload the bitwise operators, but not the arithmetic operators, for the `iostate` type in the preceding example. The `sign` type does not need any additional operators; the comparison operators work just fine by implicitly converting `sign` values to integers. Other enumerations might call for overloading `++` and `--` operators (similar to the `succ` and `pred` functions in Pascal). [Example 2-10](#) shows how operators can be overloaded for enumerations.

### *Example 2-10. Overloading operators for enumerations*

```
// Explicitly cast to int, to avoid infinite recursion.
inline iostate operator|(iostate a, iostate b) {
    return iostate(int(a) | int(b));
}
inline iostate& operator|=(iostate& a, iostate b) {
    a = a | b;
    return a;
}
// Repeat for &, ^, and ~.

int main( )
{
    iostate err = goodbit;
    if (error( ))
        err |= badbit;
}
```

## **POD Types**

POD is short for Plain Old Data. The fundamental types and enumerated types are POD types, as are pointers and arrays of POD types. You can declare a POD class, which is a class or struct that uses only POD types for its nonstatic data members. A POD union is a union of POD types.

POD types are special in several ways:

- A POD object can be copied byte for byte and retain its value. In particular, a POD object can be safely copied to an array of char or unsigned char, and when copied back, it retains its original value. A POD object can also be copied by calling memcopy; the copy has the same value as the original.
- A local POD object without an initializer is uninitialized, that is, its value is undefined. Similarly, a POD type in a new expression without an initializer is uninitialized. (A non-POD class is initialized in these situations by calling its default constructor.) When initialized with an empty initializer, a POD object is initialized to 0, false, or a null pointer.
- A goto statement can safely branch across declarations of uninitialized POD objects. (See [Chapter 4](#) for more information.)
- A POD object can be initialized using an aggregate initializer. (See [Section 2.6.3](#) later in this chapter for details.)

See [Chapter 6](#) for more information about POD types, especially POD classes.

## typedef Declarations

A typedef declares a synonym for an existing type. Syntactically, a typedef is like declaring a variable, with the type name taking the place of the variable name, and with a typedef keyword out in front. More precisely, typedef is a specifier in a declaration, and it must be combined with type specifiers and optional const and volatile qualifiers (called *cv-qualifiers*), but no storage class specifiers. A list of declarators follow the specifiers. (See the next section, [Section 2.6](#), for more information

about *cv* qualifiers, storage class specifiers, and declarators.)

The declarator of a typedef declaration is similar to that for an object declaration, except you cannot have an initializer. The following are some examples of typedef declarations:

```
typedef unsigned int uint;  
typedef long int *long_ptr;  
typedef double matrix[3][3];  
typedef void (*thunk)( );  
typedef signed char schar;
```

By convention, the typedef keyword appears before the type specifiers. Syntactically, typedef behaves as a storage class specifier (see [Section 2.6.1](#) later in this chapter for more information about storage class specifiers) and can be mixed in any order with other type specifiers. For example, the following typedefs are identical and valid:

```
typedef unsigned long ulong; // Conventional  
long typedef int unsigned ulong; // Valid, but strange
```

A typedef is especially helpful with complicated declarations, such as function pointer declarations and template instantiations. They help the author who must concoct the declarations, and they help the reader who must later tease apart the morass of parentheses, asterisks, and angle brackets. The standard library uses them frequently. (See also [Example 2-11](#).)

```
typedef std::basic_string<char, std::char_traits<char> > string;
```

A typedef does not create a new type the way `class` and `enum` do. It simply declares a new name for an existing type. Therefore, function declarations for which the parameters differ only as typedefs are not actually different declarations. The two function declarations in the following example are really two declarations of the same function:

```
typedef unsigned int uint;  
uint func(uint); // Two declarations of the
```

```
unsigned func(unsigned);    // same function
```

Similarly, because you cannot overload an operator on fundamental types, you cannot overload an operator on typedef synonyms for fundamental types. For example, both the following attempts to overload + result in an error:

```
int operator+(int, int);    // Error
typedef int integer;
integer operator+(integer, integer); // Error
```

C programmers are accustomed to declaring typedefs for struct, union, and enum declarations, but such typedefs are not necessary in C++. In C, the struct, union, and enum names are separate from the type names, but in C++, the declaration of a struct, union, class, or enum adds the type to the type names. Nonetheless, such a typedef is harmless. The following example shows typedef being used to create the synonym point for the struct point:

```
struct point { int x, y; }
typedef struct point point; // Not needed in C++, but harmless
point pt;
```

## Elaborated Type Specifier

An *elaborated type specifier* begins with an introductory keyword: class, enum, struct, typename, or union. The keyword is followed by a (possibly) qualified name of a suitable type. That is, enum is followed by the name of an enumeration, class is followed by a class name, struct by a struct name, and union by a union name. The typename keyword is used only in templates (see [Chapter 7](#)) and is followed by a name that must be a type name in a template instantiation.

A typename-elaborated type specifier is often needed in template definitions. The other elaborated type specifiers tend to be used in headers that must be compatible with C

or in type names that have been hidden by other names.  
For example:

```
enum color { black, red };  
color x; // No need for elaborated name  
enum color color( ); // Function hides color  
enum color c = color( ); // Elaborated name is needed here
```

# Object Declarations

An *object* in C++ is a region of storage with a type, a value, and possibly a name. In traditional object-oriented programming, "object" means an instance of a class, but in C++ the definition is slightly broader to include instances of any data type.

An object (variable or constant) declaration has two parts: a series of *specifiers* and a list of comma-separated *declarators*. Each declarator has a name and an optional *initializer*.

## Specifiers

Each declaration begins with a series of specifiers. The series can contain a storage class, `const` and `volatile` qualifiers, and the object's type, in any order.

### Storage class specifiers

A storage class specifier can specify scope linkage and lifetime. The storage class is optional. For function parameters and local variables in a function, the default storage class specifier is `auto`. For declarations at namespace scope, the default is usually an object with static lifetime and external linkage. C++ has no explicit storage class for such a declaration. (See [Section 2.6.4](#) later in this chapter and [Section 2.4](#) earlier in this chapter for more information.) If you use a storage class specifier, you must choose only one of the following:

`auto`

Denotes an automatic variable—that is, a variable with a lifetime limited to the block in which the variable is declared. The `auto` specifier is the default for function

parameters and local variables, which are the only kinds of declarations for which it can be used, so it is rarely used explicitly.

#### extern

Denotes an object with external linkage, which might be defined in a different source file. Function parameters cannot be extern.

#### mutable

Denotes a data member that can be modified even if the containing object is const. See [Chapter 6](#) for more information.

#### register

Denotes an automatic variable with a hint to the compiler that the variable should be stored in a fast register. Many modern compilers ignore the register storage class because the compilers are better than people at determining which variables belong in registers.

#### static

Denotes a variable with a static lifetime and internal linkage. Function parameters cannot be static.

### **const and volatile qualifiers**

The `const` and `volatile` specifiers are optional. You can use either one, neither, or both in any order. The `const` and `volatile` keywords can be used in other parts of a declaration, so they are often referred to by the more general term *qualifiers*; for brevity, they are often referred to as *cv-qualifiers*.

#### const

Denotes an object that cannot be modified. A `const` object cannot ordinarily be the target of an assignment. You cannot call a non-`const` member function of a `const` object.

## `volatile`

Denotes an object whose value might change unexpectedly. The compiler is prevented from performing optimizations that depend on the value not changing. For example, a variable that is tied to a hardware register should be `volatile`.

## **Type specifiers**

Every object must have a type in the form of one or more type specifiers. The type specifiers can be any of the following:

- The name of a class, enumeration, or typedef
- An elaborated name
- A series of fundamental type specifiers
- A class definition
- An enumerated type definition

Enumerated and fundamental types are described earlier in this chapter in [Section 2.5](#). Class types are covered in [Chapter 6](#). The `typename` keyword is covered in [Chapter 7](#).

## **Using specifiers**

Specifiers can appear in any order, but the convention is to list the storage class first, followed by the type specifiers, followed by `cv`-qualifiers.

```
extern long int const mask; // Conventional
int extern const long mask; // Valid, but strange
```



Many programmers prefer a different order: storage class, *cv*-qualifiers, type specifiers. More and more are learning to put the *cv*-qualifiers last, though. See the examples under [Section 2.6.2.2](#) later in this chapter to find out why.

The convention for types that require multiple keywords is to place the base type last and the modifiers first:

```
unsigned long int x; // Conventional
int unsigned long y; // Valid, but strange
long double a;      // Conventional
double long b;      // Valid, but strange
```

You can define a class or enumerated type in the same declaration as an object declaration:

```
enum color { red, black } node_color;
```

However, the custom is to define the type separately, then use the type name in a separate object declaration:

```
enum color { red, black };
color node_color;
```

## Declarators

A declarator declares a single name within a declaration. In a declaration, the initial specifiers apply to all the declarators in the declaration, but each declarator's modifiers apply only to that declarator. (See [Section 2.6.2.2](#) in this section for examples of where this distinction is crucial.) A declarator contains the name being declared, additional type information (for pointers, references, and arrays), and an optional initializer. Use commas to separate multiple declarators in a declaration. For example:

```
int plain_int, array_of_int[42], *pointer_to_int;
```

## Arrays

An array is declared with a constant size specified in square brackets. The array size is fixed for the lifetime of the object and cannot change. (For an array-like container

whose size can change at runtime, see `<vector>` in [Chapter 13](#).) To declare a multidimensional array, use a separate set of square brackets for each dimension:

```
int point[2];
double matrix[3][4]; // A 3 × 4 matrix
```

You can omit the array size if there is an initializer; the number of initial values determines the size. In a multidimensional array, you can omit only the first (leftmost) size:

```
int data[] = { 42, 10, 3, 4 }; // data[4]
int identity[][3] = { { 1,0,0 }, {0,1,0}, {0,0,1} }; //
identity[3][3]
char str[] = "hello"; // str[6], with trailing \0
```

In a multidimensional array, all elements are stored contiguously, with the rightmost index varying the fastest (usually called *row major order*).

When a function parameter is an array, the array's size is ignored, and the type is actually a pointer type, which is the subject of the next section. For a multidimensional array used as a function parameter, the first dimension is ignored, so the type is a pointer to an array. Because the first dimension is ignored in a function parameter, it is usually omitted, leaving empty square brackets:

```
long sum(long data[], size_t n);
double chi_sq(double stat[][2]);
```

## Pointers

A pointer object stores the address of another object. A pointer is declared with a leading asterisk (\*), optionally followed by *cv*-qualifiers, then the object name, and finally an optional initializer.

When reading and writing pointer declarations, be sure to keep track of the *cv*-qualifiers. The *cv*-qualifiers in the declarator apply to the pointer object, and the *cv*-qualifiers in the declaration's specifiers apply to the type of the

pointer's target. For example, in the following declaration, the `const` is in the specifier, so the pointer `p` is a pointer to a `const int`. The pointer object is modifiable, but you cannot change the `int` that it points to. This kind of pointer is usually called a pointer to `const`.

```
int i, j;
int const *p = &i;
p = &j; // OK
*p = 42; // Error
```

When the *cv*-qualifier is part of the declarator, it modifies the pointer object. Thus, in the following example, the pointer `p` is `const` and hence not modifiable, but it points to a plain `int`, which can be modified. This kind of pointer is usually called a `const` pointer.

```
int i, j;
int * const p = &i;
p = &j; // Error
*p = 42; // OK
```

You can have pointers to pointers, as deep as you want, in which each level of pointer has its own *cv*-qualifiers. The easiest way to read a complicated pointer declaration is to find the declarator, work your way from the inside to the outside, and then from right to left. In this situation, it is best to put *cv*-qualifiers after the type specifiers. For example:

```
int x;
int *p; // Pointer to int
int * const cp = &x; // const pointer to int
int const * pc; // Pointer to const int
int const * const cpc = &x; // const pointer to const int
int *pa[10]; // Array of 10 pointers to int
int **pp; // Pointer to pointer to int
```

When a function parameter is declared with an array type, the actual type is a pointer, and at runtime the address of the first element of the array is passed to the function. You can use array syntax, but the size is ignored. For example, the following two declarations mean the same thing:

```
int sum(int data[], size_t n);
int sum(int *data, size_t n);
```

When using array notation for a function parameter, you can omit only the first dimension. For example, the following is valid:

```
void transpose(double matrix[][3]);
```

but the following is not valid. If the compiler does not know the number of columns, it does not know how to lay out the memory for `matrix` or compute array indices.

```
void transpose(double matrix[][]);
```

A useful convention is to use array syntax when declaring parameters that are used in an array-like fashion—that is, the parameter itself does not change, or it is dereferenced with the `[]` operator. Use pointer syntax for parameters that are used in pointer-like fashion—that is, the parameter value changes, or it is dereferenced with the unary `*` operator.

## Function pointers

A function pointer is declared with an asterisk (`*`) and the function signature (parameter types and optional names). The declaration's specifiers form the function's return type. The name and asterisk must be enclosed in parentheses, so the asterisk is not interpreted as part of the return type. An optional exception specification can follow the signature. See [Chapter 5](#) for more information about function signatures and exception specifications.

```
void (*fp)(int); // fp is pointer to function that takes an int
parameter
// and returns void.
void print(int);
fp = print;
```

A declaration of an object with a function pointer type can be hard to read, so typically you declare the type separately

with a typedef declaration, and then declare the object using the typedef name:

```
typedef void (*Function)(int);
Function fp;
fp = print;
```

[Example 2-11](#) shows a declaration of an array of 10 function pointers, in which the functions return `int*` and take two parameters: a function pointer (taking an `int*` and returning an `int*`) and an integer. The declaration is almost unreadable without using typedef declarations for each part of the puzzle.

*Example 2-11. Simplifying function pointer declarations with typedef*

```
// Array of 10 function pointers
int* (*fp[10])(int*(*)(int*),
int);

// Declare a type for pointer to int.
typedef int* int_ptr;
// Declare a function pointer type for a function that takes an
int_ptr parameter
// and returns an int_ptr.
typedef int_ptr (*int_ptr_func)(int_ptr);
// Declare a function pointer type for a function that returns
int_ptr and takes
// two parameters: the first of type int_ptr and the second of
type int.
typedef int_ptr (*func_ptr)(int_ptr_func, int);
// Declare an array of 10 func_ptrs.
func_ptr fp[10];
```

## Member pointers

Pointers to members (data and functions) work differently from other pointers. The syntax for declaring a pointer to a nonstatic data member or a nonstatic member function requires a class name and scope operator before the asterisk. Pointers to members can never be cast to ordinary pointers, and vice versa. You cannot declare a reference to a member. (See [Chapter 3](#) for information about expressions that dereference pointers to members.) A

pointer to a static member is an ordinary pointer, not a member pointer. The following are some simple examples of member pointers:

```
struct simple {
    int data;
    int func(int);
};
int simple::* p = &simple::data;
int (simple::*fp)(int) = &simple::func;
simple s;
s.*p = (s.*fp)(42);
```

## References

A reference is a synonym for an object or function. A reference is declared just like a pointer, but with an ampersand (&) instead of an asterisk (\*). A local or global reference declaration must have an initializer that specifies the target of the reference. Data members and function parameters, however, do not have initializers. You cannot declare a reference of a reference, a reference to a class member, a pointer to a reference, an array of references, or a cv-qualified reference. For example:

```
int x;
int &r = x;           // Reference to int
int& const rc = x;   // Error: no cv qualified references
int &&rr;             // Error: no reference of reference
int& ra[10];         // Error: no arrays of reference
int*&* rp = &r;      // Error: no pointer to reference
int* p = &x;         // Pointer to int
int*&* pr = p;       // OK: reference to pointer
```

A reference, unlike a pointer, cannot be made to refer to a different object at runtime. Assignments to a reference are just like assignments to the referenced object.

Because a reference cannot have cv-qualifiers, there is no such thing as a const reference. Instead, a reference to const is often called a const reference for the sake of brevity.

References are often used to bind names to temporary objects, implement call-by-reference for function parameters, and optimize call-by-value for large function parameters. The `divide` function in [Example 2-12](#) demonstrates the first two uses. The standard library has the `div` function, which divides two integers and returns the quotient and remainder in a `struct`. Instead of copying the structure to a local object, `divide` binds the return value to a reference to `const`, thereby avoiding an unnecessary copy of the return value. Furthermore, suppose that you would rather have `divide` return the results as arguments. The function parameters `quo` and `rem` are references; when the `divide` function is called, they are bound to the function arguments, `q` and `r`, in `main`. When `divide` assigns to `quo`, it actually stores the value in `q`, so when `divide` returns, `main` has the quotient and remainder.

*Example 2-12. Returning results in function arguments*

```
#include <cstdlib>
#include <iostream>
#include <ostream>

void divide(long num, long den, long& quo, long& rem)
{
    const std::ldiv_t& result = std::div(num, den);
    quo = result.quot;
    rem = result.rem;
}

int main( )
{
    long q, r;
    divide(42, 5, q, r);
    std::cout << q << " remainder " << r << '\n';
}
```

The other common use of references is to use a `const` reference for function parameters, especially for large objects. Function arguments are passed by value in C++, which requires copying the arguments. The copy operation

can be costly for a large object, so passing a reference to a large object yields better performance than passing the large object itself. The reference parameter is bound to the actual argument, avoiding the unnecessary copy. If the function modifies the object, it would violate the call-by-value convention, so you should declare the reference `const`, which prevents the function from modifying the object. In this way, call-by-value semantics are preserved, and the performance of call-by-reference is improved. The standard library often uses this idiom. For example, `operator<<` for `std::string` uses a `const` reference to the string to avoid making unnecessary copies of the string. (See `<string>` in [Chapter 13](#) for details.)

If a function parameter is a non-`const` reference, the argument must be an lvalue. A `const` reference, however, can bind to an rvalue, which permits temporary objects to be passed to the function, which is another characteristic of call-by-value. (See [Chapter 3](#) for the definitions of "lvalue" and "rvalue.")

A reference must be initialized so it refers to an object. If a data member is a reference, it must be initialized in the constructor's initializer list (see [Chapter 6](#)). Function parameters that are references are initialized in the function call, binding each reference parameter to its corresponding actual argument. All other reference definitions must have an initializer. (An extern declaration is not a definition, so it doesn't take an initializer.)

A `const` reference can be initialized to refer to a temporary object. For example, if a function takes a `const` reference to a `float` as a parameter, you can pass an integer as an argument. The compiler converts the integer to `float`, saves the `float` value as an unnamed temporary object, and passes that temporary object as the function argument. The `const` reference is initialized to refer to the temporary



object. After the function returns, the temporary object is destroyed:

```
void do_stuff(const float& f);
do_stuff(42);
// Equivalent to:
{
    const float unnamed = 42;
    do_stuff(unnamed);
}
```

The restrictions on a reference, especially to a reference of a reference, pose an additional challenge for template authors. For example, you cannot store references in a container because a number of container functions explicitly declare their parameters as references to the container's value type. (Try using `std::vector<int&>` with your compiler, and see what happens. You should see a lot of error messages.)

Instead, you can write a wrapper template, call it `rvector<typename T>`, and specialize the template (`rvector<T&>`) so references are stored as pointers, but all the access functions hide the differences. This approach requires you to duplicate the entire template, which is tedious. Instead, you can encapsulate the specialization in a traits template called `Ref<>` (refer to [Chapter 7](#) for more information about templates and specializations, and to [Chapter 8](#) for more information about traits), as shown in [Example 2-13](#).

### *Example 2-13. Encapsulating reference traits*

```
// Ref type trait encapsulates reference type, and mapping to and
// from the type
// for use in a container.
template<typename T>
struct Ref {
    typedef T value_type;
    typedef T& reference;
    typedef const T& const_reference;
    typedef T* pointer;
    typedef const T* const_pointer;
```

```

typedef T container_type;
static reference from_container(reference x) { return x; }
static const_reference from_container(const_reference x)
                                     { return x; }
static reference to_container(reference x) { return x; }
};

template<typename T>
struct Ref<T&> {
    typedef T value_type;
    typedef T& reference;
    typedef const T& const_reference;
    typedef T* pointer;
    typedef const T* const_pointer;
    typedef T* container_type;
    static reference from_container(pointer x) { return *x; }
    static const_reference from_container(const_pointer x)
                                     { return *x; }
    static pointer to_container(reference x) { return &x; }
};

// rvector<> is similar to vector<>, but allows references by
// storing references
// as pointers.
template<typename T, typename A=std::allocator<T> >
class rvector {
    typedef typename Ref<T>::container_type container_type;
    typedef typename std::vector<container_type> vector_type;
public:
    typedef typename Ref<T>::value_type value_type;
    typedef typename Ref<T>::reference reference;
    typedef typename Ref<T>::const_reference const_reference;
    typedef typename vector_type::size_type size_type;
    . . . // Other typedefs are similar.
    class iterator { ... }; // Wraps a vector<>::iterator
    class const_iterator { ... };
    . . . // Constructors pass arguments to v.
    iterator begin( ) { return iterator(v.begin( )); }
    iterator end( ) { return iterator(v.end( )); }
    void push_back(typename Ref<T>::reference x) {
        v.push_back(Ref<T>::to_container(x));
    }
    reference at(size_type n) {
        return Ref<T>::from_container(v.at(n));
    }
    reference front( ) {
        return Ref<T>::from_container(v.front( ));
    }
}

```

```

const_reference front( ) const {
    return Ref<T>::from_container(v.front( ));
}
. . . // Other members are similar.
private:
    vector_type v;
};

```

## Initializers

An initializer supplies an initial value for an object being declared. You must supply an initializer for the definition of a reference or const object. An initializer is optional for other object definitions. An initializer is not allowed for most data members within a class definition, but an exception is made for static const data members of integral or enumerated type. Initializers are also not allowed for extern declarations and function parameters. (Default arguments for function parameters can look like initializers. See [Chapter 5](#) for details.)

The two forms of initializers are *assignment-like* and *function-like*. (In the C++ standard, assignment-like is called copy initialization, and function-like is called direct initialization.) An assignment-like initializer starts with an equal sign, which is followed by an expression or a list of comma-separated expressions in curly braces. A function-like initializer is a list of one or more comma-separated expressions in parentheses. Note that these initializers look like assignment statements or function calls, respectively, but they are not. They are initializers. The difference is particularly important for classes (see [Chapter 6](#)). The following are some examples of initializers:

```

int x = 42; // Initializes x with the value 42
int y(42); // Initializes y with the value 42
int z = { 42 }; // Initializes z with the value 42
int w[4] = { 1, 2, 3, 4 }; // Initializes an array
std::complex<double> c(2.0, 3.0); // Calls complex constructor

```

When initializing a scalar value, the form is irrelevant. The initial value is converted to the desired type using the usual conversion rules (as described in [Chapter 3](#)).

Without an initializer, all non-POD class-type objects are initialized by calling their default constructors. (See [Chapter 6](#) for more information about POD and non-POD classes.) All other objects with static lifetimes are initialized to `0`; objects with automatic lifetimes are left uninitialized. (See [Section 2.6.4](#) later in this chapter.) An uninitialized reference or `const` object is an error.

## Function-like Initializers

You must use a function-like initializer when constructing an object whose constructor takes two or more arguments, or when calling an explicit constructor. The usual rules for resolving overloaded functions apply to the choice of overloaded constructors. (See [Chapter 5](#) for more information about overloading and [Chapter 6](#) for more information about constructors.) For example:

```
struct point {
    point(int x, int y);
    explicit point(int x);
    point( );
    ...
};

point p1(42, 10); // Invokes point::point(int x, int y);
point p2(24);    // Invokes point::point(int x);
point p3;       // Invokes point::point( );
```

Empty parentheses cannot be used as an initializer in an object's declaration, but can be used in other initialization contexts (namely, a constructor initializer list or as a value in an expression). If the type is a class type, the default constructor is called; otherwise, the object is initialized to `0`. [Example 2-14](#) shows an empty initializer. No matter what type `T` is, the `wrapper<>` template can rely on `T( )` to be a meaningful default value.

### Example 2-14. An empty initializer

```
template<typename T>
struct wrapper {
    wrapper( ) : value_(T( )) {}
    explicit wrapper(const T& v) : value_(v) {}
private:
    T value_;
};

wrapper<int> i;    // Initializes i with int( ), or zero
enum color { black, red, green, blue };
wrapper<color> c; // Initializes c with color( ), or black
wrapper<bool> b;  // Initializes b with bool( ), or false
wrapper<point> p; // Initializes p with point( )
```

## Assignment-like initializers

In an assignment-like initializer, if the object is of class type, the value to the right of the equal sign is converted to a temporary object of the desired type, and the first object is constructed by calling its copy constructor.

The generic term for an array or simple class is *aggregate* because it aggregates multiple values into a single object. "Simple" in this case means the class does not have any of the following:

- User-declared constructors
- Private or protected nonstatic data members
- Base classes
- Virtual functions

To initialize an aggregate, you can supply multiple values in curly braces, as described in the following sections. A POD object is a special kind of an aggregate. (See [Section 2.5.3](#) earlier in this chapter for more information about POD types; see also [Chapter 6](#) for information about POD classes.)

To initialize an aggregate of class type, supply an initial value for each nonstatic data member separated by commas and enclosed in curly braces. For nested objects, use nested curly braces. Values are associated with members in the order of the members' declarations. More values than members results in an error. If there are fewer values than members, the members without values are initialized by calling each member's default constructor or initializing the members to 0.

An initializer list can be empty, which means all members are initialized to their defaults, which is different from omitting the initializer entirely. The latter causes all members to be left uninitialized. The following example shows several different initializers for class-type aggregates:

```
struct point { double x, y, z; };
point origin = { };          // All members initialized to 0.0
point unknown;              // Uninitialized, value is not known
point pt = { 1, 2, 3 };     // pt.x==1.0, pt.y==2.0, pt.z==3.0
struct line { point p1, p2; };
line vec = { { }, { 1 } };  // vec.p1 is all zero.
                          // vec.p2.x==1.0, vec.p2.y==0.0, vec.p2.z==0.0
```

Only the first member of a union can be initialized:

```
union u { int value; unsigned char bytes[sizeof(int)]; };
u x = 42; // Initializes x.value
```

## **I n i t i a l i z i n g a r r a y s**

Initialize elements of an array with values separated by commas and enclosed in curly braces. Multidimensional arrays can be initialized by nesting sets of curly braces. An error results if there are more values than elements in the array; if an initializer has fewer values than elements in the array, the remaining elements in the array are initialized to zero values (default constructors or 0). If the declarator omits the array size, the size is determined by counting the number of values in the initializer.

An array initializer can be empty, which forces all elements to be initialized to 0. If the initializer is empty, the array size must be specified. Omitting the initializer entirely causes all elements of the array to be uninitialized (except non-POD types, which are initialized with their default constructors).

In the following example the size of `vec` is set to 3 because its initializer contains three elements, and the elements of zero are initialized to 0s because an empty initializer is used:

```
int vec[] = { 1, 2, 3 }; // Array of three elements
                        // vec[0]==1 ... vec[2]==3
int zero[4] = { }; // Initialize to all zeros.
```

When initializing a multidimensional array, you can flatten the curly braces and initialize elements of the array in row major order (last index varies the fastest). For example, both `id1` and `id2` end up having the same values in their corresponding elements:

```
// Initialize id1 and id2 to the identity matrix.
int id1[3][3] = { { 1 }, { 0, 1 }, { 0, 0, 1 } };
int id2[3][3] = { 1, 0, 0, 0, 1, 0, 0, 0, 1 };
```

An array of `char` or `wchar_t` is special because you can initialize such arrays with a string literal. Remember that every string literal has an implicit null character at the end. For example, the following two `char` declarations are equivalent, as are the two `wchar_t` declarations:

```
// The following two declarations are equivalent.
char str1[] = "Hello";
char str2[] = { 'H', 'e', 'l', 'l', 'o', '\0' };

wchar_t ws1[] = L"Hello";
wchar_t ws2[] = { L'H', L'e', L'l', L'l', L'o', L'\0' };
```

The last expression in an initializer list can be followed by a comma. This is convenient when you are maintaining software and find that you often need to change the order

of items in an initializer list. You don't need to treat the last element differently from the other elements.

```
const std::string keywords[] = {  
    "and",  
    "asm",  
    ...  
    "while",  
    "xor",  
};
```

Because the last item has a trailing comma, you can easily select the entire line containing "xor" and move it to a different location in the list, and you don't need to worry about fixing up the commas afterward.

## **I n i t i a l i z i n g s c a l a r s**

You can initialize any scalar object with a single value in curly braces, but you cannot omit the value the way you can with a single-element array:

```
int x = { 42 };
```

## **Object Lifetime**

Every object has a *lifetime*, that is, the duration from when the memory for the object is allocated and the object is initialized to when the object is destroyed and the memory is released. Object lifetimes fall into three categories:

### Automatic

Objects are local to a function body or a nested block within a function body. The object is created when execution reaches the object's declaration, and the object is destroyed when execution leaves the block.

### Static

Objects can be local (with the `static` storage class specifier) or global (at namespace scope). Static objects



are constructed at most once and destroyed only if they are successfully constructed. Local static objects are constructed when execution reaches the object's declaration. Global objects are constructed when the program starts but before main is entered. Static objects are destroyed in the opposite order of their construction. For more information, see "The main Function" in [Chapter 5](#).

## Dynamic

Objects created with new expressions are dynamic. Their lifetimes extend until the delete expression is invoked on the objects' addresses. See [Chapter 3](#) for more information about the new and delete expressions.

## Namespaces

A *namespace* is a named scope. By grouping related declarations in a namespace, you can avoid name collisions with declarations in other namespaces. For example, suppose you are writing a word processor, and you use packages that others have written, including a screen layout package, an equation typesetting package, and an exact-arithmetic package for computing printed positions to high accuracy with fixed-point numbers.

The equation package has a class called `fraction`, which represents built-up fractions in an equation; the arithmetic package has a class called `fraction`, for computing with exact rational numbers; and the layout package has a class called `fraction` for laying out fractional regions of a page. Without namespaces, all three names would collide, and you would not be able to use more than one of the three packages in a single program.

With namespaces, each class can reside in a separate namespace—for example, `layout::fraction`, `eqn::fraction`, and `math::fraction`.

### Tip

C++ namespaces are similar to Java packages, with a key difference: in Java, classes in the same package have additional access rights to each other; in C++, namespaces confer no special access privileges.

## Namespace Definitions

Define a namespace with the `namespace` keyword followed by an optional identifier (the namespace name) and zero or more declarations in curly braces. Namespace declarations

can be discontinuous, even in separate source files or headers. The namespace scope is the accumulation of all definitions of the same namespace that the compiler has seen at the time it looks up a given name in the namespace. Namespaces can be nested. [Example 2-15](#) shows a sample namespace definition.

### *Example 2-15. Defining a namespace*

```
// The initial declaration
namespace numeric {
    class rational { ... }
    template<typename charT, typename traits>
    basic_ostream<charT,traits>& operator<<(
        basic_ostream<charT,traits>& out, const rational& r);
}

. . .

// This is a second definition. It adds an operator to the
// namespace.
namespace numeric {
    rational operator+(const rational&, const rational&);
}

// The definition of operator+ can appear inside or outside the
// namespace
// definition. If it is outside, the name must be qualified with
// the scope
// operator.
numeric::rational numeric::operator+(const rational& r1,
                                     const rational& r2)
{
    . . .
}

int main( )
{
    using numeric::rational;
    rational a, b;
    std::cout << a + b << '\n';
}
```

You can define a namespace without a name, in which case the compiler uses a unique, internal name. Thus, each

source file's unnamed namespace is separate from the unnamed namespace in every other source file.

You can define an unnamed namespace nested within a named namespace (and vice versa). The compiler generates a unique, private name for the unnamed namespace in each unique scope. As with a named namespace, you can use multiple namespace definitions to compose the unnamed namespace, as shown in [Example 2-16](#).

### *Example 2-16. Unnamed namespaces*

```
#include <iostream>
#include <ostream>

namespace {
    int i = 10;
}

namespace {
    int j;          // Same unnamed namespace
    namespace X {
        int i = 20; // Hides i in outer, unnamed namespace
    }
    namespace Y = X;
    int f( ) { return i; }
}

namespace X {
    int i = 30;
    // X::unnamed is different namespace than ::unnamed.
    namespace {
        int i = 40; // Hides ::X::i, but is inaccessible outside the
unnamed
                    // namespace
        int f( ) { return i; }
    }
}

int main( )
{
    int i = X::i; // ambiguous: unnamed::X or ::X?
    std::cout << ::X::f( ) << '\n'; // Prints 40
    std::cout << Y::i << '\n';      // Prints 20
    std::cout << f( ) << '\n';      // Prints 10
}
```

The advantage of using an unnamed namespace is that you are guaranteed that all names declared in it can never clash with names in other source files. The disadvantage is that you cannot use the scope operator (`::`) to qualify identifiers in an unnamed namespace, so you must avoid name collisions within the same source file.

### Tip

C programmers are accustomed to using global static declarations for names that are private to a source file. You can do the same in C++, but it is better to use an unnamed namespace because a namespace can contain any kind of declaration (including classes, enumerations, and templates), whereas static declarations are limited to functions and objects.

Declarations of static objects and functions at namespace scope are deprecated in C++.

Declarations outside of all namespaces, functions, and classes are implicitly declared in the global namespace. A program has a single global namespace, which is shared by all source files that are compiled and linked into the program. Declarations in the global namespace are typically referred to as *global declarations*. Global names can be accessed directly using the global scope operator (the unary `::`), as described earlier in [Section 2.3.1](#).

## Namespace Aliases

A namespace alias is a synonym for an existing namespace. You can use an alias name to qualify names (with the `::` operator) in using declarations and directives, but not in namespace definitions. [Example 2-17](#) shows some alias examples.

*Example 2-17. Namespace aliases*

```

namespace original {
    int f( );
}namespace ns = original;    // Alias

int ns::f( ) { return 42; } // OK
using ns::f;                // OK

int g( ) { return f( ); }

namespace ns { // Error: cannot use alias here
    int h( );
}

```

A namespace alias can provide an abbreviation for an otherwise unwieldy namespace name. The long name might incorporate a full organization name, deeply nested namespaces, or version numbers:

```

namespace tempest_software_inc {
    namespace version_1 { ... }
    namespace version_2 { ... }
}
namespace tempest_1 = tempest_software_inc::version_1;
namespace tempest_2 = tempest_software_inc::version_2;

```

## using Declarations

A using declaration imports a name from one namespace and adds it to the namespace that contains the using declaration. The imported name is a synonym for the original name. Only the declared name is added to the target namespace, which means using an enumerated type does not bring with it all the enumerated literals. If you want to use all the literals, each one requires its own using declaration.

Because a name that you reference in a using declaration is added to the current namespace, it might hide names in outer scopes. A using declaration can also interfere with local declarations of the same name.

[Example 2-18](#) shows some examples of using declarations.

### *Example 2-18. using declarations*

```
namespace numeric {
    class fraction { . . . };
    fraction operator+(int, const fraction&);
    fraction operator+(const fraction&, int);
    fraction operator+(const fraction&, const fraction&);
}

namespace eqn {
    class fraction { . . . };
    fraction operator+(int, const fraction&);
    fraction operator+(const fraction&, int);
    fraction operator+(const fraction&, const fraction&);
}

int main( )
{
    numeric::fraction nf;
    eqn::fraction qf;

    nf = nf + 1;           // OK: calls numeric::operator+
    qf = 1 + qf;          // OK: calls eqn::operator+
    nf = nf + qf;         // Error: no operator+

    using numeric::fraction;
    fraction f;           // f is numeric::fraction
    f = nf + 2;           // OK
    f = qf;               // Error: type mismatch
    using eqn::fraction; // Error: like trying to declare
                        // fraction twice in the same scope

    if (f > 0) {
        using eqn::fraction; // OK: hides outer fraction
        fraction f;         // OK: hides outer f
        f = qf;             // OK: same types
        f = nf;             // Error: type mismatch
    }
    int fraction;         // Error: name fraction in use
}
```

You can copy names from one namespace to another with a using declaration. Suppose you refactor a program and realize that the `numeric::fraction` class has all the functionality you need in the equation package. You decide to use `numeric::fraction` instead of `eqn::fraction`, but

you want to keep the eqn interface the same. So you insert `using numeric::fraction;` in the eqn namespace.

Incorporating a name into a namespace with a `using` declaration is not quite the same as declaring the name normally. The new name is just a synonym for the original name in its original namespace. When the compiler searches namespaces under argument-dependent name lookup, it searches the original namespace. [Example 2-19](#) shows how the results can be surprising if you are not aware of the `using` declaration. The eqn namespace declares `operator<<` to print a fraction, but `fraction` is declared in the numeric namespace. Although eqn uses `numeric::fraction`, when the compiler sees the use of `operator<<`, it looks in only the numeric namespace, and never finds `operator<<`.

### *Example 2-19. Creating synonym declarations with using declarations*

```
namespace eqn {
    using numeric::fraction;
    // Big, ugly declaration for ostream << fraction
    template<typename charT, typename traits>
    basic_ostream<charT,traits>& operator<<(
        basic_ostream<charT,traits>& out, const fraction& f)
    {
        out << f.numerator( ) << '/' << f.denominator( );
        return out;
    }
}

int main( )
{
    eqn::fraction qf;
    numeric::fraction nf;
    nf + qf;           // OK because the types are the same
    std::cout << qf; // Error: numeric namespace is searched for
                    // operator<<, but
                    // not eqn namespace
}
```



The using declaration can also be used within a class. You can add names to a derived class from a base class, possibly changing their accessibility. For example, a derived class can promote a protected member to public visibility. Another use of using declarations is for private inheritance, promoting specific members to protected or public visibility. For example, the standard container classes are not designed for public inheritance. Nonetheless, in a few cases, it is possible to derive from them successfully. [Example 2-20](#) shows a crude way to implement a container type to represent a fixed-size array. The array class template derives from `std::vector` using private inheritance. A series of using declarations make selected members of `std::vector` public and keep those members that are meaningless for a fixed-size container, such as `insert`, private.

*Example 2-20. Importing members with using declarations*

```
template<typename T>
class array: private std::vector<T>
{
public:
    typedef T value_type;
    using std::vector<T>::size_type;
    using std::vector<T>::difference_type;
    using std::vector<T>::iterator;
    using std::vector<T>::const_iterator;
    using std::vector<T>::reverse_iterator;
    using std::vector<T>::const_reverse_iterator;

    array(std::size_t n, const T& x = T( )) : std::vector<T>(n, x)
    {}
    using std::vector<T>::at;
    using std::vector<T>::back;
    using std::vector<T>::begin;
    using std::vector<T>::empty;
    using std::vector<T>::end;
    using std::vector<T>::front;
    using std::vector<T>::operator[];
    using std::vector<T>::rbegin;
    using std::vector<T>::rend;
```

```
using std::vector<T>::size;
};
```

See [Chapter 6](#) for more information about using declarations in class definitions.

## using Directives

A using directive adds a namespace to the list of scopes that is used when the compiler searches for a name's declaration. Unlike a using declaration, no names are added to the current namespace. Instead, the used namespace is added to the list of namespaces to search right after the innermost namespace that contains both the current and used namespaces. (Usually, the containing namespace is the global namespace.) The using directive is transitive, so if namespace A uses namespace B, and namespace B uses namespace C, a name search in A also searches C. [Example 2-21](#) illustrates the using directive.

*Example 2-21. The using directive*

```
#include <iostream>
#include <ostream>

namespace A {
    int x = 10;
}
namespace B {
    int y = 20;
}
namespace C {
    int z = 30;
    using namespace B;
}
namespace D {
    int z = 40;
    using namespace B;           // Harmless but pointless because D::y
    hides B::y
    int y = 50;
}

int main( )
{
```

```

int x = 60;
using namespace A;           // Does not introduce names, so there
is no conflict                // with x

using namespace C;

using namespace std;        // To save typing std::cout repeatedly

cout << x << '\n';          // Prints 60 (local x)
cout << y << '\n';          // Prints 20
cout << C::y << '\n';       // Prints 20
cout << D::y << '\n';       // Prints 50

using namespace D;
cout << y << '\n'; // Error: y is ambiguous. It can be found in
D::y and C's
                        // use of B::y.
}

```

## How to Use Namespaces

Namespaces have no runtime cost. Don't be afraid to use them, especially in large projects in which many people contribute code and might accidentally devise conflicting names. The following are some additional tips and suggestions for using namespaces:

- When you define a class in a namespace, be sure to declare all associated operators and functions in the same namespace.
- To make namespaces easier to use, keep namespace names short, or use aliases to craft short synonyms for longer names.
- Never place a using directive in a header. It can create name collisions for any user of the header.
- Keep using directives local to functions to save typing and enhance clarity.
- Use using namespace std outside functions only for tiny programs or for backward compatibility in legacy

projects.

## Chapter 3. Expressions

An expression combines literals, names, operators, and symbols to express or compute a value, or to achieve a side effect. This chapter describes the rules for writing and understanding expressions, including information on lvalues and rvalues, type conversion, constant expressions, and how C++ evaluates expressions. It contains detailed descriptions of all the operators and other forms of expressions.

The syntax descriptions in this chapter are informal. See [Chapter 12](#) for a precise BNF grammar.

### Lvalues and Rvalues

Lvalues and rvalues are fundamental to C++ expressions. Put simply, an *lvalue* is an object reference and an *rvalue* is a value. The difference between lvalues and rvalues plays a role in the writing and understanding of expressions.

An lvalue is an expression that yields an object reference, such as a variable name, an array subscript reference, a dereferenced pointer, or a function call that returns a reference. An lvalue always has a defined region of storage, so you can take its address.

An rvalue is an expression that is not an lvalue. Examples of rvalues include literals, the results of most operators, and function calls that return nonreferences. An rvalue does not necessarily have any storage associated with it.

Strictly speaking, a function is an lvalue, but the only uses for it are to use it in calling the function, or determining the function's address. Most of the time, the term lvalue means object lvalue, and this book follows that convention.

C++ borrows the term lvalue from C, where only an lvalue can be used on the left side of an assignment statement. The term rvalue is a logical counterpart for an expression that can be used only on the righthand side of an assignment. For example:

```
#define rvalue 42
int lvalue;
lvalue = rvalue;
```

In C++, these simple rules are no longer true, but the names remain because they are close to the truth. The most significant departure from traditional C is that an lvalue in C++ might be const, in which case it cannot be the target of an assignment. (C has since evolved and now has const lvalues.)

The built-in assignment operators require an lvalue as their lefthand operand. The built-in address (&) operator also requires an lvalue operand, as do the increment (++) and decrement (- -) operators. Other operators require rvalues. The rules are not as strict for user-defined operators. Any object, including an rvalue, can be used to call member functions, including overloaded =, &, ++, and - - operators.

Some other rules for lvalues and rvalues are:

- An array is an lvalue, but an address is an rvalue.
- The built-in array subscript ([ ]), dereference (unary \*), assignment (=, +=, etc.), increment (++), and decrement (- -) operators produce lvalues. Other built-in operators produce rvalues.
- A type cast to a reference type produces an lvalue; other casts produce rvalues.
- A function call (including overloaded operators) that returns a reference returns an lvalue; otherwise, it returns an rvalue.

- An lvalue is converted implicitly to an rvalue when necessary, but an rvalue cannot be implicitly converted to an lvalue.

[Example 3-1](#) shows several different kinds of lvalues and rvalues.

### *Example 3-1. Lvalues and rvalues*

```
class number {
public:
    number(int i = 0) : value(i) {}
    operator int( ) const { return value; }
    number& operator=(const number& n);
private:
    int value;
};

number operator+(const number& x, const number& y);

int main( )
{
    number a[10], b(42);
    number* p;
    a;           // lvalue
    a[0];        // lvalue
    &a[0];       // rvalue
    *a;          // lvalue
    p;           // lvalue
    *p;          // lvalue
    10;          // rvalue
    number(10); // rvalue
    a[0] + b;    // rvalue
    b = a[0];    // lvalue
}
```

## Type Conversions

In an arithmetic expression, binary operators require operands of the same type. If this is not the case, the type of one operand must be converted to match that of the other operand. When calling a function, the argument type must match the parameter type; if it doesn't, the argument is converted so its type matches. C++ has cast operators, which let you define a type conversion explicitly, or you can let the compiler automatically convert the type for you. This section presents the rules for automatic type conversion.

## Arithmetic Types

An *arithmetic type* is a fundamental integral or floating-point type: `bool`, `char`, `signed char`, `unsigned char`, `int`, `short`, `long`, `unsigned int`, `unsigned short`, `unsigned long`, `float`, `double`, or `long double`. Some operations are permitted only on arithmetic types, pointer types, enumerations, class types, or some combination of types. The description of each operator tells you which types the operator supports.

## Type Promotion

*Type promotion* is an automatic type conversion that applies only to arithmetic types, converting a "smaller" type to a "larger" type while preserving the original value. Contrast promotions with other automatic type conversions (described in later subsections), which can lose information. Promotions involve either integral or floating-point values. The rules for integral promotion are:

- A "small" integral rvalue is converted to an `int` if the type `int` can represent all of the values of the source



type; otherwise, it is converted to unsigned int. A "small" value is an integral bit-field (see [Chapter 6](#)) whose size is smaller than an int, or a value with one of the following types: char, signed char, unsigned char, short int, unsigned short int.

- An rvalue whose type is wchar\_t or an enumeration (including bit-fields with enumeration type) is converted to the first type that can hold all the values of the source type. The type is one of: int, unsigned int, long, or unsigned long.
- An rvalue of type bool can be converted to an int; the value true becomes 1, and the value false becomes 0.

One floating-point promotion rule is defined:

- An rvalue of type float can be converted to type double.

## Arithmetic Type Conversions

An arithmetic type conversion is an automatic type conversion that the compiler applies to the operands of the built-in arithmetic and comparison operators. For the arithmetic operators, the result type is the same as the operand type.



The arithmetic type conversions try to preserve the original values as much as possible. However, they do not always succeed. For example, `-1 / 1u` does not produce `-1` as a result because `-1` has type `int`, which is converted to type `unsigned int`, yielding an implementation-defined arithmetic value. On a 32-bit, two's-complement system, the result is `4294967295u`.

The rules for arithmetic type conversion are applied in the following order:

1. If one operand has type `long double`, the other is converted to `long double`.
2. Otherwise, if one operand is `double`, the other is converted to `double`.
3. Otherwise, if one operand is `float`, the other is converted to `float`.
4. Otherwise, integral promotions are performed (see the section [Section 3.2.2](#)).
5. After integral promotion, if one operand is `unsigned long`, the other is converted to `unsigned long`.
6. Otherwise, if one operand is `long`, and the other is `unsigned int`, type conversion occurs as follows:
  1. If all values of type `unsigned int` fit in a `long int`, the `unsigned int` is converted to `long int`.
  2. Otherwise, both operands are converted to `unsigned long`.
7. Otherwise, if one operand is `long`, the other is converted to `long`.
8. Otherwise, if one operand is `unsigned`, the other is converted to `unsigned`.
9. Otherwise, both operands are `int`.

## Implicit Numeric Conversions

An arithmetic value can be converted implicitly to a different arithmetic value, even if that conversion loses information. These implicit conversions can take place in assignments, initializers, arguments to function calls, returning values from a function, and template arguments.

Do not confuse these conversions with the arithmetic type conversions described earlier, which can take place in any arithmetic operation.

The basic rule is that any arithmetic type can be converted to any other. If the destination type cannot hold the converted value, the behavior is undefined. When assigning a floating-point number to an integer, the floating-point value is truncated toward zero by discarding the fractional part. When converting an integer or enumerated value to a floating-point number, if the integer value falls between two floating-point values (that is, the integer is not representable in the floating-point format), the implementation chooses one of the two neighboring values.

A value can be converted to a class type if the class has a suitable constructor that does not use the `explicit` specifier. A value of class type can be converted to a non-class type if the source class has a type conversion operator of the target type. When converting a class type to another class type, the compiler considers the constructors for the target class and the type conversion operators for the source class.

## **Lvalue Conversions**

Lvalue conversions automatically convert an lvalue to an rvalue for contexts in which an rvalue is required. The need to convert an lvalue to an rvalue is usually transparent, and these conversions are listed only for the sake of completeness:

- An array lvalue can be converted to a pointer rvalue, pointing to the first element of the array.
- A function lvalue can be converted to an rvalue pointer to the function.

- Any other lvalue can be converted to an rvalue with the same value. If the type is not a class type, *cv*-qualifiers are removed from the type. Thus, a `const int` lvalue is converted to an `int` rvalue.

## Conversion to bool

Arithmetic, enumeration, and pointer values can be converted to `bool`, in which null pointers and zero arithmetic values are false, and everything else is true. A common idiom is to test whether a pointer is a null pointer:

```
if (ptr)
    ... // Do something with *ptr.
else
    ... // Pointer is null

if (! ptr)
    // Another way to test if ptr is null
```

This idiom is so common that some classes have operators to convert an object to `void*`, so `void*` can be converted to `bool`. (Any pointer type would do, but `void*` is used because the pointer is not meant to be dereferenced, only converted to `bool`.) The `void*` operator is used instead of a direct conversion to `bool` to avoid automatic promotion to an integer. For example, `basic_ios` defines operator `void* ( )` to return a null pointer if the `iostream`'s `failbit` is set. (See `<ios>` in [Chapter 13](#) for details.)

## Type Casts

A *type cast* is an explicit type conversion. C++ offers several different ways to cast an expression to a different type. The different ways expose the evolution of the language. The six forms of type cast expressions are:

- `( type ) expr`
- `type ( expr )`

- `const_cast< type >( expr )`
- `dynamic_cast< type >( expr )`
- `reinterpret_cast< type >( expr )`
- `static_cast< type >( expr )`

The first form was inherited from C; the second form was added in the early days of C++ and has the same meaning as the first, but with a slightly different syntax. The final four forms supplant the two older forms and are preferred because they are more specific, thus reducing the risk of error. All type cast expressions are described in detail later in this chapter; the first form is covered in the section [Section 3.5.4](#), and the other five forms are covered in [Section 3.5.2](#). The remainder of this section briefly discusses when to use each form.

If you simply want to force an expression to be `const`, or to remove a `const` (e.g., prior to passing a pointer to a legacy library that is not written in C++ and does not know about `const` declarations), use `const_cast<>`. (You can also change the `volatile` qualifier, but that is less common.)

If you have a pointer or reference to an object whose declared type is a base class, but you need to obtain a derived class pointer or reference, you should use `dynamic_cast<>`. The dynamic cast performs a runtime check to make sure the cast is valid. A dynamic cast works only with polymorphic classes, which have at least one virtual function.

The common uses of `static_cast<>` are to force one enumeration to a different enumerated type, force an arithmetic type to a different type, or force a particular conversion from a class-type object that has multiple type conversion operators. The simpler type casts are sometimes used in these cases, for the sake of brevity. For example, `sqrt(float(i))` can be easier to read than

`sqrt(static_cast<float>(i))`). Use `static_cast` to reverse any implicit type conversion. For example, C++ automatically converts an enum to an integer; use `static_cast` if you want to convert an integer to an enum. You can use `static_cast` to cast a pointer to and from `void*`.

A `reinterpret_cast<>` is reserved for potentially unsafe type casts, such as converting one type of pointer to another. You can also convert a pointer to an integer or vice versa. For example, an internal debugging package might record debug log files. The logger might convert pointers to integers using `reinterpret_cast<>` and print the integers using a specific format.

If you try to perform the wrong kind of cast with one of the template-like cast operators, the compiler informs you of your mistake. For example, if you use `static_cast<>` and accidentally cast away `const`-ness, the compiler complains. Or if you use `static_cast<>` where you should have used `reinterpret_cast<>`, the compiler complains. The short forms of casting do not provide this extra level of error-checking because one form must suffice for the several different kinds of type casts.

When you see a type cast, you should read it as a warning that something unusual is happening. The longer forms of type casts provide additional clues about the programmer's intent, and help the compiler enforce that intent.

## Constant Expressions

A constant expression is an expression that can be evaluated at compile time. Constants of integral or enumerated type are required in several different situations, such as array bounds, enumerator values, and case labels. Null pointer constants are a special case of integral constants.

## Integral Constant Expressions

An integral constant expression is an expression that can be evaluated at compile time, and whose type is integral or an enumeration. The situations that require integral constant expressions include array bounds, enumerator values, case labels, bit-field sizes, static member initializers, and value template arguments. The compiler must be able to evaluate the expression at compile time, so you can use only literals, enumerators, `const` objects that have constant initializers, integral or enumerated template parameters, `sizeof` expressions, and constant addresses. The address of a static lvalue object is a constant address, as is the address of a function. A string literal, being a static array of characters, is also a constant address.

An integral static `const` data member can be initialized in the class definition if the initializer is a constant integral or enumerated expression. The member can then be used as a constant expression elsewhere in the class definition. For example:

```
template<typename T, size_t size>
class array {
public:
    static const size_t SIZE = size;
    ...
private:
```

```
    T data[SIZE];  
};
```

See [Chapter 6](#) for more information about static data members.

## Null Pointers

A constant expression with a value of `0` can be a *null pointer constant*. A null pointer constant can be converted to a *null pointer value*. The colloquial term null pointer almost always means null pointer value.



A null pointer value has an implementation-defined bit pattern. Many implementations use all zero bits, but some do not. Thus, the null pointer constant is not representative of the bits that make up a null pointer value, but serves only as a mnemonic for the programmer, much like `= 0` does for a pure virtual function ([Chapter 6](#)).

When you assign a null pointer constant to a pointer-type variable, the compiler converts the null pointer constant to a null pointer value of the appropriate type. Similarly, when comparing pointers, the compiler ensures that the comparison is meaningful. In particular, a null pointer value is never equal to any valid pointer value, is always equal to another null pointer value, and is always equal to a null pointer constant. A null pointer value, when converted to `bool`, is `false`. A pointer that is initialized with an empty initializer is initialized to a null pointer value.

The `NULL` macro, defined in `<cstdlib>` and other headers (see [Chapter 13](#)), expands to a null pointer constant. Using `NULL` instead of `0` can be a helpful reminder to the reader or maintainer of a program, especially when typedefs obscure the nature of a type:



```
Token tok1 = 0;    // Is Token a pointer or an integer?  
Token tok2 = NULL; // Token is probably a pointer.
```

Dereferencing a null pointer results in undefined behavior.

## Expression Evaluation

At its most fundamental level, the execution of a C++ program is the successive evaluation of expressions, under the control of statements ([Chapter 4](#)), in which some expressions can produce side effects. Any expression might have one or more of the following side effects:

- Accessing a volatile object
- Modifying an object
- Calling a function in the standard library
- Calling any other function that has side effects

## Sequence Points

During the execution of a program, there are well-defined points in time called *sequence points*, at which the side effects have all been completed for expressions that have been evaluated, and no side effects have been started for any unevaluated expression. Between sequence points, the compiler is free to reorder expressions in any way that preserves the original semantics. The same term also refers to the positions in the source code that produce sequence points when the code executes. You can usually ignore the details of sequence points, but when you are using global or volatile objects, it is important that you know exactly when it is safe to access those objects. That time is after a sequence point. Also, any expression that modifies a scalar object more than once between sequence points, or that examines a scalar object's value after modifying it, yields undefined behavior. This rule often bites the unwary programmer who uses the increment and decrement operators. For example:

```

int i = 0;
i = ++i - ++i;           // Error: undefined behavior
printf("%d,%d", ++i, ++i); // Error: undefined behavior
i = 3, ++i, i++;        // OK: i == 5

```

There are sequence points in the following positions:

- At the end of every expression that is not a subexpression. Such an expression might be used in an expression statement, in an initializer, as a condition in an `if` statement, etc.
- After evaluating all function arguments but before calling the function.
- When a function returns: after copying the return value from the function call (if any), but before evaluating any other expressions outside the function.
- After evaluating the first expression (`expr1`) in each of the following expressions, provided they use the built-in operators and not overloaded operators:
  - `expr1 && expr2`
  - `expr1 || expr2`
  - `expr1 ? expr2 : expr3`
  - `expr1 , expr2`

## Order of Evaluation

In general, the order in which operands are evaluated is unspecified, so you should never write code that depends on a particular order. For example, in the expression `f( ) / g( )`, `f( )` might be called first, or `g( )` might be called first. The difference can be significant when the functions have side effects. [Example 3-2](#) shows a contrived situation in which a program prints 2 if `g( )` is called first, or 1 if `f( )` is called first.

*Example 3-2. Demonstrating order of evaluation*

```

#include <iostream>
#include <ostream>

int x = 1;

int f( )
{
    x = 2;
    return x;
}

int g( )
{
    return x;
}

int main( )
{
    std::cout << f( ) / g( ) << '\n';
}

```

A simpler example follows. The increment of `i` can happen before or after the assignment, so `i` might be 2 or 3.

```

int i = 1;
i = i++ + 1; // Value of i is unspecified

```

In a function call, all arguments are evaluated before the function is called. As you might expect, the order in which the arguments are evaluated is unspecified.

## Short-Circuit Evaluation

The logical operators (`&&` and `||`) perform *short-circuit evaluation*. The left operand is evaluated, and if the expression result can be known at that point, the right operand is *not* evaluated:

```

if (false && f( )) ... // f( ) is never called.
if (true || f( )) ... // f( ) is never called.

```

If the logical operator is overloaded, however, it cannot perform short-circuit evaluation. Like any other function, all the arguments are evaluated before the function is

called. For this reason, you should avoid overloading the `&&` and `||` operators.

## Expression Rules

C++ has the usual unary operators such as logical negation (!a), binary operators such as addition (a+b), and even a ternary operator (a?b:c). Unlike many other languages, an array subscript is also an operator (a[b]), and a function call is an n-ary operator (e.g., a(b, c, d)).

Every operator has a *precedence*. Operators with higher precedence are grouped so that they are logically evaluated before operators with lower precedence. (Note that precedence determines how the compiler parses the expression, not necessarily the actual order of computation. For example, in the expression a( ) + b( ) \* c( ), the multiplication has higher precedence, but a( ) might be called first.)

Some operators group from left to right. For example, the expression x / y / z is equivalent to (x / y) / z. Other operators group right to left, as in x = y = z, which is equivalent to x = (y = z). The order of grouping is called the operator's *associativity*.

When reading C++ expressions, you must be aware of the precedence and associativity of the operators involved. For example, \*ptr++ is read as \*(ptr++) because the postfix ++ operator has higher precedence than the unary \* operator.

[Table 3-1](#) summarizes the syntax, precedence, and associativity of each kind of expression. The subsections that follow describe the kinds of expressions in depth; each subsection covers a single precedence group.

*Table 3-1. Expression syntax and associativity, grouped by precedence*

Group	Associativity	Expression
-------	---------------	------------

Group	Associativity	Expression
Primary (highest precedence)	Left-to-right	<pre> literal this(   expr ) name ::name class-or-namespace :: name </pre>
Postfix	Left-to-right	<pre> pointer [   expr ] expr (   expr, ... ) type (   expr, ... ) object.member  pointer -&gt; member  cast_keyword &lt; type &gt;(   expr )  typeid(   expr )  typeid(   type ) lvalue ++ lvalue -- </pre>

Group	Associativity	Expression
Unary	Right-to-left	<pre> ++ lvalue -- lvalue ~ expr compl expr ! expr not expr + expr - expr * pointer &amp; lvalue sizeof expr sizeof( type ) new-expr delete-expr </pre>
Cast	Right-to-left	<pre> ( type ) expr </pre>



Group	Associativity	Expression
Pointer to Member	Left-to-right	$object$ $. *$ $expr$ $pointer$ $-> *$ $expr$
Multiplicative	Left-to-right	$expr$ $*$ $expr$ $expr$ $/$ $expr$ $expr$ $\%$ $expr$
Additive	Left-to-right	$expr$ $+$ $expr$ $expr$ $-$ $expr$
Shift	Left-to-right	$expr$ $<<$ $expr$ $expr$ $>>$ $expr$

Group	Associativity	Expression
Relational	Left-to-right	$expr < expr$ $expr > expr$ $expr <= expr$ $expr >= expr$
Equality	Left-to-right	$expr == expr$ $expr != expr$ $not\_eq\ expr$
Bitwise And	Left-to-right	$expr \& expr$ $bitand\ expr$
Bitwise Exclusive Or	Left-to-right	$expr \wedge expr$ $xor\ expr$

Group	Associativity	Expression
Bitwise Inclusive Or	Left-to-right	$expr$ $ $ $expr$  $expr$ bitor $expr$
Logical And	Left-to-right	$expr$ $\&\&$ $expr$  $expr$ and $expr$
Logical Or	Left-to-right	$expr$ $  $ $expr$  $expr$ or $expr$
Conditional	Right-to-left	$expr$ $?$ $expr : expr$

Group	Associativity	Expression
Assignment	Right-to-left	$lvalue$ $=$ $expr$  $lvalue$ $op$ $=$ $expr$ $throw$ $expr$ $throw$
Comma (lowest precedence)	Left-to-right	$expr , expr$

## Primary Expressions

A primary expression is the basic building block for more complex expressions. It is an expression in parentheses, a literal, or a name (possibly qualified). The various forms of primary expressions are:

### *literal*

A constant value. String literals (being arrays of `const char` or `const wchar_t`) are lvalues. All other literals are rvalues. (See [Chapter 1](#).)

### *this*

Refers to the target object in a nonstatic member function. Its type is a pointer to the class type; its value is an rvalue.

### *( expression )*

Has the type and value of *expression*.

### *unqualified-name*

Names an entity according to the name lookup rules in [Chapter 2](#). The expression result is the entity itself, and the expression type is the entity's type. The result is an lvalue if the entity is an object, data member, or function. The following are the various kinds of unqualified identifiers:

#### *identifier*

Names an object, function, member, type, or namespace. The name is looked up according to the rules in [Chapter 2](#). The type is the declared type of the entity. If the entity is an object, data member, or function, the value is an lvalue.

#### *operator symbol*

Names an operator. See [Chapter 5](#) for more information.

#### *template-name < optional-template-args >*

Names a template instance. See [Chapter 7](#) for more information.

#### *operator type*

Names a type conversion operator. The *type* is a type specifier, possibly with one or more pointer symbols in the declarator. (See [Chapter 2](#) for details about type specifiers and declarators.)

#### *~ class-name*

Names the destructor for *class-name*.

### *qualified-name*

Uses the scope operator to qualify an identifier, operator, or destructor. The qualified name can be in the global scope or in the scope of a class or namespace:

*:: identifier*

Names a global identifier. The type is the declared type of the entity. If the entity is an object, data member, or function, the value is an lvalue; otherwise, it is an rvalue.

*:: operator symbol*

Names a global operator. Note that type conversion operators must be member functions, so there cannot be a global type conversion operator. See [Chapter 5](#) for more information.

*nested-name :: unqualified-name, nested-name ::  
template unqualified-name*

Names an entity in a class or namespace scope. The *nested-name* can be a class or namespace name, or it can have the form *class-or-namespace-name :: nested-name* or *class-name :: template nested-name*. Use the `template` keyword when instantiating a template. See [Chapter 7](#) for information about template members.

*:: nested-name :: unqualified-name, :: nested-name  
:: template unqualified-name*

Names an entity in a class or namespace scope. The first (left-most) class or namespace name is looked up in the global scope.

In the rest of this chapter, the syntax element *name-expr* refers to a qualified or unqualified name, as described in this section. In particular, a *name-expr* can be used to the right of the `.` or `->` in a postfix expression.

[Example 3-3](#) shows some primary expressions.

*Example 3-3. Primary expressions*

```

namespace ns {
    int x;
    class cls {
    public:
        cls(int);
        ~cls( );
    };
}
int x;

3.14159          // Literal
(2 + 3 * 4)      // Parenthesized expression
x                // Unqualified identifier
ns::x            // Qualified identifier
ns::cls::cls     // Qualified constructor
operator*        // Unqualified operator name

```

## Postfix Expressions

A postfix expression is an expression that uses postfix syntax (operator follows the operand) with some exceptions that just happen to have the same precedence. The postfix expressions are:

*pointer* [ *expr* ]

Returns an element of an array. The subscript operator requires a pointer as the left operand. An array is implicitly converted to a pointer. The right operand is converted to an integer, and the expression is evaluated as  $*((\textit{pointer}) + (\textit{expr}))$ . If the array index is out of bounds, the behavior is undefined. The result is an lvalue whose type is the base type of *pointer*.

*expr* ( *optional-expr-list* )

Calls a function. The function call operator requires one of the following as a left operand: a function name, an expression that returns a function pointer, or an expression that returns an object that has a function call operator. (An operator name is the same as a function name in this case.) The *optional-expr-list* is a

comma-separated list of zero or more assignment expressions. (See [Section 3.5.17](#) later in this chapter.) All the expressions in the expression list are evaluated, and then the function is called. The result type is the return type of the function. If the return type is a reference type, the result is an lvalue; otherwise, the result is an rvalue. If the return type is void, no value is returned. See [Chapter 5](#) for more information about functions.

*simple-type-specifier ( optional-expr-list )*

Performs type conversion or construction. A *simple-type-specifier* is a single name: a fundamental type or a qualified name of a class, an enumeration, or a typedef. The result is an instance of the specified type, initialized as follows:

- If the expression list is empty, the result is an rvalue that is initialized with a default constructor, or that is initialized to 0. (See [Chapter 2](#) for details about initialization.)
- If the expression list contains one expression, that value is cast to the desired type in the same manner as a cast expression—that is, ( *type* ) *expr*. If the type is a reference, the result is an lvalue; otherwise, it is an rvalue.
- If the expression list contains more than one expression, the type must be a class, and the expression list is passed as arguments to a suitable constructor to create an instance of the class, which is returned as an rvalue.

*object . name-expr*

Returns a member of an object. The name can be qualified to refer to a name in a base class (see [Chapter 2](#)). The return type is the type of *name-expr*. The return



value depends on whether *name-expr* is a data member, member function, or enumerator:

- If *name-expr* names a static data member, the member is returned as an lvalue.
- If *name-expr* names a nonstatic data member, the result is an lvalue only if *object* is an lvalue. If *name-expr* is declared mutable, the result is not const even if *object* is const; otherwise, the result is const if either *object* or *name-expr* is const. Similarly, the result is volatile if either *object* or *name-expr* is volatile.
- If *name-expr* names a member function, the usual rules for overload resolution apply (see [Chapter 5](#)). If the function is a static member function, the result is an lvalue. You can take the function's address (with &) or call the function.
- If the function is a nonstatic member function, it must be used in a function call—for example, `obj.memfun(arg)`.
- If *name-expr* is an enumerator, the result is an rvalue.

*pointer* -> *name-expr*

Returns `*(pointer)`.*name-expr*.

*lvalue* ++

Increments *lvalue* and returns its value prior to incrementing (as an rvalue). The type of *lvalue* must be arithmetic or a pointer. The new value is *lvalue* + 1.

If *lvalue* has type `bool`, the new value is always `true`. This `bool`-specific behavior is deprecated.

*lvalue* --

Decrements *lvalue* and returns its value prior to decrementing (as an rvalue). The type must be arithmetic or a pointer and cannot be `bool`. The new value is *lvalue* - 1.

`const_cast< type >( expr )`

Casts *expr* to *type*. If *type* is a reference, the result is an lvalue; otherwise, it is an rvalue. The new type must match the type of *expr*, but the `const` and `volatile` qualifiers can be changed.

A `const_cast` that removes a `const` qualifier is generally a bad idea. Nonetheless, it is sometimes necessary to cast away `const`-ness, especially when passing pointers to legacy libraries.

See [Chapter 6](#) for a discussion of the mutable modifier, which lets you modify a data member of a `const` object.

`dynamic_cast< type >( expr )`

Casts a base class pointer or reference *expr* to a derived class *type*. A runtime check is performed to make sure the true class of *expr* is *type* or a class derived from *type*. The class must be polymorphic, that is, have at least one virtual function. The base class can be virtual. A `dynamic_cast<>` cannot cast away `cv`-qualifiers. The cast works as follows:

- If *type* is `void*`, the return value is a pointer to the most-derived object that *expr* points to. The type does not have to be polymorphic in this case.
- If *expr* is a pointer, *type* must be a pointer type. If the type of *expr* does not match *type* (is not the same as *type* or derived from *type*), a null pointer value is returned. Otherwise, the value of *expr* cast to *type* is returned. If *expr* is a null pointer, a null pointer is returned.

- If *expr* is an object and *type* is a reference, *expr* is cast to *type*. If the type of *expr* does not match, a `bad_cast` exception is thrown.
- You can also cast from a derived class to a base class, which is the same as an ordinary implicit conversion. The type does not have to be polymorphic in this case.

[Example 3-4](#) shows some uses of `dynamic_cast<>`.

*Example 3-4. Using `dynamic_cast<>`*

```
#include <iostream>
#include <ostream>

class base {
public:
    virtual ~base( ) {}
};

class derived : public base {};
class most_derived : public derived {};
class other : public base {};

int main( )
{
    base* b = new derived;
    dynamic_cast<most_derived*>(b); // Null pointer
    dynamic_cast<derived&>(*b);    // OK
    dynamic_cast<other*>(b);       // Null pointer

    derived* d = new most_derived;
    b = d;
    b = dynamic_cast<base*>(d);    // OK, but dynamic_cast<>
                                   // is unnecessary.
}
```



`reinterpret_cast< type >( expr )`

Casts *expr* to *type*. When using `reinterpret_cast<>`, no conversion functions or constructors are called. Casting

to a reference yields an lvalue; otherwise, it yields an rvalue. Only the following conversions are allowed:

- A pointer can be converted to an integer. The integer must be large enough to hold the pointer's value. Which integer type you should use is implementation-defined, as is the mapping from pointer to integer.
- An integer or enumerated value can be converted to a pointer. The mapping is implementation-defined. Casting from a pointer to an integer back to the original pointer type yields the original pointer value, provided the integer type is large enough.
- Casting an integer constant of value 0 to a pointer always produces a null pointer value. Casting any other integer expression of value 0 to a pointer produces an implementation-defined pointer, which may or may not be a null pointer.
- A function pointer can be converted to a function pointer of a different type. Calling such a function results in undefined behavior. Converting back to the original pointer type produces the original pointer value.
- An object pointer can be converted to an object pointer of a different type. Using such an object results in undefined behavior. Converting back to the original pointer type produces the original pointer value.
- A pointer to a member can be converted to a pointer to a different member. Using the pointer to a member has undefined behavior, except that casting a pointer to a data member to a different pointer to a data member and back to the original type produces the original value, and casting a pointer to a member

function to a different member function and back produces the original value.

- A null pointer constant or value can be converted to a null pointer of the target type.
- A reference can be cast in the same manner as a pointer (except that the pointer is dereferenced). That is, casting `reinterpret_cast<T&>(x)` is just like casting `*reinterpret_cast<T*>(&x)`.

A `reinterpret_cast` has the following restrictions:

- A function pointer cannot be converted to an object pointer, or an object pointer to a function pointer.
- A member function pointer cannot be converted to a data member pointer, or a data member pointer to a member function pointer.
- A member pointer cannot be converted to a nonmember pointer, or a nonmember pointer to a member pointer.
- The target type must not cast away *cv*-qualifiers.

The need for `reinterpret_cast<>` is rare in an ordinary application.

[Example 3-5](#) shows some uses of `reinterpret_cast<>`. The first use reinterprets the representation of a float as an int to show the underlying representation of a float. This use is implementation-dependent and requires `sizeof(int)` to be greater than or equal to `sizeof(float)`. The second use is a simple conversion from a function pointer to an integer to print the address in a specific format. It requires that an int be large enough to hold a function pointer.

*Example 3-5. Using `reinterpret_cast<>`*

```

#include <cassert>
#include <iomanip>
#include <iostream>
#include <ostream>

int foo( )
{
    return 0;
}

int main( )
{
    using namespace std;
    float pi = 3.1415926535897;
    int    ipi;

    // Print numbers in pretty hexadecimal.
    cout << setfill('0') << showbase << hex << internal;

    // Show the representation of a floating-point number.
    assert(sizeof(int) == sizeof(float));
    ipi = reinterpret_cast<int&>(pi);
    cout << "pi bits=" << setw(10) << ipi << '\n';

    // Show the address of foo( ).
    cout << "&foo=" << setw(10) <<
        reinterpret_cast<int>(&foo) << '\n';
}

```

`static_cast< type >( expr )`

Casts *expr* to *type* using a standard or user-defined conversion, as though you declared a temporary *type tmp* (*expr*) and used the value of *tmp* in place of the cast expression. The result is an lvalue if *type* is a reference; otherwise the result is an rvalue. A `static_cast<>` cannot cast away cv-qualifiers. The rules of permitted conversions are:

- The *type* can be void, in which case the result of *expr* is discarded.
- A base-class lvalue can be cast to a reference to a derived class, provided a standard conversion exists from a derived-class pointer to the base class. The

base class must not be virtual. If *expr* is not actually a subobject of *type*, the behavior is undefined. (See `dynamic_cast<>` to learn how to make error-checking safe.)

- A base-class pointer can be converted to a derived-class pointer in the manner described for class references.
- A pointer to a member of a derived class can be converted to a pointer to a member of a base class if there is a standard conversion in the other direction. The target class (or an ancestor class) must contain the original member.
- Standard arithmetic conversions can work in reverse—for example, `long` can be cast to `short`. Integers can be cast to enumerations.
- Enumerations can be cast to other enumerations.
- A `void` pointer can be converted to any object pointer. Converting a pointer to `void*` (with the same *cv*-qualifiers as the original pointer) and back produces the original pointer value.

[Example 3-6](#) shows some uses of `static_cast<>` .

*Example 3-6. Using `static_cast<>`*

```
#include <iostream>
#include <ostream>

class base {};
class derived : public base {};
class other : public base {};

enum color { red, black };
enum logical { no, yes, maybe };

int main( )
{
    base* b = new derived;
    static_cast<derived*>(*b); // OK
```

```

    static_cast<other*>(b);    // Undefined behavior

    derived* d = new derived;
    b = d;
    b = static_cast<base*>(d); // OK, but unnecessary

    color c = static_cast<color>(yes);

    int i = 65;
    std::cout << static_cast<char>(i);
}

```

`typeid( expr )`

Returns type information for the type of *expr* without evaluating *expr*. The type information is an lvalue of type `const std::type_info` (or an implementation-defined type that derives from `type_info`). See `<typeinfo>` in [Chapter 13](#) for information about this class.

If *expr* is an lvalue of a polymorphic type (a class with at least one virtual function), the type information is for the most-derived class of *expr*. If *expr* is a dereference of a null pointer, `bad_typeid` is thrown.

If *expr* is not an lvalue, or the type is not polymorphic, the type information is for the static type of *expr*.

`typeid( type )`

Returns the type information for *type* as described for `typeid( expr )`. [Example 3-7](#) shows some uses of `typeid`.

*Example 3-7. Using typeid*

```

#include <iostream>
#include <ostream>
#include <typeinfo>

class base {
public:
    virtual ~base( ) {}
};

```



```

class derived : public base {};
enum color { red, black };

// The actual output is implementation-defined, but should
// reflect the types
// shown in the comments.
int main( )
{
    base* b = new derived;
    std::cout << typeid(*b).name( ) << '\n'; // Derived
    std::cout << typeid(base).name( ) << '\n'; // Base
    derived* d = new derived;
    std::cout << typeid(*d).name( ) << '\n'; // Derived
    std::cout << typeid(derived).name( ) << '\n'; // Derived
    std::cout << typeid(red).name( ) << '\n'; // Color
    std::cout << typeid(color).name( ) << '\n'; // Color
}

```

## Unary Expressions

A unary expression uses a unary prefix operator:

**++ *lvalue***

Increments *lvalue*, which must be of arithmetic or pointer type, and returns the new value as an lvalue. The expression ++*x* is equivalent to *x* += 1, unless *x* is of type `bool`, in which case the expression ++*x* is equivalent to *x*=`true`. The special-case behavior for `bool` is deprecated.

**-- *lvalue***

Decrements *lvalue*, which must be of arithmetic or pointer type (not `bool`), and returns the new value as an lvalue. The expression --*x* is equivalent to *x* -= 1.

**\* *pointer***

Dereferences *pointer* and returns an lvalue for the object that *pointer* points to. If *pointer* has type *T*\*, the expression has type *T* (preserving any *cv*-qualifiers).

**& *lvalue*, & *qualified-name***

Returns the address of *lvalue* or *qualified-name*. If *lvalue* has type *T* or if *qualified-name* is a static member of type *T*, the result is the object's address, which is an rvalue of type pointer to *T*. If *qualified-name* is a nonstatic member of class *C*, the result type is a pointer to a member of class *C*.

Note that a pointer to a member is formed only by applying the & operand to a qualified name. Even in the scope of a class, &*unqualified-name* is not a pointer to a member, but an ordinary pointer to an object.

You cannot take the address of a bit-field. To take the address of an overloaded function, the context must make it clear which function you mean. [Example 3-8](#) shows uses of the & operator.

### *Example 3-8. Using the & operator*

```
class demo
{
public:
    int x;
    static int y;
    int get_x( ) { return x; }
};

int demo::y = 10;

int add(int a, int b) { return a + b; }
double add(double a, double b) { return a + b; }

int main( )
{
    demo d;
    int demo::*p;
    int (demo::*func)( );
    int *i;
    int local = 42;
    int *ptr = &local;

    p = &demo::x;
    i = &demo::y;
    func = &demo::get_x;
```

```

d.*p = *ptr;
*i = (d.*func)( );

int (*add)(int, int);
add = &add;
d.*p = add(42, *i);

return d.y;
}

```

+ *expr*

Returns *expr*, which must have arithmetic, enumeration, or pointer type. The usual type promotions take place, and the result type is the promoted type. The result is an rvalue.

- *expr*

Negates *expr*, which must have arithmetic or enumerated type. The usual type promotions take place, and the result type is the promoted type. The result is an rvalue. If the type is unsigned, the result is  $2^n - expr$ , in which  $n$  is the number of bits in the result type.

~ *expr*, compl *expr*

Returns the bitwise complement of *expr*, which must have integral or enumeration type. The type is promoted according to the usual rules, and the result type is the promoted type. The result is an rvalue. Every zero bit in *expr* is converted to a one bit, and every one bit is converted to a zero bit.

In the ambiguous case of  $\sim C( )$ , in which  $C$  is a class name,  $\sim C( )$  is interpreted as the complement operator, not the destructor of  $C$ . If  $C$  does not have an overloaded operator  $\sim$  or an implicit conversion to an integral or enumerated type, the expression  $\sim C( )$  results in an error. To force a reference to  $C$ 's destructor, use a member reference (e.g.,  $this \rightarrow \sim C( )$ ) or a qualified name (e.g.,  $C::\sim C( )$ ).

`! expr, not expr`

Returns the logical negation of *expr* after converting it to `bool`. The result is an rvalue of type `bool`. If *expr* is true, the result is false; if *expr* is false, the result is true.

`sizeof expr, sizeof ( type )`

Returns the size in bytes of *type* or the type of *expr* (without evaluating *expr*). By definition, `sizeof(char) == 1`. You cannot take the size of a bit-field, a function, or an incomplete type. The size of a reference is the size of the referenced type.

The `sizeof` operator always returns a value greater than zero for a class or object of class type. The size of a base-class subobject within a derived-class object can be zero, so the compiler is not necessarily wasting space. You can see this in [Example 3-9](#), which shows that the size of the derived class is the same as the size of the base class. The expression result is an rvalue of type `size_t`. (See `<cstdlib>` in [Chapter 13](#).)

*Example 3-9. Using the sizeof operator*

```
#include <iostream>
#include <ostream>

class base {};
class derived : public base {};

int main( )
{
    // The values actually printed depend on the implementation,
    // but many
    // common implementations print the values shown.
    using namespace std;
    cout << sizeof(base) << '\n';           // Prints 1
    cout << sizeof(derived) << '\n';       // Prints 1
    base b[3];
    cout << sizeof(b) << '\n';             // Prints 3
    derived d[5];
```

```
    cout << sizeof(d) << '\n';           // Prints 5
}
```

*new type, new type ( optional-expr-list ), new ( expr-list ) type, new ( expr-list ) type ( optional-expr-list )*

Allocates and initializes a dynamic object or array of objects. The new expression first calls an allocation function (operator new) to allocate memory. It then constructs the object in the allocated memory. A class can provide its own allocation function by overriding operator new as a member function. Otherwise, a global operator new is called. (See <new> in [Chapter 13](#) for the standard allocation functions.)

The parts of a new expression are:

*new, ::new*

The new keyword can be prefixed with the global scope operator to call the global operator new as the allocation function.

*( expr-list )*

An expression list in parentheses is called the *placement*. The placement is optional, but if it is used, it must contain at least one expression. If present, the expressions are passed directly to the allocation function without further interpretation.

The standard library defines two placement new functions, which are discussed in [Chapter 13](#). You can also write your own overloaded operator new functions for other forms of placement new. The first operand to operator new is always the amount of memory to allocate followed by the placement parameters.

*type*

The type to allocate. It has the following form (optionally in parentheses):

*type-specifiers*

*ptr-operators*  
*dimensions*

(See [Chapter 2](#) for information about type specifiers.) The *ptr-operators* are optional and can be \* or & for pointers or references. The array *dimensions* are optional. All dimensions of an array except the first must be constant integral expressions (enclosed in square brackets). The first dimension can be any integral expression.

The compiler reads the longest sequence of declarators as the *type*, even if it results in a syntax error. If *type* contains parentheses (e.g., a function pointer) or if you want to force a particular type declaration, surround *type* with parentheses. For example, (new int[ *n* ])[2] allocates an array of *n* integers and extracts the element at index 2. Without the parentheses, new int[ *n* ][2] allocates a two-dimensional array of int.

( *optional-expr-list* )

An optional initializer that follows the usual rules for initializers. (See [Chapter 2](#) for more information about initializers.)

If the expression list has a single expression, and a single object is allocated, the expression is the object's initial value.

If multiple expressions are in the initializer, the type must be a class type, and the expression list is passed to a suitable constructor, which is found by the usual rules for resolving overloaded functions ([Chapter 5](#)).

An array cannot have an initializer. If the base type is a POD type, the array elements are uninitialized; otherwise, they are initialized by calling the default constructor for each array element.

See [Chapter 6](#) for a comparison of POD and non-POD types.

The allocation function (operator `new`) is called with at least one argument: the number of bytes to allocate (of type `size_t`). If a placement is used, the placement arguments are passed as additional arguments to the allocation function to the right of the size. If the allocation function cannot fulfill the request, it typically throws `std::bad_alloc`. However, if you pass `std::nothrow` as the placement argument, the standard allocation function returns a null pointer as an error indicator instead of throwing `bad_alloc`.



Allocating an array is different from allocating a scalar. The allocation function is `operator new[]`. The requested size is the number of elements in the array times the size of each element. An implementation is free to request additional memory for its own use, perhaps to store the number of elements in the array. The amount of additional memory is implementation-defined. Even if the array size is zero, the returned pointer is not null. The allocated memory is aligned to the most restrictive boundary for any type. (More precisely, the allocation function must return a pointer that is aligned for any type, and the `new` expression simply uses that pointer.) Thus, you can, for example, allocate an array of `char` and use the memory to store any object. The standard containers often do this. See `<memory>` in [Chapter 13](#) for algorithms that work with uninitialized memory.

If an exception is thrown during initialization, the memory is freed by calling a corresponding deallocation function (corresponding to the equivalent delete expression). If placement new is used, a placement delete operator with the same additional parameters is the deallocation function, if such a function exists; otherwise, no deallocation function is called.

The following are some examples of new expressions:

```
int n = 10;                // Note that n is not const
new int                    // Pointer to uninitialized int
new int( )                 // Pointer to int,
                           // initialized to 0
new int[n]                 // n uninitialized ints
new (int*)                 // Pointer to uninitialized
                           // pointer to int
new (int (*[n]))( )       // n function
                           // pointers
    typedef int (*int_func)( );
new int_func[n];           // n function pointers
new (int*[n][4])           // n*4 array of pointers to int
new complex<int>(42)       // Pointer to a complex object
new complex<int>[5]        // Five default-initialized
                           // complex objects
```

*delete pointer, delete[] pointer*

Destroys and frees a dynamic object or array of objects and returns void. The actual memory deallocation is performed by a deallocation function. A class can provide its own deallocation function by overriding operator delete. A plain delete expression looks up the deallocation function first in the class (if the *pointer* type is a pointer to a class type) and, if it is not found, in the global scope. Use the global scope operator to look only in the global scope. (See <new> in [Chapter 13](#) for the default deallocation functions.)

To free an array, you must use delete[]. To free a scalar, you must use delete. If you make a mistake, the results are undefined. Note that the compiler cannot generally



help you avoid mistakes because a pointer to a scalar cannot be distinguished from a pointer to an array. (Some libraries are more forgiving of this error than others.)

The *pointer* expression is evaluated once. If *pointer* is a pointer to a class type, the scalar form calls the object's destructor first, and the array form calls the destructor for each element of the array. The value of *pointer* is then passed to the deallocation function as a `void*`. If the expression's static type does not match the object's dynamic type, the static class must have a virtual destructor, or else the behavior is undefined. See [Chapter 6](#) for more information.

You can delete a pointer to a `const` object. It is also safe to delete a null pointer value, in which case the deallocation function does nothing.

## Cast Expressions

A cast expression performs an explicit type conversion. The cast expression is a holdover from C. In C++, the preferred cast syntax uses one of the explicit cast operators (described in [Section 3.5.2](#) earlier in this chapter). The C-style casts are still used for their brevity, however.

( *type* ) *expr*

The C-style cast converts *expr* to *type* using one or more template-like type conversions. If *type* is a reference, the result is an lvalue; otherwise the result is an rvalue. The following type conversions are tried in order. The first one that is syntactically allowed is the one used, even if the expression is not permitted semantically.

1. `const_cast< type >( expr )`
2. `static_cast< type >( expr )`

3. `const_cast< type >( static_cast< type1 >( expr ))`
4. `reinterpret_cast< type >( expr )`
5. `const_cast< type >( reinterpret_cast< type1 >( expr ))`

The type *type1* is the same as *type*, but its *cv*-qualifiers are changed to match the type of *expr*. Thus, the C-style type cast can mix a `const` cast with a `static` or `reinterpret` cast. A C-style cast can also cast to an otherwise inaccessible base class (see [Chapter 6](#) for information about accessibility). That is, you can cast a derived class to an inaccessible base class, cast a pointer to a member of a derived class to a pointer to a member of an inaccessible base class, or cast from an inaccessible base class to an accessible derived class.

## Pointer-to-Member Expressions

A pointer-to-member expression takes an object or a pointer to an object as the lefthand operand and a pointer-to-member as the righthand operand, and it binds the pointer-to-member to the object. The result can be a data member or a member function. A member function can be used only to call the function. [Example 3-8](#) shows some uses of pointers to members. The pointer-to-member operator has the following syntax:

*object* .\* *expr*

Binds *expr* to *object*, in which *expr* is a pointer-to-member of class *C*, and the type of *object* is *C* or a class derived from *C*. The result is an lvalue if *object* is an lvalue and *expr* points to a data member; otherwise, the result is an rvalue. The type of the result is determined

by the type of *expr*. The behavior is undefined if *expr* is a null pointer-to-member.

*pointer* ->\* *expr*

Binds *expr* to the object that *pointer* points to, in which *expr* is a pointer-to-member of class *C*, and the type of *object* is *C* or a class derived from *C*. The result is an lvalue if *expr* points to a data member. The type of the result is determined by the type of *expr*. The behavior is undefined if *pointer* is null or if *expr* is a null pointer-to-member.

If *expr* points to a virtual function, the usual rules apply. That is, the actual function called is that of the most-derived type of \**pointer* or object. See [Chapter 6](#) for more information about virtual functions.

## Multiplicative Expressions



A multiplicative expression is used for multiplication, division, and remainders. The multiplicative operators require arithmetic or enumeration types; the usual conversions are performed, and an rvalue is returned. If the result is too large, the behavior is undefined (except for unsigned types, for which arithmetic is performed modulo the integer size; see [Chapter 1](#) for details). Many C++ implementations ignore integer overflow. The multiplicative operators have the following syntax:

*expr1* \* *expr2*

Performs multiplication.

*expr1* / *expr2*

Performs division. If the divisor is zero, the behavior is undefined.



*expr1 % expr2*

Returns the remainder of dividing *expr1* by *expr2*. The operands must have integral or enumerated types. If *expr2* is 0, the behavior is undefined; otherwise, the value is such that  $(a/b)*b + a\%b == a$ . If both operands are nonnegative, the result is nonnegative; otherwise, the sign of the result is implementation-defined.

## Additive Expressions



An additive expression is used for addition and subtraction. The additive operators require arithmetic, enumerated, or pointer types. The usual conversions are performed, and an rvalue is returned. If the result of an additive expression is too large, the behavior is undefined (except for unsigned types, for which arithmetic is performed modulo the integer size; see [Chapter 1](#) for details). Many C++ implementations ignore integer overflow. The additive operators have the following syntax:

*expr + expr*

Performs addition. If one operand has a pointer type, the other must have an integral or enumeration type. The result is a pointer to the same array, but with the index offset by *N* positions (*N* can be positive, negative, or 0), in which *N* is the integral operand. The resulting pointer must be within the bounds of the array or point to one past the end of the array; otherwise, the behavior is undefined. Note that a pointer to any object can be treated as a one-element array.

*expr - expr*

Performs subtraction. If both operands have arithmetic or enumeration types, the usual promotions apply and the result is the difference of the operands.

If both operands are pointers, they must point to elements of the same array, or to one element past the end of the array. The result has type `ptrdiff_t` (declared in `<cstdlib>`) and is equal to the difference of the indices of the two objects.

If the left operand is a pointer and the right operand has integral or enumeration type, the result is the same as *expr1 - expr2*.

## Shift Expressions

A shift expression shifts the bits of the left operand by an amount specified by the right operand. The operands must be have integral or enumerated types; both types are promoted to integral types. The result type is the promoted type of the left operand.



The result of a shift operation is undefined if the right operand is negative or is larger than the number of bits in the left operand.

The shift operators have the following syntax:

*expr1 << expr2*

Performs a left shift of *expr1* by *expr2* bits. Vacated bits are filled with zeros. If *expr1* is unsigned, the result is equal to multiplying *expr1* by 2 raised to *expr2* (modulo the integer size; see [Chapter 1](#) for more information about unsigned integer arithmetic).



*expr1* >> *expr2*

Performs a right shift of *expr1* by *expr2* bits. If the *expr1* is unsigned, or if it is signed and has a positive value, vacated bits are filled with zeros. The result is equal to dividing *expr1* by 2 raised to *expr2* (truncated to an integer). If *expr1* has a signed type and negative value, the result is implementation-defined.

The standard library overloads the shift operators for the I/O stream class templates. As with any overloaded function, the syntax (precedence and associativity) remains the same. Only the runtime behavior is different. See [Chapter 9](#) for more information.

## Relational Expressions

A relational expression compares two values for relative order. Relational operators have higher precedence than equality operators, so the following two expressions are equivalent:

```
a < b == c > d
(a < b) == (c > d)
```

The result of a relational expression is an rvalue of type `bool`. The operands must have arithmetic, enumeration, or pointer types. For arithmetic and enumeration types, the usual conversions are performed, and the resulting values are compared.

When comparing pointers, the pointers must have the same types (after the usual conversions and ignoring *cv*-qualification), one must be a pointer to `void`, or one operand must be a null pointer constant. If the pointer types are compatible, they are compared as follows (*L* is the left operand and *R* is the right operand):

- If  $L$  and  $R$  point to the same function, the same object, or one element past the end of the same array, or if both are null pointers, they are considered to be equal to each other. That is,  $L \leq R$  and  $L \geq R$  are true, and  $L < R$  and  $L > R$  are false.



If  $L$  and  $R$  point to different objects (that are not members of a common object and not elements of the same array), different functions, or if only one is a null pointer, the result of the relational operators depends on the implementation.

- If  $L$  and  $R$  point to data members (as ordinary pointers, not pointers-to-members) within the same object (members of the same object, elements of arrays that are data members, and so on, recursively applied), and if the members are not separated by an access specifier label, and if the object is not a union, then  $L > R$  if the member to which  $L$  points is declared later than the member to which  $R$  points. Similarly,  $L < R$  if the member to which  $L$  points is declared earlier than the member to which  $R$  points. If the members are separated by an access specifier label, the results are implementation-dependent.
- If  $L$  and  $R$  point to data members (as ordinary pointers, not pointers-to-members) of the same union, they are considered to be equal to each other.
- If  $L$  and  $R$  point to elements of the same array, including one element past the end of the same array, the pointer with the higher subscript is larger.

[Example 3-10](#) shows some pointer comparisons.

*Example 3-10. Comparing pointers*

```

#include <iostream>
#include <ostream>

struct Demo {
    int x;
    int y;
};

union U {
    int a;
    double b;
    char c[5];
    Demo d;
};

int main( )
{
    Demo demo[10];
    std::cout << std::boolalpha;
    // Everything prints "true".
    std::cout << (&demo[0] < &demo[2]) << '\n';
    std::cout << (&demo[0] == demo) << '\n';
    std::cout << (&demo[10] > &demo[9]) << '\n';
    std::cout << (&demo[0].x < &demo[0].y) << '\n';

    U u;
    std::cout << (&u.d == static_cast<void*>(u.c)) << '\n';
    std::cout << (&u.a == static_cast<void*>(&u.b)) << '\n';
}

```

The relational operators have the following syntax:

*expr1* < *expr2*

Returns true if *expr1* is less than *expr2*

*expr1* > *expr2*

Returns true if *expr1* is greater than *expr2*

*expr1* <= *expr2*

Returns true if *expr1* is less than or equal to *expr2*

*expr1* >= *expr2*

Returns true if *expr1* is greater than or equal to *expr2*



## Equality Expressions

An equality expression compares two values to see if they are equal or different. Equality operators have lower precedence than relational operators, so the following two expressions are equivalent:

```
a < b == c > d
(a < b) == (c > d)
```

The result of an equality expression is an rvalue of type `bool`. The operands must have arithmetic, enumeration, or pointer types. For arithmetic and enumeration types, the usual conversions are performed, and the resulting values are compared.

### Warning

Note that comparing the results of floating-point values for equality rarely gives the result you want. Instead, you probably want a fuzzy comparison that allows for floating-point imprecision.

When comparing pointers, the pointers must have the same type (after the usual conversions). The pointers are equal if any of the following conditions hold, and are not equal if none of the conditions hold:

- Both pointers are null pointers.
- Both object pointers point to the same object.
- Both object pointers point to one element past the end of the same array.
- Both function pointers point to the same function.
- Both member pointers point to the same member of the same most-derived object.

- Both member pointers point to any data members of the same union.

The equality operators have the following syntax:

*expr == expr*

Returns true if the operands are equal.

*expr != expr, expr not\_eq expr*

Returns false if the operands are equal.

## Bitwise And Expressions

A bitwise and expression performs and on its operands' bits. The bitwise and expression is permitted only on integral types after the usual arithmetic conversions. The bitwise and operator has the following syntax:

*expr & expr, expr bitand expr*

Performs bitwise and of the operands. Each bit of the result is 1 if the corresponding bits of the operands are both 1; otherwise, the result bit is 0.

## Bitwise Exclusive Or Expressions

A bitwise exclusive or expression performs exclusive or on its operands' bits. The bitwise exclusive or expression is permitted only on integral types after the usual arithmetic conversions. The exclusive or operator has the following syntax:

*expr ^ expr, expr xor expr*

Performs bitwise exclusive or of the operands. Each bit of the result is 1 if the corresponding bits of the operands are not equal; otherwise, the result bit is 0.

## Bitwise Inclusive Or Expressions

A bitwise inclusive or expression performs inclusive or on its operands' bits. The bitwise inclusive or expression is permitted only on integral types after the usual arithmetic conversions. The bitwise inclusive or operator has the following syntax:

*expr | expr, expr bitor expr*

Performs bitwise inclusive or of the operands. Each bit of the result is 0 if the corresponding bits of the operands are both 0; otherwise, the result bit is 1.

## Logical And Expressions

A logical and expression implicitly converts its operands to type `bool`. The result has type `bool`: `true` if both operands are `true`; otherwise, it is `false`. The logical and operator is a short-circuit operator, so the second operand is evaluated only if the first evaluates to `true`. The logical and operator has the following syntax:

*expr && expr, expr and expr*

Performs the logical and of the operands.

## Logical Or Expressions

A logical or expression implicitly converts its operands to type `bool`. The result has type `bool`: `false` if both operands are `false`; otherwise, it is `true`. The logical or operator is a short-circuit operator, so the second operand is evaluated only if the first evaluates to `false`. The logical or operator has the following syntax:

*expr || expr, expr or expr*

Performs the logical or of the operands.

## Conditional Expressions

A conditional expression is like an if statement in an expression:

*condition ? true-expr : false-expr*

The first operand is converted to `bool`. If the value is true, only the second operand is evaluated; if it is false, only the third operand is evaluated. The result of the conditional expression is the result of the second or third operand, whichever is evaluated.

The type of the expression depends on the types of the second and third operands: if both operands are lvalues of the same type, the result is an lvalue of that type. If one operand is an lvalue of type  $T$ , and the other operand is an lvalue that can be implicitly converted to a reference to  $T$ , the result is an lvalue of type  $T$ . An error results if the conversion is ambiguous—that is, the second operand can be converted to the type of the third just as readily as the third can be converted to the type of the second. Otherwise, the result is an rvalue, and the type is determined as follows:

- If both operands have the same type, that is the type of the result.
- If one operand is a throw expression, the result type is that of the other operand.
- If one type  $S$  can be implicitly converted to the other type  $T$ ,  $T$  is the result type. An error results if each operand can be implicitly converted to the other type.

## Assignment Expressions

An assignment expression assigns its right operand to its left operand. In addition to plain assignment ( $x = y$ ), there are several other assignment operators, each of which is a shorthand for an arithmetic operation and an assignment. The left operand of an assignment expression must be a modifiable lvalue. The result is an lvalue: the left operand after the assignment is complete.

Assignment operators have the following syntax:

*lvalue = expr*

Assigns *expr* to *lvalue*. If the left operand has class type, a suitable assignment operator is called. (See [Chapter 6](#) for more information.) If the left operand is not of class type, and the operands have different types, the right operand is converted to the type of the left operand.

*lvalue += expr, lvalue -= expr*

The assignment operator  $x \text{ op} = y$  is shorthand for  $x = x \text{ op} y$ , except that  $x$  is evaluated only once. The type of  $x$  must be an arithmetic type or a pointer type. The usual conversions are performed for *op*, and the result is converted to the type of  $x$ .

*lvalue \*= expr, lvalue /= expr, lvalue %= expr, lvalue <<= expr, lvalue >>= expr, lvalue &= expr, lvalue and\_eq expr, lvalue ^= expr, lvalue xor\_eq expr, lvalue |= expr, lvalue or\_eq expr*

The assignment operator  $x \text{ op} = y$  is shorthand for  $x = x \text{ op} y$ , except that  $x$  is evaluated only once. The type of  $x$  must be arithmetic (unlike the additive assignment operators, which allow pointers because the additive operators allow pointer operands). The usual conversions are performed for *op*, and the result is converted to the type of  $x$ .

`throw expr, throw`

Throws an exception. The first form throws *expr* as the exception object. The second form rethrows the current exception object. If a program uses the second form, and no exception is being handled, `throw` calls `terminate()`. (See `<exception>` in [Chapter 13](#) for more information about `terminate`.)

You can throw any expression, but the convention is to throw an object of type `exception` or of a derived class, especially a class derived from one of the standard exception classes (see `<stdexcept>` in [Chapter 13](#)).



The `throw` expression initializes a temporary exception object, and that object is thrown to the exception handler, which can copy the exception object in its `catch` clause. When the handler finishes, the object is destroyed. An implementation can optimize away the extra copy and initialize the `catch` object directly with the exception *expr*. Even if the object is never copied, the class must have a copy constructor.

See the `try` statement in [Chapter 4](#) for a description of how exceptions are handled and for examples of `throw` expressions. See also [Chapter 5](#) for information about `throw` specifications in function declarations.

## Comma Expressions

A comma expression serializes expression evaluation. The comma operator evaluates its left operand, discards the result, and then evaluates the right operand. The result is the value of the right operand, and the type is the type of the right operand. If the right operand is an lvalue, the result is an lvalue; otherwise, the result is an rvalue.

Typically, the left operand is evaluated for a side effect, such as an assignment or a function call.

The comma operator can be overloaded, in which case the serial nature is lost. Both operands are evaluated, and the operator function is called. Overloading the comma operator is seldom done.

The syntax of the comma operator is:

*expr1* , *expr2*

Evaluates *expr1*, then *expr2*, and returns *expr2*. To use a comma expression as a function argument or in other contexts in which a comma is significant, enclose the comma expression in parentheses:

```
sin((angle = 0.0, angle + 1.0));
```

## Chapter 4. Statements

Statements define and control what a program does. This chapter describes the syntax and rules for C++ statements: expressions, loops, selection, and control. The statement syntax rules apply recursively, and wherever a statement is called for, you can use (almost) any of the statements in this chapter.

The syntax descriptions in this chapter are informal. See [Chapter 12](#) for a precise BNF grammar.

### Expression Statements

An expression statement computes an expression, such as a function call or assignment. The expression result is discarded, so the expression is typically evaluated for its side effects. (See [Chapter 3](#) for details about expressions.) The statement syntax is simply an optional expression followed by a semicolon:

```
    expr ;
```

or:

```
    ;
```

A statement with no expression is called a *null statement*. Null statements are most often used for loops when no code is needed in the loop body.

Here are several examples of expression statements:

```
42;           // Valid but pointless
cout << 42;   // More typical
x = y * z;    // Remember that assignment is an expression
;            // Null statement
```



## Declarations

A declaration can appear anywhere a statement appears, and certain statements permit additional declarations within those statements.

Declarations made in a substatement (of a selection or loop statement) are limited in scope to the substatement, even if the substatement is not a compound statement. For example, the following statement:

```
while ( test( ) )  
    int x = init( );
```

is equivalent to:

```
while ( test( ) ) {  
    int x = init( );  
}
```

The first example uses a declaration as the entire loop body, and the second uses a compound statement (enclosing the loop body in curly braces). In both cases, though, the scope of `x` is limited to the body of the `while` loop.

## Declaration Statements

A simple declaration can appear wherever a statement can be used. You can declare an object, a type, or a namespace alias. You can also write a `using` declaration or `using` directive. You can declare a function, but not define a function, although there is rarely any reason to declare a function locally. You cannot define a namespace or declare a template.

In traditional C programming, declarations appear at the start of each block or compound statement. In C++ (and in the C99 standard), declarations can appear anywhere a statement can, which means you can declare variables

close to where they are used. [Example 4-1](#) shows examples of how declarations can be mixed with statements.

### *Example 4-1. Mixing declarations and statements*

```
#include <cctype>
#include <cstddef>
#include <iomanip>
#include <iostream>
#include <ostream>
#include <string>

// Count lines, words, and characters in the standard input.
int main( )
{ unsigned long num_lines, num_words, num_chars;

  num_lines = num_words = num_chars = 0;

  using namespace std;
  string line;
  while (getline(cin, line)) {
    ++num_lines;
    num_chars += line.length( ) + 1;

    bool in_word = false;
    for (size_t i = 0; char c = line[i]; ++i)
      if (isspace(static_cast<unsigned char>(c))) {
        if (in_word)
          ++num_words;
        in_word = false;
      } else if (! in_word)
        in_word = true;
    if (in_word)
      ++num_words;
  }
  cout << right <<
    setw(10) << num_lines <<
    setw(10) << num_words <<
    setw(10) << num_chars << '\n';
}
```

Sometimes a construct can look like an expression, statement or a declaration. These ambiguities are resolved in favor of declarations. [Example 4-2](#) shows some declarations that look like they might be expressions.

### *Example 4-2. Declarations that seem like expressions*

```

class cls {
public:
    cls( );
    cls(int x);
};

int x;
int* y = &x;

int main( )
{
    // The following are unambiguously expressions, constructing
    instances of cls.  cls(int(x));
    cls(*static_cast<int*>(y));

    // Without the redundant casts, though, they would look like
    declarations, not
    // expressions.
    cls(x);    // Declares a variable x
    cls(*y);   // Declares a pointer y
}

```

## Condition Declarations

The for, if, switch, and while statements permit a declaration within each statement's condition:

```
if (int x = test_this(a, b)) { cout << x; }
```

If the condition contains a declaration, the scope of the declared name extends to the end of the entire statement. In particular, a name declared in the condition of an if statement is visible in the else part of the statement. In a loop, the condition object is created and destroyed for each iteration of the loop. The name cannot be redeclared in the immediate substatement (but can be redeclared in a nested statement). The name is not visible in subsequent statements. For example:

```

if (derived* d = dynamic_cast<derived*>(b)) {
    d->derived_only_func( );
} else {
    assert(d == NULL); // Same d as above
    double d;         // Error: can't redeclare d
    if (d == 0)

```

```
    int d;           // Valid: inner block
}
cout << d;         // Invalid: d no longer in scope
```

Like the `if`, `switch`, and `while` statements, a `for` loop permits a declaration in its condition. Unlike those other statements, it also allows a declaration to initialize the loop. Both declarations are in the same scope. See [Section 4.5.2](#) for details.

The syntax for a condition declaration is:

```
type-specifiers
    declarator = expression
```

This syntax is similar to the syntax of a simple object declaration. In this case, the initializer is required (without it, the condition would not have a value), and only one declarator is permitted. See [Chapter 2](#) for more information about type specifiers and declarators.

## Compound Statements

A compound statement is a sequence of zero or more statements enclosed within curly braces. Compound statements are frequently used in selection and loop statements. They enable you to write loop bodies that are more than one statement long, among other things. A compound statement is sometimes called a *block*.

Here is the syntax of a compound statement:

```
{statement ... }
```

or:

```
{ }
```

A compound statement can be used as a single statement wherever a statement is called for. A compound statement delimits a declarative scope, that is, any name declared in the compound statement is not visible outside the statement, and names in the statement can hide names from outside the statement. The lifetime of any automatic object declared in a compound statement is confined to that statement. All such objects are destroyed when execution leaves the compound statement for any reason (e.g., branching outside the statement, execution reaches the end of the statement, or an exception is thrown).

Compound statements are most often used as the bodies of selection and loop statements, but you can also stick a compound statement in the middle of a compound statement to restrict the scope of variables that are local to the compound statement. See the examples in this chapter for uses of compound statements.

# Selections

A selection statement chooses one branch from multiple possibilities and executes the statements in that branch. The two selection statements supported by C++ are `if` and `switch`.

## if Statements

An `if` statement has one of the following forms:

```
if ( condition ) statement
```

or:

```
if ( condition ) statement else statement
```

In the first form, the *condition* is converted to a `bool`, and if it is true, the *statement* is executed. Otherwise, the statement is skipped, and execution continues with the subsequent statement.

The second form chooses one of two alternative code paths. The *condition* is converted to a `bool`, and if it is true, the first *statement* is executed. Otherwise, the second *statement* is executed. Declarations in the first statement are not visible in the second. If `if` statements are nested, the `else` part binds with the closest, preceding `if` statement. [Example 4-3](#) shows nested `if` statements.

### *Example 4-3. Nested if statements*

```
if (c1)
    if (c2)
        cout << "c1 and c2 are true\n";
    else
        cout << "c1 is true, but c2 is false\n";
else if (c2)
    cout << "c1 is false, but c2 is true\n";
else
    cout << "c1 and c2 are false\n";
```

## switch Statements

A switch statement chooses one execution path from among many alternatives. The syntax is:

```
switch (condition ) statement
```

The *condition* must have an integral or enumerated type, or be of a class type in which the class has a single conversion function to an integral or enumerated type. The *condition* is evaluated once. Its value is compared against the case labels in the *statement*. If a case label matches the *condition*, execution continues with the statement immediately after the case label. If no case matches the *condition*, execution continues after the default label, if one is present. If there is no default label, the switch's *statement* is skipped and execution continues with the subsequent statement.

The *statement* part of a switch statement is typically a compound statement, in which every substatement has one or more case labels or a default label. The syntax for case and default labels is:

```
case constant-expression : statement  
default : statement
```

The *constant-expression* must have an integral or enumerated type. The value is implicitly converted to the type of the *condition*. In a single switch statement, all case *constant-expressions* must have different values. A single statement in the switch's substatement can have multiple case labels, and a single switch statement can have any number of cases.

## Tip

In C++, like C and Java, but unlike most other languages, a case or default label does not affect control flow. Execution continues from one case to the next, which is known as "falling through" to the next case. Use the break statement (described later in this chapter) to exit from a switch statement.

There can be at most one default case in a switch statement; it can appear anywhere in the statement. (The default case does not have to be last, as in some languages.)

By convention, case and default labels appear at the top level of the switch's substatement. They can appear in nested statements, but that makes the statements hard to read. In nested switch statements, case and default labels apply only to the innermost switch statement.

You should not define any objects in the switch's substatement unless they are enclosed in a nested compound statement. When the switch dispatches control to a particular case, it might jump over the declaration. Jumping over a declaration results in undefined behavior unless the declared object has POD type and is not initialized.

[Example 4-4](#) shows sample switch statements.

### *Example 4-4. Switch statements*

```
enum color { black, red, green, yellow, blue,
            magenta, cyan, white };
color get_pixel(unsigned r, unsigned c) { ... }

void demo( )
{
    using std::cout;
    int r = ...
    int c = ...  switch (get_pixel(r, c))
```



```

{
    cout << "this is never executed, but it is valid\n";
    case black:
        cout << "no color\n";
        break; // Don't forget the break statements!
    case red: case green: case blue:
        cout << "primary\n";
        break; // Omitting break is a common mistake.
    default:
        cout << "mixed\n";
        switch (get_pixel(r+1, c+1))
        case white: {
            const int x = 0;
            cout << " white\n"; // This case is private to the inner
switch
                                // statement.

            c = x;
        }
        if (r > 0)
            // If the color is yellow, the switch branches directly
to here. For
            // colors other than red, green, blue, and black,
execution jumps to the
            // default label and arrives here if r > 0.
            case yellow:
                cout << " yellow or r > 0\n";

                break; // A break after the last case is not necessary, but
is a good idea
                    // if you add a case later.
        }
}
}

```

The outer switch statement has one case for black; another case for red, green, or blue; a default case; and a final case for yellow. The first two branches are conventional, and each ends with a break statement to exit the switch.

The case for yellow is unusual and hard to read because it is buried in an if statement. The same statement can be reached from the default case when the if condition is true (that is,  $r > 0$ ).

The default case has a nested switch, which has a single case, white. The inner switch statement is atypical because it does not have a compound statement as its body.

# Loops

A loop statement allows you to execute a statement multiple times. Two loop statements, `for` and `while`, test at the top of the loop. The `do` statement tests at the bottom, thereby ensuring that the loop body executes at least once. This section describes the loop statements. See [Section 4.6](#) for additional statements that affect or control loop execution.

A loop statement can declare variables in the scope of the loop's substatement. Every time the loop iterates, it reenters the substatement scope. This means objects that are declared in the substatement (and in the loop's condition) are created and destroyed with every loop iteration.

## while Statements

A `while` statement repeats a statement while a condition is true. The syntax is:

```
while (condition ) statement
```

The *condition* is evaluated and converted to `bool`. If the value is `true`, *statement* is executed, and the loop repeats. If the value is `false`, the loop finishes and execution continues with the subsequent statement. Thus, if *condition* is `false` the first time it is evaluated, the *statement* is never executed.

A `continue` statement in a `while` loop branches to the top of the loop, and execution continues with the evaluation of *condition*. A `break` statement exits immediately.

A `while` loop is typically used for unbounded iteration, that is, when you don't know beforehand how many times a loop

will iterate. [Example 4-5](#) shows how the `while` loop can be used to control I/O.

### *Example 4-5. Controlling I/O with the while loop*

```
#include <algorithm>
#include <iostream>
#include <iterator>
#include <ostream>
#include <string>
#include <vector>

// Sort lines of text.
int main( )
{
    using namespace std;
    string line;
    vector<string> data; while (getline(cin, line))
        data.push_back(line);
    sort(data.begin( ), data.end( ));
    copy(data.begin( ), data.end( ),
        ostream_iterator<string>(cout, "\n"));
}
```

## for Statements

A `for` loop is a generalization of the traditional counted loop that appears in most programming languages. The syntax for a `for` loop is:

```
for (init ; condition ; iterate-expr ) statement
```

The *init*, *condition*, and *iterate-expr* parts are optional. The *init* part of the `for` statement can be an expression or a simple declaration. The *init* part offers more flexibility than a condition. While a condition can declare only one name, the *init* part can declare multiple names. The syntax for the *init* part is:

```
specifier-list declarator-list
```

or:

```
expression
```

As with the *condition*, the scope of the *init* part extends to the end of *statement*. The *init*, *condition*, and *iterate-expr* parts are in the same scope as the loop body. See [Chapter 2](#) for more information about specifiers and declarators.

The for loop starts by executing the *init* part, if present. It then evaluates the *condition* (just like a while loop). If the *condition* is true, the loop executes *statement*. It then evaluates *iterate-expr* and reevaluates *condition*. This process continues while *condition* is true. If *condition* is false, the loop finishes and execution continues with the subsequent statement. Thus, the *init* part is evaluated exactly once. The *condition* is evaluated at least once. The *statement* might be executed zero times.

The most common use of a for loop is to count a bounded loop, although its flexibility makes it useful for unbounded loops, too, as you can see in [Example 4-6](#).

#### *Example 4-6. Multiple uses of for loops*

```
// One way to implement the for_each standard algorithm
template<typename InIter, typename Function>
Function for_each(InIter first, InIter last, Function f)
{   for ( ; first != last; ++first)
    f(*first);
    return f;
}

// One way to implement the generate_n standard algorithm
template<typename OutIter, typename Size, typename Generator>
void generate_n(OutIter first, Size n, Generator gen)
{
    for (Size i = 0; i < n; ++i, ++first)
        *first = gen;
}
```

A continue statement in a for loop branches to the top of the loop, and execution continues by evaluating the *iterate-expr* and then the *condition*. A break statement exits the loop without evaluating the *iterate-expr*. See

[Section 4.6](#) later in this chapter for examples of break and continue.

The *init*, *condition*, and *interate-expr* parts are all optional. If the *init* or *interate-expr* parts are omitted, nothing happens to initialize the loop or after the statement executes. If the *condition* is omitted, it defaults to true.

## do Statements

The do statement is like a while statement, except that it tests at the end of the loop body. The syntax is:

```
dostatement while ( expression ) ;
```

The *statement* is executed, then the *expression* is evaluated and converted to bool. If the value is true, the *statement* is repeated and the *expression* is checked again. If the *expression* is false, the loop finishes and execution continues with the subsequent statement. Thus, *statement* is always executed at least once.

A continue statement in a do loop jumps to the end of *statement*, and execution continues with the evaluation and test of *expression*. A break statement exits immediately.

# Control Statements

Control statements change execution from its normal sequence. When execution leaves a scope, all automatic objects that were created in that scope are destroyed. (see [Chapter 2](#) for a discussion of automatic and other object lifetimes)

C++ supports the following control statements:

**break;**

A break statement can be used only in the body of a loop or switch statement. It terminates the loop or switch statement and transfers execution to the statement immediately following the loop or switch.

In a nested loop or switch, the break applies only to the innermost statement. To break out of multiple loops and switches, you must use a goto statement or redesign the block to avoid nested loops and switches (by factoring the inner statement into a separate function, for example). [Example 4-7](#) shows a simple use of break.

*Example 4-7. Using break to exit a loop*

```
// One way to implement the find_if standard algorithm.
template<typename InIter, typename Predicate>
InIter find_if(InIter first, InIter last, Predicate pred)
{
    for ( ; first != last; ++first)
        if (pred(*first))      break;

    return first;
}
```

**continue;**

A continue statement can be used only in the body of a loop. It causes the loop to skip the remainder of its body and immediately retest its condition prior to reiterating

(if the condition is true). In a for loop, the *iterate-expr* is evaluated before testing the condition. [Example 4-8](#) shows how `continue` is used in a loop.

#### *Example 4-8. Using `continue` in a loop*

```
#include <cmath>
#include <iostream>
#include <istream>
#include <limits>
#include <ostream>

int main( )
{
    using std::cin;
    using std::cout;
    while(true) {
        cout << "Enter a number: ";
        double x;
        cin >> x;
        if (cin.eof( ) || cin.bad( ))
            // Input error: exit
            break;
        else if (cin.fail( )) {
            // Invalid input: skip the rest of the line
            cin.clear( );
            cin.ignore(std::numeric_limits<int>::max( ), '\n');
continue;
        }
        cout << "sqrt(" << x << ")=" << std::sqrt(x) << std::endl;
    }
}
```

`goto identifier ;`

The `goto` statement transfers control to the statement that has *identifier* as a label. The `goto` statement and the labeled statement must be in the same function. Jumping into a block is usually a bad idea. In particular, if the jump bypasses the declaration of an object, the results are undefined unless the object has POD type and no initializer. (See [Chapter 2](#) for information about POD types and initializers.) [Example 4-9](#) shows some uses of `goto` statements.



### Example 4-9. goto statements

```
#include <iostream>
#include <ostream>

int main(int argc, char* argv[])
{
    int matrix[4][5];
    for (int i = 0; i < 4; ++i)
        for (int j = 0; j < 5; ++j)
            if (! (std::cin >> matrix[i][j]))           goto error;
                                                         goto end;

                                                         error:
    std::cerr << "Need 20 values for the matrix\n";
    end:
    return 0;
}
```

return ;, return *expr* ;

The return statement transfers execution out of a function to the caller. The first form does not return a value, so it should be used only in functions of type void, in constructors, and in destructors. The latter form cannot be used in constructors and destructors. In a function of type void, you can use the second form, but only if *expr* has type void. See [Chapter 5](#) for more information about returning from functions.



The value of *expr* is converted to the function's return type and returned to the caller. The compiler is free to construct a temporary object and copy *expr* when returning. Some compilers optimize away the extra copy.

If execution reaches the last statement of a function without executing a return statement, an implicit return; is assumed. If the function has a non-void return type, the behavior is undefined.

The main function is special. If it ends without a return statement, `return 0;` is assumed.

*identifier : statement*

Any statement can have a label. A label is used only as a target of a `goto` statement. Label identifiers must be unique within a function; the scope of a label is the function in which it is declared. Label identifiers do not conflict with any other identifiers.

A statement can have multiple labels, including case and default labels.

## Handling Exceptions

An *exception* interrupts the normal flow of control in a program. An exception *throws* an object from the point where the exception is raised to the point where the exception is handled. In between those points, function calls are aborted, and local scopes are abruptly exited. Local and temporary objects in those scopes are destroyed.

This section discusses exception handling in general and the try statement in particular. For information about other aspects of exception handling, see [Chapter 3](#), which describes the throw expression, and [Chapter 5](#), which describes function try blocks and throw specifications. See also <exception> and <stdexcept> in [Chapter 13](#) for information about the standard exception classes and related functions.

### Tip

As the name implies, an exception is used for exceptional circumstances, such as indexing a vector with an index that is out of bounds. The intention in C++ is that try statements should have near zero cost, but throwing and handling an exception can be expensive in terms of performance.

A try statement establishes a local scope plus exception handlers. A try statement begins with the try keyword followed by a compound statement and one or more catch handlers. Each catch handler has an exception type in parentheses followed by a compound statement. The exception type can be a sequence of type specifiers followed by an optional declarator, or it can be an ellipsis:

```
trycompound-statement  
catch(type
```

```
        declarator) compound-statement  
catch(type) compound-statement  
catch(...) compound-statement
```

Each compound statement (after `try` and for each exception handler) forms a separate local scope, following the rules for any other compound statement.

Every time a local scope is entered—for example, by calling a function or by entering a compound statement or other statement that establishes a scope—you can think of the scope as pushing an entry on an execution stack. When execution leaves the scope, the scope is popped from the stack. Each `try` statement is also pushed on the stack.

When an exception is thrown, the execution stack is *unwound*. Local scopes are popped from the stack (destroying any automatic objects that were declared in each scope) until a `try` statement is found. The `try` statement is popped from the exception stack. Then the type of the exception object is compared with each exception handler. If a match is found, the exception object is copied to the locally-declared exception object, and execution transfers to the compound statement for that handler. After the handler finishes, the exception object is destroyed, and execution continues with the statement that follows the `try` statement (and its handlers).

If no matching handler is found, the stack continues to be popped until the next `try` statement is reached. Its handlers are compared with the exception type, and the process repeats until a matching handler is found. If no `try` statement has a matching handler, and the entire execution stack is popped, the `terminate` function is called. (See `<exception>` in [Chapter 13](#).)

An exception handler can throw an exception, in which case, the usual rules apply for unwinding the stack. Note that the `try` statement has already been popped from the

stack, so the same `try` statement cannot catch the exception. The handler can throw a new exception or rethrow the existing exception (in which case, the exception object is not freed yet).

A catch handler can use an ellipsis to mean it catches any type of exception. When an ellipsis is used, the handler must be last in the list of exception handlers.

Handlers are checked in order. The *type* in each catch clause is compared with the type *T* of the exception object according to the following rules:

- If *type* is the same as *T* or *T* & (not considering *cv*-qualifiers), the handler matches.
- If *type* is a base class of *T*, the handler matches.
- If *type* is a pointer that can be converted to *T* using a standard conversion ([Chapter 3](#))—e.g., *type* is `void*`—the handler matches.
- An ellipsis (`...`) always matches.

Thus, it is important to put the most specific exceptions first, and put base classes later. If you use an ellipsis, it must be last.

A common idiom is to throw an exception object and catch a reference. This avoids unnecessary copies of the exception object.

[Example 4-10](#) shows a typical use of a `try` statement. It also shows a `try` function block, which is covered in [Chapter 5](#).

*Example 4-10. Throwing and catching exceptions*

```
#include <cstdlib>
#include <fstream>
#include <iostream>
#include <numeric>
#include <ostream>
```

```

#include <sstream>
#include <stdexcept>
#include <string>
#include <vector>

class bad_data : public std::out_of_range {
public:
    bad_data(int value, int min, int max)
        : std::out_of_range(make_what(value, min, max))
    {}
private:
    std::string make_what(int value, int min, int max);
};

std::string bad_data::make_what(int value, int min, int max)
{
    std::ostringstream out;
    out << "Invalid datum, " << value << ", must be in [" <<
        min << ", " << max << " ]";
    return out.str( );
}

// Read a set of numbers from an input stream. Verify that all
// data is within the
// defined boundaries. Throw bad_data if any data is invalid. If
// an exception is
// thrown, tmp's destructor is automatically called (if it has a
// destructor).
template<typename T, typename charT, typename traits>
void getdata(std::basic_istream<charT,traits>& in,
             std::vector<T>& data, T min, T max)
{
    T tmp;
    while (in >> tmp)
    {
        if (tmp < min)          throw bad_data(tmp, min, max);
        else if (tmp > max)
            throw bad_data(tmp, min, max);
        else
            data.push_back(tmp);
    }
}

// Arbitrary precision integer
class bigint {
public:
    bigint( );
    ~bigint( );
};

```

```

    ...
};

int main(int argc, char** argv)
{
    using namespace std;
    if (argc < 2) {
        cerr << "usage: " << argv[0] << " FILE\n";
        return EXIT_FAILURE;
    }

    vector<bigint> data;
    ifstream in(argv[1]);
    if (! in) {
        perror(argv[1]);
        return EXIT_FAILURE;
    }
    try {
        getdata(in, data, bigint(0), bigint(10));
    } catch (const bad_data& ex) {
        cerr << argv[1] << ": " << ex.what( ) << '\n';
        return EXIT_FAILURE;
    } catch(...) {
        std::cerr << "fatal error: unknown exception\n";
        return EXIT_FAILURE;
    }

    if (data.size( ) == 0)
        cout << "no data\n";
    else {
        bigint sum(accumulate(data.begin(),data.end( ),bigint( )));
        std::cout << "avg=" << sum / data.size( ) << '\n';
    }
}

```

## Chapter 5. Functions

Every C++ program has at least one function (`main`), and all but the most trivial programs define additional functions. The C++ standard library provides numerous functions that your program can call. This chapter discusses how to declare, define, and call functions. A function *declaration* tells the compiler about a function's name, return type, and parameters. A function *definition* also provides the body of the function.

See [Chapter 3](#) for more information about function call expressions and [Chapter 4](#) for information about statements, which make up function bodies. This chapter presents information that is common to all kinds of functions. For characteristics that are unique to member functions, see [Chapter 6](#), and for information that pertains specifically to function templates, see [Chapter 7](#).

The syntax descriptions in this chapter are informal. See [Chapter 12](#) for a precise BNF grammar.

### Function Declarations

A function declaration tells the compiler about a function name and how to call the function. The actual body of the function can be defined separately (described later in this chapter). A function declaration has the following parts:

```
type
    name ( parameters ) cv-qualifiers
    except-spec ;
```

The *parameters*, *cv-qualifiers*, and *except-spec* are optional. The *type* is required, except for constructors, destructors, and type conversion operators. The *name* is the function name. (Each of these parts is described later in



this chapter.) [Example 5-1](#) shows a variety of function declarations.

### *Example 5-1. Declaring functions*

```
// Function named "add", which returns type int, and takes two
// parameters, each
// of type int. The names a and b are optional.int add(int a, int
b);

// Function named "print", which takes a const string reference,
// and does not
// return anything. The function is expanded inline.
inline void print(const std::string& str)
{
    std::cout << str;
}

// Function named "test", which takes two floating-point
// arguments and returns an
// enumeration. This function does not throw any exceptions.
enum fptest { less=-1, equal, greater };
fptest test(double, double) throw( );

class demo {
public:
    // Member function named "value", which returns type int, takes
    // no arguments,
    // and is const, which means it can be called for a const
    // object.
    int value( ) const;
    // Function named "~demo", that is, a destructor, that is
    // virtual. Constructors
    // and destructors do not have return types. Destructors do not
    // take arguments.
    virtual ~demo( );
    // Inline, overloaded, const type conversion operator
    operator bool( ) const { return value( ) != 0; }
};
```

## Return Type

The *type* in a function declaration is the function's return type. It is a series of type specifiers (see [Chapter 2](#)) with pointer and reference operators. You can mix any of the

following function specifiers freely with the type specifiers, but the convention is to list function specifiers before other type specifiers:

### explicit

Applies only to constructors. An explicit constructor cannot be used in an implicit type conversion. See [Chapter 6](#) for details.

### inline

Tells the compiler to expand the function body at the point of the function call. An inline function must be defined in every file in which it is used, and the definition must be identical in every file. If the function body contains a local static object, including string literals, every expansion of the function in every file refers to a common object.

A function definition in a class definition is an inline function definition, even without the use of the `inline` specifier.



The `inline` specifier is a hint to the compiler, and the compiler is free to ignore the hint. Most compilers impose a variety of restrictions on which functions can be expanded inline. The restrictions vary from one compiler to another. For example, most compilers cannot expand a recursive function.

Inline functions are most often used for extremely simple functions. For example, all standard containers have a member function `empty`, which returns `true` if the container is empty. Some containers might implement the function as the following `inline` function:

```
inline bool empty( ) const { return size( ) != 0; }
```

## virtual

Applies only to nonstatic member functions. A virtual function's definition is bound to the function call at runtime instead of at compile time. See [Chapter 6](#) for details.

If a function's return type is `void`, no value is returned to the caller. The function does not need a return statement. The form `return;` is permitted, or the return expression must have type `void`.

If the return type is anything other than `void`, every return statement must have an expression, the type of which can be implicitly converted to the function's return type. See [Chapter 4](#) for more information about the return statement.

A function's return type cannot be an array or function type, but it can be a pointer or reference to an array or function.

In a function declaration (but not a definition), you can use the extern storage class specifier. Declarations and definitions can have linkage specifications, e.g., `extern "C"`. See [Chapter 2](#) for more information about storage class and linkage.

A friend specifier can be used to declare a friend function. See [Chapter 6](#) for more information.

Note that the return type is not considered in overload resolution. See [Section 5.3](#) later in this chapter for details.

## Parameters

Function parameters are optional. If a function takes no parameters, you can leave the parentheses empty or use the keyword `void`. If a function requires parameters, the

parameter list is comma-separated, in which each parameter is a simple declaration of the following form:

```
type-specifiers  
    declarator = expr
```

*expr* is optional; if it is omitted, the = symbol is also omitted. The *type-specifiers* allow for the optional register and auto storage class specifiers and pointer and reference operators. See [Chapter 2](#) for more information about declarators and specifiers.

### Tip

In C, a function that takes no arguments requires the `void` keyword, but in C++, `void` is optional. Thus, `void` appears most often in headers that must be used in both C and C++ programs. For example:

```
#ifdef __cplusplus  
    #define EXTERN extern "C"  
#else  
    #define EXTERN extern  
#endif  
EXTERN int get_version(void);
```

In other situations, you can use whichever style you prefer.

You can omit the parameter name from the *declarator*. In a function declaration, the name is important only to the human reader. In a function definition, a nameless parameter cannot be used in the function body. For example, suppose a graphics package defines a variety of shape classes, all deriving from a common base class, `shape`. Among the operations permitted on a `shape` is `scale`, which takes two arguments: the scaling amounts in the x and y directions. Also, suppose that the square shape

(unlike `rectangle`) heeds only the `x` scale factor. The `square::scale` function might be written as:

```
void square::scale(double xscale, double)
{
    this->size *= xscale;
}
```

A parameter can have *cv*-qualifiers (`const` and `volatile`, as discussed in [Chapter 2](#)). The qualifiers have their usual meaning in the function body. The qualifiers and storage class specifier are not part of the function type and do not participate in overload resolution.

## Default Arguments

A parameter can have a default argument (separated from the declarator by an equal sign). Only the right-most parameters can have default arguments. If any given parameter does not have a default argument, all parameters to its left cannot have default arguments. The default argument can be any expression. (If you want to use a comma operator, enclose the expression in parentheses.)

In a function call, arguments that have default values can be omitted, starting from the right. For each omitted argument, the default argument is substituted in the function call. Each default argument is implicitly converted to the parameter type, applying the same rules that govern initializers in declarations. The default argument expressions are evaluated every time the argument is needed in a function call. Names used in the default arguments are looked up at the point of declaration, not the point of use, as shown in [Example 5-2](#).

*Example 5-2. Declaring and using default arguments*

```
#include <iostream>
#include <ostream>

namespace demo {
```

```

int f( )
{
    return 20;
}

int f( )
{
    return 10;
}

// The default argument for y is always the global f( ), even if
// a different f( )
// is visible where func( ) is called.
int func(int x, int y = f( ))
{
    return x + y;
}

int main( )
{
    using demo::f;
    std::cout << f( ) << '\n';        // Prints 20
    std::cout << func(32) << '\n';    // Prints 42
}

```

Default arguments are cumulative in multiple declarations of the same function in the same scope. Later declarations can provide default arguments for additional parameters, in which case the declaration must omit the default arguments for parameters that already have default arguments, as shown in [Example 5-3](#).

### *Example 5-3. Accumulating default arguments*

```

void func(int x, int y);
void func(int x, int y = 10);
void func(int x = 20, int y);
void other( )
{
    func( ); // Same as func(20, 10)
}

```

Different scopes can have different default arguments. For example, the source file in which a function is defined might have different default arguments from those used in

function declarations where the function is used. However, most of the time, different default arguments suggests programmer errors.

In a derived class, an overridden virtual function can have different default arguments than its counterpart in the base class. The default argument is chosen at compile time, based on the object's static type. Thus, the default arguments are typically those of the base class, even if the function actually called is from the derived class. To avoid confusion, it is best to avoid default arguments with virtual functions, or make sure they are the same for all overridden functions. (See [Chapter 6](#) for more information about virtual functions.)

In a member function declaration, you cannot use a nonstatic data member as a default argument unless it is the member of a specific object. If you want to use the value of a data member as the default value for a parameter, use an overloaded function, as shown in [Example 5-4](#). (See [Section 5.3](#) for more on overloaded functions.)

#### *Example 5-4. Default arguments in member functions*

```
class example {
public:
    void func(int x, int y = data_); // Error

    // Achieve the desired effect with overloaded functions.
    void func(int x, int y);
    void func(int x)          { func(x, data_); }
private:
    int data_;
};
```

## **Variable Number of Arguments**

The last parameter in a function declaration can be an ellipsis (...), which permits a variable number of arguments to be passed to the function. The comma that

separates the next-to-last parameter from the ellipsis is optional. However, if portability with C is important, be sure to include the comma. (See `<cstdarg>` in [Chapter 13](#) to learn how to access the additional arguments.) You can use an ellipsis as the sole parameter in a function, but there is no mechanism in standard C++ to access the arguments from the function body. Such a declaration might be used for an external function, however.

## cv-qualifiers

Only nonstatic member functions (but not constructors or destructors) can have *cv*-qualifiers (`const` and `volatile`). They are optional, and if used in a member function declaration, apply to the implicit object parameter of the member function (`this`). You can use `const`, `volatile`, neither, or both in any order. Place *cv*-qualifiers after the closing parenthesis of the function parameters and before the exception specification. The qualifiers are part of the function type and participate in overload resolution, so you can have multiple functions with the same name and parameters, but with different qualifiers (but only if you do not also have a static member function of the same name and parameters; see [Section 5.3](#) later in this chapter for details).

A pointer-to-member function and a function typedef can also have *cv*-qualifiers. Only a top-level typedef can have *cv*-qualifiers; you cannot declare a typedef that combines a function typedef and a qualifier.

*cv*-qualifiers are most often used to declare `const` member functions. These functions can be called for a `const` object. In general, member functions that do not change `*this` should be declared `const`. (See [Chapter 6](#) for more information on how *cv*-qualifiers affect member functions.) [Example 5-5](#) shows some simple uses of qualifiers.

---



### Example 5-5. Using qualifiers with member functions

```
class point
{
public:
    point(int x, int y) : x_(x), y_(y) {}
    int x( )const    { return x_; }
    int y( )        const    { return y_; }
    double abs( ) const    { return sqrt(double(x( ))*x( ) + y(
)*y( )); }
    void offset(const point& p) {
        // Cannot be const because offset( ) modifies x_ and y_
        x_ += p.x( );
        y_ += p.y( );
    }
private:
    int x_, y_;
};
```

## Exception Specifications

An exception specification tells the compiler which exceptions a function can throw. Exception specifications are optional in a function declaration and are rarely used. The syntax is:

```
throw (type-list )
```

The *type-list* is optional. The exception specification follows the function header and cv-qualifiers. If present, it is a comma-separated list of type names. (See [Chapter 2](#) for details about type names.) Each type name is an exception type that the function can throw. If the function throws an exception that is not listed in the exception specification, the unexpected function is called. If the function declaration does not have an exception specification, the function can throw any exception.

The default implementation of unexpected calls terminate to terminate the program. You can set your own unexpected handler, which must call terminate or throw an exception. If your handler throws an exception that is

not listed in the function's exception specification, `bad_exception` is thrown. If `bad_exception` is not listed in the function's exception specification, `terminate` is called. In other words, if there is an exception specification, only exceptions of the listed types (or derived from one of the listed types) can be thrown from the function, or else the program terminates. See `<exception>` in [Chapter 13](#) for details.

An overridden virtual function must have an exception specification that lists only types that are also listed in the base-class exception specifications. In particular, if the base-class function does not throw any exceptions, the derived class function must not throw any exceptions.

An exception specification most often marks functions that do not throw exceptions at all (`throw( )`). [Example 5-6](#) shows various uses of exception specifications.

#### *Example 5-6. Declaring exception specifications*

```
class base {
public:
    virtual void f( ) throw( );
    virtual void g( ); // Can throw anything
    virtual void h( ) throw(std::string);
};

class derived : public base {
public:
    virtual void f( ) throw( ); // OK: same as base
    virtual void g( ) throw(int); // OK: subset of base
    virtual void h( ) throw(int); // Error: int not in base
};

class more : public derived {
public:
    virtual void f( ); // Error: can throw anything
    virtual void g( ) throw( ); // OK
};

// Function does not throw any exceptions
int noproblem(int x, int y) throw( )
try
```

```
{
    dostuff(x);
    dostuff(y);
    return 1;
}
catch(...)
{
    return 0;
}

derived* downcast(base* b) throw(std::bad_cast)
{
    return dynamic_cast<derived*>(b);
}
```

### **Ti p**

Java programmers should note two significant differences between C++ and Java with respect to exception specifications:

- The exception specification is introduced by `throw`, not `throws`.
- The correctness of the exception specification is checked at runtime, not at compile time.

## Function Definitions

Every function that a program uses must be defined exactly once in the program, except for inline functions. (Function templates are a little different; see [Chapter 7](#) for details.) An inline function must be defined in every source file that uses the function. This section discusses function definitions and their relationship to their declarations.

## Declarations and Definitions

In a source file, every function must be declared or defined before it is used. For functions defined in libraries or other source files, the convention is to declare the function in a header (.h or .hpp) file, and the source file where the function is called must `#include` the header file. Every function that is used in the program must have a definition.

Inline functions must be defined in every source file in which they are used. This is typically accomplished by defining them in a header file, which you must `#include` in each source file that calls the inline function. Every definition of an inline function must be identical.

## Function Types and Signatures

A *function type* includes the language linkage, return type, parameter types, and *cv*-qualifiers. Note that for each parameter, only its type is significant; its name, storage class, and *cv*-qualifiers are not part of the function type. Exception specifications are not part of a function's type.

A single source file can have multiple declarations of the same function (that is, functions with the same type), even if those declarations differ in other ways. For example, one declaration can declare a parameter `const` and another

might declare it `volatile`. Because *cv*-qualifiers on parameters do not affect a function's type, both declarations are equivalent. (Parameter qualifiers matter only in the function definition, not the declaration.)

[Example 5-7](#) shows several declarations and one definition, all for a single function.

### *Example 5-7. Declaring and defining functions*

```
// Three declarations of the same function type
int add(const int a, const int b);
int add(int x, volatile int);
int add(signed int, int signed);

// Definition of the function. The parameter qualifiers in the
// declarations are
// irrelevant. Only those in the definition matter.
int add(int x, int y)
{
    return x + y;
}
```

Array and pointer parameter types are also equivalent. In a function body, a parameter with an array type is actually a pointer, not an array. The first (leftmost) size in a multidimensional array is ignored and can be omitted. Similarly, a parameter of function type is the same as a parameter of function pointer type. The rules for function types apply recursively to function and function pointer parameters. Thus, a parameter with a type that is a pointer to a function that takes an `int` is equivalent to one with a type that is a pointer to a function that takes a `const int` parameter. [Example 5-8](#) illustrates equivalent pointer types.

### *Example 5-8. Equivalent pointer types*

```
int first(int const array[10]); // Size is ignored
int first(int const array[]);  // Equivalent declaration
int first(int const *array);   // Equivalent declaration
int first(int const *array)    // Definition
{
    return array[0];
}
```

```

}

int apply(int func(int), int arg);
int apply(int func(int const), int);           // Equivalent
int apply(int (*func)(int), int arg);         // Equivalent
int apply(int (*func)(int const), int arg)    // Definition
{
    return func(arg);
}

```

Because typedef declarations do not create new types, but only synonyms for existing types, parameters that differ only in their use of typedef types are equivalent, as shown in [Example 5-9](#).

*Example 5-9. Using typedef in equivalent parameter types*

```

typedef int INT;
typedef int const CINT;
void func(int);
void func(INT);           // Equivalent
void func(INT const);    // Equivalent
void func(CINT);         // Equivalent
void func(signed int i)  // Definition
{
    std::cout << i;
}

```

You can declare a function type from a typedef, but you cannot use the typedef in a function definition. This usage is uncommon. For example:

```

typedef int func(int, int);
func add;           // Declares int add(int, int);
int add(int a, int b) // Cannot use "func add" here
{
    return a + b;
}

```

Be careful to distinguish between a function type and a function pointer type. A function can be implicitly converted to a function pointer and can be called using a function pointer. A function pointer, however, cannot be used to declare a function. The following is an example of the implicit conversion of a function, `add`, to a function

pointer, and the use of a function pointer, `a`, to call the function:

```
typedef func* funcptr; // Pointer-to-function type
funcptr a = add;      // Pointer-to-function object
int i = a(1, 2);      // Call the function that a points to.
```

A *function signature* is similar to a function type, but it ignores the return type. Overload resolution relies on function signatures to determine which overloaded function to call. See [Section 5.3](#) later in this chapter.

## Function Bodies and try Blocks

A function definition consists of a function declaration followed by a function body. The function body has one of two forms:

### *compound-statement*

Executed when the function is called. When execution reaches a return statement or the end of the compound statement, the function returns.

### *try ctor-initializers compound-statement handlers*

Sets up a try block that surrounds the constructor initializers and function body. If an exception is thrown from any of the *ctor-initializers*, it is handled by the handlers in the same manner as any exception thrown from the compound statement. Thus, this form is typically used only with constructors that have initializers. See [Chapter 6](#) for more information about constructor initializers.

The constructor initializers in the second form are optional. Without them, the form can be used for any function. It is equivalent to having a try statement around the entire function body. [Example 5-10](#) shows a try function body used in this way.

*Example 5-10. A try function body*

```
void func1( )
try {
    // Statements
} catch (...) {
    // Handler
}

void func2( )
{
    try {
        // Statements
    } catch(...) {
        // Handler
    }
}
```



## Function Overloading

A single function name can have multiple declarations. If those declarations specify different function signatures, the function name is *overloaded*. A function call to an overloaded function requires the compiler to resolve the overloaded name and decide which function to call. The compiler uses the argument types (but not the return type) to resolve calls to overloaded functions. [Example 5-11](#) shows simple examples of two overloaded functions named `sqrt`: one for `int` arguments and the other for `double` arguments. The rest of this section explains the rules for overloading and resolution.

### *Example 5-11. Overloaded functions*

```
int sqrt(int);
double sqrt(double);

int main( )
{
    std::cout << sqrt(3) << '\n';    // sqrt(int)
    std::cout << sqrt(3.14) << '\n'; // sqrt(double)
}
```

## Declaring Overloaded Functions

Whenever you declare more than one function with the same name in the same scope, you are overloading the function name. The function can be an ordinary function, member function, constructor, or overloaded operator.

Overloaded functions must differ in their parameter lists: they must have a different number of parameters, or the parameter types must be different. Refer to [Section 5.2.2](#) earlier in this chapter for information on equivalent parameter types.

Default arguments are not considered when declaring overloaded functions (but are important when resolving a call to an overloaded function, as described in [Section 5.3.2](#)). [Example 5-12](#) shows overloaded constructors and member functions for the point class, and overloaded declarations for the add function.

*Example 5-12. Overloaded constructors, member functions, and function declarations*

```
class point {
public:
    point(int x = 0);           // Overloaded constructors
    point(int x, int y);
    point(const point& pt);
    // x and y are overloaded.
    int x( ) const { return x_; }
    void x(int newx) { x_ = newx; }
    int y( ) const { return y_; }
    void y(int newy) { y_ = newy; }
private:
    int x_, y_;
};

// add is overloaded.
int  add(int, int);
int  add(int);
double add(double, double);
double add(double);
long  add(long, long);
long  add(long);

// The following declarations are not overloaded, but are
// redeclarations of the
// add functions.
int  add(int a = 0, int b = 0);
long add(signed long, long signed);

long add(int, int); // Error: cannot overload on return type
```

## Calling Overloaded Functions

A function call expression ([Chapter 3](#)) must identify the function being called, based on its name and arguments.

For example, the simple expression  $f(x)$  might be a call to a function named  $f$ , a function call through a function pointer named  $f$ , a function template named  $f$ , the construction of an object of class  $f$ , a conversion of  $x$  to a type named  $f$ , or the invocation of the function call operator of an object named  $f$ . In each situation, the compiler might use different rules for interpreting  $f$  and  $x$ .

The compiler first uses the function name and context to create a list of *candidate functions*. The number and types of the arguments are used to select the best match from the candidates, and that function is the one that is called. If no match is found, the compiler reports an error. If more than one match ties for "best," the compiler reports an ambiguity error.

For example, the C++ standard library declares three different `sqrt` functions (see `<cmath>` in [Chapter 13](#)):

```
float sqrt(float);
double sqrt(double);
long double sqrt(long double);
```

Suppose you add another function named `sqrt`, such as the following, to apply it to each element of an array :

```
void sqrt(const double data[], size_t count);
```

In a function call to `sqrt` (e.g., `sqrt(x)`), the compiler first uses the ordinary rules of name lookup ([Chapter 2](#)) to find the first suitable object or function named `sqrt`. Suppose you used `using namespace std;` to import the standard `sqrt` functions into the same namespace as your `sqrt` function. The compiler would collect all the overloaded `sqrt` functions as the candidate list, which has four elements in this case (the three original functions plus the array version). The list is then pruned to eliminate functions with the wrong number of arguments. In this case, the array version of the function is eliminated because the expression has only one argument. Finally, the

type of `x` is used to determine which function to call. If there is an exact match, the corresponding function is called. If `x` is, for example, an integer, the compiler reports an error because the three floating-point `sqrt` functions look equally good. If `x` has class type, and the class has a conversion function to one of the floating-point types, the compiler implicitly calls that conversion and then calls the corresponding function:

```
struct FixedPoint {
    ...
    operator long double( );
};
void demo( )
{
    FixedPoint x;
    std::cout << sqrt(x) << '\n'; // Prints a long double
}
```

If a candidate is a function template, the compiler deduces the argument types (see [Chapter 7](#)), and from that point on, the template instance is treated as a normal function.

The rules for creating the candidate list and argument list depend on the context of the function call. The argument list can also depend on context: when choosing an overloaded function to call, a member function's class is treated as an argument type. More precisely, member functions have an *implicit object parameter* whose type is a reference to  $T$ , in which  $T$  is the class that defines the member function. Any qualifiers on the member function also qualify the type (that is, the object type for a `const` function is `const T &`). In the function body, the type of `this` is a pointer to qualified  $T$ . (See [Chapter 6](#) for more information about this.)

A call to a member function applies to a specific object, which is the *implicit object argument*. When calling a member function with the `->` or `.` operator, the implicit object argument is the left operand. For an overloaded

binary operator (such as operator<<), the implicit object argument is the left operand. For a unary operator, the implicit object argument is the operand. When calling an unqualified function inside a nonstatic member function, the implicit object argument is *\*this*. The implicit object argument is considered the first argument in the argument list. Unlike normal arguments, implicit type conversions do not take place for the implicit object argument.

## Candidate Functions

This section describes how the compiler creates its list of candidate functions. [Table 5-1](#) summarizes the various categories of function calls, and the subsequent subsections provide the details.

*Table 5-1. Function calls and candidate lists*

Cat egor y	Funct i on cal l synt ax	Candi dat e funct i ons
Qualified member function call	<i>expr . name ( args )</i> <i>expr -&gt; name ( args )</i>	Member functions
Unqualified function call	<i>name ( args )</i> <i>expr ( args )</i>	Member functions Nonmember functions
Operator	<i>expr</i> <i>op</i> <i>expr</i>	Member functions Nonmember functions
Function-like initialization	<i>type</i> <i>name ( args )</i>	Constructors
Assignment-like initialization	<i>type</i> <i>name = expr</i>	Constructors Type conversion operators

Category	Function call syntax	Candidate functions
Conversion initialization	<i>type</i> <i>name = expr</i>	Type conversion operators
Reference bound to initialization	<i>type&amp; name = expr</i>	Type conversion operators

## Qualified member function call

An expression that uses the `.` or `->` operator to call a function must call a member function. The function name is looked up in the class of the left operand and in its base classes (using the name lookup rules listed in [Chapter 2](#)). The search starts in the most-derived class. If that class declares a member function with the desired name, the search stops, and the candidate functions are all the member functions with the same name in that class. If no matching member function is found, the search continues with the immediate base classes. The search stops when a matching name is found.

In other words, a derived class cannot overload a function that is declared in a base class. Instead, if the derived class has a function with the same name, the derived class hides the name that would be inherited from the base class (or the derived class might override a virtual function; see [Chapter 6](#)). Insert a `using` declaration in the derived class if you want the compiler to consider the base class functions as candidates, as shown in [Example 5-13](#).

### *Example 5-13. Overloading inherited functions*

```
#include <iostream>
#include <ostream>

class base {
public:
    void f(int) { std::cout << "f(int)\n"; }
    void g(int) { std::cout << "g(int)\n"; }
```

```

};

class derived : public base {
public:
    void f(double) { std::cout << "f(double)\n"; }
    void g(double) { std::cout << "g(double)\n"; }
    using base::g; // g(int) and g(double) are visible.
};

int main( )
{
    derived d;d.f(3); // Calls derived::f(double)
    d.g(42); // Calls base::g(int)
}

```

If a class has multiple immediate base classes, overload resolution must find a name in only one of the base classes. If functions with the desired name are found in multiple immediate base classes, the compiler reports an ambiguity error. To resolve this ambiguity, use the scope operator (::) to qualify the function name in the derived class, as shown in [Example 5-14](#).

#### *Example 5-14. Avoiding ambiguous base-class overloads*

```

struct base1 {
    void func(int);
};
struct base2 {
    void func(double);
};
struct derived : base1, base2 {
    // Call to overloaded func is ambiguous.
    void demo1(long x) { func(x); }
    // Qualify the name to resolve ambiguity.
    void demo2(long x) { base2::func(x); }
};

```

## **Unqualified function call**

An ordinary-looking function call can be a nonmember function, a member function, an object of class type, a type name, or a variable of type pointer-to-function. For a variable of type pointer-to-function, overload resolution

takes place when a value is assigned to the variable (discussed later in this chapter). In the other cases, the usual name lookup rules apply when finding the candidate functions.

A function call in a member function searches first for matching member functions in the same class or an ancestor class. The search for a match begins in the class that is declaring the member function. If a match is found, candidate functions are taken from that class. If no matches are found, the search continues with ancestor classes, following the same rules as for qualified member function calls. If no matches are found in any ancestor classes, the namespace of the class is searched for nonmember functions. [Example 5-15](#) shows how a matching member function in a base class precludes finding a better match in the global namespace.

#### *Example 5-15. Finding candidate member functions*

```
#include <iostream>
#include <ostream>

void proc(int x)
{
    std::cout << "proc(int:" << x << ")\n";
}

class base {
public:
    void f(int) { std::cout << "f(int)\n"; }
    void g(int) { std::cout << "g(int)\n"; }
    void proc(double) { std::cout << "base::proc(double)\n"; }
};

class derived : public base {
public:
    void f(double) { std::cout << "f(double)\n"; }
    void g(double x) {
        std::cout << "g(double)\n";
        proc(42); // Calls base::proc(double), not ::proc(int)
    }
    using base::g;
};
```



```

};

// Declared after derived, so call to proc( ) inside g( ) never
// sees this proc( ).
void proc(double x)
{
    std::cout << "proc(double:" << x << ")\n";
}

int main( )
{
    derived d;d.g(3.14159); // Calls g(double)
}

```

If a function call expression resolves to an object of class type, the class must have a function call operator or a conversion operator, in which the conversion is to a function type: pointer-to-function, reference-to-function, or reference-to-pointer-to function.

A conversion operator is rarely used. The compiler constructs a wrapper function so that the conversion function is the first argument. The conversion type followed by the types of the actual arguments is the new list of argument types to use in overload resolution. In other words, all of a class's function-type conversion operators participate in overload resolution.

[Example 5-16](#) shows how a class-type object is used as the left operand of a function call.

*Example 5-16. Calling a class-type object as a function*

```

#include <iostream>
#include <ostream>

typedef void (*strproc)(const char*);

void print(const char* str)
{
    std::cout << "const char*:" << str << '\n';
}

void print(int x)
{

```

```

    std::cout << "int:" << x << '\n';
}

void print(double x)
{
    std::cout << "double:" << x << '\n';
}

struct simple
{
    void operator( )(int x)    { print(x); } // print(int)
    void operator( )(double x) { print(x); } // print(double)
};

typedef void (*intfunc)(int);
typedef void (*dblfunc)(double);

struct indirect
{
    operator intfunc( ) { return print; } // print(int)
    operator dblfunc( ) { return print; } // print(double)
    operator strproc( ) { return print; } // print(const char*)
};

int main( )
{
    simple sim;
    indirect ind; sim(42); // Prints "int:42"
    sim.operator( )(42); // Prints "int:42"
    sim(42.0); // Prints "double:42"
    ind(42); // Prints "int:42"
    ind.operator intfunc( )(42); // Prints "int:42"
    ind(42.0); // Prints "double:42"
    ind("forty-two"); // Prints "const char*:forty-two"
}

```

## Oper at or

The function for an overloaded operator is chosen according to the usual rules for resolving overloaded functions, in which the operator's operands are the function's arguments. You can overload operators only if at least one operand has a user-defined type. If all operands have built-in types, an operator has its built-in meaning.

If the left operand of an operator has class type, the operator can be a nonstatic member function or a nonmember function. Otherwise, the function must be a nonmember function. The operator function name is formed from the keyword operator followed by the operator symbol, e.g., operator[], operator++, or operator-. Unary operators can be member functions with no arguments or nonmember functions of one argument. Binary operators are member functions of one argument or nonmember functions of two arguments. Postfix increment and decrement operators are different. They are implemented as binary operators for which the right operand is an int. (See [Chapter 3](#) for more information.)

The -> operator is also different. Although it is a binary operator, it is treated as a unary operator. It must be implemented as a member function, so the function takes no arguments. It returns a pointer or another object that overloads the -> operator. Ultimately, the overloaded operators must resolve to a pointer of class type, to which the built-in -> operator can be applied.

The candidate functions for the overloaded operator include all member, nonmember, and built-in candidate functions. Member functions do not take precedence over nonmember functions.

[Example 5-17](#) shows how operator functions are different from named functions. In particular, it shows how operators are resolved by considering member functions and global functions, whereas named member functions take precedence over named global functions.

*Example 5-17. Calling overloaded operators*

```
class demo
{
public:
    demo(int v) : value_(v) {}
```

```

demo add(const demo& d) const;
demo sub(const demo& d) const;
demo mul(const demo& d) const;
demo operator+(const demo& d) const;
demo operator-(const demo& d) const;
demo operator*(const demo& d) const;
operator int( ) const { return value_; }
private:
    int value_;
};

// Silly examples, but illustrative
demo add(const demo& a) { return a; }
demo mul(const demo& a) { return a; }
demo div(const demo& a) { return a; }

demo operator+(const demo& a, const demo& b)
{
    return a.operator+(b); // Force use of member function.
}

demo demo::add(const demo& d)
const
{return *this + d; // Error: calls ::operator+( ) or
demo::operator+( )?
}

demo demo::sub(const demo& d) const
{
    return this->operator-(d); // Member operator
}

demo demo::mul(const demo& d) const
{
    return ::operator*(*this, d); // Global operator
}

demo demo::operator+(const demo& d) const
{
    return demo(int(*this) + int(d));
}

demo demo::operator-(const demo& d) const
{
    return sub(d); // Calls demo::sub (recurses infinitely)
}

demo demo::operator*(const demo& d) const

```

```
{
    return ::mul(d); // Scopes operator to call global mul( )
}
```

## Function-like initialization

An object can be initialized using function-like syntax (see [Chapter 2](#)). When an object of class type is so initialized, the candidate functions are constructors of the named class. The same syntax applies to conversion initialization covered later in this chapter. [Example 5-18](#) shows function-like initializers that call an overloaded constructor.

### *Example 5-18. Calling an overloaded constructor*

```
class point {
public:
    point( ) : x_(0), y_(0) {}
    point(int x) : x_(x), y_(0) {}
    point(int x, int y): x_(x), y_(y) {}
private:
    int x_, y_;
};point p1(42);
point p2(4, 2);
point p3(p1);
```

## Assignment-like initialization

An object can be initialized using assignment-like syntax (see [Chapter 2](#)). The candidate functions for "T x = i" are single-argument, non-explicit constructors of T. If i has class type, the candidates are conversion functions that convert i to type T, as shown in [Example 5-19](#). The compiler is free to call T's copy constructor to copy a temporary T object to x. Even if the compiler optimizes away this extra copy, T must have an accessible copy constructor.

### *Example 5-19. Resolving assignment-like initialization*

```
class point {
public:
    point( ) : x_(0), y_(0) {}
    point(int x) : x_(x), y_(0) {}
```

```

    point(int x, int y) : x_(x), y_(y) {}
    int x( ) const { return x_; }
    int y( ) const { return y_; }
private:
    int x_, y_;
};

class dot {
public:
    dot(int x, int y) : center_(point(x, y)) {}
    dot(const point& center) : center_(center) {}
    operator point( ) const { return center_; }
private:
    point center_;
}; point p1 = 3; // Invokes point(int) constructor
point p2 = 4.2; // Converts 4.2 to 4, and invokes point(int)
dot d1 = p1; // Invokes dot(const point&) constructor
point p3 = d1; // Invokes dot::operator point( ) and implicit
// point(const point&) copy constructor

```

## Conversion Initialization

An object that does not have class type can be initialized with an object of class type. The candidate functions are the conversion functions of that class type. [Example 5-20](#) shows two cases of conversion initialization. In the first, the object `c` is initialized with an `rbnode<int>` object, `n`, and the compiler calls `operator color( )`. The second invokes `operator T( )`, in which `T` is `int`, to initialize `i` to the value 42.

*Example 5-20. Initializing non-class-type objects by calling conversion functions*

```

enum color { red, black };

template<typename T>
class rbnode {
public:
    rbnode(const T& value, color c);
    operator color( ) const { return color_; }
    operator T( ) const { return value_; }
private:
    T value_;
    color color_;
}

```

```
};
rbnode<int> n(42, black);color c = n;
int i = n;
```

## Reference bound to conversion

Similar to initializing an object of non-class type with an expression of class type, you can initialize a reference to an lvalue that results from a conversion function and initialize a const reference to an rvalue. Most conversion operators do not return lvalues, so the const reference version of this rule is used more often. [Example 5-21](#) shows some examples of binding references to conversions. The object `c1` is bound to a temporary copy of `n.color_`, and `i1` is bound directly to `n.value_`.

### *Example 5-21. Binding references to conversion lvalues*

```
enum color { red, black };

template<typename T>
class rbnode {
public:
    rbnode(const T& value, color c);
    operator color( ) { return color_; }
    operator T&( ) { return value_; }
private:
    T value_;
    color color_;
};

rbnode<int> n(42, black);const color& c1 = n;
int& i1 = n;
```

## Addresses of Overloaded Functions

When taking the address of a function, the compiler does not have the benefit of an argument list to help resolve overloading. Instead, it collects a list of candidate functions and picks one based on the type required by the context. The context can be an initializer, an assignment, a function

argument, the return value of a function, or an explicit type cast.

Of the potentially matching functions, nontemplate functions are better than template functions, and more-specific template functions are better than less-specific template functions. There must be exactly one best function, or else the compiler reports an error. [Example 5-22](#) shows simple examples of resolving the addresses of overloaded functions.

*Example 5-22. Taking the address of an overloaded function*

```
int abs(int);
long abs(long);
double abs(double);

template<typename T>
T abs(T x);

template<typename T, typename U>
T abs(T x, U = U( ));

int main( )
{
    int (*intfunc)(int) = &abs;           // abs(int)
    double (*dblfunc)(double) = &abs;    // abs(double)
    float (*fltfunc)(float) = &abs;      // abs<float>
    short (*shrtfunc1)(short, int) = &abs; // abs<short,int>
    short (*shrtfunc2)(short) = &abs;    // abs<short>
}
```

## Best Overloaded Function

Once the compiler has found the list of candidate functions, it must choose the "best" match from the list. The candidate list is first pruned by removing functions with the wrong number of arguments. Then the remaining functions are checked to determine how to convert the actual arguments to the desired parameter types. The function with the simplest conversions wins. This section discusses these two steps.



## Pruning the candidate list

Once the compiler has assembled its list of candidate functions, it prunes the list by removing functions that have the wrong number of arguments. If the function call has  $n$  argument expressions, a function is kept in the list if any of the following apply:

- It has exactly  $n$  parameters.
- It has fewer than  $n$  parameters followed by an ellipsis parameter.
- It has more than  $n$  parameters, and the extra parameters have default arguments.

Also, each actual argument must be convertible (applying the rules described in the next section) to the parameter type.

### Overloading versus Default Arguments

Overloading functions can have the same effect as using default arguments. The question is when to use overloading and when to use default arguments. As usual, the answer is, "It depends."

If a default argument is complicated, you are probably better off using overloading. With default arguments, the complicated code would be duplicated at every function call. With overloading, you can concentrate it at a single point.

Ask yourself whether it makes sense to omit arguments, starting from the right. If so, you might be able to use default arguments. If some arguments cannot be omitted singly, you might need to use overloading. In the following example, omitting  $y$  without also omitting  $x$  is unusual:

```
void circle(float radius, float x, float y);  
void circle(float radius); // x=0, y=0
```

If a function is complicated, you might want to use default arguments so you can write the function only once. Even if you use overloading, write a single base function and let the overloaded functions call the base function:

```
class shape {  
public:  
    // Using default arguments keeps the code simple because  
    // there is  
    // only one constructor.  
    explicit shape(color c = black, size s = 1);  
    ...  
};
```

## Choosing the best function

Of the remaining candidate functions, the "best" function is the one that is called. If multiple functions are tied for best, the compiler reports an error. The best function is found by comparing the implicit type conversion sequences needed for each argument type (and possibly an implicit object argument). As described in the next section, some sequences are better than others. When comparing two functions, the better function is the one with the better conversion sequences:

- If two functions A and B have argument conversion sequences that are equally good, but A has at least one argument with a better sequence, A is the better function.
- If all the argument conversion sequences are equally good, a nontemplate function is better than a template function.

- If the conversion sequences for two template functions are equally good, a more-specific template function is better than a less-specific template function.

## Argument conversion sequences

An actual argument is implicitly converted to the desired parameter type by undergoing a series of transformations. Basically, a better conversion is one in which the argument type is closer to the parameter type and, therefore, the type undergoes fewer transformations.

A *standard conversion sequence* is a sequence of built-in transformations that are based on automatic type conversions ([Chapter 3](#)). A sequence is built from at most one of each of three kinds of transformations. The quality of a standard conversion sequence is dictated by its worst transformation, in which the value-preserving transformations are better than those that might discard information. The transformations are as follows, ordered from best to worst:

### Lvalue transformation and qualification adjustment

An lvalue transformation is the conversion of an array to a pointer, an lvalue to an rvalue, or a function to a function pointer. A qualification adjustment is adding `const` or `volatile` qualifiers to match the parameter. These transformations are just as good as no transformation at all:

```
void func1(const int& x);  
void func2(const char* str);  
func1(42);  
func2("Hello");
```

### Promotion

The built-in type promotions are listed in [Chapter 3](#). In general, type promotions are from smaller types to larger types, so information cannot be lost:

```
void func3(long l);  
func3(42);
```

## Conversion

A built-in type conversion is one that might lose information, such as converting a number to a smaller type, or a floating point to an integer. Certain conversions are better than others:

- A conversion of a pointer to `bool` is worse than other conversions.
- Converting a derived-class pointer to a base-class pointer is better than converting the pointer to `void*`.
- Converting a derived-class pointer to `void*` is better than converting a base-class pointer to `void*`.
- When converting classes, class pointers, and pointers to members, shorter distances in a class hierarchy are better than longer distances; for example, if A inherits from B and B inherits from C, `A*` to `B*` is better than `A*` to `C*`.

The following are some examples of conversion transformations:

```
void func4(bool b);  
func4(3.14);  
void func5(void* ptr);  
func5(&x);
```

*A user-defined conversion sequence* has up to three parts: a standard conversion sequence followed by a single, user-defined conversion (constructor or type conversion operator) followed by another standard conversion sequence.

A standard conversion sequence is better than a user-defined conversion sequence. A user-defined conversion

sequence is better than matching an ellipsis parameter.

If one sequence of transformations is a subsequence of another, the shorter sequence is better than the longer one.

Among sequences that differ only by qualification, the one with fewer qualification adjustments is better than one with more adjustments.

[Example 5-23](#) shows how overloaded functions are chosen. Notice that type promotions are preferred to the conversion to Num, even though Num has an exact type conversion to long. Also note how an unsigned long cannot be promoted, so it must undergo a built-in conversion. The compiler has no preference of int, long or double, so the conversion is ambiguous and, therefore, results in an error. The same applies to long double; even though you might consider the conversion to double to be "better" than conversion to int, the rules of overloading say otherwise. The final example results in an error because there is no matching function. An array of wchar\_t cannot be converted to any of the types used as func parameters.

### *Example 5-23. Resolving overloaded functions*

```
#include <string>

void func(double);
void func(long);
void func(int);
void func(const std::string&);

class Num {
public:
    Num(int i) : num_(i) {}
    operator long( ) const { return num_; }
private:
    int num_;
};

int main( )
{
    short n = 42;
```

```
func(n);           // func(int)
func(Num(n));     // func(long)
func(true);      // func(int)
func(3.1415f);   // func(double)
func("string");  // func(string);
std::string s("string");
func(s);         // func(string);
func(3.14159L);  // Error: ambiguous
func(42UL);     // Error: ambiguous
func(L"widestr"); // Error: no match
}
```

# Operator Overloading

For user-defined types (classes and enumerations), you can define alternate behavior for the C++ operators. This is called overloading the operators. You cannot define new operators, and not all operators can be overloaded. [Table 5-2](#) lists all the operators and indicates which can be overloaded. For these, it shows whether the overload must be a member function. Overloaded operators that are implemented as member functions must be nonstatic member functions.

*Table 5-2. Operators and overloading*

Operator	Meaning	Overloading permitted?	Must be member function?
+	Addition, unary plus	yes	no
&	Address of	yes	no
[]	Array subscript	yes	yes
&=	Assign bitwise and	yes	no
^=	Assign bitwise exclusive or	yes	no
=	Assign bitwise or	yes	no
-=	Assign difference	yes	no
<<=	Assign left shift	yes	no
=	Assignment	yes	yes
*=	Assign product	yes	no

Operator	Meaning	Overloading permitted?	Must be member function?
/=	Assign quotient	yes	no
%=	Assign remainder	yes	no
>>=	Assign right shift	yes	no
+=	Assign sum	yes	no
&	Bitwise and	yes	no
~	Bitwise complement	yes	no
^	Bitwise exclusive or	yes	no
	Bitwise or	yes	no
? :	Conditional	no	N/A
new	Create dynamic object	yes	no
new[]	Create dynamic array	yes	no
--	Decrement	yes	no
delete	Destroy dynamic object	yes	no
delete[]	Destroy dynamic array	yes	no
/	Division	yes	no
==	Equal	yes	no
( )	Function call	yes	yes



Operator	Meaning	Overloading permitted?	Must be member function?
>	Greater than	yes	no
>=	Greater than or equal	yes	no
++	Increment	yes	no
<<	Left shift	yes	no
<	Less than	yes	no
<=	Less than or equal	yes	no
&&	Logical and	yes	no
!	Logical complement	yes	no
	Logical or	yes	no
.*	Member reference	no	N/A
->*	Member reference	yes	yes
.	Member reference	no	N/A
->	Member reference	yes	yes
*	Multiplication, dereference	yes	no
!=	Not equal	yes	no
%	Remainder	yes	no
>>	Right shift	yes	no

Operator	Meaning	Overloading permitted?	Must be member function?
::	Scope	no	N/A
,	Serial evaluation	yes	no
-	Subtraction, negation	yes	no
<i>type</i>	Type conversion	yes	yes

An overloaded operator is a function in which the function name has the form operator followed by the operator symbol, or in the case of a type conversion member function, a list of type specifiers (with pointer, reference, and array operators). For example, the following code declares functions to overload the ! and && operators:

```
enum logical { no, maybe, yes };
logical operator !(logical x);
logical operator &&(logical a, logical b);
```

Some overloaded operators must be member functions, and others can be member or nonmember functions (as shown in [Table 5-2](#)). When you define a member function, the object is always the lefthand operand. A unary operator, therefore, takes no arguments because the one operand is the object itself. Likewise, a binary operator takes one argument: the righthand operand; the lefthand operand is the object. For a nonmember function, a unary operator takes one argument and a binary operator takes two arguments: the first argument is the lefthand operand, and the second is the righthand operand.

Use overloaded operators as you would built-in operators. You can also use function notation, in which the function name is operator followed by the operator symbol, but this usage is uncommon. You can use the function notation with

built-in operators, too, but such usage is extremely uncommon. For example:

```
operator-(42, 10) // Same as 42 - 10
operator-(33)     // Same as -33
```

The usual rules for resolving overloaded functions applies to overloaded operators. The only difference is that the built-in operators are added to the list of candidates along with the user-defined operators. Remember that you cannot overload operators when all the operands have fundamental types. At least one operand must have a user-defined type (class or enumeration).

### Defining Commutative Operators

When overloading a binary operator, consider whether the operator should be commutative ( $a + b$  is the same as  $b + a$ ). If that is the case, you might need to define two overloaded operators:

```
enum priority { idle, low, normal, high };
// Adding an integer to a priority commutes. For example:
//  setpriority(priority( ) + 2);
//  setpriority(2 + priority( ));
priority operator+(priority p, int i);
priority operator+(int i, priority p);
// Subtracting an integer from a priority does not commute.
For
// example:
//  setpriority(priority( ) - 1);
priority operator-(priority p, int i);
```

## Short-Circuit Operators

A key difference between the overloaded operators and the built-in operators is that the logical `&&` and `||` operators are *short-circuit operators*. If the expression result is known by evaluating only the left operand, the right operand is never

evaluated. For overloaded operators, all operands are evaluated before a function is called, so short-circuit evaluation is impossible.

In other words, you cannot tell whether the `&&` and `||` operators perform short-circuit evaluation by merely glancing at them. You must study the types of the operands. It is safest, therefore, never to overload these operators. If the operators are never overloaded, you know that they are always short-circuit operators.

## Comma Operator

You can overload the comma operator (`,`), but you will rarely have any reason to do so. If you do, however, you change its semantics in a subtle way. The built-in comma operator has a sequence point (see [Chapter 3](#)) between its operands, so you know that the lefthand expression is completely evaluated before the righthand expression. If you overload the operator, you lose that guarantee. The ordinary rules apply, so the operands might be evaluated in any order.

## Increment and Decrement

When overloading the increment and decrement operators, remember that they have two forms: prefix and postfix. To distinguish between the two forms, the postfix form takes an additional `int` parameter. The compiler always passes `0` as the additional argument. [Example 5-24](#) shows one way to overload the increment operator. (Decrement is analogous.)

*Example 5-24. Overloading the increment operator*

```
enum status { stopped, running, waiting };
status& operator++(status& s) { // Prefix
    if (s != waiting)
        s = status(s + 1);
```

```

    return s;
}
status operator++(status& s, int) {    // Postfix
    status rtn = s;
    ++s;
    return rtn;
}
int main( )
{
    status s(stopped);
    ++s;                                // Calls operator++(s);
    s++;                                // Calls operator++(s, 0);
}

```

## Member Reference

The `->` operator is different from the other operators. Although you use it as a binary operator, you overload it as a unary operator. It must be implemented as a member function, so the function takes no arguments. It must return one of the following:

- An object of class type, for which the type overloads the `->` operator
- A pointer to a class type

A chain of `->` operators is followed until it ends with a pointer to a class type. The actual right operand must name a member of that class. The `->` operator is most often overloaded to implement a smart pointer. See the `auto_ptr<>` template in the `<memory>` section of [Chapter 13](#) for an example.

## Function Call

The function call operator (`operator( )`) takes any number of arguments. It must be implemented as a member function. To invoke the operator, use an object of class type as the "name" of the function. Pass the arguments as you

would any other function arguments. With a simple variable of class type, the syntax looks like an ordinary function call.

An object that overloads the function call operator is often called a *functor*. Functors are typically used with the standard algorithms to better encapsulate functionality. Some algorithms, such as `for_each`, also permit the functor to store state information, which cannot be done with a plain function. Comparison functions for the associative containers are easier to implement as functors. See `<algorithm>` and `<functional>` in [Chapter 13](#) for examples.

## **operator new and operator delete**

A new expression ([Chapter 3](#)) calls `operator new` to allocate memory, and a delete expression calls `operator delete`. (A `new[]` expression calls `operator new[]`, and a `delete[]` expression calls `operator delete[]`. For the sake of simplicity, whenever I refer to a new expression, I mean a new expression or `new[]` expression. Similarly, `operator new` refers to `operator new` and `operator new[]`. Ditto for delete.)

You can overload `operator new` and `operator delete` in the global scope or as members of a class. In the global scope, the functions must not be static, nor can you declare them in a namespace. When you define these operators as member functions, the functions are always static, even if you omit the `static` specifier. If you do not overload the global operators, the C++ library provides an implementation for you. (See `<new>` in [Chapter 13](#).) If you do not overload the operators for a class, the global operators are used.

If a class overloads the `operator new` and `operator delete` functions, the corresponding operator functions are called for new and delete expressions involving that class. When

overloading operator new or operator delete as member functions or with placement arguments, you can call the global operator, as shown in [Example 5-25](#).

*Example 5-25. Overloading operator new and operator delete*

```
#include <cstddef>
#include <iostream>
#include <memory>
#include <new>
#include <ostream>

class demo
{
public:
    static void* operator new(std::size_t size)
        throw (std::bad_alloc)
    {
        std::cout << "demo::new\n";
        if (instance == 0)
            instance = ::new demo;
        ++count;
        return instance;
    }
    static void operator delete(void* p) {
        std::cout << "demo::delete\n";
        if (--count == 0) {
            ::delete instance;
            instance = 0;
        }
    }

    static demo* make( ) { return new demo( ); }

private:
    demo( ) {}
    demo(const demo&);
    static demo* instance;
    static std::size_t count;
};

demo* demo::instance;
std::size_t demo::count;

int main( )
{
```

```
std::auto_ptr<demo> s1(demo::make( ));  
std::auto_ptr<demo> s2(demo::make( ));  
return s1.get( ) == s2.get( );  
}
```

The first parameter to operator new has type size\_t and is the amount of memory to allocate. The function returns a pointer to the allocated memory as a void\*. Additional parameters are allowed for placement new functions (see [Chapter 3](#)). The first parameter to operator delete is a void\* pointer to the memory; the function returns void. Additional parameters are allowed for placement delete. See <new> in [Chapter 13](#) for more information about overloaded operator new and operator delete.

When you overload the operator new and operator delete functions, you will probably want to overload the scalar (operator new) and array versions (operator new[]) of the operator. The scalar and array versions often behave identically, but you have the option of making them behave differently. Note that the compiler initializes the objects, so your allocation or deallocation function does not need to know the number of objects being allocated or freed.

If you overload operator new, you should probably also overload operator delete. In the case of placement new, the corresponding placement delete function is called if an exception is thrown while constructing the newly allocated object or array. Without a corresponding placement delete function, no operator delete function is called. This is the only time a placement delete function is called. A delete expression always calls the plain, single-argument form of operator delete.

## Type Conversion

A class can declare type conversion operators to convert class-type objects to other types. The operator functions must be nonstatic member functions. The name of each



operator is the desired type, which can be a series of type specifiers with pointer, reference, and array operators, but cannot be a function or array type:

```
class bigint {
public:
    operator long( ); // Convert object to type long.
    operator unsigned long( );
    operator const char*( ); // Return a string representation
    ...
};
```

# The main Function



Every program must have a function in the global namespace called `main`, which is the main program. This function must return type `int`. The C++ environment calls `main`; your program must never call `main`. The `main` function cannot be declared `inline` or `static`. It can be called with no arguments or with two arguments:

```
int main( )
```

or:

```
int main(int argc, char* argv[])
```

The `argc` parameter is the number of command-line arguments, and `argv` is an array of pointers to the command-line arguments, which are null-terminated character strings. By definition, `argv[argc]` is a null pointer. The first element of the array (`argv[0]`) is the program name or an empty string.

## Tip



An implementation is required to support at least two forms of `main`: one that takes no parameters, and one that takes the two parameters for command-line arguments. An implementation can support other forms of `main`. The standard recommends that the first two parameters should be the same as they are in the standard, and that additional parameters follow them.



Static objects at namespace scope can have constant or dynamic initial values, those of POD type with constant values are initialized by constant data before the program starts, and those with dynamic values are initialized by code when the program begins. When, exactly, the objects are initialized is implementation-defined. It might happen before `main` is called, or it might be after.

You should avoid writing code that depends on the order in which static objects are initialized. If you cannot avoid it, you can work around the problem by defining a class that performs the required initialization and defining a static instance of your class. For example, you can guarantee that the standard I/O stream objects are created early so they can be used in the constructor of a static object. See `<ios>` in [Chapter 13](#) for more information and an example.

A local static object is initialized when execution first reaches its declaration. If the function is never called, or if execution never reaches the declaration, the object is never initialized.

When `main` returns or when `exit` is called (see `<cstdlib>` in [Chapter 13](#)), static objects are destroyed in the reverse order of their construction, and the program terminates. All local, static objects are also destroyed. If a function that contains a local, static object is called during the destruction of static objects, the behavior is undefined.



The value returned from `main` is passed to the host environment. You can return `0` or `EXIT_SUCCESS` (declared in `<cstdlib>`) to indicate success, or `EXIT_FAILURE` to tell the environment that the program failed. Other values are implementation-defined. Some environments ignore the value returned from `main`; others rely on the value.

## Chapter 6. Classes

C++ supports object-oriented programming with a traditional, statically-typed, class-based object model. That is, a class defines the behavior and the state of objects that are instances of the class. Classes can inherit behavior from ancestor classes. Virtual functions implement type polymorphism, that is, a variable can be declared as a pointer or reference to a base class, and at runtime it can take on the value of any class derived from that base class. C++ also supports C-style structures and unions using the same mechanism as classes. This chapter describes classes in all their glory. For information about class templates, see [Chapter 7](#). Member functions are a special kind of function; see [Chapter 5](#) for general information about functions.

The syntax descriptions in this chapter are informal. See [Chapter 12](#) for a precise BNF grammar.

### Class Definitions

A class definition starts with the `class`, `struct`, or `union` keyword. The difference between a `class` and a `struct` is the default access level. (See [Section 6.5](#) later in this chapter for details.) A `union` is like a `struct` for which the storage of all the data members overlap, so you can use only one data member at a time. (See [Section 6.1.3](#) later in this section for details.)

A class definition can list any number of base classes, some or all of which might be virtual. (See [Section 6.4](#) later in this chapter for information about base classes and virtual base classes.)

In the class definition are declarations for data members (called instance variables or fields in some other

languages), member functions (sometimes called methods), and nested types.

A class definition defines a scope, and the class members are declared in the class scope. The class name itself is added to the class scope, so a class cannot have any members, nested types, or enumerators that have the same name as the class. As with any other declaration, the class name is also added to the scope in which the class is declared.

You can declare a name as a class, struct, or union without providing its full definition. This incomplete class declaration lets you use the class name in pointers and references but not in any context in which the full definition is needed. You need a complete class definition when declaring a nonpointer or nonreference object, when using members of the class, and so on. An incomplete class declaration has a number of uses:

### Forward class declarations

If two classes refer to each other, you can declare one as an incomplete class, then provide the full definitions of both:

```
class graphics_context;
class bitmap {
    ...
    void draw(graphics_context*);
    ...
};
class graphics_context {
    ...
    bitblt(const bitmap&);
    ...
};
```

### Opaque types

Sometimes a class uses hidden helper classes. The primary class can hold a pointer to an incomplete helper

class, and you do not need to publish the definition of the helper class. Instead, the definition might be hidden in a library. This is sometimes called a *pimpl* (for a number of reasons, one of which is "pointer to implementation").

For example:

```
class bigint {
public:
    bigint( );
    ~bigint( );
    bigint(const bigint&);
    bigint& operator=(const bigint&);
    ...
private:
    class bigint_impl;
    std::auto_ptr<bigint_impl> pImpl_;
};
```

For a complete discussion of the pimpl idiom, see *More Exceptional C++*, by Herb Sutter (Addison-Wesley).

[Example 6-1](#) shows several different class definitions.

#### *Example 6-1. Class definitions*

```
#include <string>
struct point {
    double x, y;
};

class shape {
public:
    shape( );
    virtual ~shape( );
    virtual void draw( );
};

class circle : public shape {
public:
    circle(const point& center, double radius);
    point center( ) const;
    double radius( ) const;
    void move_to(const point& new_center);
    void resize(double new_radius);
    virtual void draw( );
private:
    point center_;
```

```
    double radius_;
};

class integer {
public:
    typedef int value_type;
    int value( ) const;
    void set_value(int value);
    std::string to_string( ) const;
private:
    int value_;
};

class real {
public:
    typedef double value_type;
    double value( ) const;
    void set_value(double value);
    std::string to_string( ) const;
private:
    double value_;
};

union number {
    number(int value);
    number(double value);
    integer i;
    real r;
};
```

## Tip

In C, a struct or union name is not automatically a type name, as it is in C++. A variable declaration in C, for example, would need to be elaborated as `struct demo x` instead of simply `demo x`. There is no harm in using the fully elaborated form in C++, and this is often done in headers that are meant to be portable between C and C++.

There are times, even in C++, when the fully elaborated form is needed, such as when a function or object declaration hides a class name. In this case, use the elaborated form to refer to the class name. For example, the POSIX API has a structure named `stat`, which is hidden by a function `stat`:

```
extern "C" int stat(const char* filename, struct stat* st);
int main(int argc, char* argv[])
{
    struct stat st;
    if (argc > 1 && stat(argv[1], &st) == 0) show_stats(st);
}
```

See [Chapter 2](#) for more information about elaborated type specifiers.

## Plain Old Data

Some programming languages differentiate between records and classes. Typically, a record is a simple storage container that lacks the more complex features of a class (inheritance, virtual functions, etc.). In C++, classes serve both purposes, but you can do things with simple classes (called POD, for plain old data) that you cannot do with complicated classes.



Basically, a POD class is a structure that is compatible with C. More precisely, a POD class or union does not have any of the following:

- User-defined constructors
- User-defined destructor
- User-defined copy assignment operator
- Virtual functions
- Base classes
- Private or protected nonstatic members
- Nonstatic data members that are references

Also, all nonstatic data members must have POD type. A POD type is a fundamental type, an enumerated type, a POD class or union, or a pointer to or array of POD types

Unlike C structures, a POD class can have static data members, nonvirtual functions, and nested types (members that do not affect the data layout in an object).

POD classes are often used when compatibility with C is required. In that case, you should avoid using any access specifier labels because they can alter the layout of data members within an object. (A POD class cannot have private or protected members, but you can have multiple `public:` access specifier labels and still have a class that meets the standard definition of a POD class.)

[Example 6-2](#) shows POD types (`point1` and `info`) and non-POD types (`point2` and `employee`).

*Example 6-2. Declaring POD and non-POD types*

```
struct point1 {                                // POD
    int x, y;
};

class point2 {                                  // Not POD
public:
```

```

    point2(int x, int y);
private:
    int x, y;
};

struct info {                                // POD
    static const int max_size = 50;
    char name[max_size];
    bool is_name_valid( ) const;
    bool operator<(const info& i); // Compare names.
};

struct employee : info {                    // Not POD
    int salary;
};

```

The virtue of a POD object is that it is just a contiguous area of storage that stores some value or values. Thus, it can be copied byte for byte and retain its value. In particular, a POD object can be safely copied to an array of char or unsigned char, and when copied back, it retains its original value. A POD object can also be copied by calling memcopy; the copy has the same value as the original.

A local POD object without an initializer is left uninitialized, but a non-POD object is initialized by calling its default constructor. Similarly, a POD type in a new expression is uninitialized, but a new non-POD object is initialized by calling its default constructor. If you supply an empty initializer to a new expression or other expression that constructs a POD object, the POD object is initialized to 0.

A POD class can contain padding between data members, but no padding appears before the first member. Therefore, a pointer to the POD object can be converted (with reinterpret\_cast<>) into a pointer to the first element.

A goto statement can safely branch into a block, skipping over declarations of uninitialized POD objects. A goto that skips any other declaration in the block results in undefined behavior. (See [Chapter 2](#) for more information

about initializing POD objects, and [Chapter 4](#) for more information about the goto statement.)

## Trivial Classes

A *trivial class* is another form of restricted class (or union). It cannot have any of the following:

- User-defined constructors
- User-defined destructor
- User-defined copy assignment operator
- Virtual functions
- Virtual base classes

Also, all base classes must be trivial, and all nonstatic data members must be trivial or have non-class type. Unlike POD classes, a trivial class can have base classes, private and protected members, and members with reference type.

Trivial classes are important only because members of a union must be trivial. Fundamental types are trivial, as are pointers to and arrays of trivial types.

## Unions

A union is like a struct, but with the following restrictions:

- It cannot have base classes.
- It cannot be a base class.
- It cannot have virtual functions.
- It cannot have static data members.
- Its data members cannot be references.
- All of its data members must be trivial.

An object of union type has enough memory to store the largest member, and all data members share that memory. In other words, a union can have a value in only one data member at a time. It is your responsibility to keep track of which data member is "active."

A union can be declared without a name (an *anonymous* union), in which case it must have only nonstatic data members (no member functions, no nested types). The members of an anonymous union are effectively added to the scope in which the union is declared. That scope must not declare any identifiers with the same names as the union's members. In this way, an anonymous union reduces the nesting that is needed to get to the union members, as shown in [Example 6-3](#).

### *Example 6-3. An anonymous union*

```
struct node {
    enum kind { integer, real, string } kind;
    union {
        int intval;
        double realval;
        *char strval[8];
    };
};

node* makeint(int i)
{
    node* rtn = new node;
    rtn->kind = node::integer;
    rtn->intval = i;
    return rtn;
}
```

## Tip

Unions are used less often in C++ than in C. Two common uses for unions in C have been supplanted with other alternatives in C++:

- Unions can be used to peek inside the data representation for certain types, but a `reinterpret_cast<>` can do the same thing.
- Unions were sometimes used to save memory when storing different kinds of data in a structure, but inheritance is safer and offers more flexibility.

## Local Classes

You can declare a local class, that is, a class definition that is local to a block in a function body. A local class has several restrictions when compared to nonlocal classes:

- A local class cannot have static data members.
- Member functions must be defined inline in the class definition.
- You cannot refer to nonstatic objects from within a local class, but you can refer to local static objects, local enumerators, and functions that are declared locally.
- A local class cannot be used as a template argument, so you cannot use a local functor with the standard algorithms.

Local classes are not used often. [Example 6-4](#) shows one use of a local class.

### *Example 6-4. A local class*

```
// Take a string and break it up into tokens, storing the tokens
// in a vector.
void get_tokens(std::vector<std::string>& tokens, const
std::string& str)
```

```
{
class tokenizer {
public:
    tokenizer(const std::string& str) : in_(str) {}
    bool next( )
    {
        return in_ >> token_;
    }
    std::string token( ) const { return token_; }
private:
    std::istringstream in_;
    std::string token_;
};

tokens.clear( );
tokenizer t(str);
while (t.next( ))
    tokens.push_back(t.token( ));
}
```

## Data Members

A class can have static and nonstatic data members. The static storage class specifies a static data member; with no storage class, a data member is nonstatic. No other storage class specifier is allowed. Every object has its own copy of the class's nonstatic data members, and they share a single copy of each static data member. A data member can also be declared with *cv*-qualifiers. See [Chapter 2](#) for more information about storage class specifiers and *cv*-qualifiers.

You declare data members as you would local variables in a function, except you cannot usually supply initializers. Instead, nonstatic data members are initialized in the class's constructors. See [Section 6.3](#) later in this chapter for details.

Data members are typically declared at the `private` access level. See [Section 6.5](#) later in this chapter for details.

## Data Layout



Nonstatic data members are organized so that members declared later have higher addresses than those declared earlier. Access specifier labels, however, can interrupt the order, and the relative order of members separated by an access specifier label is implementation-defined. Writing code that depends on the layout of a class is usually a bad idea, but when interfacing with external code, it is sometimes unavoidable. See [Section 6.1.1](#) earlier in this chapter for more information.

The layout of base-class subobjects within a derived-class object is also unspecified. If a base class appears more than

once in the inheritance graph, it will have multiple copies of its data members in the derived-class object, unless the inheritance is virtual. See [Section 6.4](#) later in this chapter for more information.

Individual data members might have specific alignment requirements, so the compiler can insert padding between members. Thus, adjacent member declarations might not have adjacent addresses.

Even if a class has no nonstatic data members, the size of the object is guaranteed to be greater than zero. If the class is used as a base class, however, the compiler can optimize the base-class subobject so it has a size of 0 within the derived class.

## Mutable Data Members

A nonstatic data member can be declared with the `mutable` type specifier. A mutable member can be changed even when the object is `const`.

Mutable members are often used to implement a class that can have objects that are logically constant, even if they are not physical constants. For example, a class can implement a private cache with a mutable data member. Suppose you are writing a class, `bigint`, that implements very large integers. The `to_double` member function computes an approximation of the value as a floating-point number. Instead of computing this value each time `to_double` is called, the `bigint` class saves the value after the first call and returns the cached value for the second and subsequent calls. When calling `to_double` for a `const bigint`, you still want to be able to modify the `bigint` object to cache the floating-point value, so the cache is declared mutable, as shown in [Example 6-5](#).

*Example 6-5. Caching a value in a mutable data member*



```

class bigint {
public:
    bigint( ) : has_approx_(false) { ... }
    double to_double( )
    const
    {
        if (! has_approx_) {
            approx_ = as_double( );
            has_approx_ = true;
        }
        return approx_;
    }
    ...
private:
    double as_double( ) const;
    mutable double approx_;
    mutable bool has_approx_;
    ...
};

```

## Bit-Fields

A nonstatic data member can be a *bit-field*, which is a sequence of bits packed into an object. The declarator for a bit-field uses a colon followed by an integral constant expression, which specifies the number of bits in the field. [Example 6-6](#) shows the layout of the control word for an Intel x87 floating-point processor. (Whether this struct definition accurately maps to the actual control word layout is implementation-defined, but using the x87 control word is inherently nonportable, anyway.)

### *Example 6-6. Using bit-fields for a control word*

```

// Intel x87 FPU control word.
struct fpu_control {
    enum precision_control { single, double_prec=2, extended };
    enum rounding_control { nearest, down, up, truncate };
    int
        : 4; // Reserved
    rounding_control round_ctl : 2;
    precision_control prec_ctl : 2;
    int
        : 2; // Reserved
    bool precision : 1;
};

```

```
bool        underflow    : 1;
bool        overflow     : 1;
bool        zero_divide  : 1;
bool        denormal     : 1;
bool        invalid_op   : 1;
};
```

Use a bit-field as you would any other data member, but with the following caveats:

- You cannot take the address of a bit-field.
- You cannot bind a bit-field to a non-const reference.
- When you bind a bit-field to a const reference, the compiler creates a temporary object and binds that to the reference.
- Whether a bit field is signed or unsigned is implementation defined unless you explicitly declare it with the `signed` or `unsigned` specifier. Thus, the bit-field `int bf : 1;` might take the values `0` and `1` or `-1` and `0` (two's complement), or even `-0` and `+0` (signed magnitude).
- A bit-field size can be larger than the declared type of the member, in which case the excess bits are used as padding and are not part of the member's value.
- Order and alignment of bit-fields are implementation-defined.

A bit-field without a name cannot be used or referred to, and is typically used for padding. A nameless bit-field can have size `0`, which pads the object so the next bit-field is aligned on a natural memory boundary.

If you need to work with large bit-fields or treat a set of bits as an object, consider using the `bitset` template. (See `<bitset>` in [Chapter 13](#).)

## Static Data Members

A static data member is similar to an object declared at namespace scope; the class name assumes the role of the namespace name. In other words, there is a single copy of each static data member regardless of the number of instances of the class. Derived classes also share the single static data member. In some languages, static data members are called class variables.

The lifetime of a static data member is similar to the lifetime of a global static object. See [Chapter 2](#) for details.

Local classes and anonymous unions cannot have static data members. A static data member cannot be mutable.

The member declaration in the class definition is just a declaration. You must supply a definition elsewhere in the program. The definition can have an initializer.

Only an integral or enumerated static data member can have an initializer in the class definition. The initializer must be a constant integral expression. The value of the member can be used as a constant elsewhere in the class definition. The definition of the member must then omit the initializer. This feature is often used to define the maximum size of data member arrays, as shown in [Example 6-7](#).

#### *Example 6-7. Declaring static data members*

```
// filename.h
class filename
{
public:
    static const int max_length = 1024;
    static const char path_separator = '/';
    static filename current_directory;
    filename(const char* str);
    ...
private:
    char name_[max_length];
};

// filename.cpp
// Definitions for the static data members and member functions
```

```
const int filename::max_length;
const char filename::path_separator;
filename filename::current_directory(".");

filename::filename(const char* str)
{
    strncpy(name_, str, max_length-1);
}
...
```

## Member Functions

Member functions implement the behavior of a class. Member functions can be defined within the class definition or separately. You can use the `inline` function specifier and either the `static` or `virtual` (but not both) specifier. (See [Chapter 2](#) for more about function specifiers.) Defining a member function within the class definition declares the function inline, even if you do not use the `inline` specifier.

A nonstatic member function can have `const`, `volatile`, or both function qualifiers. Qualifiers appear after the function parameters and before the exception specification. Function qualifiers are discussed in the next section, [Section 6.3.2](#).

[Example 6-8](#) shows various member function declarations and definitions.

### *Example 6-8. Declaring and defining member functions*

```
#include <cmath>
#include <iostream>
#include <istream>
#include <ostream>

class point {
public:
    typedef double value_type;
    // Constructors are special member functions.
    explicit point(value_type x = 0.0, value_type y = 0.0);
    value_type x( ) const { return x_; }
    value_type y( ) const { return y_; }
    void x(value_type x) { x_ = x; }
    void y(value_type y) { y_ = y; }

    value_type distance( ) const;
    bool operator==(const point& pt) const;

    inline static point origin( );
private:
    value_type x_, y_;
```

```

};

point::point(value_type x, value_type y)
: x_(x), y_(y)
{}

point::value_type point::distance( )
const
{
    return std::sqrt(x() * x() + y( ) * y( ));
}

bool point::operator==(const point& pt)
const
{
    return x() == pt.x() && y( ) == pt.y( );
}

inline point point::origin( )
{
    return point( );
}

int main( )
{
    point p1;
    point::value_type n;
    std::cin >> n;
    p1.x(n);
    std::cin >> n;
    p1.y(n);
    if (p1 == point::origin( ))
        std::cout << "Origin\n";
    else
        std::cout << p1.distance( ) << '\n';
}

```

When defining a function inside a class definition, the entire function definition is in the scope of the class. Thus, name lookup for the return type and all parameter types looks first in the class, then in all base classes, and then in the surrounding namespaces. (See [Chapter 2](#) for more information about name lookup.)

If the function definition is outside the class definition, the function return type is at namespace scope. Thus, if the

type is a nested type in the class, you must fully qualify the type name. The function name must also be qualified so the compiler knows which class contains the function. After the compiler has seen the class name that qualifies the function name, it enters the class scope. Parameter types are then looked up in the class scope.

## Pointers-to-Members

The address of a static member is no different than the address of an ordinary function or object at namespace scope. The address of a nonstatic member that is taken in a member function with an unqualified name is also an ordinary address. The address of a nonstatic member taken with a qualified member name (e.g., `&cls::mem`), however, is quite different from any other kind of address. The address even has a special name: *pointer-to-member*. You cannot use `reinterpret_cast<>` to cast a pointer-to-member to or from an ordinary data pointer or function pointer.

A pointer-to-member does not necessarily point to a particular function or object. Instead, think of it as a handle that keeps track of which function or data member you have selected but does not refer to a specific object. To use a pointer-to-member, you must use the `.*` or `->*` operator, which requires an object or pointer as the lefthand operand and the pointer-to-member as the righthand operand. The operator looks up the member in the object and then obtains the actual data member or member function. (See [Chapter 3](#) for more information about pointer-to-member expressions.) [Example 6-9](#) shows one use of a pointer-to-member.

### *Example 6-9. A pointer-to-member*

```
#include <iostream>
#include <ostream>
```

```

class base {
public:
    base(int i) : x_(i) {}
    virtual ~base( ) {}
    virtual void func( ) { std::cout << "base::func( )\n"; }
private:
    int x_;
};

class derived : public base {
public:
    derived(int i) : base(i) {}
    virtual void func( ) { std::cout << "derived::func( )\n"; }
};

int main( )
{
    base *b = new derived(42);    void (base::*fp)( ) = &base::func;
                                (b->*fp)( );                // Prints
    derived::func( )
}

```

## this Pointer

A nonstatic member function can be called only for an object of its class or for a derived class. The object is implicitly passed as a hidden parameter to the function, and the function can refer to the object by using the `this` keyword, which represents an rvalue pointer to the object. That is, if the object has type `T`, `this` is an rvalue of static type `T*`. In a call to a virtual function, the object's dynamic type might not match the static type of `this`.

Static member functions do not have `this` pointers. (See the next section, [Section 6.3.3](#).)

If the function is qualified with `const` or `volatile`, the same qualifiers apply to `this` within the member function. In other words, within a `const` member function of class `T`, `this` has type `const T*`. A `const` function, therefore, cannot



modify its nonstatic data members (except those declared with the `mutable` specifier).

Within a member function, you can refer to data members and member functions using just their names, and the compiler looks up the unqualified names in the class, in its ancestor classes, and in namespace scopes. (See [Chapter 2](#) for details.) You can force the compiler to look only in the class and ancestor classes by explicitly using `this` to refer to the members—for example, `this->data` and `this->func ( )`. (When using templates, an explicit member reference can reduce name lookup problems; see [Chapter 7](#) for details.)

[Example 6-10](#) shows some typical uses of `this`.

*Example 6-10. The `this` keyword*

```
class point {
public:
    point(int x=0, int y=0) : x_(x), y_(y) {}
    int x( ) const { return this->x_; }
    int y( ) const { return y_; }
    void x(int x) { this->x_ = x; }
    void y(int y) { y_ = y; }
    bool operator==(const point& that) {
        return this->x( ) == that.x( ) &&
            this->y( ) == that.y( );
    }
    void move_to(context*) const { context->move_to(*this); }
    void draw_to(context*) const { context->draw_to(*this); }
    void get_position(context*) { context->getpos(this); }
private:
    int x_, y_;
};
```

## Static Member Functions

A static member function is like a function declared at namespace scope; the class name assumes the role of the namespace name. Within the scope of the class, you can call the function using an unqualified name. Outside the

class scope, you must qualify the function name with the class name (e.g., `cls::member`) or by calling the function as a named member of an object (e.g., `obj.member`). In the latter case, the object is not passed as an implicit parameter (as it would be to a nonstatic member function), but serves only to identify the class scope in which the function name is looked up.

Static member functions have the following restrictions:

- They do not have this pointers.
- They cannot be `virtual`.
- They cannot have `const` or `volatile` qualifiers.
- They cannot refer to nonstatic members, except as members of a specific object (using the `.` or `->` operator).
- A class cannot have a static and nonstatic member function with the same name and parameters.

[Example 6-11](#) shows some uses of static member functions.

*Example 6-11. Static member functions*

```
class point {
public:
    point(int x, int y);

    static point origin( ) { return point(0, 0); }

    // Error: calls non-static function, x( )
    static bool is_zero( ) { return x( ) == 0; }

    int x( ) const;
    int y( ) const;
private:
    int x_, y_;
};
```

## Constructors

Constructors and destructors are special forms of member functions. A constructor is used to initialize an object, and a destructor is used to finalize an object.

## Declaring constructors

A constructor's name is the same as the class name. If you use typedef to create a synonym for the class name, you cannot use the typedef name as the constructor name:

```
struct point {
    point(int x, int y);
    ...
typedef point p;

p::point(int x, int y)    { . . . } // OK
point::point(int x, int y) { . . . } // OK
p::p(int x, int y)       { . . . } // Error
point::p(int x, int y)   { . . . } // Error
```

A constructor cannot have a return type, and you cannot return a value. That is, you must use a plain `return;` or return an expression of type `void`. A constructor cannot have `const` or `volatile` qualifiers, and it cannot be `virtual` or `static`.

Constructors can initialize data members with a list of member initializers, which appear after the function header but before the function body. A colon separates the header from the comma-separated initializers. Each initializer names a data member, an immediate base class, or a virtual base class. The value of the initializer follows in parentheses:

```
class-name(parameters)
: member(expr), base-class(expr-list), ...
compound-statement
```

You can also use a function try block, which wraps the constructor initializers as part of the try block. (See [Chapter 4](#) for more information about try blocks and exception handlers.) The syntax is:

```
class-name(parameters)
try
  : member(expr), base-class(expr-list), ...
  compound-statement
  exception-handlers
```

You can initialize a base class by invoking any of its constructors, passing zero or more expressions as arguments. See [Section 6.4](#) later in this chapter for more information.

### **Tip**

The order in which initializers appear in the initializer list is irrelevant. The order of initialization is determined by their declarations according to the following rules:

- Virtual base classes are initialized first, in order of declaration in a depth-first traversal of the inheritance graph. Each virtual base class is initialized exactly once for the most-derived object.
- Nonvirtual, direct base classes are initialized next, in order of declaration.
- Nonstatic data members are initialized next, in order of declaration.
- Finally, the body of the constructor is executed.

Note that each base class is constructed according to the same rules. This causes the root of the inheritance tree to be initialized first, and the most-derived class to be initialized last.

A constructor can be declared with the `explicit` specifier. An explicit constructor is never called for an implicit type conversion, but only for function-like initialization in a declaration and for explicit type casts. See the next section for more information.

## Special constructors

Two kinds of constructors are special, so special they have their own names: *default constructors* and *copy constructors*.

A default constructor can be called with no arguments. It might be declared with no parameters, default arguments for every parameter, or an ellipsis as the sole parameter. (The last case is uncommon.) Default constructors are called when you construct an array of objects, when you omit the initializer for a scalar object, or when you use an empty initializer in an expression:

```
point corners[2];
point p;
point *ptr = new point( );
```

(Remember not to use an empty initializer in a declaration. For example, `point p( );` declares a function named `p`, not a default-initialized point object. See [Chapter 2](#) for details.)

A copy constructor can take a single argument whose type is a reference to the class type. (Additional parameters, if any, must have default arguments.) The reference is usually `const`, but it does not have to be. The copy constructor is called to copy objects of the class type when passing objects to functions and when returning objects from functions.

The following example shows a default and a copy constructor:

```
struct point {
    point( );                // Default constructor
    point(const point& pt);  // Copy constructor
    ...
};
```

Default and copy constructors are so important, the compiler might generate them for you. See [Section 6.3.6](#) later in this chapter for details.

## Calling constructors

Constructors are special because you never call them directly. Instead, the compiler calls the constructor as part of the code it generates to create and initialize an object of class type. Objects can be created in several different ways:

- With automatic variables and constants in a function body:

```
void demo( )
{
    point p1(1, 2);
    const point origin;
    ...
}
```

- With static objects that are local or global:

```
point p2(3, 4);
void demo( )
{
    static point p3(5, 6);
    ...
}
```

- Dynamically, with new expressions:

```
point *p = new point(7, 8);
```

- With temporary objects in expressions:

```
set_bounds(point(left, top), point(right, bottom));
```

- With function parameters and return values:

```
point offset(point p) { p.x(p.x( ) + 2); return p; }
point x(1,2);
x = offset(x); // Calls point(const point&) twice
```

- With implicit type conversions:

```
point p = 2; // Invokes point(int=0, int=0), then
point(const point&)
```

- With constructor initializers:

```
derived::derived(int x, int y) : base(x, y) {}
```

In each case, memory is set aside for the object, then the constructor is called. The constructor is responsible for

initializing the object's members. Members that are listed in the initializer list are initialized as requested: scalar and pointer types are initialized with their given values; class-type objects are initialized by invoking the appropriate constructors. Class-type objects that are not listed are initialized by calling their default constructors; other types are left uninitialized. Every const member and reference-type member must have an initializer (or you must be willing to accept the default constructor for a class-type const member). Within the body of the constructor, you are free to change any non-const member. Many simple constructors are written with empty bodies:

```
struct point {
    point(int x, int y) : x_(x), y_(y) {}
    point( )           : x_(0), y_(0) {}
    point(double radius, double angle) {
        x = radius * cos(angle);
        y = radius * sin(angle);
    }
    ...
}
```



If a constructor is declared with the `explicit` specifier, it can never be called for an implicit type conversion. An implicit type conversion is commonly used during assignment-like construction. In the following example, two constructors are called: `complex(1.0, 0.0)` is implicitly called to construct a nameless temporary object, then the copy constructor is called to copy the temporary object into `p`. (The compiler is allowed to optimize away the copy and initialize `p` directly. Most modern compilers perform this optimization, but `point` must have a copy constructor even if it is never called.)

```
struct complex {
    complex(double re, double im = 0.0);
    ...
};
complex z = 1;
```

Implicit type conversions can produce some surprising results by calling conversion constructors when you least expect it. In the following example, a nameless temporary complex object is created as `complex(42.0, 0.0)`, and this temporary object is bound to the `z` parameter in the call to `add2`.

```
complex add2(const complex& z) {
    z.real(z.real( ) + 2);
    return p;
}
z = add2(42.0);
```

To avoid unexpected constructor calls, declare a constructor with the `explicit` specifier if that constructor can be called with a single argument. Calls to an explicit constructor require function-like construction or an explicit type cast:

```
struct complex {
    explicit complex(double re, double im = 0.0);
    ...
};
complex z = 1; // Error
complex z(1); // OK
add2(42); // Error
add2(static_cast<complex>(42)); // OK
add2(p); // OK
```

## Throwing exceptions during construction

A constructor can also use a function try block as its function body. The `try` keyword comes before the initializers. If an exception is thrown during the initialization of any member, the corresponding catch blocks are examined for matching handlers in the usual manner. (See [Chapter 4](#) for information about try blocks and exception handlers.) [Example 6-12](#) shows a constructor with a function try block.

*Example 6-12. Catching exceptions in constructors*



```

struct array {
    // Initialize the array with all the items in the range (first,
    last). If an
    // exception is thrown while copying the items, be sure to
    clean up by
    // destroying the items that have been added so far.
    template<typename InIter>
    array(InIter first, InIter last)
    try
        : data_(0), size_(0)
    {
        for (InIter iter(first); iter != last; ++iter)
            push_back(*iter);
    } catch (...) {
        clear( );
    }
    ...
private:
    int* data_;
    std::size_t size_;
};

```

Regardless of whether the function body is a plain compound statement or a function try block, the compiler keeps track of which base-class subobjects and which members have been initialized. If an exception is thrown during initialization, the destructors of only those objects that have been fully constructed are called.

If the object is being created as part of a new expression, and an exception is thrown, the object's memory is deallocated by calling the appropriate deallocation function. If the object is being created with a placement new operator, the corresponding placement delete operator is called—that is, the delete function that takes the same additional parameters as the placement new operator. If no matching placement delete is found, no deallocation takes place. (See [Chapter 3](#) for more information on new and delete expressions.)

[Example 6-13](#) shows a silly class that allocates two dynamic arrays in its constructor. Suppose the first allocation fails (member `str_`). In that case, the compiler

knows that `str_` has not been initialized and so does not call the `auto_ptr<>` destructor, nor does the compiler call the destructor for the `wstr_` member. If, however, the allocation for `str_` succeeds and fails for `wstr_`, the compiler calls the `auto_ptr<>` destructor for `str_` to free the memory that had been allocated successfully. Note that if `auto_ptr<>` were not used, the class would have a memory leak because the compiler would not be able to free the memory for `str_`. (Pointers don't have destructors.)

### *Example 6-13. Exception-safe constructor*

```
struct silly {
    silly(std::size_t n)
        : size_(n), str_(new char[n+1]), wstr_(new wchar_t[n+1])
    {}
    silly(const silly& that);
    silly& operator=(silly that);
private:
    std::size_t size_;
    std::auto_ptr<char> str_;
    std::auto_ptr<wchar_t> wstr_;
};
```

## **Destructors**

A destructor finalizes an object when the object is destroyed. Typically, finalization frees memory, releases resources (such as open files), and so on. A destructor is a special member function. It has no return type, no arguments, and cannot return a value. A destructor can be virtual, it can be inline, but it cannot be static, const, or volatile. The name of a destructor is a tilde (~) followed by the class name. A class has only one destructor (although it may have several constructors).

Just as constructors automatically construct base-class objects, a destructor automatically calls destructors for nonstatic data members, direct base classes, and virtual

base classes after the derived-class destructor body returns. Member and base-class destructors are called in reverse order of their constructors. If you do not write a destructor, the compiler does so for you. (See the next section, [Section 6.3.6](#).) Any class that has virtual functions should have a virtual destructor, as explained in [Section 6.4](#) later in this chapter.

You can call a destructor explicitly, but there is rarely any reason to do so. The compiler calls destructors automatically when objects are destroyed. Dynamically allocated objects are destroyed by `delete` expressions. Static objects are destroyed when a program terminates. Other named objects are destroyed automatically when they go out of scope. Temporary, unnamed objects are destroyed automatically at the end of the expression that created them.

When you write a destructor, make sure that it never throws an exception. When an exception causes the stack to be unwound (see [Chapter 5](#)), objects are automatically destroyed. If a destructor throws an exception during stack unwinding, the program terminates immediately. (See `<exception>` in [Chapter 13](#).)

[Example 6-14](#) shows how a simple destructor is declared.

*Example 6-14. Declaring a destructor*

```
class string {
public:
    string(std::size_t n, char c = '\0')
        : str_(new char[n+1]), size_(n) {
        std::fill(str_, str_+size_+1, c);
    }
    string( ) : str_(new char[1]), size_(0) {
        str_[0] = '\0';
    } ~string( ) { delete[] str_; }
    const char* c_str( ) const { return str_; }
    std::size_t size( ) const { return size_; }
private:
    std::size_t size_;
```

```
char* str_;  
};
```

## Implicit Member Functions

The compiler implicitly declares and defines default and copy constructors, the copy assignment operator, and the destructor in certain circumstances. All these implicit functions are inline public members. This section describes what each function does and the circumstances under which the function is implicitly declared and defined. [Example 6-15](#) explicitly shows how each implicit member function is defined.

### *Example 6-15. Implicit member functions*

```
// These classes show explicitly what the implicit member  
// functions would be.  
class base {  
public:  
    base( ) {}  
    base(const base& that) : m1_(that.m1_), m2_(that.m2_) {}  
    base& operator=(const base& that) {  
        this->m1_ = that.m1_;  
        this->m2_ = that.m2_;  
        return *this;  
    }  
    ~base( ) {}  
private:  
    int m1_;  
    char* m2_;  
};  
  
class demo {  
public:  
    demo( ) {} // Default constructs three base objects  
    demo(demo& d) {} // Copies three base objects in ary_[]  
    demo& operator=(const demo& that) {  
        this->ary_[0] = that.ary_[0];  
        this->ary_[1] = that.ary_[1];  
        this->ary_[2] = that.ary_[2];  
        return *this;  
    }  
    ~demo( ) {} // Default destructs three base objects  
private:
```

```

    base ary_[3];
};

class derived : public base {
public:
    derived( ) : base( ) {} // Constructs m3_[]
    derived(derived& that) : base(that) {} // Copies m3_[]
    derived& operator=(const derived& that) {
        static_cast<base&>(*this) =
            static_cast<const base&>(that);
        this->m3_[0] = that.m3_[0];
        this->m3_[1] = that.m3_[1];
        this->m3_[2] = that.m3_[2];
        return *this;
    }
    ~derived( ) {} // Calls ~base( ), destructs 3 demo objects
private:
    demo m3_[3];
};

```

For classes whose data members have fundamental, class, or enumeration types, the implicit functions are often adequate. The most common case in which you must implement these functions explicitly is when an object manages pointers to memory that the object controls. In this case, a copy constructor or copy assignment operator must not blindly copy a member that is a pointer, which results in two pointers to the same memory. Instead, you should allocate a new pointer and copy the contents. In such cases, you will often find yourself providing the copy constructor, copy assignment operator, and destructor.

### **Tip**

A useful guideline is that if you write one of the three special functions (copy constructor, copy assignment operator, or destructor), you will probably need to write all three.

If you want to store objects in a standard container, you must make sure it has a copy constructor and a copy

assignment operator (implicit or explicit). See [Chapter 10](#) for information about the standard containers.

### **I m p l i c i t   d e f a u l t   c o n s t r u c t o r**

The compiler implicitly declares a default constructor if a class has no user-defined constructors. The implicit default constructor calls the default constructor for all base classes and for all nonstatic data members of class type. Other nonstatic data members are left uninitialized. In other words, the behavior is the same as if you wrote the default constructor with no initializers and an empty function body.

### **I m p l i c i t   c o p y   c o n s t r u c t o r**

The compiler implicitly declares a copy constructor if a class has no copy constructor. If every direct base class and virtual base class has a copy constructor that takes a `const` reference parameter, and every nonstatic data member has a copy constructor with a `const` reference parameter, the implicit copy constructor also takes a `const` reference parameter. Otherwise, the implicit copy constructor takes a plain reference parameter.

In other words, the compiler tries to declare the implicit copy constructor so it takes a `const` reference parameter. If it cannot because the implicit function would end up calling an inherited or member copy constructor that does not take a `const` parameter, the compiler gives up and declares a copy constructor that takes a non-`const` parameter.

An implicit copy constructor calls the copy constructor for each direct base class and each virtual base class and then performs a member-by-member copy of all nonstatic data members. It calls the copy constructor for each member of class type and copies the values for members of nonclass type.

### **I m p l i c i t   c o p y   a s s i g n m e n t   o p e r a t o r**

The compiler implicitly declares a copy assignment operator (operator=) if a class does not have one. If every direct base class and virtual base class has a copy assignment operator that takes a const reference parameter, and if every nonstatic data member has a copy assignment operator with a const reference parameter, the implicit copy assignment operator also takes a const reference parameter. Otherwise, the implicit copy assignment operator takes a plain reference parameter.

In other words, the compiler tries to declare the implicit copy assignment operator, so it takes a const reference parameter. If it cannot because the implicit function would end up calling an inherited or member copy assignment operator that does not take a const parameter, the compiler gives up and declares a copy assignment operator that takes a non-const parameter.

An implicit copy assignment operator calls the copy assignment operator for each direct base class and each virtual base class and then performs a member-by-member assignment of all nonstatic data members. It calls the copy assignment operator for each member of class type and assigns the values for members of nonclass type.

### **I m p l i c i t   d e s t r u c t o r**

The compiler declares an implicit destructor if the programmer does not provide one. If a base class has a virtual destructor, the implicit destructor is also virtual. The implicit destructor is like a programmer-supplied destructor with an empty function body.

# Inheritance

A class can inherit from zero or more *base classes*. A class with at least one base class is said to be a *derived class*. A derived class inherits all the data members and member functions of all of its base classes and all of their base classes, and so on. A class's immediate base classes are called *direct base classes*. Their base classes are *indirect base classes*. The complete set of direct and indirect base classes is sometimes called the *ancestor classes*.

A class can derive directly from any number of base classes. The base-class names follow a colon and are separated by commas. Each class name can be prefaced by an access specifier (described later in this chapter). The same class cannot be listed more than once as a direct base class, but it can appear more than once in the inheritance graph. For example, `derived3` in the following code has `base2` twice in its inheritance tree, once as a direct base class, and once as an indirect base class (through `derived2`):

```
class base1 { ... };
class derived1 : public base1 { ... };
class base2 { ... }
class derived2 : public derived1, public base2 { ... }
class derived3 : protected derived2, private base2 { ... }
```

A derived class can access the members that it inherits from an ancestor class, provided the members are not private. (See [Section 6.5](#) later in this chapter for details.) To look up a name in class scope, the compiler looks first in the class itself, then in direct base classes, then in their direct base classes, and so on. See [Chapter 2](#) for more information about name lookup.

To resolve overloaded functions, the compiler finds the first class with a matching name and then searches for



overloads in that class. [Chapter 5](#) lists the complete rules for overload resolution.

An object with a derived-class type can usually be converted to a base class, in which case the object is *sliced*. The members of the derived class are removed, and only the base class members remain:

```
struct file {
    std::string name;
};
struct directory : file {
    std::vector<file*> entries;
};
directory d;
file f;
f = d; // Only d.name is copied to f; entries are lost.
```

Slicing usually arises from a programming error. Instead, you should probably use a pointer or reference to cast from a derived class to a base class. In that case, the derived-class identity and members are preserved. This distinction is crucial when using virtual functions, as described in the next section. For example:

```
directory* dp = new directory;
file* fp;
fp = dp; // Keeps entries and identity as a directory object
```

As you can see in the previous examples, the compiler implicitly converts a derived class to a base class. You can also use an explicit cast operator, but if the base class is virtual, you must use `dynamic_cast<>`, not `static_cast<>`. See [Chapter 3](#) for more information about cast expressions.

If a base class is ambiguous (see [Section 6.4.5](#) later in this chapter) or inaccessible (see [Section 6.5](#) later in this chapter), you cannot slice a derived class to the base class, nor can you cast a pointer or reference to the inaccessible base class (an old-style cast allows it but such use is not recommended).

## Virtual Functions

A nonstatic member function can be declared with the `virtual` function specifier, and is then known as a *virtual function*. A virtual function can be *overridden* in a derived class. To override a virtual function, declare it in a derived class with the same name and parameter types. The return type is usually the same but does not have to be identical (as described in the next section, [Section 6.4.2](#)). The `virtual` specifier is optional in the derived class but is recommended as a hint to the human reader. A constructor cannot be virtual, but a destructor can be. A virtual function cannot be static.

A class that has at least one virtual function is *polymorphic*. This form of polymorphism is more precisely known as *type polymorphism*. (C++ also supports *parametric polymorphism* with templates; see [Chapter 7](#).) Most programmers mean type polymorphism when they talk about object-oriented programming.

When calling virtual functions, you must distinguish between the declared type of the object, pointer, or reference and the actual type at runtime. The declared type is often called the static type, and the actual type is the dynamic type. For example:

```
struct base {
    virtual void func( );
};
struct derived : base {
    virtual void func( ); // Overload
};
base* b = new derived;    // Static type of b is base*.
                        // Dynamic type is derived*.
b->func( );              // Calls dynamic::func( )
```

When any function is called, the compiler uses the static type to pick a function signature. If the function is virtual, the compiler generates a virtual function call. Then, at

runtime, the object's dynamic type determines which function is actually called—namely, the function in the most-derived class that overrides the virtual function. This is known as a *polymorphic function call*.

### **Dispatching Virtual Functions**

Virtual functions are most commonly implemented using *virtual function tables*, or *vtables*. Each class that declares at least one virtual function has a hidden data member (e.g., `__vtbl`). The `__vtbl` member points to an array of function pointers. Every derived class has a copy of the table. Every instance of a class shares a common table. Each entry in the table points to a function in a base class, or if the function is overridden, the entry points to the derived class function. Any new virtual functions that the derived class declares are added at the end of the table.

When an object is created, the compiler sets its `__vtbl` pointer to the vtable for its dynamic class. A call to a virtual function is compiled into an index into the table and into a call to the function at that index. Note that the `dynamic_cast<>` operator can use the same mechanism to identify the dynamic type of the object.

Multiple inheritance complicates matters slightly, yet the basic concept remains the same: indirection through a table of pointers.

Compilers do not have to use vtables, but they are used so widely, the term "vtable" has entered the common parlance of many C++ programmers.

An object's dynamic type can differ from its static type only if the object is accessed via a pointer or reference. Thus, to call a virtual function, you typically access the target object

via a pointer (e.g., `ptr->func( )`). Inside a member function, if you call a virtual member function using its unqualified name, that is the same as calling the function via `this->`, so the function is called virtually. If a nonpointer, nonreference object calls a virtual function, the compiler knows that the static type and dynamic type always match, so it can save itself the lookup time and call the function in a nonvirtual manner.

[Example 6-16](#) shows a variety of virtual functions for implementing a simple calculator. A parser constructs a parse tree of `expr` nodes, in which each node can be a literal value or an operator. The operator nodes point to operand nodes, and an operand node, in turn, can be any kind of `expr` node. The virtual `evaluate` function evaluates the expression in the parse tree, returning a `double` result. Each kind of node knows how to evaluate itself. For example, a node can return a literal value or add the values that result from evaluating two operands.

### *Example 6-16. Declaring and using virtual functions*

```
class expr {
public:
    virtual ~expr( ) {}
    virtual double evaluate( ) const = 0;
    std::string as_string( ) const {
        std::ostringstream out;
        print(out);
        return out.str( );
    }
    virtual void print(std::ostream& out) const {}
    virtual int precedence( ) const = 0;
    template<typename charT, typename traits>
    static std::auto_ptr<expr> parse(
        std::basic_istream<charT,traits>& in);
};

// cout << *expr prints any kind of expression because expr-
// >print( ) is virtual.
template<typename charT, typename traits>
std::basic_ostream<charT,traits>&
operator<<(std::basic_ostream<charT,traits>& out, const expr&
```

```

e)
{
    e.print(out);
    return out;
}

class literal : public expr {
public:
    literal(double value) : value_(value) {}
    virtual double evaluate( ) const { return value_; }
    virtual void print(std::ostream& out) const {
        out << value_;
    }
    virtual int precedence( ) const { return 1; }
private:
    double value_;
};

// Abstract base class for all binary operators
class binop : public expr {
public:
    binop(std::auto_ptr<expr> left, std::auto_ptr<expr> right)
        : left_(left), right_(right) {}
    virtual double evaluate( ) const {
        return eval(left_->evaluate( ), right_->evaluate( ));
    }
    virtual void print(std::ostream& out) const {
        if (left_->precedence( ) > precedence( ))
            out << '(' << *left_ << ')';
        else
            out << *left_;

        out << op( );

        if (right_->precedence( ) > precedence( ))
            out << '(' << *right_ << ')';
        else
            out << *right_;
    }
    // Reminder that derived classes must override precedence
    virtual int precedence( ) const = 0;
protected:
    virtual double eval(double left, double right) const = 0;
    virtual const char* op( ) const = 0;
private:
    // No copying allowed (to avoid messing up auto_ptr<>s)
    binop(const binop&);
    void operator=(const binop&);

```

```

    std::auto_ptr<expr> left_;
    std::auto_ptr<expr> right_;
};

// Example binary operator.
class plus : public binop {
public:
    plus(std::auto_ptr<expr> left, std::auto_ptr<expr> right)
        : binop(left, right) {}
    virtual int precedence( ) const { return 3; }
protected:
    virtual double eval(double left, double right) const {
        return left + right;
    }
    virtual const char* op( ) const { return "+"; }
};

int main( )
{
    while(std::cin) {
        std::auto_ptr<expr> e(expr::parse(std::cin));
        std::cout << *e << '\n';
        std::cout << e->evaluate( ) << '\n';
    }
}

```

Sometimes you do not want to take advantage of the virtualness of a function. Instead, you may want to call the function as it is defined in a specific base class. In such a case, qualify the function name with the base-class name, which tells the compiler to call that class's definition of the function:

```

literal* e(new literal(2.0));
e->print(std::cout);           // Calls literal::print
e->expr::print(std::cout);     // Calls expr::print

```

A class that has at least one virtual function should also have a virtual destructor. If a delete expression (see [Chapter 3](#)) deletes a polymorphic pointer (for which the dynamic type does not match the static type), the static class must have a virtual destructor. Otherwise, the behavior is undefined.

## Covariant Return Types

The return type of an overriding virtual function must be the same as that of the base function, or it must be *covariant*. In a derived class, a covariant return type is a pointer or reference to a class type that derives from the return type used in the base class. Note that the return type classes do not necessarily have to match the classes that contain the functions, but they often do. The return type in the derived class can have additional `const` or `volatile` qualifiers that are not present in the base-class return type.

In a function call, the actual return value is implicitly cast to the static type used in the function call. [Example 6-17](#) shows one typical use of covariant types.

### *Example 6-17. Covariant return types*

```
struct shape {
    virtual shape* clone( ) = 0;
};
struct circle : shape {
    virtual circle* clone( ) {
        return new circle(*this);
    }
    double radius( ) const { return radius_; }
    void radius(double r) { radius_ = r; }
private:
    double radius_;
    point center_;
};
struct square : shape {
    virtual square* clone( ) {
        return new square(*this);
    }
private:
    double size_;
    point corners_[4];
};

circle unit_circle;

circle* big_circle(double r)
```

```

{
    circle* result = unit_circle.clone( );
    result->radius(r);
    return result;
}

int main( )
{
    shape* s = big_circle(42.0);
    shape* t = s->clone( );
    delete t;
    delete s;
}

```

## Pure Virtual Functions

A virtual function can be declared with the pure specifier (=0) after the function header. Such a function is a *pure virtual function* (sometimes called an *abstract function*). The syntax for a pure specifier requires the symbols = 0. You cannot use an expression that evaluates to 0.

Even though a function is declared pure, you can still provide a function definition (but not in the class definition). A definition for a pure virtual function allows a derived class to call the inherited function without forcing the programmer to know which functions are pure. A pure destructor must have a definition because a derived-class destructor always calls the base-class destructor.

A derived class can override a pure virtual function and provide a body for it, override it and declare it pure again, or simply inherit the pure function. See the next section, [Section 6.4.4](#).

[Example 6-18](#) shows some typical uses of pure virtual functions. A base class, `shape`, defines several pure virtual functions (`clone`, `debug`, `draw`, and `num_sides`). The `shape` class has no behavior of its own, so its functions are pure virtual.



### Example 6-18. Pure virtual functions

```
class shape {
public:
    virtual ~shape( );
    virtual void draw(graphics* context) = 0;
    virtual size_t num_sides( ) const = 0;
    virtual shape* clone( ) const = 0;
    virtual void debug(ostream& out) const = 0;
};

class circle : public shape {
public:
    circle(double r) : radius_(r) {}
    virtual void draw(graphics* context);
    virtual size_t num_sides( ) const { return 0; }
    virtual circle* clone( ) const { return new
circle(radius( )); }
    virtual void debug(ostream& out) const {
        shape::debug(out);
        out << "radius=" << radius_ << '\n';
    }
    double radius( ) const { return radius_; }
private:
    double radius_;
};

class filled_circle : public circle {
public:
    filled_circle(double r, ::color c) : circle(r), color_(c) {}
    virtual filled_circle* clone( ) const {
        return new filled_circle (radius( ), color( ));
    }
    virtual void draw(graphics* context);
    virtual void debug(ostream& out) const {
        circle::debug(out);
        out << "color=" << color_ << '\n';
    }
    ::color color( ) const { return color_;}
private:
    color color_;
};

void shape::debug(ostream& out)
const
{}
```

Even though `shape::debug` is pure, it has a function body. Derived classes must override `shape::debug`, but they can also call it, which permits uniform implementation of the various debug functions. In other words, every implementation of `debug` starts by calling the base class `debug`. Classes that inherit directly from `shape` do not need to implement `debug` differently from classes that inherit indirectly.

## Abstract Classes

An *abstract class* declares at least one pure virtual function or inherits a pure virtual function without overriding it. A *concrete class* has no pure virtual functions (or all inherited pure functions are overridden). You cannot create an object whose type is an abstract class. Instead, you must create objects of concrete type. In other words, a concrete class that inherits from an abstract class must override every pure virtual function.

Abstract classes can be used to define a pure interface class, that is, a class with all pure virtual functions and no nonstatic data members. Java and Delphi programmers recognize this style of programming because it is the only way these languages support multiple inheritance. [Example 6-19](#) shows how interface classes might be used in C++.

*Example 6-19. Using abstract classes as an interface specification*

```
struct Runnable {
    virtual void run( ) = 0;
};

struct Hashable {
    virtual size_t hash( ) = 0;
};

class Thread : public Runnable, public Hashable {
public:
```

```

Thread( ) { start_thread(*this); }
Thread(const Runnable& thread) { start_thread(thread); }
virtual void run( );
virtual size_t hash( ) const { return thread_id( ); }
size_t thread_id( ) const;
...
private:
    static void start_thread(const Runnable&);
};

// Derived classes can override run to do something useful.
void Thread::run( )
{}

```

## Multiple Inheritance

A class can derive from more than one base class. You cannot name the same class more than once as a direct base class, but a class can be used more than once as an indirect base class, or once as a direct base class and one or more times as an indirect base class. Some programmers speak of inheritance trees or hierarchies, but with multiple base classes, the organization of inheritance is a directed acyclic graph, not a tree. Thus, C++ programmers sometimes speak of *inheritance graphs*.

If multiple base classes declare the same name, the derived class must qualify references to the name or else the compiler reports an ambiguity error:

```

struct base1 { int n; };
struct base2 { int n; };
struct derived : base1, base2 {
    int get_n( ) { return base1::n; } // Plain n is an error.
};

```

Objects of a derived-class type contain separate subobjects for each instance of every base class to store the base class's nonstatic data members. To refer to a member of a particular subobject, qualify its name with the name of its base class. Static members, nested types, and enumerators are shared among all instances of a repeated base class, so

they can be used without qualification (unless the derived class hides a name with its own declaration), as shown in [Example 6-20](#).

*Example 6-20. Multiple inheritance*

```
struct base1 {
    int n;
};
struct base2 {
    enum color { black, red };
    int n;
};
struct base3 : base2 {
    int n; // Hides base2::n
};
struct derived : base1, base2, base3 {
    color get_color( ); // OK: unambiguous use of base2::color
    int get_n( ) { return n; } // Error: ambiguous
    int get_n1( ) { return base2::n; } // Error: which base2?
    int get_n2( ) { return base3::n; } // OK
    int get_n3( ) { // OK: another way to get to a specific
member n
        base3& b3 = *this;
        base2& b2 = b3;
        return b2.n;
    }
};
```

A well-designed inheritance graph avoids problems with ambiguities by ensuring that names are unique throughout the graph, and that a derived class inherits from each base class no more than once. Sometimes, however, a base class must be repeated in an inheritance graph. [Figure 6-1](#) illustrates the organization of multiple base classes, modeled after the standard I/O stream classes. Because `basic_iostream` derives from `basic_istream` and from `basic_ostream`, it inherits two sets of flags, two sets of buffers, and so on, even though it should have only one set of each. This problem is solved by virtual base classes, as explained in the next section.

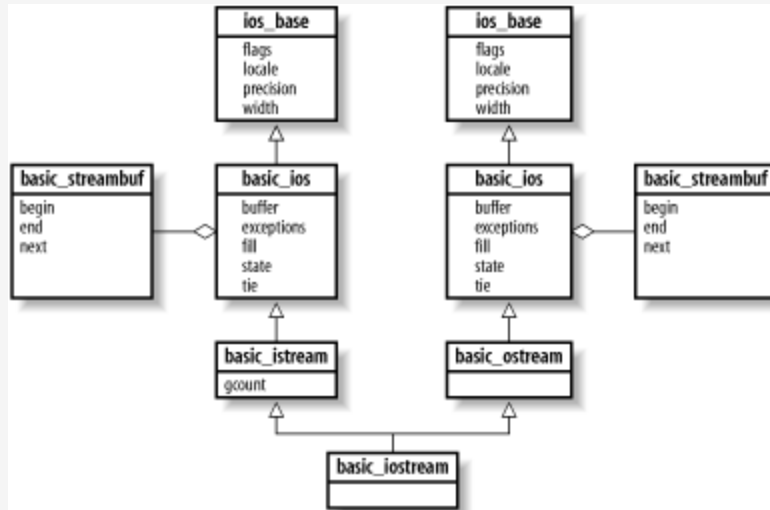


Figure 6-1. Deriving from multiple base classes

## Virtual Inheritance

A base class can be declared with the `virtual` specifier, which is important when a class has multiple base classes. Ordinarily, each time a class is named in an inheritance graph, the most-derived class is organized so that an object has a distinct subobject for each occurrence of each base class. Virtual base classes, however, are combined into a single subobject. That is, each nonvirtual base class results in a distinct subobject, and each virtual base class gets a single subobject, no matter how many times that base class appears in the inheritance graph.

[Figure 6-2](#) illustrates virtual base classes as they are used in the standard I/O streams. Because `basic_istream` and `basic_ostream` derive virtually from `basic_ios`, the `basic_iostream` class inherits a single copy of `basic_ios`, with its single copy of `flags`, `buffer`, and so on.

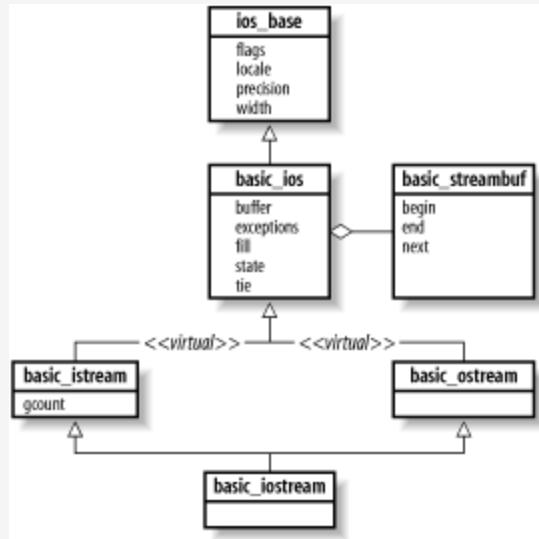


Figure 6-2. Virtual base classes

If a base class is used as a virtual and nonvirtual base class in an inheritance graph, all the virtual instances are shared, and all the nonvirtual instances are separate from the virtual instances. This situation is rare, and usually indicates a programming error.

When constructing a class that has virtual base classes, the most-derived class's constructors must initialize all the virtual base classes. If a constructor does not explicitly initialize a virtual base class, the virtual base class's default constructor is used. Initializers for the virtual base classes are ignored when all base class constructors run.

### Warning

You should design your virtual base classes so they require only the default constructor. Otherwise, you impose a burden on every derived class to initialize the virtual base class properly. Any change to the parameters of the virtual base class constructor necessitates a corresponding change to every constructor for every class that derives, directly or indirectly, from the virtual base class.

[Example 6-21](#) shows the skeletons of the declarations for the standard I/O streams, as illustrated in [Figure 6-2](#).

*Example 6-21. Virtual inheritance in the standard I/O classes*

```
class ios_base;

template<typename charT, typename traits>
class basic_ios : public ios_base;

template<typename charT, typename traits>
class basic_istream : public virtual basic_ios<charT,traits>;

template<typename charT, typename traits>
class basic_ostream : public virtual basic_ios<charT,traits>;

template<typename charT, typename traits>
class basic_iostream : public basic_istream<charT,traits>,
                      public basic_ostream<charT,traits>;
```

## Polymorphic Constructors and Destructors

When an object is constructed, the class of the object being created is called the *most-derived class*. When a constructor body runs, the class identity and virtual functions are those of the class whose constructor is running, which is not necessarily the most-derived class. If you call a pure virtual function, the behavior is undefined.

Destructors run in the opposite order of constructors. In the body of a destructor, the class identity and virtual functions are those of the destructor's class, not the most-derived class. As mentioned earlier, any class that has a virtual function should also have a virtual destructor.

[Example 6-22](#) illustrates how an object's identity changes as it is being constructed.

*Example 6-22. Constructing and destroying an object*

```
#include <iostream>
#include <ostream>
```

```
struct base {
    base( ) { whoami( ); }
    ~base( ) { whoami( ); }
    virtual void whoami( ) { std::cout << "I am base\n"; }
};

struct derived : base {
    derived( ) { whoami( ); }
    ~derived( ) { whoami( ); }
    virtual void whoami( ) { std::cout << "I am derived\n"; }
};

struct most_derived : virtual derived {
    most_derived( ) { whoami( ); }
    ~most_derived( ) { whoami( ); }
    virtual void whoami( )
    {
        std::cout << "I am most_derived\n";
    }
};

int main( )
{
    most_derived md;
    // Prints:
    // I am base
    // I am derived
    // I am most_derived
    // I am most_derived
    // I am derived
    // I am base
}
```



## Access Specifiers

Access specifiers restrict who can access a member. You can use an access specifier before a base-class name in a class definition and have access specifier labels within a class definition. The access specifiers are:

`public`

Anyone can access a public member.

`protected`

Only the class, derived classes, and friends can access protected members.

`private`

Only the class and friends can access private members.

In a class definition, the default access for members and base classes is `private`. In a struct definition, the default is `public`. That is the only difference between a class and a struct, although by convention, some programmers use struct only for POD classes and use class for all other classes.

The access level of a base class affects which members of the base class are accessible to users of a derived class (not the derived class's access to the base class). The access level caps the accessibility of inherited members. In other words, private inheritance makes all inherited members private in the derived class. Protected inheritance reduces the accessibility of public members in the base class to protected in the derived class. Public inheritance leaves all accessibility as it is in the base class.

The access level of a base class also limits type casts and conversions. With public inheritance, you can cast from a

derived class to a base class in any function. Only derived classes and friends can cast to a protected base class. For private inheritance, only the class that inherits directly and friends can cast to the base class.

Access level labels in a class definition apply to all data members, member functions, and nested types. A label remains in effect until a new access label is seen or the class definition ends.



Access specifier labels affect the layout of nonstatic data members. In the absence of access specifier labels, nonstatic data members are at increasing addresses within an object in the order of declaration. When separated by access specifier labels, however, the order is implementation-defined.

When looking up names and resolving overloaded functions, the access level is not considered. The access level is checked only after a name has been found and overloaded functions have been resolved, and if the level does not permit access, the compiler issues an error message. [Example 6-23](#) shows how the compiler ignores the access level when resolving names and overloading.

*Example 6-23. The access level and overloading*

```
class base {
public:
    void func(double);
protected:
    void func(long);
private:
    void func(int);
};

class demo : public base {
public:
    demo( )    { func(42L); } // Calls base::func(long)
    void f( ) { func(42); }  // Error: func(int) is private
```

```

};

class closed : private demo {
public:
    closed() { f( ); } // OK: f( ) is accessible from closed
};

int main( )
{
    demo d;
    d.func(42L); // Error: func(long) accessibly only from base and
demo
    d.func(42); // Error: func(int) is private

    closed c;
    c.f( ); // Error: private inheritance makes demo::f( )
private in closed.
}

```

A derived class can change the accessibility of inherited members with using declarations. A derived class can restore the accessibility of a member whose access was reduced by protected or private inheritance, or it can increase the accessibility of an inherited protected member. [Example 6-24](#) shows how using declarations work. (You can omit the using keyword, leaving only the qualified name, but such usage is deprecated.)

*Example 6-24. Adjusting accessibility with using declarations*

```

struct base { // struct is public by default.
    int x;
protected:
    int y;
};
// Private inheritance makes x and y private in derived1.
class derived1 : private base {
public:
    using base::x; // x is now public.
};
// public inheritance means x is public and y is protected in
derived2.
class derived2 : public base {
public:
    using base::y; // y is now public.
}

```

```

private:
    using base::x; // Pointless: x is still public
};

int main( )
{
    base b;
    derived1 d1;
    derived2 d2;
    b.x = 0;           // OK: x is public in base
    b.y = 42;         // Error: y is protected in base
    d1.x = b.x;       // OK: x is public in derived1
    d2.x = d1.x;      // OK: x is public in derived2
    d1.y = 42;        // Error: y is private in derived1
    d2.y = b.x;       // OK: y is public in derived2
    return d2.y;
}

```

A member's accessibility depends on how it is accessed. In particular, a protected, nonstatic member can be accessed only through its own class or a derived class, but not through a base class. [Example 6-25](#) demonstrates this principle.

### *Example 6-25. Accessing a protected member*

```

class base {
protected:
    int m;
};

struct derived : base {
    void reset( ) {
        this->m = 0; // OK
    }
    void set(base* b) {
        b->m = 0;    // Error: cannot refer to m through base
    }
    bool operator==(const derived& d) {
        return this->m == d.m; // OK
    }
};

```

## Friends

A *friend* is permitted full access to private and protected members. A friend can be a function, function template, or member function, or a class or class template, in which case the entire class and all of its members are friends.

Use the friend specifier to declare a friend in the class granting friendship. Note that friendship is given, not taken. In other words, if class A contains the declaration `friend class B;`, class B can access the private members of A, but A has no special privileges to access B (unless B declares A as a friend).

By convention, the friend specifier is usually first, although it can appear in any order with other function and type specifiers. The friend declaration can appear anywhere in the class; the access level is not relevant.

You cannot use a storage class specifier in a friend declaration. Instead, you should declare the function before the class definition (with the storage class, but without the friend specifier), then redeclare the function in the class definition (with the friend specifier and without the storage class). The function retains its original linkage. If the friend declaration is the first declaration of a function, the function gets external linkage. (See [Chapter 2](#) for more information about storage classes and linkage.) For example:

```
class demo;
static void func(demo& d);
class demo {
    friend void func(demo&);
    ...
}
```

Friendship is not transitive—that is, the friend of my friend is not my friend (unless I declare so in a separate friend declaration)—nor is a nested class a friend just because the

outer class is a friend. (See the next section, [Section 6.7](#), for more information.)

Friendship is not inherited. If a base class is a friend, derived classes do not get any special privileges.

You cannot define a class in a friend declaration, but you can define a function, provided the class granting friendship is not local to a block. The function body is in the class scope, which affects name lookup (see [Chapter 2](#)). The friend function is automatically inline. Usually, friend functions are declared, not defined, in the class.

A declaration or definition of a friend function does not make that function a member of the class and does not introduce the name into the class scope. [Example 6-26](#) shows several different kinds of friends.

#### *Example 6-26. Friend functions and classes*

```
#include <iterator>

// Simple container for singly-linked lists
template<typename T>
class slist {
    // Private type for a link (node) in the list
    template<typename U>
    struct link {
        link* next;
        U value;
    };
    typedef link<T> link_type;

    // Base class for iterator and const_iterator. Keeps track of
    // current node, and
    // previous node to support erase( ).
    class iter_base :
    public std::iterator<std::forward_iterator_tag, T> {
    protected: friend class slist; // So slist can construct
    iterators
        iter_base(slist::link_type* prv, slist::link_type* node);
        slist::link_type* node_;
        slist::link_type* prev_;
    };
};
```

```

public:
    typedef T value_type;
    typedef std::size_t size_type;

    class iterator : public iter_base {
        // Members omitted for brevity . . .
    private:
        friend class slist; // So slist can call constructor
        iterator(slist::link_type* prev, slist::link_type* node)
            : iter_base(prev, node) {}
    };

    friend class iter_base; // So iter_base can use link_type
    friend class iterator; // So iterator can use link_type
    template<typename U>
    friend void swap(slist<U>& a, slist<U>& b);

    iterator begin( )      { return iterator(0, head_); }
    iterator end( )       { return iterator(0, 0); }
private:
    link_type* head_;
    size_type count_;
};

template<typename T>
slist<T>::iter_base::iter_base(slist::link_type* prev,
                             slist::link_type* node)
: prev_(prev), node_(node)
{}

// Swap two lists in constant time by swapping members.
template<typename T>
void swap(slist<T>& a, slist<T>& b)
{
    typename slist<T>::link_type* tmp_head = a.head_;
    typename slist<T>::size_type tmp_count = a.count_;
    a.head_ = b.head_;
    a.count_ = b.count_;
    b.head_ = tmp_head;
    b.count_ = tmp_count;
}

```

## Nested Types

You can declare and define types within a class definition. Accessing these nested types is similar to using types declared within a namespace; the class name serves as the namespace name.

### Tip

Java programmers have several different kinds of nested classes to choose from, but C++ has only one. A C++ nested class is like a static member class in Java. You can construct the other forms of nested classes by adding the appropriate declarations and support code. For example:

```
class outer {
public:
    friend inner;    // Implicit in Java
    class inner {
    friend outer; // Implicit in Java
    public:
        // Inner member class keeps track of outer instance.
        inner(outer& out) : out_(out) {}
        int get( ) const { return out_.x( ); }
    private:
        outer& out_;
    }
    int x( );
    int y( ) { inner i(*this); return i.get( ); }
};
```

Nested types obey access specifiers, so private types are usable only by the class and its friends; protected types are usable by the class, derived classes, and friends. Public types are usable anywhere. To use a nested type outside the containing class and derived classes, it must be qualified with the class name.



Nested enumerations add every enumerator to the class scope:

```
class Graphics {
public:
    enum color { black, red, green, yellow, blue, magenta, cyan,
white };
    color pixel(int row, int col) const;
    ...
};
Graphics::color background = Graphics::white;
```

A typedef declaration can also be nested in a class. This technique is often used for traits (see [Chapter 8](#)):

```
template<typename T, std::size_t N>
class array {
public:
    typedef T value_type;
    typedef T* pointer_type;
    typedef T& reference_type;
    typedef std::size_t size_type;
    ...
};
```

A nested-class definition is like an ordinary class definition, except the name is local to the outer class. The inner class has no special access privileges to the outer class, and vice versa. A nested class can be declared within its outer class definition and later defined at namespace scope. In that case, the nested name must be qualified with the outer class name. Often, nested classes are fully defined within their outer classes:

```
struct outer {
    struct inner1 {
        int foo( ) { return 10; }
    };
    class inner2;
};
struct outer::inner2 {
    int foo( );
};
int outer::inner2::foo( ) {
    outer::inner2 x;
```

```
    return x.foo( );  
}
```

## Chapter 7. Templates

C++ is more than a mere object-oriented programming language. The full power of C++ is seen when programming with templates. Templates lie at the heart of the standard library: strings, I/O streams, containers, iterators, algorithms, and more.

A *template* is a pattern for creating classes or functions as instances of the template at compile time, similar to the manner in which a class is a pattern for creating objects as instances of the class at runtime. A template takes one or more parameters, and when you instantiate a class or function template, you must supply arguments for the parameters. The classes and functions can have different behaviors or implementations, depending on the arguments. This style of programming is often called *generic programming*. Templates can also be used to select code at compile time, affect the behavior of generated code, and set policies. This style of programming is often known as *template metaprogramming*.

Programming with templates is unlike traditional object-oriented programming. Object-oriented programming centers around type polymorphism (requiring classes, objects, and virtual functions). Template-based programming centers around *parametric polymorphism*, in which a function or class is defined independently of its parameters (which can be values, types, or even other templates).

This chapter describes the syntax and semantics of declaring, specializing, instantiating, and using templates. See [Chapter 8](#) for information about some typical uses of templates in the standard library. See [Appendix B](#) for information about template-oriented projects.

The syntax descriptions in this chapter are informal. See [Chapter 12](#) for a precise BNF grammar.

## Overview of Templates

A template declaration can be a function declaration, function definition, class declaration, or class definition. The template declaration takes one or more parameters, which can be values, types, or class templates.

In most cases, you can use a template simply by naming the template and providing arguments for the template parameters: constant expressions, types, or template references. This is known as *instantiating* the template. You can instantiate a template at its point of use, or declare a separate instantiation as a class or function declaration. An instance of a function template creates a function; an instance of a class template creates a class and all of its members.

A template lets you define a class or function once for a wide range of template arguments, but sometimes you need to customize a template for particular arguments. This is known as *specializing* the template. A template specialization, as its name implies, is a special case of the template pattern. When you instantiate a template, the compiler uses the template arguments to pick a specialization, or if no specialization matches the arguments, the original template declaration. A specialization can be total or partial. A *total specialization* specifies values for all of the template arguments, and a *partial specialization* specifies only some of the template arguments.

The terminology used in this book reflects the terminology that many C++ programmers have adopted, even though that terminology differs slightly from that used in the C++

standard. In the standard, "specialization" means an instance of a template. "Instantiation" also refers to an instance of a template. When you declare a special case of a template, that is known as *explicit specialization*.

Many C++ programmers prefer to keep specialization and instantiation as separate concepts and separate terms, and I have adopted the simpler terminology for this book.

Note that a template declaration defines only the pattern. A specialization defines a pattern that applies to a specific set of template arguments. Only by instantiating a template do you declare or define a function or class. When you instantiate a template, the compiler uses the template arguments to pick which pattern to instantiate: a specialization or the main template.

Writing a template is more difficult than writing a non-template class or function. The template can be instantiated in almost any context, and the context can affect how the template definition is interpreted. Name lookup rules are more complicated for templates than for non-templates.

[Example 7-1](#) shows several different kinds of templates and their uses.

### *Example 7-1. Declaring and using templates*

```
#include <cmath>
#include <complex>
#include <iostream>
#include <ostream>

// Template declaration of pointtemplate<typename T>
class point {
public:
    typedef T value_type;
    point(const T& x, const T& y) : x_(x), y_(y) {}
    point() : x_(T( )), y_(T( )) {}
    T x( ) const { return x_; }
    T& x( ) { return x_; }
    void x(const T& new_x) { x_ = new_x; }
```

```

    T y( )          const { return y_; }
    T& y( )         { return y_; }
    void y(const T& new_y) { y_ = new_y; }
private:
    T x_, y_;
};

// Instantiate point<>.
typedef point<std::complex<double> > strange;
strange s(std::complex<double>(1, 2),
          std::complex<double>(3, 4));

// Specialize point<int> to use call-by-value instead of const
references.
template<>
class point<int> {
public:
    typedef int value_type;
    point(int x, int y) : x_(x), y_(y) {}
    point( ) : x_(0), y_(0) {}
    int x( )          const { return x_; }
    int& x( )         { return x_; }
    void x(int new_x) { x_ = new_x; }
    int y( )          const { return y_; }
    int& y( )         { return y_; }
    void y(int new_y) { y_ = new_y; }
private:
    int x_, y_;
};

// Instance of the specialized point<int>
point<int> p(42, 0);
// Instance of the general point<>, using long as the template
argument
point<long> p(42, 0);

// Function template
template<typename T>
T abs(T x)
{
    return x < 0 ? -x : x;
}

namespace {
    // Explicit instantiation
    const int abs_char_min1 = abs<int>(CHAR_MIN);
    // Implicit instantiation
    const int abs_char_min2 = abs(CHAR_MIN);
}

```

```

}

// Overload abs( ) with another function template.
template<typename floatT, typename T>
floatT abs(const point<T>& p)
{
    return std::sqrt(static_cast<floatT>(p.x( )*p.x( ) +
                                         p.y( )*p.y( )));
}

int main( )
{
    point<double> p;
    // Call instance of function template. Compiler deduces second
    template
    // argument (double) from the type of p.
    double x = abs<long double>(p);
    std::cout << x << '\n'; // prints 0
    std::cout << abs_char_min1 << '\n';
}

```

# Template Declarations

A template declaration begins with a template header followed by a function declaration or definition, or a class declaration or definition. Template declarations can appear only at namespace or class scope. The template name must be unique in its scope (except for overloaded functions).

The template header starts with the `template` keyword followed by the template parameters enclosed in angle brackets (`<>`). Multiple parameters are separated by commas. The syntax is:

```
template <parameter-list > declaration
```

There are three kinds of template parameters: values, types, and class templates. Similar to a function parameter, a template parameter has an optional name and an optional default argument.

Function templates cannot have default template arguments. If a class template has member definitions that are outside the class definition, only the class template takes default arguments; the individual member definitions do not. If a default argument is present, it is preceded by an equal sign. Only the rightmost parameters can have default arguments. If any parameter has a default argument, all parameters to its right must also have default arguments.

[Example 7-2](#) shows valid and invalid member definitions.

## *Example 7-2. Defining members of a class template*

```
// OK: default argument for template parameter A
template<typename T, typename A = std::allocator<T> >
class hashset {
    bool empty( ) const;
    size_t size( ) const;
    ...
};
// Error: do not use default argument here
template<typename T, typename A = std::allocator<T> >
```



```

bool hashset<T,A>::empty( ) { return size( ) == 0; }

// OK
template<typename T, typename A>
size_t hashset<T,A>::size( ) { ... }

```

Each template header defines the template name and template parameters. The scope of a parameter name extends from its declaration to the end of the declaration or definition of the class or function. A parameter name can be used in subsequent template parameters in the same template header (such as `std::allocator<T>` in [Example 7-2](#)). The template parameter name must be unique in the template declaration and cannot be redeclared in its scope. If a class template has separate definitions for its members, each member definition is free to use different names for the template parameters. (See [Section 7.4.1](#) later in this chapter for more information.)

There are three kinds of template parameters:

### *Value template parameter*

Declared in the same manner as a function parameter:

```

type-specifiers declarator
                                type-specifiers declarator = expr

```

The type must be an integral, enumeration, pointer, reference, or pointer-to-member type. When the template is instantiated, the argument must be a constant integral or enumeration expression, the address of a named object or function with external linkage, or the address of a member:

```

template<unsigned Size>
struct name {
    // ...
    unsigned size( ) const { return Size; }
private:
    char name_[Size+1];
};
name<100> n;

```

Note that a string literal is an unnamed object with internal linkage, so you cannot use it as a template argument:

```
template<const char* Str> void print(const char* s = Str);
print<"default">( ); // Error
const char def[] = "default";
print<def>( );      // OK
```

The *type-specifiers* can be elaborated type specifiers that start with `typename`, that is, `typename` followed by a qualified type name. (If `typename` is followed by a plain identifier, the template parameter is a type parameter, as described later.) For more information about this use of `typename`, see [Section 7.8](#) later in this chapter.

```
template<typename list<int>::value_type value>
int silly( ) { return value; }
```

### *Type template parameter*

Introduced with the keyword `typename` followed by the optional parameter name:

```
typename identifier
typename identifier = type
```

The `class` keyword can be used in place of `typename` and has the same meaning in this context. (A useful convention is to use `class` when the argument must be a class and `typename` when the argument can be any type.) When the template is instantiated, the argument must be a type (that is, a list of type specifiers with optional pointer, reference, array, and function operators). In the following example, the `point` template is instantiated with `unsigned long int` as the template argument:

```
template<typename T> struct point {
    T x, y;
};
point<unsigned long int> pt;
```

If `typename` is followed by a qualified type name instead of a plain identifier, it declares a value parameter of that

type, as described earlier.

### *Template template parameter*

Must be a class template. It has the form of a template declaration:

```
template <parameter-list > class identifier  
template < parameter-list > class identifier = template-id
```

When the template is instantiated, the argument must be a class template:

```
// Erase all occurrences of item from a sequence container.  
template<template<typename T, typename A> class C, typename T,  
        typename A>  
void erase(C<T,A>& c, const T& item)  
{  
    c.erase(std::remove(c.begin(), c.end( ), item), c.end( ));  
}  
...  
list<int> l;  
...  
erase(l, 42);
```

To use a template declaration, you must create an instance of the template, either explicitly (by naming the template and enclosing a list of template arguments in angle brackets) or implicitly (by letting the compiler deduce the template arguments from context) for a function template. In either case, the compiler must know about a template declaration before the template is used. Typically, a template is declared in an `#include` file or header. The header declares the function or class template and possibly provides the definition of the function template or the definitions of all the members of the class template. See [Section 7.10](#) later in this chapter for details about the files that declare and define templates.

A template instance must provide an argument for each template parameter. If a class template has fewer arguments than parameters, the remaining parameters

must have default arguments, which are used for the template instance.

If a function template has fewer arguments than parameters, the remaining arguments are deduced from the context of the function call (as explained in the next section). If the arguments cannot be deduced, the compiler reports an error. See [Section 7.7](#) later in this chapter for more information.

# Function Templates

A function template defines a pattern for any number of functions whose definitions depend on the template parameters. You can overload a function template with a non-template function or with other function templates. You can even have a function template and a non-template function with the same name and parameters.

Function templates are used throughout the standard library. The best-known function templates are the standard algorithms (such as `copy`, `sort`, and `for_each`).

Declare a function template using a template declaration header followed by a function declaration or definition. Default template arguments are not allowed in a function template declaration or definition. In the following example, `round` is a template declaration for a function that rounds a number of type `T` to `N` decimal places (see example 7-14 for the definition of `round`), and `min` is a template definition for a function that returns the minimum of two objects:

```
// Round off a floating-point value of type T to N digits.
template<unsigned N, typename T> T round(T x);
// Return the minimum of a and b.
template<typename T> T min(T a, T b)
{ return a < b ? a : b; }
```

To use a function template, you must specify a template instance by providing arguments for each template parameter. You can do this explicitly, by listing arguments inside angle brackets:

```
long x = min<long>(10, 20);
```

However, it is more common to let the compiler deduce the template argument types from the types of the function arguments:

```
int x = min(10, 20); // Calls min<int>( )
```

See [Section 7.3.2](#) later in this chapter for details.

## Function Signatures

The signature of a function template includes the template parameters even if those parameters are never used in the function's return type or parameter types. Thus, instances of different template specializations produce distinct functions. The following example has two function templates, both named `func`. Because their template parameters differ, the two lines declare two separate function templates:

```
template<typename T> void func(T x = T( ));  
template<typename T, typename U> void func(T x = U( ));
```

Typically, such function templates are instantiated in separate files; otherwise, you must specify the template arguments explicitly to avoid overload conflicts.

The function signature also includes expressions that use any of the template parameters. In the following example, two function templates overload `func`. The function parameter type is an instance of the `demo` class template. The template parameter for `demo` is an expression that depends on `func`'s template parameter (`T` or `i`):

```
template<int x> struct demo {};  
template<typename T> void func(demo<sizeof(T)>);  
template<int i> void func(demo<i+42>);
```

Just as you can have multiple declarations (but not definitions) of a single function, you can declare the same function template multiple times using expressions that differ only by the template parameter names. Even if the declarations are in separate files, the compiler and linker ensure that the program ends up with a single copy of each template instance.

If the expressions differ by more than just the parameter names, they result in distinct functions, so you can have distinct function template definitions. If the argument expressions result in the same value, the program is invalid, although the compiler is not required to issue an error message:

```
template<int i> void func(demo<i+42>);  
template<int x> void func(demo<x+42>) {} // OK  
template<int i> void func(demo<i+41>) {} // OK  
template<int i> void func(demo<i+40+2>) {} // Error
```

Follow the standard procedure of declaring templates in header files and using those headers in the source files where they are needed, thereby ensuring that every source file sees the same template declarations. Then you will not have to be concerned about different source files seeing declarations that differ only by the form of the template argument expressions. The compiler tries to ensure that the program ends up with a single copy of functions that are meant to be the same, but distinct copies of those that are meant to be distinct.

## Deducing Argument Types

Most uses of function templates do not explicitly specify all of the template arguments, letting the compiler deduce the template arguments from the function call. You can provide explicit arguments for the leftmost template parameters, leaving the rightmost parameters for the compiler to deduce. If you omit all the template arguments, you can also choose to omit the angle brackets.

In the following discussion, make sure you distinguish between the parameters and arguments of a template and those of a function. Argument type deduction is the process of determining the template arguments, given a function call or other use of a function (such as taking the address of a function).

The basic principle is that a function call has a list of function arguments, in which each function argument has a type. The function argument types must be matched to function parameter types; in the case of a function template, the function parameter types can depend on the template parameters. The ultimate goal, therefore, is to determine the appropriate template arguments that let the compiler match the function parameter types with the given function argument types.

Usually, the template parameters are type or template parameters, but in a few cases, value parameters can also be used (as the argument for another template or as the size of an array).

The function parameter type can depend on the template parameter in the following ways, in which the template parameter is  $v$  for a value parameter,  $T$  for a type parameter, or  $TT$  for a template template parameter:

- $T$ ,  $T^*$ ,  $T\&$ ,  $\text{const } T$ ,  $\text{volatile } T$ ,  $\text{const volatile } T$
- Array of  $T$
- Array of size  $v$
- Template  $TT$ , or any class template whose template argument is  $T$  or  $v$
- Function pointer in which  $T$  is the return type or the type of any function parameter
- Pointer to a data member of  $T$ , or to a data member whose type is  $T$
- Pointer to a member function that returns  $T$ , to a member function of  $T$ , or to a member function in which  $T$  is the type of any function parameter

Types can be composed from any of these forms. Thus, type deduction can occur for a struct in which the types of its



members are also deduced, or a function's return type and parameter types can be deduced.

A single function argument can result in the deduction of multiple template parameters, or multiple function arguments might depend on a single template parameter.

If the compiler cannot deduce all the template arguments of a function template, you must provide explicit arguments. [Example 7-3](#) shows several different ways the compiler can deduce template arguments.

### *Example 7-3. Deducing template arguments*

```
#include <algorithm>
#include <cstdint>
#include <iostream>
#include <list>
#include <ostream>
#include <set>

// Simple function template. Easy to deduce T.
template<typename T>
T min(T a, T b)
{
    return a < b ? a : b;
}

template<typename T, std::size_t Size>
struct array
{
    T value[Size];
};

// Deduce the row size of a 2D array.
template<typename T, std::size_t Size>
std::size_t dim2(T [][][Size])
{
    return Size;
}

// Overload the function size( ) to return the number of
elements in an array<>
// or in a standard container.
template<typename T, std::size_t Size>
std::size_t size(const array<T,Size>&)
```

```

{
    return Size;
}

template<template<typename T, typename A> class container,
        typename T, typename A>
typename container<T,A>::size_type size(container<T,A>& c)
{
    return c.size( );
}

template<template<typename T, typename C, typename A>
        class container,
        typename T, typename C, typename A>
typename container<T,C,A>::size_type
size(container<T,C,A>& c)
{
    return c.size( );
}

// More complicated function template. Easy to deduce Src, but
// impossible to
// deduce Dst.
template<typename Dst, typename Src>
Dst cast(Src s)
{
    return s;
}

int main( )
{min(10, 100);          // min<int>
  min(10, 1000L);     // Error: T cannot be int and long

  array<int,100> data1;
  // Deduce T=int and SIZE=100 from type of data1.
  std::cout << size(data1) << '\n';

  int y[10][20];
  std::cout << dim2(y) << '\n'; // Prints 20

  std::list<int> lst;
  lst.push_back(10);
  lst.push_back(20);
  lst.push_back(10);
  std::cout << size(lst) << '\n'; // Prints 3

  std::set<char> s;
  const char text[] = "hello";

```

```
std::copy(text, text+5, std::inserter(s, s.begin( )));  
std::cout << size(s) << '\n'; // Prints 4  
  
// Can deduce Src=int, but not return type, so explicitly set  
Dst=char  
char c = cast<char>(42);  
}
```

## Overloading

Function templates can be overloaded with other function templates and with non-template functions. Templates introduce some complexities, however, beyond those of ordinary overloading (as covered in [Chapter 5](#)).

As with any function call, the compiler collects a list of functions that are candidates for overload resolution. If the function name matches that of a function template, the compiler tries to determine a set of template arguments, deduced or explicit. If the compiler cannot find any suitable template arguments, the template is not considered for overload resolution. If the compiler can find a set of template arguments, they are used to instantiate the function template, and the template instance is added to the list of candidate functions. Thus, each template contributes at most one function to the list of overload candidates. The best match is then chosen from the list of template and non-template functions according to the normal rules of overload resolution: non-template functions are preferred over template functions, and more specialized template functions are preferred over less specialized template functions. See [Chapter 5](#) for more details on overload resolution.

[Example 7-4](#) shows how templates affect overload resolution. Sometimes, template instantiation issues can be subtle. The last call to `print` results in an error, not because the compiler cannot instantiate `print`, but because

the instantiation contains an error, namely, that operator << is not defined for `std::vector<int>`.

#### *Example 7-4. Overloading functions and function templates*

```
#include <cctype>
#include <iomanip>
#include <iostream>
#include <ostream>
#include <string>
#include <vector>

template<typename T>
void print(const T& x)
{
    std::cout << x;
}

void print(char c)
{
    if (std::isprint(c))
        std::cout << c;
    else
        std::cout << std::hex << std::setfill('0')
                    << "\\x" << std::setw(2) << int(c) << "\\'"
                    << std::dec;
}

template<>
void print(const long& x)
{
    std::cout << x << 'L';
}

int main( )
{
    print(42L);    // Calls print<long> specialization
    print('x');   // Calls print(char) overload
    print('\\0'); // Calls print(char) overload
    print(42);    // Calls print<int>
    print(std::string("y")); // Calls print<string>
    print( );     // Error: argument required
    print(std::vector<int>( )); // Error
}
```

## Operators

You can create operator templates in the same manner as function templates. The usual rules for deducing argument types apply as described earlier in this section. If you want to specify the arguments explicitly, you can use the operator keyword:

```
template<typename T>
bigint operator+(const bigint& x, const T& y);
bigint a;
a = a + 42; // Argument type deduction
a = operator+<long>(a, 42); // Explicit argument type
```

When using operator<, be sure to leave a space between the operator symbol and the angle brackets that surround the template argument, or else the compiler will think you are using the left-shift operator:

```
operator<<long>(a, 42); // Error: looks like left-shift
operator< <long>(a, 42); // OK
```

Type conversion operators can use ordinary type cast syntax, but if you insist on using an explicit operator keyword, you cannot use the usual syntax for specifying the template argument. Instead, the target type specifies the template argument type:

```
struct demo {
    template<typename T>
    operator T*( ) const { return 0; }
};
demo d;
char* p = d.operator int*( ); // Error: illegal
cast
char* c = d; // d.operator char*( ); // OK
```

# Class Templates

A class template defines a pattern for any number of classes whose definitions depend on the template parameters. The compiler treats every member function of the class as a function template with the same parameters as the class template.

Class templates are used throughout the standard library for containers (`list<>`, `map<>`, etc.), complex numbers (`complex<>`), and even strings (`basic_string<>`) and I/O (`basic_istream<>`, etc.).

The basic form of a class template is a template declaration header followed by a class declaration or definition:

```
template<typename T>
struct point {
    T x, y;
};
```

To use the class template, supply an argument for each template parameter (or let the compiler substitute a default argument). Use the template name and arguments the way you would a class name.

```
point<int> pt = { 42, 10 };
typedef point<double> dpoint;
dpoint dp = { 3.14159, 2.71828 };
```

In member definitions that are separate from the class definition, you must declare the template using the same template parameter types in the same order, but without any default arguments. (You can change the names of the template parameters, but be sure to avoid name collisions with base classes and their members. See [Section 7.8](#) later in this chapter for more information.) Declare the member definitions the way you would any other definition, except that the class name is a template name (with arguments),

and the definition is preceded by a template declaration header.

In the class scope (in the class definition, or in the definition of a class member), the bare class name is shorthand for the full template name with arguments. Thus, you can use the bare identifier as the constructor name, after the tilde in a destructor name, in parameter lists, and so on. Outside the class scope, you must supply template arguments when using the template name. You can also use the template name with different arguments to specify a different template instance when declaring a member template. The following example shows how the template name can be used in different ways:

```
template<typename T> struct point {
    point(T x, T y);    // Or point<T>(T x, T y);
    ~point<T>( );     // Or ~point( );
    template<typename U>
    point(const point<U>& that);
    ...
};
```

[Example 7-5](#) shows a class template with several different ways to define its member function templates.

### *Example 7-5. Defining class templates*

```
template<typename T, typename U = int>
class demo {
public:
    demo(T t, U u);
    ~demo( );
    static T data;
    class inner; // inner is a plain class, not a template.
};

template<typename T, typename U>
demo<T,U>::demo(T, U)    // Or demo<T,U>::demo<T,U>(T, U)
{}

template<typename T, typename U>
demo<T,U>::~~demo<T,U>( ) // Or demo<T,U>::~~demo( )
{}
```

```

template<typename U, typename T> // Allowed, but confusing
U demo<U,T>::data;

template<typename T, typename U>
class demo<T,U>::inner {
public:
    inner(T, demo<T,U>& outer);
private:
    demo<T,U>& outer;
};

// The class name is "demo<T,U>::inner", and the constructor name
is "inner".
template<typename T, typename U>
demo<T,U>::inner::inner(T, demo<T,U>& outer)
: outer(outer)
{}

```

The template defines only the pattern; the template must be instantiated to declare or define a class. See the sections [Section 7.5](#) and [Section 7.7](#) later in this chapter for more information.

## Member Templates

A *member template* is a template inside a class template or non-template class that has its own template declaration header, with its own template parameters. It can be a member function or a nested type. Member templates have the following restrictions:

- Local classes cannot have member templates.
- A member function template cannot be a destructor.
- A member function template cannot be virtual.
- If a base class has a virtual function, a member function template of the same name and parameters in a derived class does not override that function.

Type conversion functions have additional rules because they cannot be instantiated using the normal template



instantiation or specialization syntax:

- A using declaration cannot refer to a type conversion template in a base class.
- Use ordinary type conversion syntax to instantiate a type conversion template:

```
template<typename T>
struct silly {
    template<typename U> operator U*( ) { return 0; }
};
...
silly<int> s, t(42);
if (static_cast<void*>(s))
    std::cout << "not null\n";
```

A class can have a non-template member function with the same name as a member function template. The usual overloading rules apply ([Chapter 5](#)), which means the compiler usually prefers the non-template function to the function template. If you want to make sure the function template is called, you can explicitly instantiate the template:

```
template<typename T>
struct demo {
    template<typename U> void func(U);
    void func(int);
};

demo<int> d;
d.func(42);          // Calls func(int)
d.func<int>(42);    // Calls func<int>
```

## Friends

A friend (of a class template or an ordinary class) can be a template, a specialization of a template, or a non-template function or class. A friend declaration cannot be a partial specialization. If the friend is a template, all instances of the template are friends. If the class granting friendship is a template, all instances of the template grant friendship. A

specialization of a template can be a friend, in which case instances of only that specialization are granted friendship:

```
template<typename T> class buddy {};  
template<typename T> class special {};  
class demo {  
    // All instances of buddy<> are friends.  
    template<typename T> friend class buddy;  
    // special<int> is a friend, but not special<char>, etc.  
    friend class special<int>;  
};
```

When you use the containing class template as a function parameter, you probably want friend functions to be templates also:

```
template<typename T>  
class outer {  
    friend void func1(outer& o);    // Wrong  
  
    template<typename U>  
    friend void func2(outer<U>& o); // Right  
};
```

A friend function can be defined in a class template, in which case every instance of the class template defines the function. The function definition is compiled even if it is not used.

You cannot declare a friend template in a local class.

If a friend declaration is a specialization of a function template, you cannot specify any default function arguments, and you cannot use the `inline` specifier. These items can be used in the original function template declaration, however.

All the rules for function templates apply to friend function templates. (See [Section 7.3](#) earlier in this chapter for details.) [Example 7-6](#) shows examples of friend declarations and templates.

*Example 7-6. Declaring friend templates and friends of templates*

```

#include <algorithm>
#include <cstddef>
#include <iostream>
#include <iterator>
#include <ostream>

template<typename T, std::size_t Size>
class array {
public:
    typedef T value_type;

    class iterator_base :
        std::iterator<std::random_access_iterator_tag, T> {
    public:
        typedef T value_type;
        typedef std::size_t size_type;
        typedef std::ptrdiff_t distance_type;        friend inline bool
operator==(const iterator_base& x,
            const iterator_base& y)
        {
            return x.ptr_ == y.ptr_;
        }
friend inline bool operator!=(const iterator_base& x,
            const iterator_base& y)
        {
            return ! (x == y);
        }
friend inline ptrdiff_t operator-(const iterator_base& x,
            const iterator_base& y)
        {
            return x.ptr_ - y.ptr_;
        }

    protected:
        iterator_base(const iterator_base& that)
            : ptr_(that.ptr_) {}
        iterator_base(T* ptr) : ptr_(ptr) {}
        iterator_base(const array& a, std::size_t i)
            : ptr_(a.data_ + i) {}

        T* ptr_;
};

friend class iterator_base;

class iterator : public iterator_base {
public:
    iterator(const iterator& that) : iterator_base(that) {}

```

```

T& operator*( ) const { return *this->ptr_; }
iterator& operator++( ) { ++this->ptr_; return *this; }
T* operator->( ) const { return this->ptr_; }

private:
friend class array;
iterator(const array& a, std::size_t i = 0)
    : iterator_base(a, i) {}
iterator(T* ptr) : iterator_base(ptr) {}

friend inline iterator operator+(iterator iter, int off)
{
    return iterator(iter.ptr_ + off);
}
};

array( ) : data_(new T[Size]) {}
array(const array& that);
~array( ) { delete[] data_; }
iterator begin( ) { return iterator(*this); }
const_iterator begin( ) const { return iterator(*this); }
iterator end( );
const_iterator end( ) const;
T& operator[](std::size_t i);
T operator[](std::size_t i);

template<typename U, std::size_t USize>
friend void swap(array<U,USize>& a, array<U,USize>& b);

private:
    T* data_;
};

template<typename T, std::size_t Size>
void swap(array<T,Size>& a, array<T,Size>& b)
{
    T* tmp = a.data_;
    a.data_ = b.data_;
    b.data_ = tmp;
}

```

## Specialization

A template declares a set of functions or classes from a single declaration, which is a powerful tool for writing code once and having it work in a multitude of situations. A single, one-size-fits-all approach is not always appropriate, however. You can therefore specialize a template for specific values of one or more template arguments and provide a completely different definition for that special case.

You can specialize a class template, a function template, or a member of a class template. Declare a specialization with an empty set of template parameters in the template header followed by a declaration that provides arguments for all the template parameters:

```
template<>declaration
```

Specializations must appear after the declaration of the primary template, and they must be declared in the same namespace as the primary template. You can also have a *partial specialization*, which provides some of the template arguments. (See the next section, [Section 7.6](#).)

When a template is instantiated, if the template arguments of the instantiation match the arguments used in a specialization, the compiler instantiates the matching specialization. Otherwise, it instantiates the primary template. See [Section 7.7](#) later in this chapter for details.

To specialize a function template, the function return type and parameter types must match the primary template declaration. The function name can be followed by the specialized template arguments (enclosed in angle brackets and separated by commas), or you can omit them and let the compiler deduce the specialized arguments. See [Section 7.3.2](#) earlier in this chapter for details.

If you need to change the function parameters, you can overload the function, but that requires an entirely new function or function template declaration, not a specialization. [Example 7-7](#) shows two specializations of a function template.

### *Example 7-7. Specializing a function template*

```
// Primary template declaration
template<typename T> T root(T x, T y)
{
    return pow(x, 1.0/y);
}

// Specialization in which T is deduced to be long
template<> long root(long x, long y)
{
    if (y == 2)
        return sqrt((long double)x);
    else
        return pow((long double)x, 1.0l/y);
}

// Specialization for explicit T=int
template<> int root<int>(int x, int y)
{
    if (y == 2)
        return sqrt(double(x));
    else
        return pow(double(x), 1.0/y);
}

// Overload with a different function template, such as
valarray<T>.
template<template<typename T> class C, typename T>
C<T> root(const C<T>& x, const C<T>& y)
{
    return pow(x, C<T>(1.0)/y);
}
```

A specialization of a class template requires an entirely new class declaration or definition. The members of the specialized class template can be entirely different from the members of the primary class template. If you want to

specialize only some members, you can specialize just those members instead of specializing the entire class.

In a specialization of a static data member, you must supply an explicit initializer to define the data member. Without an initializer, the member template declaration is just a declaration of the data member, not a definition. If you want to use the static data member, you must make sure it is defined:

```
template<typename T>
struct demo {
    static T data;
};
template<> int demo<int>::data;    // Declaration
template<> int demo<int>::data = 42; // Definition
```

Declare specializations before the template is used. (See [Section 7.7](#) later in this chapter for information on how templates are "used.") A program must have a single specialization for a given set of template arguments.

[Example 7-8](#) shows some specializations of a class template, `type_traits`, which exposes a few attributes, or traits, of a type. The primary template describes the default traits, and the template is specialized for specific types, such as `int`. See [Chapter 8](#) for more information about traits.

### *Example 7-8. Specializing a class template*

```
#include <iostream>
#include <ostream>
#include <sstream>
#include <string>
#include <typeinfo>

template<typename T>
struct type_traits
{
    typedef T base_type;
    enum { is_fundamental = 0 };
    enum { is_integer = 0 };
    enum { is_float = 0 };
};
```

```

enum { is_pointer = 0 };
enum { is_reference = 0 };
static std::string to_string(const T&);
};

template<typename T>
std::string type_traits<T>::to_string(const T& x)
{
    return typeid(x).name( );
}

// Specialize entire class for each fundamental type. template<>
struct type_traits<int>
{
    typedef int base_type;
    enum { is_fundamental = 1 };
    enum { is_integer = 1 };
    enum { is_float = 0 };
    enum { is_pointer = 0 };
    enum { is_reference = 0 };
    static std::string to_string(const int& x) {
        std::ostringstream out;
        out << x;
        return out.str( );
    }
};

struct point
{
    int x, y;
};

// Specialize only the to_string( ) member function for type
point.
template<>
std::string type_traits<point>::to_string(const point& p)
{
    std::ostringstream out;
    out << '(' << p.x << ', ' << p.y << ')';
    return out.str( );
}

int main( )
{
    using std::cout;
    cout << type_traits<point>::is_fundamental << '\n';
    cout << type_traits<point>::is_pointer << '\n';
    cout << type_traits<point>::to_string(point( )) << '\n';
}

```



```
cout << type_traits<int>::is_fundamental << '\n';  
cout << type_traits<int>::is_pointer << '\n';  
cout << type_traits<int>::is_integer << '\n';  
cout << type_traits<int>::to_string(42) << '\n';  
}
```

## Partial Specialization

You can choose to specialize only some of the parameters of a class template. This is known as partial specialization. Note that function templates cannot be partially specialized; use overloading to achieve the same effect.

A partial specialization is declared with a template header that contains one or more template parameters. (With no template parameters, it is a total specialization, not a partial specialization; see the previous section for details.) The class declaration that follows the template header must supply an argument for each parameter of the primary template (or rely on default arguments), and the arguments can depend on the template parameters of the partial specialization. For example, suppose you have a simple `pair` class, similar to the one in the standard library, but you want to allow one member of the `pair` to be `void`. In this case, an object can be made smaller and store only one item:

```
template<typename T, typename U>
struct pair {
    T first;
    U second;
};
template<typename X> struct pair<X, void> {
    X first;
};
```

The partial specialization of a class is a distinct template and must provide a complete class definition. You cannot partially specialize a member of a class template, only the entire class template.

[Example 7-9](#) shows partial specializations of the `type_traits` template from [Example 7-8](#). The first partial specialization applies to all pointer types. It sets the `is_pointer` member to 1, for example. The second partial

specialization applies to all const types. It sets most members to their values obtained from the non-const template instance. Thus, `type_traits<const int*>::base_type` is plain `int`.

*Example 7-9. Partially specializing a class template*

```
// Specialize for all pointer types.template<typename T>
struct type_traits<T*>
{
    typedef T base_type;
    enum { is_fundamental = type_traits<T>::is_fundamental };
    enum { is_integer = 0 };
    enum { is_float = 0 };
    enum { is_pointer = 1 };
    enum { is_reference = 0 };
    static std::string to_string(const T& x) {
        std::ostringstream out;
        out << x;
        return out.str( );
    }
};

// Specialize for const types, so base_type refers to the non-
const type.
template<typename T>
struct type_traits<const T>
{
    typedef T base_type;
    typedef type_traits<base_type> base_type_traits;
    enum { is_fundamental = base_type_traits::is_fundamental };
    enum { is_integer = base_type_traits::is_integer };
    enum { is_float = base_type_traits::is_float };
    enum { is_pointer = base_type_traits::is_pointer };
    enum { is_reference = base_type_traits::is_reference };
    static std::string to_string(const base_type& x) {
        return type_traits<base_type>::to_string(x);
    }
};
```

## Instantiation

Template declarations and specializations describe the form that a class or function can take but do not create any actual classes or functions. To do these things, you must instantiate a template.

Most often, you implicitly instantiate a template by using it in a function call, object declaration, or similar context. The template instance requires a template argument for each template parameter. The arguments can be explicit (enclosed in angle brackets and separated by commas) or implicit (default arguments for class templates or deduced arguments for function templates):

```
template<typename T> T plus(T a, T b) { return a + b; }
template<typename T = int> struct wrapper {
    T value;
};
int x = plus(10, 20); // Instantiate plus<int>.
wrapper<double> x;    // Instantiate wrapper<double>.
wrapper<> w;          // Instantiate wrapper<int>.
```

A class member expression (. or -> operator) that names an instance of a member function template with template arguments must use the `template` keyword before the member name. Similarly, a qualified name of a member template must use `template` before the member name. Otherwise, the `<` symbol that introduces the template argument list is interpreted as the less-than operator:

```
class bigfloat {
    template<unsigned N> double round( );
    template<unsigned N> static double round(double);
};

bigfloat f;
std::cout << f.template round<2>( ) << '\n';
std::cout << bigfloat::template round<8>(PI) << '\n';
```



Instantiating a class template does not necessarily instantiate all the members. Members of a class template are instantiated only when they are needed. A static data member is "needed" if the program refers to the member. A member function is "needed" if the function is called, has its address taken, or participates in overload resolution. An implementation is permitted to define every virtual function in a class template, and most implementations do. Each instance of a class template has its own, separate copy of the class's static data members.

If a class template has any specializations, the compiler must choose which specialization to instantiate. It starts by choosing which partial specializations match the template arguments. A partial specialization matches if the specialized template arguments can be deduced from the actual template arguments. See [Section 7.3.2](#) earlier in this chapter.

If no specializations match, the compiler instantiates the primary template. If one specialization matches, that one is instantiated. Otherwise, the best match is chosen, in which "best" means most specific. An error results if two or more specializations are tied for best. Template specialization A is at least as specific as template specialization B if the parameters of A can be deduced from B's template parameter list. No implicit conversion takes place when comparing specializations. The following example shows several instances of the demo class template. The template instantiations are instances of the primary template or one of the three partial specializations:

```
template<typename T, typename U, int N> class demo {};  
template<typename T, int N> class demo<T, int, N> {};  
template<typename T> class demo<T, T, 0> {};  
template<typename T> class demo<T*, T, 1> {};  
demo<int, char, 1> w;      // Primary template  
demo<char, int, 10> x;    // First specialization  
demo<char, char, 0> y;    // Second specialization
```

```
demo<char, char, 1> z;    // Primary template
demo<char*,char, 1> p;   // Third template
demo<char*,char, 0> q;   // Primary template
demo<char*,char*, 0> r;  // Second template
demo<int*, int, 1> s;    // Error: ambiguous
```

In addition to the obvious points at which a function or class is needed (e.g., calling the function, taking the address of the function, declaring an object whose type is the class, or casting an object to the class), there are more subtle points of instantiation. For example, a template can be instantiated if it is used in the value of a default function argument, and the default argument is needed in a function call.



A function template is needed if an instance of the template is needed for overload resolution. If a class template participates in function overloading (that is, as part of a function parameter type), but the template does not need to be instantiated to resolve the overload (e.g., because the function parameter is a pointer or reference to the class template), the class template may or may not be instantiated. The standard leaves this up to the implementation.

If a template declaration contains an error, the compiler might diagnose the error when it compiles the template or when it instantiates the template. If a program does not instantiate the template, you might find that you can compile a program successfully using one compiler (which reports errors only when a template is instantiated), but not a different compiler (which reports errors when a template is declared).

You can explicitly instantiate a template. Use a bare template keyword followed by a declaration that supplies the template arguments for the template:

```
template long int plus<long int>(long int a, long int b);
template short plus(short a, short b); // Deduce plus<short>.
template struct wrapper<unsigned const char*>;
```

The class, function, or member template must be defined before it can be instantiated. If you are instantiating a specialization, the template specialization must appear before the instantiation in the source. The behavior is undefined if a template is specialized after it has been instantiated for the same template arguments. [Example 7-10](#) illustrates template instantiation.

### *Example 7-10. Instantiating templates*

```
#include <iomanip>
#include <iostream>
#include <ostream>
const double pi = 3.1415926535897932;
// Function template
template<typename T> T sqr(T x)
{
    return x * x;
}

// Class template
template<typename T>
class circle
{
public:
    circle(T r) : radius_(r) {}
    // sqr<> is instantiated when circle<> is instantiated.
    T area( ) const { return pi * sqr(radius_); }
    T radius( ) const { return radius_; }
private:
    T radius_;
};

// Function template
template<typename T>
void print(T obj)
{
    std::cout << obj << '\n';
}

// Overload the function template with another template.
template<typename T>
```

```

void print(const circle<T>& c)
{
    std::cout << "circle(" << c.radius( ) << ")\n";
} template int sqr<int>(int); // Explicit instantiation

// Explicit instantiation of circle<double> and implicit
instantiation of
// sqr<double>
template class circle<double>;

// Error: after instantiation of sqr<double>, illegal to
specialize it
template<> double sqr(double x)
{
    return x * x;
}

int main( )
{
    using namespace std;
    // No circle<> instance is needed yet, even to resolve
overloaded print
    // function.
    print(42);
    for (int i = 0; i < 10; ++i)
        // Implicit instantiation of sqr<int>
        cout << setw(2) << i << setw(4) << sqr(i) << '\n';
    // Instantiation of circle<float> and therefore sqr<float>
circle<float> unit(1.0f);
    // Implicit instantiation of print<circle<float> >
print(unit);
}

```



## Name Lookup

Templates introduce a new wrinkle to name lookup. (See [Chapter 2](#) for the non-template rules of name lookup.) When compiling a template, the compiler distinguishes between names that depend on the template parameters (called *dependent names*) and those that do not (*nondependent names*). Nondependent names are looked up normally when the template is declared. Dependent names, on the other hand, must wait until the template is instantiated, when the template arguments are bound to the parameters. Only then can the compiler know what those names truly mean. This is sometimes known as *two-phase lookup*.

## Dependent Names

This section describes dependent names, and the following section describes what the compiler does with them.

A dependent name can have different meanings in different template instantiations. In particular, a function is dependent if any of its arguments are type-dependent. An operator has a dependent name if any of its operands are type-dependent.

A *dependent type* is a type that can change meaning if a template parameter changes. The following are dependent types:

- The name of a type template parameter or template template parameter:

```
template<typename T> struct demo {T dependent; }
```

- A template instance with template arguments that are dependent:

```
template<typename T> class templ {};  
template<typename U> class demo {templ<U> dependent; }
```

- A nested class in a dependent class template (such as the class template that contains the nested class):

```
template<typename T> class demo {class  
    dependent {};
```

- An array that has a base type that is a dependent type:

```
template<typename T> class demo {T dep[1]; }
```

- An array with a size that is value-dependent (defined later in this section):

```
template<typename T> class demo {int dep[sizeof(T)]; }
```

- Pointers and references to dependent types or functions (that is, functions whose return types or parameter types are dependent or whose default arguments are dependent):

```
template<typename T> class demo {T& x; T (*func)( ); }
```

- A class, struct, or union that depends on a template parameter for a base class or member (note that a non-template class nested in a class template is always dependent):

```
template<typename T> class demo {class  
    nested { T x; }; };
```

- A const- or volatile-qualified version of a dependent type:

```
template<typename T> class demo {const  
    T x; };
```

- A qualified name, in which any qualifier is the name of a dependent type:

```
template<typename T> struct outer { struct inner {} };  
template<typename T> class demo {outer<T>::inner dep; };
```

*A type-dependent expression* has a dependent type. It can be any of the following:

- `this`, if the class type is dependent
- A qualified or unqualified name if its type is dependent
- A cast to a dependent type
- A new expression, creating an object of dependent type
- Any expression that is built from at least one dependent subexpression, except when the result type is not dependent, as in the following:
  - A `sizeof` or `typeid` expression
  - A member reference (with the `.` or `->` operators)
  - A `throw` expression
  - A `delete` expression

A constant expression can also be *value-dependent* if its type is dependent or if any subexpression is value-dependent. An identifier is value-dependent in the following cases:

- When it is the name of an object with a dependent type
- When it is the name of a value template parameter
- When it is a constant of integral or enumerated type, and its initial value is a value-dependent expression

A `sizeof` expression is value-dependent only if its operand is type-dependent. A cast expression is dependent if its operand is a dependent expression. [Example 7-11](#) shows a variety of dependent types and expressions.

### *Example 7-11. Dependent names*

```
template<typename T> struct base {
    typedef T value_type; // value_type is dependent.
    void func(T*);        // func is dependent.
    void proc(int);       // proc is nondependent.

    class inner {         // inner is dependent.
        int x;           // x is nondependent.
    };
};
```

```

};
template<unsigned N>
class data {          // data is dependent.
    int array[N];    // array is dependent.
};

class demo : inner { // demo is dependent.
    char y[sizeof(T)]; // y is dependent.
};
};

int main( )
{
    base<int> b;
}

```

## Resolving Names

When writing a template declaration or definition, you should use qualified names as much as possible. Use member expressions to refer to data members and member functions (e.g., `this->data`). If a name is a bare identifier, the name lookup rules are different from the rules for non-templates.

The compiler looks up unqualified, nondependent names at the point where the template is declared or defined. Dependent base classes are not searched (because, at the point of declaration, the compiler does not know anything about the instantiation base class). This can give rise to surprising results. In the following example, the `get_x` member function does not see `base<T>::x`, so it returns the global `x` instead:

```

template<typename T> struct base {
    double x;
};
int x;
template<typename T>
struct derived : base<T> {
    int get_x( ) const { return x; } // Returns ::x
};

```

Dependent names are looked up twice: first in the context of the declaration and later in the context of the instantiation. In particular, when performing argument-dependent name lookup ([Chapter 2](#)), the compiler searches the declaration and instantiation namespaces for the function argument types.

Essentially, the instantiation context is the innermost namespace scope that encloses the template instantiation. For example, a template instance at global scope has the global scope as its instantiation context. The context for an instance that is local to a function is the namespace where the function is defined. Thus, the instantiation context never includes local declarations, so dependent names are never looked up in the local scope.

A function template can have multiple instantiation points for the same template arguments in a single source file. A class template can have multiple instantiations for the same template arguments in multiple source files. If the different contexts for the different instantiations result in different definitions of the templates for the same template arguments, the behavior is undefined. The best way to avoid this undefined behavior is to avoid using unqualified dependent names.

[Example 7-12](#) shows several ways dependent name lookup affects a template. In particular, note that `iterator_base` can refer to its members without qualification or member expressions. However, the derived classes, such as `iterator`, must use `this->` or qualify the member name with the base class because the base class is not searched for unqualified names. The `print` member function is also interesting. It prints the array by using an `ostream_iterator`, which calls `operator<<` to print each element. The name `operator<<` is dependent, so it is not looked up when the template is declared, but when the

template is instantiated. At that time, the compiler knows the template argument, `big::integer`, so it also knows to search the `big` namespace for the right overloaded `operator<<`.

### *Example 7-12. Resolving dependent names*

```
#include <algorithm>
#include <iostream>
#include <iterator>
#include <ostream>
#include <stdexcept>

template<unsigned Size, typename T>
class array {
    template<unsigned Sz, typename U>
        friend class array<Sz,U>::iterator;
    template<unsigned Sz, typename U>
        friend class array<Sz,U>::const_iterator;
    class iterator_base {
    public:
        iterator_base& operator++( ) {
            ++ptr;
            return *this;
        }
        T operator*( ) const { check( ); return *ptr; }
    protected:
        iterator_base(T* s, T* p) : start(s), ptr(p) {}
        void check( ) const {
            if (ptr >= start + Size)
                throw std::out_of_range("iterator out of range");
        }
        T* ptr;
        T* start;
    };
    public:
        array( ): data(new T[Size]) {}
        class iterator : public iterator_base,
            public std::iterator<std::random_access_iterator_tag,T>
        {
        public:
            iterator(T* s, T* p) : iterator_base(s, p) {}
            operator const_iterator( ) const { return
const_iterator(this->start, this->ptr);
            }
            T& operator*( ) {
                iterator_base::check( );
            }
        };
};
```

```

        return *this->ptr;
    }
};
iterator begin( ) { return iterator(data, data); }
iterator end( ) { return iterator(data, data + Size); }
template<typename charT, typename traits>
void print(std::basic_ostream<charT,traits>& out)
const
{
    std::copy(begin( ), end( ), std::ostream_iterator<T>(out));
}
private:
    T* data;
};

namespace big {
class integer {
public:
    integer(int x = 0) : x_(x) {}
    operator int( ) const { return x_; }
private:
    int x_; // Actual big integer implementation is left as an
exercise.
};

template<typename charT, typename traits>
std::basic_ostream<charT,traits>&
operator<<(std::basic_ostream<charT,traits>& out,
const integer& i)
{
    out << int(i);
    return out;
}
}

int main( )
{
    const array<10, big::integer> a;
    a.print(std::cout);
}

```

## Hiding Names

When using templates, several situations can arise in which template parameters hide names that would be visible in a non-template class or function. Other situations arise in

which template parameter names are hidden by other names.

If a member of a class template is defined outside of the namespace declaration that contains the class template, a template parameter name hides members of the namespace:

```
namespace ns {
    template<typename T>
    struct demo {
        demo( );
        T x;
    };
    int z;
}
template<typename z>
ns::demo::demo( )
{
    x = z( ); // Template parameter z, not ns::z
}
```

Template parameter names are hidden in the following cases:

- If a member is defined outside of its class template, members of the class or class template hide template parameter names:

```
template<typename T>
struct demo {
    T x;
    demo( );
};
template<typename x>
demo::demo( ) { x = 10; } // Member x, not parameter x
```

- In the definition of a member of a class template that lies outside the class definition, or in the definition of a class template, a base class name or the name of a member of a base class hides a template parameter if the base class is nondependent:

```
struct base {
    typedef std::size_t SZ;
```



```

};
template<int SZ>
struct derived : base {
    SZ size; // base::SZ hides template parameter.
};

```

If a base class of a class template is a dependent type, it and its members do not hide template parameters, and the base class is not searched when unqualified names are looked up in the derived-class template.

## Using Type Names

When parsing a template definition, the compiler must know which names are types and which are objects or functions. Unqualified names are resolved using the normal name lookup rules (described earlier in this section). Qualified dependent names are resolved according to a simple rule: if the name is prefaced with `typename`, it is a type; otherwise, it is not a type.

Use `typename` only in template declarations and definitions, and only with qualified names. Although `typename` is meant to be used with dependent names, you can use it with nondependent names. [Example 7-13](#) shows a typical use of `typename`.

### *Example 7-13. Using `typename` in a template definition*

```

// Erase all items from an associative container for which a
// predicate returns
// true.
template<typename C, typename Pred>
void erase_if(C& c, Pred pred)
{
    // Need typename because iterator is qualified
    for (typename C::iterator it = c.begin( ); it != c.end( );)
        if (pred(*it))
            c.erase(*it++);
        else
            ++it;
}

```

## Tricks with Templates

Template syntax permits recursion, selection, and computation. In other words, it is a full-fledged programming language, albeit a language that is hard to write and even harder to read. The full scope and power of programming with templates is beyond the scope of this book. This section presents some tips for starting your own exploration of this exciting field. [Appendix B](#) tells you about some interesting projects in this area.

Suppose you want to write a function to round off floating-point values to a fixed number of decimal places. You decide to write the function by multiplying by a power of 10, rounding off to an integer, and dividing by the same power of 10. You can hardcode the amount, (e.g., 100) in the routine, but you prefer to use a template, in which a template parameter specifies the number of decimal digits to retain. To avoid computing the same power of 10 every time the function is called, you decide to use template programming to compute the constant at compile time.

The `ipower<>` template in [Example 7-14](#) uses recursion to compute the power of any integer raised to any nonnegative integer value. Three base cases for the recursion are defined as specializations: raising any value to the 0th power is always 1, raising 0 to any power is always 0, and raising 0 to the 0th power is undefined. (As an exercise, try to define `ipower` more efficiently.) Finally, the `ipower<>` class template is used to define the `round<>` function template.

### *Example 7-14. Computing at compile time*

```
template<int x, unsigned y>
struct ipower {
    enum { value = x * ipower<x, y-1>::value };
};
```

```

template<int x>
struct ipower<x, 0> {
    enum { value = 1 };
};
template<unsigned y>
struct ipower<0, y> {
    enum { value = 0 };
};
template<> struct ipower<0, 0> {};

// Round off a floating-point value to a fixed number of digits.
template<unsigned N, typename T>
T round(T x)
{
    if (x < 0.0)
        return std::floor(x * ipower<10,N>::value - 0.5) /
            ipower<10,N>::value;
    else
        return std::floor(x * ipower<10,N>::value + 0.5) /
            ipower<10,N>::value;
}

```

In addition to compile-time computation, you can write more complicated programs that are evaluated at compile time. For example, the Boost project (described in [Appendix B](#)) uses templates to create type lists, that is, lists of types that are manipulated at compile time. Such lists can greatly simplify certain programming tasks, such as implementing type traits. [Example 7-15](#) shows a simplified version of type lists. A list is defined recursively as an empty list or a node that contains a type (head) and a list (tail). (This definition should be familiar to anyone with experience using functional programming languages.)

#### *Example 7-15. Defining type lists*

```

struct empty {};
template<typename H, typename T>
struct node {
    typedef H head;
    typedef T tail;
};

template<typename T1 = empty, typename T2 = empty,
        typename T3 = empty, typename T4 = empty,

```

```

        typename T5 = empty, typename T6 = empty,
        typename T7 = empty, typename T8 = empty,
        typename T9 = empty, typename T10 = empty,
        typename T11 = empty, typename T12 = empty
>
struct list {
    typedef node<T1, node<T2, node<T3, node<T4,
        node<T5, node<T6, node<T7, node<T8,
        node<T9, node<T10, node<T11, node<T12,
        empty
        > > > > > > > > > > type;
};

template<typename L>
struct length {
    enum { value = 1 + length<typename L::tail>::value };
};
template<>
struct length<empty> {
    enum { value = 0 };
};

template<typename L>
struct is_empty {
    enum { value = false };
};
template<>
struct is_empty<empty> {
    enum { value = true };
};
};

```

Actions on type lists are inherently recursive. Thus, to count the number of items in a type list, count the head as one, and add the length of the tail. The recursion stops at the end of the list, which is implemented as a specialization of the `length<>` template for the empty type.

An important action when using type lists is to test membership. To do this, you must be able to compare two types to see if they are the same. The `is_same_type` class template uses partial specialization to determine when two types are the same. If the two types specified as template arguments are the same, the specialization sets `value` to `true`. If the arguments are different, the primary template

is instantiated, and value is false. [Example 7-16](#) shows `is_same_type` and how it is used in the `is_member` template.

*Example 7-16. Testing membership in a type list*

```
template<typename T, typename U>
struct is_same_type {
    enum { value = false };
};
template<typename T>
struct is_same_type<T, T> {
    enum { value = true };
};

template<typename T, typename L>
struct is_member {
    enum { value = is_same_type<T, typename L::head>::value
                || is_member<T, typename L::tail>::value };
};

template<typename T>
struct is_member<T, empty> {
    enum { value = false };
};
```

Once you can define type lists and test whether a type is in a type list, you can use these templates to implement simple type traits. For example, you can create a list of the integral types and test whether a type is one of the fundamental integral types. (Testing for integral types is important when implementing standard containers, as discussed in [Chapter 10](#).) [Example 7-17](#) shows some simple uses of type lists.

*Example 7-17. Using type lists*

```
#include <iostream>
#include <ostream>

typedef list<bool, char, unsigned char, signed char,
           int, short, long, unsigned,
           unsigned long, unsigned short>::type int_types;
typedef list<float, double, long double>::type real_types;
```

```
int main( )
{
    using namespace std;
    cout << is_same_type<int,int>::value << '\n';
    cout << is_same_type<int, signed int>::value << '\n';
    cout << is_same_type<int, unsigned int>::value << '\n';
    cout << is_member<int, int_types>::value << '\n';
    cout << is_member<float, int_types>::value << '\n';
    cout << is_member<ostream, int_types>::value << '\n';
}
```

# Compiling Templates

Like an ordinary function, a function template requires a definition before the function can be called. Like an ordinary class, a class template requires a definition for each member function and static data member before they can be used. Unlike ordinary functions or members, however, templates are typically defined in every source file.

Templates are often used by placing a template declaration and all supporting definitions in a header file, (e.g., `template.h`). Then `#include` that file anywhere the template is needed. For every implicit or explicit instantiation, the compiler generates the necessary code for the template's instantiation. If multiple source files instantiate the same template with the same arguments, the compiler and linker ensure that the program contains a single copy of the instantiated functions and members.

Different compilers use different techniques to ensure that the program contains a single copy of each template instance. The following are four different approaches:

- The most common approach is to have the compiler keep track of which source files require which instantiations. When the program is linked, the compiler combines all the lists of required instantiations and compiles the template instantiations at that time. As an optimization, the compiler saves the compiled instantiations, so an instantiation that does not change does not need to be recompiled.
- Another approach is to have the compiler generate the code for all needed instantiations when it compiles each source file. The linker identifies duplicate

instantiations and ensures that the linked program gets a single copy of each instantiation.

- A third approach is to separate the template declaration from its associated definitions. If the declaration is in `template.h`, the definitions are in `template.cc`. The compiler automatically locates the definition file from the name of the declaration file. In this scenario, the compiler keeps track of which instantiations it needs and compiles each of them only once.
- Another approach uses the `export` keyword to declare and define templates. An exported template lets you define a function template or all the members of a class template in a separate source file. The template's header contains only the declarations.



As I write this, exactly one compiler supports `export`. (See the book's web site for current details.) Major compiler vendors are finally moving toward full conformance with the standard, including `export`. Nonetheless, I doubt `export` will see widespread support, at least for the next few years.

If portability is important to you, do not use `export`.

Consult your compiler's documentation to learn how it handles templates. When writing a template that can be used by multiple compilers, a common technique is to put the declaration in a header file (e.g., `template.h`), and the definitions in another file (e.g., `template.cc`), and at the end of `template.h`, use conditional directives to `#include "template.cc"` for those compilers where it is needed. Use conditional compilation to define a macro only for compilers that support `export`. [Example 7-18](#) shows this common pattern.



### *Example 7-18. Declaring and defining a template*

```
// point.h
#ifndef POINT_H
#define POINT_H

#ifdef HAS_EXPORT
#define EXPORT export
#else
#define EXPORT
#endif

EXPORT template<typename T>
class point {
public:
    point(T a, T b);
    point( );
    T x( ) const { return x_; }
    T y( ) const { return y_; }
private:
    T x_, y_;
};

#ifdef NEED_TEMPLATE_DEFINITIONS
#include "point.cc"
#endif
#endif // POINT_H

// point.cc
#include "point.h"
EXPORT template<typename T>
point<T>::point(T a, T b)
: x_(a), y_(b)
{}

EXPORT template<typename T>
point<T>::point( )
: x_(T( )), y_(T( ))
{}

// program.cc
#include "point.h"
int main( )
{
    point<float> ptf;
    point<int> pti;
    ...
}
```

If your compiler supports export, define `HAS_EXPORT`. If the compiler requires template definitions in every source file, define `NEED_TEMPLATE_DEFINITIONS`. Most compilers offer a way to define macros globally (e.g., in a project definition file, in a makefile, etc.). Another alternative is to use conditional compilation to test the predefined macros that most compilers define, and use those to set the template macros accordingly. Put these definitions in a configuration file that is included first by every other file, as shown in [Example 7-19](#).

*Example 7-19. Configuring template compilation macros*

```
// config.h
#ifndef CONFIG_H
#define CONFIG_H

#ifdef __COMO__
    #define HAS_EXPORT
    #undef NEED_TEMPLATE_DEFINITIONS
#endif
#if defined(__BORLANDC__) || defined(__GNUC__)
    #undef HAS_EXPORT
    #define NEED_TEMPLATE_DEFINITIONS
#endif
...

#endif // CONFIG_H
```

## Chapter 8. Standard Library

The previous seven chapters describe the C++ language. This and the next two chapters describe the library. The library consists of a number of headers, in which each header declares types, macros, and functions for use in C++ programs. Much of the standard library is implemented using templates, so the implementation is typically embodied entirely in header files. Nonetheless, parts of the library are separately compiled and must be linked into a C++ program, the details of which are dictated by the implementation.

This chapter presents an overview of the standard library and some of its more important components: wide and multibyte characters, traits, policies, allocators, and numerics. Other important parts of the library get their own chapters: [Chapter 9](#) introduces the input and output classes, and [Chapter 10](#) describes the containers, iterators, and algorithms that make up what is most often known as the "standard template library," or STL.

### Overview of the Standard Library

The standard library has 51 headers, each containing a set of macro, function, type, and object declarations. A header is sometimes called a header file, but that phrase is misleading. An implementation does not need to implement headers as external files, although it is often simplest to think of a header as a file that contains macro, type, and function declarations.

Almost all the names declared in the standard library are in the `std` namespace. Macros, of course, are not in any namespace, so it is important that you know which names are macros and which are not. The detailed descriptions in

[Chapter 13](#) tell you this information. The only other names that are outside the `std` namespace are the global operator `new` and operator `delete` functions, declared in the `<new>` header.



To use the standard library, you must `#include` the desired header or headers. Some implementations `#include` headers within other headers (e.g., `<set>` might `#include <iterator>`). The headers that a vendor includes within a header vary from one implementation to another, so code that compiles successfully with one vendor's environment might not compile in a different environment. The solution is to get in the habit of including all the headers you need.

### **Warning**

If you see headers of the form `<iostream.h>`, the program was probably written in the days before the C++ standard. It might also mean that your compiler and library were written before C++ was standardized. Although some popular, old compilers are still in wide use, they have newer versions that provide much better support for the C++ standard. Following the C++ standard is the best way to achieve portability across recent compilers and libraries.

You must not add any declarations to the `std` namespace, although you can specialize templates that are declared in `std`. When you add your own specialization, you must obey all the constraints specified for the template, and the specialization must depend on at least one user-defined name that has external linkage. Otherwise, the behavior is undefined.

The following are brief descriptions of the contents of each header. The headers inherited from the C standard are marked as "C header." Some of these headers have improved C++ equivalents, which are also shown. For complete descriptions of these headers, see [Chapter 13](#).

`<algorithm>`

Standard algorithms for copying, searching, sorting, and otherwise operating on iterators and containers. See [Chapter 10](#) for more information about the standard algorithms.

`<bitset>`

Class template to hold a fixed-sized sequence of bits.

`<cassert>`

Runtime assertion-checking; C header.

`<cctype>`

Character classification and case conversion; C header (see also `<locale>`).

`<cerrno>`

Error codes; C header.

`<cfloat>`

Limits of floating-point types; C header (see also `<limits>`).

`<ciso646>`

Empty header because C declarations are incorporated in the C++ language; C header.

`<climits>`

Limits of integer types; C header (see also `<limits>`).

`<locale>`

Locale-specific information; C header (see also `<locale>`).

`<cmath>`

Mathematical functions; C header.

`<complex>`

Complex numbers.

`<csetjmp>`

Nonlocal goto; C header.

`<csignal>`

Asynchronous signals; C header.

`<cstdarg>`

Macros to help implement functions that take a variable number of arguments; C header.

`<stddef>`

Miscellaneous standard definitions; C header.

`<stdio>`

Standard input and output; C header (see also `<iostream>` and related headers).

`<stdlib>`

Miscellaneous functions and related declarations; C header.

`<string>`

String-handling functions; C header (see also `<string>`).

`<ctime>`

Date and time functions and types; C header.

`<wchar>`

Wide character functions, including I/O; C header (see also `<locale>`, `<iostream>`, `<string>`, and other I/O-related headers).

`<cwctype>`

Wide character classification and case conversion; C header (see also `<locale>`).

`<deque>`

Deque (double-ended queue) standard container.

`<exception>`

Base exception class and functions related to exception-handling.

`<fstream>`

File-based stream I/O.

`<functional>`

Function objects; typically used with standard algorithms.

`<iomanip>`

I/O manipulators; used with standard I/O streams.

`<ios>`

Base class declarations for all I/O streams.

`<iosfwd>`

Forward declarations for I/O objects.

`<iostream>`

Declarations of standard I/O objects.

`<istream>`

Input streams and input/output streams.

`<iterator>`

Additional iterators for working with standard containers and algorithms. See [Chapter 10](#) for more information.

<limits>

Limits of numerical types.

<list>

Standard linked list container.

<locale>

Locale-specific information for formatting and parsing numbers, dates, times, and currency values, plus character-related functions for classifying, converting, and comparing characters and strings.

<map>

Associative map (sometimes called a dictionary) standard container.

<memory>

Allocators, algorithms for uninitialized memory, and smart pointers (`auto_ptr`).

<new>

Global operator `new` and operator `delete` and other functions related to managing dynamic memory.

<numeric>

Numerical algorithms.

<ostream>

Output streams.

<queue>

Queue and priority queue container adapters.

<set>



Associative set container.

`<sstream>`

String-based I/O streams.

`<stack>`

Stack container adapter.

`<stdexcept>`

Standard exception classes.

`<streambuf>`

Low-level stream buffers; used by high-level I/O streams.

`<string>`

Strings and wide-character strings.

`<stringstream>`

String streams that work with character arrays (see also `<sstream>`).

`<typeinfo>`

Runtime type information.

`<utility>`

Miscellaneous templates, such as `pair`, most often used with standard containers and algorithms.

`<valarray>`

Numerical arrays.

`<vector>`

Vector (array-like) standard container.

## C Library Wrappers

The C++ library includes the entire C standard library (from the 1990 C standard, plus Amendment 1), in which each C header, such as `<stdio.h>`, is wrapped as a C++ header (e.g., `<cstdio>`). Being part of the C++ standard, all types, functions, and objects are declared in the `std` namespace.

The external names are also reserved in the global namespace. Thus, proper practice is to use the names in the `std` namespace (e.g., `std::strlen`), but realize that these names are also reserved in the global namespace, so you cannot write your own `::strlen` function.

The C standard permits macros to be defined to mask function names. In the C++ wrappers for these headers, the names must be declared as functions, not macros. Thus, the C `<stdio.h>` header might contain the following:

```
extern int printf(const char* fmt, ...);
#define printf printf
```

In C++, the `printf` macro is not permitted, so the `<cstdio>` header must declare the `printf` function in the `std` namespace, so you can use it as `std::printf`.

A deprecated feature of C++ is that the C standard headers are also available as their original C names (e.g., `<stdio.h>`). When used in this fashion, their names are in the global namespace, as though a `using` declaration were applied to each name (e.g., `using std::printf`).

Otherwise, the old style headers are equivalent to the new headers. The old C header names are deprecated; new code should use the `<cstdio>`, etc., style C headers.

## Wide and Multibyte Characters

The familiar `char` type is sometimes called a *narrow character*, as opposed to `wchar_t`, which is a *wide character*. The key difference between a narrow and wide character is that a wide character can represent any single character in any character set that an implementation supports. A narrow character, on the other hand, might be too small to represent all characters, so multiple narrow `char` objects can make up a single, logical character called a *multibyte character*.

Beyond some minimal requirements for the character sets (see [Chapter 1](#)), the C++ standard is purposely open-ended and imposes few restrictions on an implementation. Some basic behavioral requirements are that conversion from a narrow character to a wide character must produce an equivalent character. Converting back to a narrow character must restore the original character. The open nature of the standard gives the compiler and library vendor wide latitude. For example, a compiler for Japanese customers might support a variety of Japanese Industrial Standard (JIS) character sets, but not any European character sets. Another vendor might support multiple ISO 8859 character sets for Western and Eastern Europe, but not any Asian multibyte character sets. Although the standard defines universal characters in terms of the Unicode (ISO/IEC 10646) standard, it does not require any support for Unicode character sets.



This section discusses some of the broad issues in dealing with wide and multibyte characters, but the details of specific characters and character sets are implementation-defined.

## Wide Characters

A program that must deal with international character sets might work entirely with wide characters. Although wide characters usually require more memory than narrow characters, they are usually easier to use. Searching for substrings in a wide string is easy because you never have the problem of matching partial characters (which can happen with multibyte characters).

A common implementation of `wchar_t` is to use Unicode UTF-32 encoding, which means each wide character is 32 bits and represents a single Unicode character. Suppose you want to declare a wide string that contains the Greek letter pi ( $\pi$ ). You can specify the string with a universal name (see [Chapter 1](#)):

```
wchar_t wpi[] = "\u03c0";
```

Using UTF-32, the string would contain `L"\x03c0"`. With a different `wchar_t` implementation, the `wpi` string would contain different values.

The standard `wstring` class supports wide strings, and all the I/O streams support wide characters (e.g., `wistream`, `wostream`).

## Multibyte Characters

A multibyte character represents a single character as a series of one or more bytes, or narrow characters. Because a single character might occupy multiple bytes, working with multibyte strings is more difficult than working with wide strings. For example, if you search a multibyte string for the character `'\x20'`, when you find a match, you must test whether the matching character is actually part of a multibyte character and is therefore not actually a match for the single character you want to find.

Consider the problem of comparing multibyte strings. Suppose you need to sort the strings in ascending order. If one string starts with the character '\xA1' and other starts with '\xB2', it seems that the first is smaller than the second and therefore should come before the second. On the other hand, these characters might be the first of multibyte character sequences, so the strings cannot be compared until you have analyzed the strings for multibyte character sequences.

Multibyte character sets abound, and a particular C++ compiler and library might support only one or just a few. Some multibyte character sets specifically support a particular language, such as the Chinese Big5 character set. The UTF-8 character set supports all Unicode characters using one to six narrow characters.

For example, consider how an implementation might encode the Greek letter pi ( $\pi$ ), which has a Unicode value of 0x03C0:

```
char pi[] = "\u03c0";
```



If the implementation's narrow character set is ISO 8859-7 (8-bit Greek), the encoding is 0xF0, so pi[] contains "\xf0". If the narrow character set is UTF-8 (8-bit Unicode), the representation is a multibyte character, and pi[] would contain "\xe0\x8f\x80". Many character sets do not have any encoding for  $\pi$ , in which case the contents of pi[] might be "?", or some other implementation-defined marker for unknown characters.

## Shift State

You can convert a multibyte character sequence to a wide character and back using the functions in <wchar>. When

performing such conversions, the library might need to keep track of state information during the conversion. This is known as the *shift state* and is stored in an `mbstate_t` object (also defined in `<wchar>`).

For example, the Japanese Industrial Standard (JIS) encodes single-byte characters and double-byte characters. A 3-byte character sequence shifts from single- to double-byte mode, and another sequence shifts back. The shift state keeps track of the current mode. The initial shift state is single-byte. Thus, the multibyte string `"\x1B$B&P\x1B(B"` represents one wide character, namely, the Greek letter pi ( $\pi$ ). The first three characters switch to double-byte mode. The next two characters encode the character, and the final three characters restore single-byte mode.

Shift states are especially important when performing I/O. By definition, file I/O uses multibyte characters. That is, a file is treated as a sequence of narrow characters. When reading a wide-character stream, the narrow characters are converted to wide characters, and when writing a wide stream, wide characters are converted back to multibyte characters. Seeking to a new position in a file might seek to a position that falls in the middle of a multibyte sequence. Therefore, a file position is required to keep track of a shift state in addition to a byte position in the file. See `<ios>` in [Chapter 13](#).

## Traits and Policies

Traits are used throughout the C++ library. A *trait* is a class or class template that characterizes a type, possibly a template parameter. At first glance, it seems that traits obscure information, hiding types and other declarations in a morass of templates. This is true, but traits are also powerful tools used in writing templates. Traits are often used to obtain information from built-in types in the same manner as user-defined types.

A *policy* is a class or class template that defines an interface as a service to other classes. Traits define type interfaces, and policies define function interfaces, so they are closely related. Sometimes, a single class template implements traits and policies.

The typical application programmer might never use traits and policies directly. Indirectly, however, they are used in the `string` class, I/O streams, the standard containers, and iterators—just about everywhere in the standard library.

## Character Traits

One of the most commonly used trait and policy templates is `char_traits<>` in the `<string>` header. The standard declares two specializations: `char_traits<char>` and `char_traits<wchar_t>`.

The rest of the C++ library uses character traits to obtain types and functions for working with characters. For example, the `basic_istream` class template takes a character type, `charT`, and a character traits type as template parameters. The default value for the traits parameter is `char_traits<charT>`, which is a set of character traits defined in `<string>`. The `basic_istream` template declares the `get( )` function, which reads a

character and returns its integer equivalent. The return type is obtained from the character traits template, specifically `int_type`.

As a policy template, `char_traits<>` provides member functions that compare characters and character arrays, copy character arrays, and so on. For example, `compare` compares two character arrays for equality. The `char_traits<char>` specialization might implement `compare` by calling `memcmp`.

At a basic level, the typical C++ programmer does not need to be concerned with the implementation of traits. Instead, you can use the `istream` and `string` classes, and everything just works. If you are curious, you can trace the declaration of, for example, `istream::int_type`:

```
istream::int_type →  
basic_istream<char>::int_type → traits::int_type  
→ char_traits<char>::int_type → int
```

As you can see, traits can be difficult to follow when you need to know the exact type of one of the types declared in a standard container.

Once you get used to them, however, you can see how valuable traits can be. Consider what happens when you change from `istream::int_type` to `wistream::int_type`:

```
wistream::int_type →  
basic_istream<wchar_t>::int_type →  
traits::int_type →  
char_traits<wchar_t>::int_type → wint_t
```

Note that the declarations of `basic_istream` and the other templates do not differ when the template parameter changes from `char` to `wchar_t`. Instead, you end up with a different template specialization for `char_traits<>`, which directs you to a different integer type.



You can implement your own character traits and policy template. For example, suppose you want to use strings that compare themselves without regard to case differences. Comparison is a policy issue, typically implemented by the `char_traits<>` template. You can define your own template that has the same trait and policy implementation, but one that implements compare to ignore case differences. Using your template, you can specialize `basic_string<>` to create a case-insensitive string class and then store those strings in sets and maps. The keys will be compared using your policy function that ignores case differences, as shown in [Example 8-1](#).

*Example 8-1. Case-insensitive character policy*

```
template<typename T> struct ci_char_traits {};template<> struct  
ci_char_traits<char> {  
    typedef char char_type;  
    typedef int int_type;  
    typedef std::streamoff off_type;  
    typedef std::streampos pos_type;  
    typedef std::mbstate_t state_type;  
  
    static void assign(char_type& dst, const char_type src) {  
        dst = src;  
    }  
    static char_type* assign(char* dst, std::size_t n, char c) {  
        return static_cast<char_type*>(std::memset(dst, n, c));  
    }  
    static bool eq(const char_type& c1, const char_type& c2) {  
        return lower(c1) == lower(c2);  
    }  
    static bool lt(const char_type& c1, const char_type& c2) {  
        return lower(c1) < lower(c2);  
    }  
    static int compare(const char_type* s1,  
const char_type* s2, std::size_t n) {  
        for (size_t i = 0; i < n; ++i)  
        {  
            char_type lc1 = lower(s1[i]);  
            char_type lc2 = lower(s2[i]);  
            if (lc1 < lc2)  
                return -1;  
            else if (lc1 > lc2)
```

```

        return 1;
    }
    return 0;
}
...
private:
    static int_type lower(char_type c) {
        return std::tolower(to_int_type(c));
    }
};

typedef std::basic_string<char, ci_char_traits<char> >
    ci_string;

void print(const std::pair<const ci_string, std::size_t>& item)
{
    std::cout << item.first << '\t' << item.second << '\n';
}

int main( )
{
    std::map<ci_string, std::size_t> count;
    ci_string word;
    while (std::cin >> word)
        ++count[word];
    std::for_each(count.begin( ), count.end( ), print);
}

```

## Iterator Traits

Traits are also useful for iterators ([Chapter 10](#)). An algorithm often needs to know the iterator category to provide specializations that optimize performance for random access iterators, for example. Traits provide a standard way to convey this information to the algorithm — namely, by using the `iterator_category` typedef. They also permit algorithms to use plain pointers as iterators.

For example, the `distance` function returns the distance between two iterators. For random access iterators, the distance can be computed by subtraction. For other iterators, the distance must be computed by incrementing an iterator and counting the number of increments needed.

[Example 8-2](#) shows a simple implementation of distance that uses the iterator traits to choose the optimized random access implementation or the slower implementation for all other input iterators.

*Example 8-2. Implementing the distance function template*

```
// Helper function, overloaded for random access iterators
template<typename InputIter>
typename std::iterator_traits<InputIter>::difference_type
compute_dist(InputIter first, InputIter last,
             std::random_access_iterator_tag)
{
    return last - first;
}

// Helper function, overloaded for all other input iterators
template<typename InputIter>
typename std::iterator_traits<InputIter>::difference_type
compute_dist(InputIter first, InputIter last,
             std::input_iterator_tag)
{
    typename std::iterator_traits<InputIter>::difference_type
        count = 0;
    while (first != last) {
        ++first;
        ++count;
    }
    return count;
}

// Main distance function, which calls the helper function, using
// the iterator
// tag to differentiate the overloaded functions.
template<typename InputIter>
typename std::iterator_traits<InputIter>::difference_type
distance(InputIter first, InputIter last)
{
    return compute_dist(first, last,
                       std::iterator_traits<InputIter>::iterator_category( ));
}
```

Being able to optimize algorithms for certain kinds of iterators is one benefit of using traits, but the real power comes from the `iterator_traits<T*>` specialization. This class permits the use of any pointer type as an iterator.

(See `<iterator>` in [Chapter 13](#) for details.) Consider how the `distance` function is called in the following example:

```
int data[] = { 10, 42, 69, 13, 100, -1 };
distance(&data[1], &data[4]);
```

The compiler infers the `InputIter` template parameter as type `int*`. The `iterator_traits<T*>` template is expanded to obtain the `iterator_category` type (`random_access_iterator_tag`) and `difference_type` (`ptrdiff_t`).

## Custom Traits

Traits can be useful whenever you are using templates. You never know what the template parameters might be. Sometimes, you want to specialize your own code according to a template parameter.

For example, all the standard sequence containers have a constructor that takes two iterators as arguments:

```
template<typename InputIterator>
list(InputIterator first, InputIterator last);
```

But take a closer look at the declaration. The author's intent is clear: that the template parameter must be an input iterator, but nothing in the declaration enforces this restriction. The compiler allows any type to be used (at least any type that can be copied).

If the `InputIterator` type actually is an input iterator, the `list` is constructed by copying all the elements in the range `[first, last)`. But if the `InputIterator` type is an integral type, the `first` argument is interpreted as a count, and the `last` argument is interpreted as an integer value, which is converted to the value type of the container; the container is then initialized with `first` copies of the `last` value, which is ordinarily the work of a different constructor. See [Chapter 10](#) for more information about these constructors.

If you need to implement your own container template, you must find a way to implement this kind of constructor. The simplest way is to define a traits template that can tell you whether a type is an integral type. [Example 8-3](#) shows one possible implementation and how it can be used by a container.

*Example 8-3. Differentiating between integral types using traits*

```
// Type trait to test whether a type is an integer.
struct is_integer_tag {};
struct is_not_integer_tag {};

// The default is that a type is not an integral type.
template<typename T>
struct is_integer {
    typedef is_not_integer_tag tag;
};

// Override the default explicitly for all integral types.
template<>
struct is_integer<int> {
    typedef is_integer_tag tag;
};
template<>
struct is_integer<short> {
    typedef is_integer_tag tag;
};
template<>
struct is_integer<unsigned> {
    typedef is_integer_tag tag;
};
// And so on for char, signed char, short, etc.

// Constructor uses the is_integer trait to distinguish integral
// from nonintegral
// types and dispatches to the correct overloaded construct
// function.
template<typename T, typename A>
template<typename InputIter>
list<T,A>::list(InputIter first, InputIter last)
{
    construct(first, last, is_integer<InputIter>::tag( ));
}
}
```

```

// The construct member functions finish the initialization of
the list. The
// integral version casts the arguments to the size and value
types.
template<typename T, typename A>
template<typename InputIter>
void list<T,A>::construct(InputIter first, InputIter last,
                          is_integer_tag)
{
    insert(begin( ), static_cast<size_type>(first),
           static_cast<T>(last));
}

// The non-integral version copies elements from the iterator
range.
template<typename T, typename A>
template<typename InputIter>
void list<T,A>::construct(InputIter first, InputIter last,
                          is_not_integer_tag)
{
    insert(begin( ), first, last);
}

```

Traits can be used to characterize any type and specialize templates for a wide variety of situations. See the Boost project (described in [Appendix B](#)) for other definitions and uses of traits.

## Allocators

An *allocator* is a policy class that defines an interface for managing dynamic memory. You already know about the `new` and `delete` expressions for allocating and freeing dynamic memory. They are simple, expressive, and useful, but the standard library does not necessarily use them internally. Instead, the standard library uses allocators, which let you provide alternative mechanisms for allocating and freeing memory.

The standard library provides a standard allocator (see `<memory>` in [Chapter 13](#)). If you don't want to use the standard allocator, you can use your own, provided it satisfies the same interface that is defined by the standard allocator.

## Using Allocators

An allocator is a simple object that manages dynamic memory, abstracting `new` and `delete` expressions. All the container class templates take an allocator template parameter and use the allocator to manage their internal memory. You can use allocators in your own container classes or wherever you want to offer flexibility to the user of your class or template.

If you do not want to bother with allocators, you don't need to. All the standard containers have a default argument for their allocator template parameters: `std::allocator`, which uses standard `new` and `delete` expressions to manage dynamic memory.

If you write a new container class template, make sure it takes an allocator parameter, as the standard containers do. Use the allocator to manage internal memory for your

container. See [Chapter 10](#) for more information about containers.

If you simply want to use an allocator to manage memory, you can do so. Your class would use the allocator to allocate and free memory, initialize and finalize objects, and take the address of an allocated object. (See `<memory>` in [Chapter 13](#) for a complete description of the allocator policy interface.) [Example 8-4](#) shows a simple class that wraps a dynamic instance of any object. It is not particularly useful, but it illustrates how a class can use an allocator. Note how the allocation of memory is separated from the construction of the object. If `allocate` fails to allocate the desired memory, it throws `bad_alloc`, so the wrapper constructor fails before it tries to construct the object. If the allocation succeeds, but the call to `construct` fails, the memory must be freed, hence the `try` statement. The destructor assumes that wrapped class is well-written and never throws an exception.

#### *Example 8-4. Wrapping a dynamic object*

```
template<typename T, typename Alloc=std::allocator<T> >
class wrapper {
public:
    typedef T value_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::reference reference;

    // Allocate and save a copy of obj.
    wrapper(const T& obj = T( ), const Alloc& a = Alloc( ))
        : alloc_(a), ptr_(0)
    {
        T* p = a.allocate(sizeof(T));
        try {
            alloc_.construct(p, obj);
        } catch(...) {
            // If the construction fails, free the memory without
            trying to finalize
            // the (uninitialized) object.
            alloc_.deallocate(p);
            throw;
        }
    }
};
```



```

    // Everything succeeded, so save the pointer.
    ptr_ = p;
}
~wrapper( )
{
    alloc_.destroy(ptr_);
    alloc_.deallocate(ptr_);
}
typename Alloc::reference operator*( ) { return *ptr_; }
value_type operator*( ) const        { return *ptr_; }
private:
    Alloc alloc_;
    typename Alloc::pointer ptr_;
};

```

## Custom Allocators

Writing a custom allocator requires care and patience. One particularly difficult point is that an implementation of the standard library is free to assume that all instances of an allocator class are equivalent, that is, allocators cannot maintain state. The standard permits this behavior without mandating it.

Thus, you need to be aware of the standard library's requirements for your implementation. Once you know the requirements, you can write a custom allocator. As a starting point, see [Chapter 13](#) (under <memory>), which implements the allocator as trivial wrappers around new and delete expressions. Other allocators might manage memory that is shared between processes or differentiate between different kinds of pointers (such as near and far pointers found on old PC operating systems). More sophisticated allocators can implement debugging or validity checks to detect programmer errors, such as memory leaks or double frees.

## Numerics

The C++ library has several headers that support numerical programming. The most basic header is `<cmath>`, which declares transcendental and other mathematical functions. This C header is expanded in C++ to declare overloaded versions of every function. For example, whereas C declares only `exp(double)`, C++ also declares `exp(float)` and `exp(long double)`.

The `<complex>` header declares a class template for complex numbers, with the specializations you would probably expect for `float`, `double`, and `long double`. Transcendental and I/O functions are also declared for complex numbers.

The `<numeric>` header declares a few algorithms (that use standard iterators) for numerical sequences.

The most interesting numerical functions are in `<valarray>`. A `valarray` is like an ordinary numerical array, but the compiler is free to make some simplifying assumptions to improve optimization. A `valarray` is not a container, so it cannot be used with standard algorithms or the standard numeric algorithms.

You can use the `complex` template as an example of how to define custom numeric types. For example, suppose you want to define a type to represent rational numbers (fractions). To use rational objects in a `valarray`, you must define the class so it behaves the same as ordinary values, such as `ints`. In other words, a custom numeric type should have the following:

- A public default constructor (e.g., `rational( )`)
- A public copy constructor (e.g., `rational(const rational&)`)

- A public destructor
- A public assignment operator (e.g., `rational& operator=(const rational&)`)
- Reasonable arithmetic and comparison operators

When you implement a class, make sure that the copy constructor and assignment operator have similar, reasonable results.

When overloading arithmetic and comparison operators, think about which operators are meaningful (e.g., most arithmetic types should have addition, subtraction, multiplication, and division, but only integral types will probably have remainder, bitwise, and shift operators). You should provide overloaded operators that accept the built-in types as operands. For example, the `complex` template defines the following functions for `operator+`:

```
template<typename T> complex<T> operator+(const complex<T>& z);
template<typename T> complex<T> operator+(const complex<T>& x,
const
                                complex<T>& y);
template<typename T> complex<T> operator+(const complex<T>& x,
const T& y);
template<typename T> complex<T> operator+(const T& x, const
complex<T>& y);
```

[Example 8-5](#) shows excerpts from a simple `rational` class template.

*Example 8-5. The rational class template for rational numbers*

```
template<typename T>
class rational
{
public:
    typedef T value_type;

    rational( )                : num_(0), den_(1) {}
    rational(value_type num)   : num_(num), den_(1) {}
    rational(value_type num, value_type den)
        : num_(num), den_(den) { reduce( ) ; }
```

```

rational(const rational& r): num_(r.num_), den_(r.den_) {}
template<typename U>
rational(const rational<U>& r)
    : num_(r.num_), den_(r.den_) { reduce( ); }

rational& operator=(const rational& r)
    { num_ = r.num_; den_ = r.den_; return *this; }
template<typename U>
rational& operator=(const rational<U>& r)
    { assign(r.numerator( ), r.denominator( )); return *this; }

void assign(value_type n, value_type d)
    { num_ = n; den_ = d; reduce( ); }

value_type numerator( ) const { return num_; }
value_type denominator( ) const { return den_; }

private:
    void reduce( );
    value_type num_;
    value_type den_;
};

// Reduce the numerator and denominator by the gcd. Make sure
// that the
// denominator is nonnegative.
template<typename T>
void rational<T>::reduce( )
{
    if (den_ < 0) {
        den_ = -den_;
        num_ = -num_;
    }
    T d = gcd(num_, den_);
    num_ /= d;
    den_ /= d;
}

// Greatest common divisor using Euclid's algorithm
template<typename T>
T gcd(T n, T d)
{
    n = abs(n);
    while (d != 0) {
        T t = n % d;
        n = d;
        d = t;
    }
}

```

```

    return n;
}

// Multiplication assignment operator. Often implemented as a
// member function,
// but there is no need to do so.
template<typename T, typename U>
rational<T>& operator*=(rational<T>& dst,
                      const rational<U>& src)
{
    dst.assign(dst.numerator( ) * src.numerator( ),
              dst.denominator( ) * src.denominator( ));
    return dst;
}

// Multiply two rational numbers.
template<typename T>
rational<T> operator*(const rational<T>& a,
                    const rational<T>& b)
{
    rational<T> result(a);
    result *= b;
    return result;
}

// Multiply rational by an integral value.
template<typename T>
rational<T> operator*(const T& a, const rational<T>& b)
{
    return rational<T>(a * b.numerator( ), b.denominator( ));
}

template<typename T>
rational<T> operator*(const rational<T>& a, const T& b)
{
    return rational<T>(b * a.numerator( ), a.denominator( ));
}

// Other arithmetic operators are similar.
// Comparison. All other comparisons can be implemented in terms
// of operator==
// and operator<.
template<typename T>
bool operator==(const rational<T>& a, const rational<T>& b)
{
    // Precondition. Both operands are reduced.
    return a.numerator( ) == b.numerator( ) &&
           a.denominator( ) == b.denominator( );
}

```

```
template<typename T>
bool operator<(const rational<T>& a, const rational<T>& b)
{
    return a.numerator( ) * b.denominator( ) <
           b.numerator( ) * a.denominator( );
}
```

Many numerical programmers find the C++ standard library to be lacking. However, the Blitz++ project is a popular, high-performance numerical library. Boost also has some numerical headers, such as a full-featured rational class template. See [Appendix B](#) for information about these and other C++ libraries.

## Chapter 9. Input and Output

C++ has a rich I/O library, which is often called *I/O streams*. This chapter presents an overview of the C++ I/O library. For details of individual classes and functions, see [Chapter 13](#). The I/O streams often use templates; refer to [Chapter 7](#) for information about templates. See also [Chapter 8](#) for information about character traits.

The standard I/O streams can also be used with standard iterator adapters. See [Chapter 10](#) for details.

### Introduction to I/O Streams

As with C and many modern languages, input and output in C++ is implemented entirely in the library. No language features specifically support I/O.

The C++ I/O library is based on a set of templates parameterized on the character type. Thus, you can read and write plain `char`-type characters, wide `wchar_t` characters, or some other, exotic character type that you might need to invent. (Read about character traits in [Chapter 8](#) first.) [Figure 9-1](#) depicts the class hierarchy. Notice that the names are of the form `basic_name`. These are the template names; the specializations have the more familiar names (e.g., `istream` specializes `basic_istream<char>`).

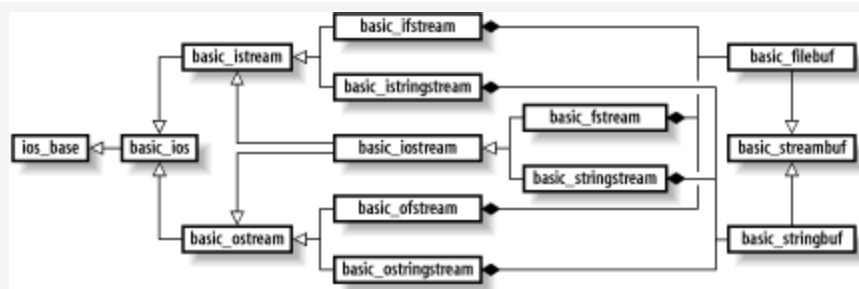


Figure 9-1. The I/O stream class hierarchy

One advantage of using inheritance in the I/O library is that the basic I/O functions are defined once in the base classes, and that interface is inherited by the derived classes and overridden when necessary. Using inheritance, you perform I/O with files as you would with strings. With only a little effort, you can derive your own I/O classes for specialized situations. (See [Example 9-6](#).) Thus, to understand I/O in C++, you must start with the base classes.

Another advantage of the I/O stream classes is that you can implement your own overloaded functions that look and behave like the standard functions. Thus, you can read and write objects of your custom classes just as easily as the fundamental types.

The `ios_base` class declares types and constants that are used throughout the I/O library. Formatting flags, I/O state flags, open modes, and seek directions are all declared in `ios_base`.

The `basic_istream` template declares input functions, and `basic_ostream` declares output functions. The `basic_iostream` template inherits input and output functions through multiple inheritance from `basic_istream` and `basic_ostream`.

The stream class templates handle high-level I/O of numbers, strings, and characters. For low-level I/O, the streams rely on *stream buffers*, which control reading and writing buffers of characters. The `basic_streambuf` template defines the stream buffer interface, and the actual behavior is implemented by derived-class templates.

I/O to and from external files is handled by `basic_fstream`, `basic_ifstream`, and `basic_ofstream`, using `basic_filebuf` as the stream buffer. These class templates are declared in `<fstream>`.



You can also treat a string as a stream using `basic_istream`, `basic_ostringstream`, and `basic_stringstream`. The stream buffer template is `basic_stringbuffer`. These class templates are declared in `<sstream>`.

The I/O library supports formatted and unformatted I/O. Unformatted I/O simply reads or writes characters or character strings without interpretation. The I/O streams have a number of functions for performing unformatted I/O.

Formatted input can skip over leading whitespace, parse text as numbers, and interpret numbers in different bases (decimal, octal, hexadecimal). Formatted output can pad fields to a desired width and write numbers as text in different bases. Formatted I/O uses a stream's *locale* to parse numeric input or format numeric output.

A locale is an object of type `locale`. It stores character attributes, formatting preferences, and related information about a particular culture or environment. This information is organized into a set of *facets*. For example, the `num_get` facet defines how numbers are read and parsed from an input stream; the `num_put` facet defines how numbers are formatted and written to an output stream; and the `numpunct` facet specifies the punctuation characters used for decimal points, minus signs, and so on. Locales are used primarily by I/O streams, but they have other uses, as well. See `<locale>` in [Chapter 13](#) for details.

To perform formatted I/O, the I/O streams overload the shift operators: left shift (`<<`) writes and right shift (`>>`) reads. Think of the shift operators as arrows pointing in the direction of data flow: output flows from an expression to the stream (`cout << expr`). Input flows from the stream to a variable (`cin >> var`). The `string`, `complex`, and other types in the standard library overload the shift operators,

so you can perform I/O with these objects just as easily as you can with the fundamental types.

## Custom I/O Operators

When you define your own classes for which I/O is meaningful, you should also override the shift operators to perform I/O in the same manner as the standard I/O operators. A common simplification is to overload the shift operators for `istream` and `ostream`, but that prevents your operators from being used with wide-character streams or streams with custom character traits. You should consider writing function templates instead, using `basic_istream` and `basic_ostream`, to take advantage of the generality that the I/O stream templates offer.

Other guidelines to follow when writing custom I/O functions are:

- Pay attention to the stream's flags, locale, and state.
- Set `failbit` for malformed input.
- Be careful with internal whitespace.
- When writing an object that has multiple parts, be careful of how you treat the stream's field width. Remember that the width is reset to 0 after each formatted output function.

[Example 9-1](#) shows an example of I/O for the `rational` class, which represents a rational number. For input, two numbers are read, separated by a slash (/). If the slash is missing, the input is malformed, so `failbit` is set. (See the `<ios>` header in [Chapter 13](#) for more information about `failbit` and the other state bits.) For output, the two parts of the rational number are formatted as a string, and the entire string is written to the output stream. Using a temporary string lets the caller set the field width and

apply the width to the entire rational number as a single entity. If the numerator and denominator were written directly to the output stream, the width would apply only to the numerator. (See [Example 9-3](#) for more information about rational.)

### *Example 9-1. Performing I/O with the rational class template*

```
// Read a rational number. The numerator and denominator must be
// written as two
// numbers (the first can be signed) separated by a slash (/).
// For example:
// "2/3", "-14/19".
template<typename T, typename charT, typename traits>
std::basic_istream<charT, traits>&
operator>>(std::basic_istream<charT, traits>& in, rational<T>& r)
{
    typename rational<T>::numerator_type n;
    typename rational<T>::denominator_type d;
    char c;

    if (! (in >> n)) return in;
    // Allow whitespace before and after the dividing '/'.
    if (! (in >> c)) return in;
    if (c != '/') {
        // Malformed input
        in.setstate(std::ios_base::failbit);
        return in;
    }
    if (! (in >> d)) return in;
    r.set(n, d);
    return in;
}

// Write a rational number as two integers separated by a slash.
// Use a string
// stream so the two numbers are written without padding, and the
// overall
// formatted string is then padded to the desired width.
template<typename T, typename charT, typename traits>
std::basic_ostream<charT, traits>&
operator<<(std::basic_ostream<charT, traits>& out, const
rational<T>& r)
{
    // Use the same flags, locale, etc. to write the
```

```

// numerator and denominator to a string stream.
std::basic_ostringstream<charT, traits> s;
s.flags(out.flags( ));
s.imbue(out.getloc( ));
s.precision(out.precision( ));
s << r.numerator( ) << '/' << r.denominator( );
// Write the string to out. The field width, padding, and
alignment are already
// set in out, so they apply to the entire rational number.
out << s.str( );
return out;
}

```

## The C Standard I/O Library

Because the C++ library includes the C library, you can use any of the C standard I/O functions, such as `fopen` and `printf`. The C standard I/O library is contained in `<cstdio>` for narrow I/O and `<cwchar>` for wide-character I/O. The C++ library inherits many attributes from the C library. For example, the C++ end-of-file marker, `char_traits<char>::eof( )`, has the same value as the C end-of-file marker, `EOF`. In most cases, however, a C++ program should use C++ I/O functions rather than C functions.

The `printf` and `scanf` functions are noteworthy for their lack of safety. Although some compilers now check the format strings and compare them with the actual arguments, many compilers do not. It is too easy to make a simple mistake. If you are lucky, your program will fail immediately. If you are unlucky, your program will appear to work on your system and fail only when it is run on your customers' systems. The following is an example of a common mistake made when using `printf` :

```

size_t s;
printf("size=%u\n", s);

```

The problem here is that the `size_t` type might be unsigned long, which means the argument `s` and the

format %u do not match. On some systems, the mismatch is harmless, but on others, wrong information will be printed, or worse.

Other functions, such as `gets` and `sprintf`, are unsafe because they write to character arrays with no way to limit the number of characters written. Without a way to prevent buffer overruns, these functions are practically useless.

Another limitation of the C I/O library is that it has little support for alternate locales. In C++, every stream can have a different locale. For example, this lets you write a program that reads a datafile in a fixed format (using the "C" locale), but prints the human-readable results in the native locale. Writing such a program using the C I/O library requires that locales be changed between each read and write call, which is much less convenient.

In spite of all these problems, some C++ programmers still use the C library. In some implementations, the C library performs better than the C++ library. Another reason is that some C++ functions, such as `printf`, have a brevity that appeals to C and C++ programmers. Compare the following examples, one using `printf` and the other using the `<<` operator, which both print the same count and mask values in the same formats:

```
unsigned long count, mask;
...
printf("count=%-9.9ld\n"
       "mask=%#-8.8lx\n", count, mask);

cout.fill('0');
cout << "count=" << right << dec << setw(9) << count <<
      "\nmask=0x" << hex << setw(8) << mask << '\n';
```

Even though the `printf` approach seems more concise and easier to understand, I recommend using the C++ I/O streams. You might think `printf` is saving you time now, but the lack of safety can present major problems in the

future. (Reread the example and imagine what would happen if count exceeded the maximum value for type long.)

Sometimes, using the C I/O library is necessary (e.g., perhaps legacy C code is being called from C++ code). To help in such situations, the standard C++ I/O objects are associated with their corresponding C FILEs. The C++ cin object is associated with the C stdin object, cout is associated with stdout, and cerr and clog are associated with stderr. You can mix C and C++ I/O functions with the standard I/O objects.

## C++ I/O Headers

This section lists the I/O-related headers in the C++ standard library, with a brief description of each. See the corresponding sections in [Chapter 13](#) for more detailed descriptions.

### <fstream>

File streams—that is, input and output using external files. Declares `basic_filebuf`, `basic_fstream`, `basic_ifstream`, `basic_ofstream`, `fstream`, `ifstream`, `ofstream`, `wfstream`, `wistream`, `wostream`, and other types.

### <iomanip>

Declares several manipulators—that is, function objects that affect an I/O stream object. A manipulator offers an alternate, often more convenient, syntax for calling member functions of an I/O stream object.

### <ios>

Base template definitions of `ios_base`, `basic_ios` and some common manipulators. All I/O streams derive from

`basic_ios`, which derives from `ios_base`.

`<iosfwd>`

Forward declarations of the standard I/O classes and templates. Judicious use of `<iosfwd>` can reduce the compile-time burden in certain situations.

`<iostream>`

Declarations of the standard I/O objects: `cin`, `cout`, etc.

`<istream>`

Declares types and functions for input-only streams (`basic_istream`, `istream`, and `wistream`), and for input and output streams (`basic_iostream`, `iostream`, and `wiostream`).

`<ostream>`

Declares types and functions for output-only streams (`basic_ostream`, `ostream`, and `wostream`).

`<sstream>`

Declares string streams (`basic_istringstream`, `basic_ostringstream`, `basic_stringbuf`, `basic_stringstream`, `istringstream`, `ostringstream`, and `stringstream`), which read from and write to strings using the stream protocol.

`<streambuf>`

Declares low-level stream buffers (`basic_streambuf`) for the I/O stream classes. Most programs do not need to deal with the stream buffers, but stick to the high-level interfaces presented by the stream classes.

`<strstream>`

Declares string streams (`istrstream`, `ostrstream`, `strstream`, `strstreambuf`), which read and write to character arrays using a stream protocol. These classes

are not template-based so they do not work for wide characters or alternative character traits. This header is deprecated.



## Text I/O

By default, an I/O stream performs text I/O, also known as formatted I/O, in which text is converted to and from numeric and other values. For output, values are converted to text, and padding, alignment, and other formatting is applied to the text. For input, formatting controls how text is converted to values and whether whitespace is skipped prior to reading an input field.



Different systems have different ways of representing the end of a line. A text I/O stream hides these details and maps all line endings in a file to the newline character ('\n'). Thus, the number of characters read from or written to a file might not match the actual file size. An implementation might require a newline at the end of the last line of a file.

To control formatting, a stream keeps track of a set of flags, a field width, and a precision. [Table 13-12](#) (in the <ios> section) lists all the formatting flags.

## Formatted Input

The formatted input functions are the overloaded `operator>>` functions. If the `skipws` flag is set (which is the default), whitespace characters (according to the locale imbued in the stream) are skipped, and input begins with the first non-whitespace character.

If reading into a string or character array, all non-whitespace characters are read into the string, ending with the first whitespace character or when `width` characters have been read (if `width > 0`), whichever comes first. The `width` is then reset to 0.

For all other types, the width is not used. To read a number from a fixed-width field, read the field into a string, then use a string stream to read the number, as shown in [Example 9-2](#).

*Example 9-2. Reading a number from a fixed-width field*

```
template<typename T, typename charT, typename traits>
std::basic_istream<charT, traits>&
    fixedread(std::basic_istream<charT, traits>& in, T& x)
{
    if (in.width( ) == 0)
        // Not fixed size, so read normally.
        in >> x;
    else {
        std::string field;
        in >> field;
        std::basic_istringstream<charT, traits> stream(field);
        if (! (stream >> x))
            in.setstate(std::ios_base::failbit);
    }
    return in;
}
```

The only other flags that affect input are `basefield` and `boolalpha`:

- The `basefield` flag determines how integers are interpreted. If `basefield` is nonzero, it specifies a fixed radix (oct, hex, or dec), or if `basefield` is 0 (the default), the input determines the radix: leading 0x or 0X for hexadecimal, leading 0 for octal, decimal otherwise.
- If the `boolalpha` flag is set, a formatted read of `bool` reads a string, which must match the names `true` or `false` (in the stream's locale, according to the `num_punct` facet). If the `boolalpha` flag is not set, a `bool` is read as a long integer, and the number is converted to `bool` using the standard rules: nonzero is `true` and 0 is `false`. By default, the flag is clear (`false`).

Floating-point numbers are accepted in fixed or scientific format.

The decimal point character is determined by the stream's locale, as is the thousands separator. Thousands separators are optional in the input stream, but if present, they must match the locale's thousands separator character and grouping rules. For example, assuming that the thousands separator is , and the grouping is for every 3 characters (grouping( ) returns "\3"), the following are three valid input examples and one invalid example:

```
1,234,567          // Valid
1,234,56           // Invalid
1234567           // Valid
1234,567          // Valid
```

When reading data that the user types, you should imbue the input stream with the user's native locale (e.g., `cin.imbue(locale(" "))`). For input that is being read from files or other sources that require portable data formats, be sure to use the "C", or classic, locale (e.g., `cin.imbue(locale::classic( ))`).

See the `num_get` facet in the <locale> section of [Chapter 13](#) for details on how numeric input is parsed and interpreted.

## Formatted Output

The formatted output functions are the overloaded `operator<<` functions. They all work similarly, using the flags and other information to format a value as a string. The string is then padded with zero or more copies of a `fill` character to achieve the desired width. The `adjustfield` flags are used to determine where the `fill` characters are added (to the left, right, or internal, with internal meaning after a sign or a leading `0x` or `0X`).

The padded string is written to the output stream, and the stream's width is reset to 0. The width is the only formatting parameter that is reset. The flags, precision, and the fill character are "sticky" and persist until they are changed explicitly.

Formatting an integer depends on the basefield (hex, oct, or dec), uppercase (0X for hexadecimal), showpos (insert a + for positive numbers), and showbase (to insert a prefix of 0x or 0X for hexadecimal or 0 for octal) flags. The defaults are decimal, lowercase, no positive sign, and no base. If the locale's numpunct facet specifies thousands grouping, thousands separators are inserted at the specified positions.

Formatting a floating-point number depends on the floatfield (fixed, scientific, or 0 for general), uppercase (E for exponent), showpoint (insert decimal point even if not needed), and showpos (insert a + for positive numbers) flags. The defaults are general, lowercase, no point unless needed, and no positive sign.

If the boolalpha flag is set, bool values are written as names (e.g., true or false, or other strings, depending on the locale). If boolalpha is not set, bool values are written as integers (described earlier). The default flag is clear (false).

When writing output for a user's immediate consumption, you should imbue the output stream with the user's native locale (e.g., `cout.imbue(locale(""))`). For output that is being written to files or other sources that require portable data formats, be sure to use the "C", or classic, locale (e.g., `cout.imbue(locale::classic( ))`).

See the `num_put` facet in the `<locale>` section of [Chapter 13](#) for details on how numeric output is formatted.

## Binary I/O

Binary, or unformatted, I/O involves characters, character arrays, or strings, which are read or written without interpretation, padding, or other adjustments. End-of-lines are not treated specially.

The unformatted read functions (such as `get` and `read` ) can read into a string or character array. The `gcount` function returns the number of characters read.

The unformatted output functions (such as `put` and `write` ) can write a string or character array. You can specify the exact number of characters to write from a character array or write all characters up to a null character (to write a C-style null-terminated string).

You can also dip down the stream buffer level to perform unformatted I/O, although this is seldom done except for bulk I/O of entire streams. See the next section for more information.

## Stream Buffers

The I/O stream classes rely on stream buffers for the low-level input and output. Most programmers ignore the stream buffer and deal only with high-level streams. You might find yourself dealing with stream buffers from the client side—for example, using stream buffers to perform low-level I/O on entire streams at once—or you might find yourself on the other side, implementing a custom stream buffer. This section discusses both of these aspects.

You can copy a file in several ways. [Example 9-3](#) shows how a C programmer might copy a stream once he has learned about templates.

### *Example 9-3. Copying streams one character at a time*

```
template<typename charT, typename traits>
void copy(std::basic_ostream<charT, traits>& out,
          std::basic_istream<charT, traits>& in)
{
    charT c;
    while (in.get(c))
        out.put(c);
}
```

After measuring the performance of this solution, the intrepid programmer might decide that copying larger buffers is the right way to go. [Example 9-4](#) shows the new approach. On my system, the new version runs roughly twice as fast as the original version. (Of course, performance measures depend highly on compiler, library, environment, and so on.)

### *Example 9-4. Copying streams with explicit buffers*

```
template<typename charT, typename traits>
void copy(std::basic_ostream<charT, traits>& out,
          std::basic_istream<charT, traits>& in)
{
    const unsigned BUFFER_SIZE = 8192;
    std::auto_ptr<charT> buffer(new charT[BUFFER_SIZE]);
```

```

while (in) {
    in.read(buffer.get( ), BUFFER_SIZE);
    out.write(buffer.get( ), in.gcount( ));
}
}

```

After reading more about the C++ standard library, the programmer might try to improve performance by delegating all the work to the stream buffer, as shown in [Example 9-5](#).

### *Example 9-5. Copying streams via stream buffers*

```

template<typename charT, typename traits>
void copy(std::basic_ostream<charT, traits>& out,
          std::basic_istream<charT, traits>& in)
{
    out << in.rdbuf( );
}

```

The version in [Example 9-5](#) runs about as fast as the version in [Example 9-4](#) but is much simpler to read and write.

Another reason to mess around with stream buffers is that you might need to write your own. Perhaps you are implementing a network I/O package. The user opens a network stream that connects to a particular port on a particular host and then performs I/O using the normal I/O streams. To implement your package, you must derive your own stream buffer class template (`basic_networkbuf`) from the `basic_streambuf` class template in `<streambuf>`.

A stream buffer is characterized by three pointers that point to the actual buffer, which is a character array. The pointers point to the beginning of the buffer, the current I/O position (that is, the next character to read or the next position for writing), and the end of the buffer. The stream buffer class manages the array and the pointers. When the array empties upon input, the stream buffer must obtain additional input from the network (the `underflow` function). When the array fills upon output, the stream buffer must

write the data to the network (the overflow function). Other functions include putting back a character after reading it, seeking, and so on. (See `<streambuf>` in [Chapter 13](#) for details about each member function.) [Example 9-6](#) shows an extremely oversimplified sketch of how the `basic_networkbuf` class template might work.

*Example 9-6. The `basic_networkbuf` class template*

```
template<typename charT,
        typename traits = std::char_traits<char> >
class basic_networkbuf :
public std::basic_streambuf<charT, traits> {
public:
    typedef charT          char_type;
    typedef traits         traits_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;

    basic_networkbuf( );
    virtual ~basic_networkbuf( );

    bool is_connected( );
    basic_networkbuf* connect(const char* hostname, int port,
        std::ios_base::openmode mode);
    basic_networkbuf* disconnect( );

protected:
    virtual std::streamsize showmanyc( );
    virtual int_type underflow( );
    virtual int_type overflow(int_type c = traits::eof( ));
    virtual pos_type seekoff(off_type offset,
        std::ios_base::seekdir dir,
        std::ios_base::openmode);
    virtual pos_type seekpos(pos_type sp,
        std::ios_base::openmode);
    virtual basic_networkbuf* setbuf(char_type* buf,
        std::streamsize size);
    virtual int sync( );

private:
    char_type* buffer_;
    std::streamsize size_;
    bool ownbuf_; // true means destructor must delete buffer_
    // network connectivity stuff...
```



```

};

// Construct initializes the buffer pointers.
template<typename charT, typename traits>
basic_networkbuf<charT,traits>::basic_networkbuf( )
: buffer_(new char_type[DEFAULT_BUFSIZ]),
  size_(DEFAULT_BUFSIZ), ownbuf_(true)
{
    this->setg(buffer_, buffer_ + size_, buffer_ + size_);
    // Leave room in the output buffer for one last character.
    this->setp(buffer_, buffer_ + size_ - 1);
}

// Return the number of characters available in the input buffer.
template<typename charT, typename traits>
std::streamsize basic_networkbuf<charT,traits>::showmanyc( )
{
    return this->egptr( ) - this->gptr( );
}

// Fill the input buffer and set up the pointers.
template<typename charT, typename traits>
typename basic_networkbuf<charT,traits>::int_type
basic_networkbuf<charT,traits>::underflow( )
{
    // Get up to size_ characters from the network, storing them in
    buffer_. Store
    // the actual number of characters read in the local variable,
    count.

    std::streamsize count;
    count = netread(buffer_, size_);
    this->setg(buffer_, buffer_, buffer_ + count);
    if (this->egptr( ) == this->gptr( ))
        return traits::eof( );
    else
        return traits::to_int_type(*this->gptr( ));
}

// The output buffer always has room for one more character, so
if c is not
// eof( ), add it to the output buffer. Then write the buffer to
the network
// connection.
template<typename charT, typename traits>
typename basic_networkbuf<charT,traits>::int_type
{
    if (c != traits::eof( )) {

```

```
        *(this->pptr( )) = c;
        this->pbump(1);
    }
    netwrite(this->pbase(), this->pptr( ) - this->pbase( ));
    // The output buffer is now empty. Make sure it has room for
one last
    // character.
    this->setp(buffer_, buffer_ + size_ - 1);
    return traits::not_eof(c);
}

// Force a buffer write.
template<typename charT, typename traits>
int basic_networkbuf<charT,traits>::sync( )
{
    overflow(traits::eof( ));
    return 0;
}
```

# Manipulators

A *manipulator* is a function object that can be used as an operand to an input or output operator to manipulate the stream. Manipulators can send additional output to a stream, read input from a stream, set flags, and more. For example, to output a zero-padded, hexadecimal integer, you can use an ostream's member functions or manipulators, whichever you prefer. [Example 9-7](#) shows both ways.

*Example 9-7. Manipulating an output stream to format a number*

```
using namespace std;

// Output a value using ostream's member functions.
cout.fill('0');
cout.width(8);
cout.setf(ios_base::internal, ios_base::adjustfield);
cout.setf(ios_base::hex, ios_base::basefield);
cout << value;

// Output the same value using manipulators.
cout << setfill('0') << setw(8) << hex << internal << value;
```

## Standard Manipulators

The standard library defines several manipulators for setting formatting flags, setting other formatting parameters, skipping whitespace, flushing output, and more. The following is a list of all the standard manipulators, grouped by header:

<ios>

Declares the manipulators that set the formatting flags:  
boolalpha , dec , fixed , hex , internal , left ,  
noboolalpha , noshowbase , noshowpoint , noshowpos ,  
noskipws , nouppercase , nounitbuf , oct , right ,

scientific , showbase , showpoint , showpos , skipws ,  
uppercase , and unitbuf

<iostream>

Declares the input manipulator: ws

<ostream>

Declares the output manipulators: endl , ends , and  
flush

<iomanip>

Declares several additional manipulators: resetioflags  
, setioflags , setbase , setfill , setprecision , and  
setw

Most manipulators are declared in the same header as the stream type they manipulate. The only time you need to `#include` an additional header is when you use a manipulator that takes an argument. These manipulators are in the `<iomanip>` header.

## Custom Manipulators

To write your own manipulator, use the standard manipulators as patterns. The easiest to use are manipulators that take no arguments. A manipulator is simply a function that takes a stream as an argument and returns the same stream. The standard streams overload `operator<<` and `operator>>` to take a pointer to such a function as an operand.

Suppose you want to write an input manipulator that skips all characters up to and including a newline. (Perhaps this manipulator is used by a command processor after reading a `//` comment sequence.) [Example 9-8](#) shows one way to write the `skipline` manipulator.

*Example 9-8. Skipping a line in an input stream*

```

template<typename charT, typename traits>
std::basic_istream<charT,traits>&
skipline(std::basic_istream<charT,traits>& in)
{
    charT c;
    while (in.get(c) && c != '\n')
        ;
    return in;
}
...
int x;
std::string next;
std::cin >> x >> skipline >> next;

```

Manipulators that take arguments are harder to write, but only slightly. You need to write some supporting infrastructure, such as additional overloaded `operator>>` or `operator<<` functions.

For example, suppose you want to parameterize your input `skipline` manipulator so it skips everything up to a caller-supplied character. This manipulator is defined as a class template, in which the constructor takes the manipulator's argument, that is, the delimiter character. You must overload `operator>>` so it recognizes your manipulator as an operand and invokes the manipulator's `operator( )`. You don't need to use `operator( )`, but this is a good choice when building a reusable infrastructure for manipulators. [Example 9-9](#) shows the new `skip` manipulator.

*Example 9-9. Writing a manipulator that takes an argument*

```

template<typename charT>
class skipper
{
public:
    typedef charT char_type;
    skipper(char_type delim) : delim_(delim) {}
    template<typename traits>
    void operator( )(std::basic_istream<charT,traits>&) const;
private:
    char_type delim_;
};

```

```

// Skip the rest of the line. The compiler deduces the traits
type from the
// stream argument.
template<typename charT>
template<typename traits>
void skipper<charT>::operator( )
(std::basic_istream<charT,traits>& stream)
const
{
    char_type c;
    while (stream.get(c) && c != delim_)
        ;
}

// Invoke the skipper manipulator.
template<typename charT, typename traits>
std::basic_istream<charT,traits>&
    const skipper<charT>& f)
{
    f(stream);
    return stream;
}

// Let the compiler deduce the character type.
template<typename charT>
skipper<charT> skip(charT c)
{
    return skipper<charT>(c);
}
...
int x;
std::string next;
std::cin >> x >> skip('\n') >> next;

```

## Errors and Exceptions

By default, I/O streams do not raise exceptions for errors. Instead, each stream keeps a mask of error bits called the *I/O state*. The state mask keeps track of formatting failures, end-of-file conditions, and miscellaneous error conditions. The `ios_base` class template defines several member functions for testing and modifying the state flags (`rdstate`, `setstate`, `fail`, etc.).

A common idiom is to read from an input stream until an input operation fails. Because this idiom is so common, the standard library makes it easy. Instead of calling `rdstate` and testing the state explicitly, you can simply treat the stream object as a Boolean value: `true` means the state is good, and `false` means the state has an error condition. Most I/O functions return the stream object, which makes the test even easier:

```
while (cin.get(c))
    cout.put(c);
```

The `basic_ios` class overloads operator `void*` to return a non-null pointer if the state is good or a null pointer for any error condition. Similarly, it overloads operator `!` to return `true` for any error condition. (As explained later in this section, an end-of-file is not an error condition.) This latter test is often used in conditional statements:

```
if (! cout)
    throw("write error");
```

The state mask has three different error bits:

### `badbit`

An unrecoverable error occurred. For example, an exception was thrown from a formatting facet, an I/O system call failed unexpectedly, and so on.

`eofbit`

An end-of-file upon input.

`failbit`

An I/O operation failed to produce any input or output. For example, when reading an integer, if the next input character is a letter, no characters can be read from the stream, which results in an input failure.

The `basic_ios` conditional operators define "failure" as when `badbit` or `failbit` is set, but not when `eofbit` is set. To understand why, consider the following canonical input pattern. During a normal program run, the input stream's state is initially zero. After reading the last item from the input stream, `eofbit` is set in the state. At this time, the state does not indicate "failure," so the program continues by processing the last input item. The next time it tries to read from the input stream, no characters are read (because `eofbit` is set), which causes the input to fail, so the stream sets `failbit`. Now a test of the input stream returns `false`, indicating failure, which exits the input loop.

Sometimes, instead of testing for failure after each I/O operation, you may want to simplify your code. You can assume that every operation succeeds and arrange for the stream to throw an exception for any failure. In addition to the state mask, every stream has an exception mask, in which the bits in the exception mask correspond to the bits in the state mask. When the state mask changes, if any bit is set in both masks, the stream throws an `ios_base::failure` exception.

For example, suppose you set the exception mask to `failbit | badbit`. Using the canonical input pattern, after reading the last item from the input stream, `eofbit` is set in the state. At this time, `rdstate( ) & exceptions( )` is still `0`, so the program continues. The next time the



program tries to read from the input stream, no characters are read, which causes the input to fail, and the stream sets `failbit`. Now `rdstate( ) & exceptions( )` returns a nonzero value, so the stream throws `ios_base::failure`.

A stream often relies on other objects (especially locale facets) to parse input or format output. If one of these other objects throws an exception, the stream catches the exception and sets `badbit`. If `badbit` is set in the `exceptions( )` mask, the original exception is rethrown.

When testing for I/O success, be sure to test for `badbit` as a special indicator of a serious failure. A simple test for `! cin` does not distinguish between different reasons for failure: `eofbit | failbit` might signal a normal end-of-file, but `failbit | badbit` might tell you that there is something seriously wrong with the input stream (e.g., a disk error). One possibility, therefore, is to set `badbit` in the `exceptions( )` mask so normal control flow deals with the normal situation of reading an end-of-file. However, more serious errors result in exceptions, as shown in [Example 9-10](#).

### *Example 9-10. Handling serious I/O errors*

```
#include <algorithm>
#include <cstdint>
#include <exception>
#include <iostream>
#include <map>
#include <string>

void print(const std::pair<std::string, std::size_t>& count)
{
    std::cout << count.first << '\t' << count.second << '\n';
}

int main( )
{
    using namespace std;

    try {
```

```
string word;
map<string, size_t> counts;
cin.exceptions(ios_base::badbit);
cout.exceptions(ios_base::badbit);
while (cin >> word)
    ++counts[word];
for_each(counts.begin( ), counts.end( ), print);
} catch(ios_base::failure& ex) {
    std::cerr << "I/O error: " << ex.what( ) << '\n';
    return 1;
} catch(exception& ex) {
    std::cerr << "Fatal error: " << ex.what( ) << '\n';
    return 2;
} catch(...) {
    std::cerr << "Total disaster.\n";
    return 3;
}
}
```

# Chapter 10. Containers, Iterators, and Algorithms

Containers (sometimes called collections) are a staple of computer programming. Every major programming language has fundamental containers, such as arrays or lists. Modern programming languages usually have an assortment of more powerful containers, such as trees, for more specialized needs.

The C++ library has a basic suite of containers (deques, lists, maps, sets, and vectors), but more important, it has *generic algorithms*, which are function templates that implement common algorithms, such as searching and sorting. The algorithms operate on *iterators*, which are an abstraction of pointers that apply to any container or other sequence.

This chapter presents an overview of the standard C++ containers, the iterators used to examine the containers, and the algorithms that can be used with the iterators. It covers how to use the standard containers, iterators, and algorithms, as well as how to write your own.

In this chapter, the mathematical notation of  $[first, last)$  is often used to denote a range. The square bracket marks an inclusive endpoint of a range, and the parenthesis marks an exclusive endpoint of a range. Thus,  $[first, last)$  means a range that extends from `first` to `last`, including `first`, but excluding `last`. Read more about ranges in "Iterators" later in this chapter.

For complete details about the various containers, iterators, and algorithms, see [Chapter 13](#).

## Containers

The fundamental purpose of a container is to store multiple objects in a single container object. Different kinds of containers have different characteristics, such as speed, size, and ease of use. The choice of container depends on the characteristics and behavior you require.

In C++, the containers are implemented as class templates, so you can store anything in a container. (Well, almost anything. The type must have *value semantics*, which means it must behave as an ordinary value, such as an `int`. Values can be copied and assigned freely. An original and its copy must compare as equal. Some containers impose additional restrictions.)

## Standard Containers

The standard containers fall into two categories: sequence and associative containers. A sequence container preserves the original order in which items were added to the container. An associative container keeps items in ascending order (you can define the order relation) to speed up searching. The standard containers are:

### `deque`

A deque (double-ended queue) is a sequence container that supports fast insertions and deletions at the beginning and end of the container. Inserting or deleting at any other position is slow, but indexing to any item is fast. Items are not stored contiguously. The header is `<deque>`.

### `list`

A list is a sequence container that supports rapid insertion or deletion at any position but does not support random access. Items are not stored contiguously. The header is `<list>`.

## map , multimap

A map (or dictionary) is an associative container that stores pairs of keys and associated values. The keys determine the order of items in the container. map requires unique keys. multimap permits duplicate keys. The header for map and multimap is <map>.

## set , multiset

A set is an associative container that stores keys in ascending order. set requires unique keys. multiset permits duplicate keys. The header for set and multiset is <set>.

## vector

A vector is a sequence container that is like an array, except that it can grow as needed. Items can be rapidly added or removed only at the end. At other positions, inserting and deleting items is slower. Items are stored contiguously. The header is <vector>.

The set and map containers perform insertions, deletions, and searches in logarithmic time, which implies a tree or tree-like implementation. Items must be kept in sorted order, so a hash table implementation is not allowed. Many programmers consider the lack of a standard hash table to be a serious omission. When the C++ standard is revised, a hash table is likely to be added. The STLport project includes hash table-based containers. See [Appendix B](#) for information about STLport.

## Container Adapters

In addition to the standard containers, the standard library has several container adapters. An adapter is a class template that uses a container for storage and provides a

restricted interface when compared with the standard containers. The standard adapters are:

### priority\_queue

A priority queue is organized so that the largest element is always the first. You can push an item onto the queue, examine the first element, or remove the first element. The header is `<queue>`.

### queue

A queue is a sequence of elements that lets you add elements at one end and remove them at the other end. This organization is commonly known as FIFO (first-in, first-out). The header is `<queue>`.

### stack

A stack is a sequence that lets you add and remove elements only at one end. This organization is commonly known as LIFO (last-in, first-out). The header is `<stack>`.

## Pseudo-Containers

The standard library has a few class templates that are similar to the standard containers but fail one or more of the requirements for a standard container:

### bitset

Represents a bitmask of arbitrary size. The size is fixed when the `bitset` is declared. There are no `bitset` iterators, so you cannot use a `bitset` with the standard algorithms. See `<bitset>` in [Chapter 13](#) for details.

### basic\_string , string , wstring

Represent character strings. The string class templates meet almost all of the requirements of a sequence container, and you can use their iterators with the

standard algorithms. Nonetheless, they fall short of meeting all the requirements of a container, such as lacking front and back member functions. The header is `<string>`.

### `valarray`

Represents an array of numeric values optimized for computational efficiency. A `valarray` lacks iterators, and as part of the optimization, the compiler is free to make assumptions that prevent `valarray` from being used with the standard algorithms. See `<valarray>` in [Chapter 13](#) for details.

### `vector<bool>`

A specialization of the `vector` template. Although `vector<>` usually meets the requirements of a standard container, the `vector<bool>` specialization does not because you cannot obtain a pointer to an element of a `vector<bool>` object. See `<vector>` in [Chapter 13](#) for details.

## Container Requirements

This section presents the rules that govern containers. Some rules apply to all the standard containers, and you can rely on the standard behavior for all C++ implementations. Other rules apply only to sequence containers, and some apply only to associative containers. If you write your own container class template, be sure to follow the same conventions and rules that apply to the standard containers.

### Member types

#### `const_iterator`

The iterator type for const values.

`const_reference`

A const lvalue type for the items stored in the container. This is typically the same as the allocator's `const_reference` type.

`difference_type`

A signed integral type denoting the difference between two iterators.

`iterator`

The iterator type.

`reference`

An lvalue type for the items stored in the container. This is typically the same as the allocator's `reference` type.

`size_type`

An unsigned integral type that can hold any nonnegative `difference_type` value.

`value_type`

The type of item stored in the container. This is typically the same as the first template parameter.

A container that supports bidirectional iterators should also define the `reverse_iterator` and `const_reverse_iterator` types.

An associative container should define `key_type` as the key type, `compare_type` as the key compare function or functor, and `value_compare` as a function or functor that compares two `value_type` objects.

Optionally, a container can declare `pointer` and `const_pointer` as synonyms of the allocator's types of the same name, and `allocator_type` for the allocator, which is typically the last template parameter.



## Member functions

Most of the standard member functions have a complexity that is constant or linear in the number of elements in the container. Some of the member functions for associative members are logarithmic in the number of elements in the container. Each of the descriptions in this section notes the complexity of the function.

A container template should have the following constructors. You do not need to write separate constructors in all cases; sometimes you can use default arguments instead of overloading. If an allocator object is supplied, it is copied to the container; otherwise, a default allocator is constructed. In each of the following descriptions, *container* is the name of the container class template.

***container*** ( ), ***container*** (allocator\_type)

Initializes the container to be empty. Complexity is constant.

***container*** (const container& that)

Initializes the container with a copy of all the items and the allocator from that. Complexity is linear.

A sequence container should have the following additional constructors:

***container*** (size\_type n, const value\_type& x),  
***container*** (size\_type n, const value\_type& x,  
allocator\_type)

Initializes the container with n copies of x. Complexity is linear with respect to n.

template<InIter> ***container*** (InIter first, InIter last), template<InIter>, ***container*** (InIter first,

InIter last, allocator\_type)

Initializes the container with copies of the items in the range [first, last). Complexity is linear with respect to last - first.

If InIter is an integral type, the container is initialized with first copies of last (converted to value\_type). Complexity is linear with respect to first.

An associative container should have the following additional constructors:

**container** (key\_compare compare), **container** (key\_compare compare, allocator\_type)

Initializes an empty container that uses compare to compare keys. Complexity is constant.

template<InIter> **container** (InIter first, InIter last, key\_compare compare), template<InIter> **container** (InIter first, InIter last, key\_compare compare, allocator\_type)

Initializes the container with copies of the items in the range [first, last), comparing keys with compare. Complexity is linear.

All containers must have a destructor:

~ **container** ( )

Calls the destructor for every object in the container, perhaps by calling clear. Complexity is linear.

Many member functions are the same for all the container templates, and when you write your own container template, you should implement the same member functions with the same behaviors. Some are specific to

sequence containers, and some to associative containers. Each container template can define additional members.

iterator **begin** ( ), const\_iterator **begin** ( ) const

Returns an iterator that points to the first item of the container. Complexity is constant.

void **clear** ( )

Erases all the items in the container. Complexity is linear.

bool **empty** ( ) const

Returns true if the container is empty (`size( ) == 0`). Complexity is constant.

iterator **end** ( ), const\_iterator **end** ( ) const

Returns an iterator that points to one past the last item of the container. Complexity is constant. (See [Section 10.2](#) later in this chapter for a discussion of what "one past the last item" means.)

**erase** (iterator p)

Erases the item that p points to. For a sequence container, erase returns an iterator that points to the item that comes immediately after the deleted item or `end( )`. Complexity depends on the container.

For an associative container, erase does not return a value. Complexity is constant (amortized over many calls).

**erase** (iterator first, last)

Erases all the items in the range `[first, last)`. For a sequence container, erase returns an iterator that points to the item that comes immediately after the last deleted item or `end( )`. Complexity depends on the container.

For an associative container, `erase` does not return a value. Complexity is logarithmic, plus last - first.

`size_type` **max\_size** ( ) const

Returns the largest number of items the container can possibly hold. Although many containers are not constrained, except by available memory and the limits of `size_type`, other container types, such as an array type, might have a fixed maximum size. Complexity is usually constant.

`container&` **operator=** (const `container&` that)

Erases all items in this container and copies all the items from that. Complexity is linear in `size( ) + that.size( )`.

`size_type` **size** ( ) const

Returns the number of items in the container. Complexity is usually constant.

`void` **swap** (const `container&` that)

Swaps the elements of this container with that. An associative container also swaps the comparison function or functions. Complexity is usually constant.

Each container should have all of its equality and relational operators defined, either as member functions or, preferably, as functions at the namespace level. Namespace-level functions offer more flexibility than member functions. For example, the compiler can use implicit type conversions on the lefthand operand but only if the function is not a member function.

A container that supports bidirectional iterators should define `rbegin` and `rend` member functions to return reverse iterators.

The following functions are optional. The standard containers provide only those functions that have constant complexity.

reference **at** (size\_type n), const\_reference **at**  
(size\_type n) const

Returns the item at index n, or throws `out_of_range` if `n >= size()`.

reference **back** ( ), const\_reference **back** ( ) const

Returns the last item in the container. Behavior is undefined if the container is empty.

reference **front** ( ), const\_reference **front** ( )  
const

Returns the first item in the container. Behavior is undefined if the container is empty.

reference **operator[]** (size\_type n), const\_reference  
**operator[]** (size\_type n)

Returns the item at index n. Behavior is undefined if `n >= size()`.

void **pop\_back** ( )

Erases the last item in the container. Behavior is undefined if the container is empty.

void **pop\_front** ( )

Erases the first item in the container. Behavior is undefined if the container is empty.

void **push\_back** (const value\_type& x)

Inserts x as the new last item in the container.

void **push\_front** (const value\_type& x)

Inserts x as the new first item in the container.

A sequence container should define the following member functions. The complexity of each depends on the container type.

iterator **insert** (iterator p, const value\_type& x)

Inserts x immediately before p and returns an iterator that points to x.

void **insert** (iterator p, size\_type n, const value\_type& x)

Inserts n copies of x before p.

template<InIter>, void **insert** (iterator p, InIter first, InIter last)

Copies the values from [first, last) and inserts them before p.

An associative container should define the following member functions. In the descriptions of the complexity, N refers to the number of elements in the container, M refers to the number of elements in the argument range (e.g., last - first), and count is the value that the function returns. Some of these member functions seem to duplicate standard algorithms (as discussed in [Section 10.3](#) later in this chapter), but the associative containers can implement them with better performance than the generic algorithms.

size\_type **count** (const key\_type& k) const

Returns the number of items equivalent to k. Complexity is  $\log N + \text{count}$ .

pair<const\_iterator, const\_iterator>, **equal\_range** (const key\_type& k) const, pair<iterator, iterator> **equal\_range** (const key\_type& k)

Returns the equivalent of `make_pair(lower_bound(k), upper_bound(k))`. Complexity is  $\log N$ .

size\_type **erase** (const key\_type& k)

Erases all the items equivalent to k. Returns the number of items erased. Complexity is  $\log N + \text{count}$ .

const\_iterator **find** (const key\_type& k) const,  
iterator **find** (const key\_type& k)

Finds an item equivalent to k and returns an iterator that points to one such item, or `end( )` if not found.

Complexity is  $\log N$ .

**insert** (const value\_type& x)

Inserts x. If the container permits duplicate keys, `insert` returns an iterator that points to the newly inserted item.

If the container requires unique keys, `insert` returns `pair<iterator, bool>`, in which the first element of the pair is an iterator that points to an item equivalent to x, and the second element is `true` if x was inserted or `false` if x was already present in the container.

Complexity is  $\log N$ .

iterator **insert** (iterator p, const value\_type& x)

Inserts x and returns an iterator that points to x. The iterator p hints to where x might belong. Complexity is  $\log N$  in general, but is constant (amortized over many calls) if the hint is correct, that is, if x is inserted immediately after p.

template<InIter>, void **insert** (InIter first, InIter last)

Copies the items from `[first, last)` and inserts each item in the container. Complexity is  $M \log(N + M)$ , but is linear if the range is already sorted.

key\_compare **key\_comp** ( ) const

Returns the key compare function or functor. Complexity is constant.

```
const_iterator lower_bound (const key_type& k)  
const_iterator lower_bound (const key_type& k)
```

Returns an iterator that points to the first item in the container that does not come before k. That is, if k is in the container, the iterator points to the position of its first occurrence; otherwise, the iterator points to the first position where k should be inserted. Complexity is log N.

```
value_compare value_comp ( ) const
```

Returns the value compare function or functor. Complexity is constant.

```
const_iterator upper_bound (const key_type& k)  
const_iterator upper_bound (const key_type& k)
```

Returns an iterator that points to the first item in the container that comes after all occurrences of k. Complexity is log N.

## **Exceptions**

The standard containers are designed to be robust in the face of exceptions. The exceptions that the containers themselves can throw are well-defined (for example, at might throw `out_of_range`), and most member functions do not throw any exceptions of their own.

If a single value is being added to a container (by calling `insert`, `push_front`, or `push_back`), and an exception is thrown, the container remains in a valid state without adding the value to the container.

When inserting more than one value, different containers have different behaviors. A `list`, for example, ensures that all items are inserted or none, that is, if an exception is thrown, the `list` is unchanged. A `map` or `set`, however, ensures only that each individual item is inserted



successfully. If an exception is thrown after inserting some of the items from a range, the destination container retains the elements that had been inserted successfully.

The `erase`, `pop_back`, and `pop_front` functions never throw exceptions.

The `swap` function throws an exception only if an associative container's `Compare` object's copy constructor or assignment operator throws an exception.

[Example 10-1](#) shows an `slist` container, which implements a singly-linked list. A singly-linked list requires slightly less memory than a doubly-linked list but offers, at best, a forward iterator, not a bidirectional iterator.

*Example 10-1. Implementing a custom container: a singly-linked list*

```
// Simple container for singly-linked lists.
template<typename T, typename Alloc = std::allocator<T> >
class slist {
    // Private type for a link (node) in the list.
    template<typename U>
    struct link {
        link* next;
        U value;
    };
    typedef link<T> link_type;

public:
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef Alloc allocator_type;
    typedef T value_type;
    typedef size_t size_type;
    typedef ptrdiff_t difference_type;

    class iterator;          // See Section 10.2 later in this chapter
    for
    class const_iterator; // the iterators.

    slist(const slist& that);
```

```

slist(const Alloc& alloc = Alloc( ));
slist(size_type n, const T& x, const Alloc& alloc=Alloc( ));
template<typename InputIter>
slist(InputIter first, InputIter last,
      const Alloc& alloc = Alloc( ));
~slist( ) { clear( ); }

slist& operator=(const slist& that);
allocator_type get_allocator( ) const { return alloc_; }

iterator begin( ) { return iterator(0, head_); }
const_iterator begin( ) const
    { return const_iterator(0, head_); }
iterator end( ) { return iterator(0, 0); }
const_iterator end( ) const { return const_iterator(0, 0); }

void pop_front( ) { erase(begin( )); }
void push_front(const T& x) { insert(begin( ), x); }
T front( ) const { return head_->value; }
T& front( ) { return head_->value; }

iterator insert(iterator p, const T& x);
void insert(iterator p, size_type n, const T& x);
template<typename InputIter>
void insert(iterator p, InputIter first, InputIter last);

iterator erase(iterator p);
iterator erase(iterator first, iterator last);

void clear( ) { erase(begin( ), end( )); }
bool empty( ) const { return size( ) == 0; }
size_type max_size( ) const
    { return std::numeric_limits<difference_type>::max( ); }
void resize(size_type sz, const T& x = T( ));
size_type size( ) const { return count_; }
void swap(slist& that);

private:
typedef typename
    allocator_type::template rebind<link_type>::other
    link_allocator_type;

link_type* newitem(const T& x, link_type* next = 0);
void delitem(link_type* item);

template<typename InputIter>
void construct(InputIter first, InputIter last,
              is_integer_tag);

```

```

template<typename InputIter>
void construct(InputIter first, InputIter last,
               is_not_integer_tag);

link_type* head_;
link_type* tail_;
size_t count_;
allocator_type alloc_;
link_allocator_type linkalloc_;
};

// Constructor. If InputIter is an integral type, the standard
// requires the
// constructor to interpret first and last as a count and value,
// and perform the
// slist(size_type, T) constructor. Use the is_integer trait to
// dispatch to the
// appropriate construct function, which does the real work.
template<typename T, typename A>
template<typename InputIter>
slist<T,A>::slist(InputIter first, InputIter last,
                 const A& alloc)
: alloc_(alloc), linkalloc_(link_allocator_type( )),
  head_(0), tail_(0), count_(0)
{
  construct(first, last, is_integer<InputIter>::tag(  ));
}

template<typename T, typename A>
template<typename InputIter>
void slist<T,A>::construct(InputIter first, InputIter last,
                          is_integer_tag)
{
  insert(begin( ), static_cast<size_type>(first),
        static_cast<T>(last));
}

template<typename T, typename A>
template<typename InputIter>
void slist<T,A>::construct(InputIter first, InputIter last,
                          is_not_integer_tag)
{
  insert(begin( ), first, last);
}

// Private function to allocate a new link node.
template<typename T, typename A>

```

```

typename slist<T,A>::link_type*
    slist<T,A>::newitem(const T& x, link_type* next)
{
    link_type* item = linkalloc_.allocate(1);
    item->next = next;
    alloc_.construct(&item->value, x);
    return item;
}

// Private function to release a link node.
template<typename T, typename A>
void slist<T,A>::delitem(link_type* item)
{
    alloc_.destroy(&item->value);
    linkalloc_.deallocate(item, 1);
}

// Basic insertion function. All insertions eventually find their
// way here.
// Inserting at the head of the list (p == begin( )) must set
// the head_member.
// Inserting at the end of the list (p == end( )) means
// appending to the list,
// which updates the tail_'s next member, and then sets tail_.
// Anywhere else in
// the list requires updating p.prev_->next. Note that inserting
// into an empty
// list looks like inserting at end( ). Return an iterator that
// points to the
// newly inserted node.
template<typename T, typename A>
typename slist<T,A>::iterator
    slist<T,A>::insert(iterator p, const T& x)
{
    // Allocate the new link before changing any pointers. If
    // newitem throws an
    // exception, the list is not affected.
    link_type* item = newitem(x, p.node_);
    if (p.node_ == 0) {
        p.prev_ = tail_;
        // At end
        if (tail_ == 0)
            head_ = tail_ = item; // Empty list
        else {
            tail_->next = item;
            tail_ = item;
        }
    }
}

```

```

else if (p.prev_ == 0)
    head_ = item;           // New head of list
else
    p.prev_ ->next = item;
p.node_ = item;
++count_;
return p;
}

// Erase the item at p. All erasures come here eventually. If
// erasing begin( ),
// update head_. If erasing the last item in the list, update
// tail_. Update the
// iterator to point to the node after the one being deleted.
template<typename T, typename A>
typename slist<T,A>::iterator slist<T,A>::erase(iterator p)
{
    link_type* item = p.node_;
    p.node_ = item->next;
    if (p.prev_ == 0)
        head_ = item->next;
    else
        p.prev_ ->next = item->next;
    if (item->next == 0)
        tail_ = p.prev_;
    --count_;
    delitem(item);
    return p;
}

// Comparison functions are straightforward.
template<typename T>
bool operator==(const slist<T>& a, const slist<T>& b)
{
    return a.size( ) == b.size( ) &&
        std::equal(a.begin(), a.end( ), b.begin( ));
}

```

## Using Containers

A container holds stuff. Naturally, you need to know how to add stuff to a container, remove stuff from a container, find stuff in a container, and so on.

### Value type requirements

Every container stores values and imposes certain restrictions on the values' types. Most important, the value must be copyable and assignable. The result of the copy constructor or assignment operator must be an exact copy of the original. (Note that you cannot store `auto_ptr<>` objects in a container because copies are not exact duplicates.)

In a sequence container, `operator==` is used to compare objects when searching. If you compare entire containers with any relational operators, the value types must also support `operator<`. All the relational operators are defined in terms of `operator==` and `operator<`.

In an associative container, values are stored in ascending order according to a comparison function or functor that you supply. The default is `std::less<>`, which uses `operator<`. Two objects A and B are considered to be equal (more precisely, *equivalent*) when `A < B` is false and `B < A` is false, so there is no need for `operator==`.

## **I n s e r t i n g v a l u e s**

To add an item to a container, call an `insert` member function. Sequence containers might also have `push_front` or `push_back` to insert an item at the beginning or end of the sequence. The `push_front` and `push_back` members exist only if they can be implemented in constant time. (Thus, for example, `vector` does not have `push_front`.)

Every container has an `insert( iter, item )` function, in which *iter* is an iterator and *item* is the item to insert. A sequence container inserts the item before the indicated position. Associative containers treat the iterator as a hint: if the item belongs immediately after the iterator's position, performance is constant instead of logarithmic.

Sequence containers have additional `insert` functions: for inserting many copies of an item at a position and for

copying a range to a given position. Associative containers have additional insert functions for inserting an item (with no positional hint) and for copying a range into the container.

## **Erasing values**

To remove an item from a container, call an erase member function. All containers have at least two erase functions: one that takes a single iterator to delete the item that the iterator points to, and another two iterators to delete every item in the range. Associative containers also have an erase function that takes a value as an argument to erase all matching items.

The standard containers are designed to be used with the standard iterators (see [Section 10.2](#) later in this chapter) and standard algorithms (see [Section 10.3](#) later in this chapter). The standard algorithms offer much more functionality than the containers' member functions, but they also have limitations. In particular, the standard algorithms cannot insert or erase items. For example, among the standard algorithms are `remove` and `remove_if`. Their names are suggestive but misleading. They do not remove anything from the container. Instead, they rearrange the elements of the container so that the items to retain are at the beginning. They return an iterator that points to the first item to be erased. Call `erase` with this iterator as the first argument and `end( )` as the second to erase the items from the container. This two-step process is needed because an iterator cannot erase anything. The only way to erase an item from a container is to call a member function of the container, and the standard algorithms do not have access to the containers, only to iterators. [Example 10-2](#) shows how to implement a generic erase function that calls `remove` and then the erase member function.

---

### *Example 10-2. Removing matching items from a sequence container*

```
// Erase all items from container c that are equal to item.
template<typename C>
void erase(C& c, const typename C::value_type& item)
{
    c.erase(std::remove(c.begin(), c.end( ), item), c.end( ));
}

template<typename C, typename Pred>
void erase_if(C& c, Pred pred)
{
    c.erase(std::remove_if(c.begin(), c.end( ), pred), c.end( ));
}

int main( )
{
    std::list<int> lst;
    ...
    // Erase all items == 20.
    erase(lst, 20);
    ...
    // Erase all items < 20.
    erase_if(lst, std::bind2nd(std::less<int>( ), 20));
    ...
}
```

## **Searching**

The standard algorithms provide several ways to search for items in a container: `adjacent_find`, `find`, `find_end`, `find_first_of`, `find_if`, `search`, and `search_n`. These algorithms essentially perform a linear search of a range. If you know exactly which item you want, you can search an associative container much faster by calling the `find` member function. For example, suppose you want to write a generic function, `contains`, that tells you whether a container contains at least one instance of an item.

[Example 10-3](#) shows one way, which relies on `find`, to implement this function.

### *Example 10-3. Determining whether a container contains an item*



```

// Need a type trait to tell us which containers are associative
and which are
// not (see Chapter 8).
struct associative_container_tag {};
struct sequence_container_tag {};

template<typename C>
struct is_associative
{};

template<typename T, typename A>
struct is_associative<std::list<T,A> >
{
    typedef sequence_container_tag tag;
};
// Ditto for vector and deque

template<typename T, typename C, typename A>
struct is_associative<std::set<T,C,A> >
{
    typedef associative_container_tag tag;
};
// Ditto for multiset, map, and multimap

template<typename C, typename T>
inline bool do_contains(const C& c, const T& item,
                       associative_container_tag)
{
    return c.end( ) != c.find(item);
}

template<typename C, typename T>
inline bool do_contains(const C& c, const T& item,
                       sequence_container_tag)
{
    return c.end() != std::find(c.begin( ), c.end( ), item);
}

// Here is the actual contains function. It dispatches to
do_contains, picking
// the appropriate overloaded function depending on the type of
the container c.
template<typename C, typename T>
bool contains(const C& c, const T& item)
{
    return do_contains(c, item, is_associative<C>::tag( ));
}

```

As you can see, iterators are important for using containers. You need them to insert at a specific position, identify an item for erasure, or specify ranges for algorithms. The next section discusses iterators in more depth.

# Iterators

An iterator is an abstraction of a pointer used for pointing into containers and other sequences. An ordinary pointer can point to different elements in an array. The ++ operator advances the pointer to the next element, and the \* operator dereferences the pointer to return a value from the array. Iterators generalize the concept so that the same operators have the same behavior for any container, even trees and lists. See the <iterator> section of [Chapter 13](#) for more details.

## Iterator Categories

There are five categories of iterators:

### Input

Permits you to read a sequence in one pass. The increment (++) operator advances to the next element, but there is no decrement operator. The dereference (\*) operator returns an rvalue, not an lvalue, so you can read elements but not modify them.

### Output

Permits you to write a sequence in one pass. The increment (++) operator advances to the next element, but there is no decrement operator. You can dereference an element only to assign a value to it. You cannot compare output iterators.

### Forward

Permits unidirectional access to a sequence. You can refer to and assign to an item as many times as you want. You can use a forward iterator wherever an input

iterator is required or wherever an output iterator is required.

### Bidirectional

Similar to a forward iterator but also supports the `--` (decrement) operator to move the iterator back one position.

### Random access

Similar to a bidirectional iterator but also supports the `[]` (subscript) operator to access any index in the sequence. Also, you can add or subtract an integer to move a random access iterator by more than one position at a time. Subtracting two random access iterators yields an integer distance between them. Thus, a random access iterator is most like a conventional pointer, and a pointer can be used as a random access iterator.

An input, forward, bidirectional, or random access iterator can be a `const_iterator`. Dereferencing a `const_iterator` yields a constant value, but otherwise behaves as described previously. See [Section 10.2.5](#) later in this chapter for details.

## Iterator Safety

The most important point to remember about iterators is that they are inherently unsafe. Like pointers, an iterator can point to a container that has been destroyed or to an element that has been erased. You can advance an iterator past the end of the container the same way a pointer can point past the end of an array. With a little care and caution, however, iterators are safe to use.

The first key to safe use of iterators is to make sure a program never dereferences an iterator that marks the end of a range. Two iterators can denote a range of values,

typically in a container. One iterator points to the start of the range and another marks the end of the range by pointing to a position one past the last element in the range. The mathematical notation of  $[first, last)$  tells you that the item that `first` points to is included in the range, but the item that `last` points to is excluded from the range.

A program must never dereference an iterator that points to one past the end of a range (e.g., `last`) because that iterator might not be valid. It might be pointing to one past the end of the elements of a container, for example.

Even a valid iterator can become invalid and therefore unsafe to use, for example if the item to which the iterator points is erased. The detailed descriptions in [Chapter 13](#) tell you this information for each container type. In general, iterators for the node-based containers (`list`, `set`, `multiset`, `map`, `multimap`) become invalid only when they point to an erased node. Iterators for the array-based containers (`deque`, `vector`) become invalid when the underlying array is reallocated, which might happen for any insertion and for some erasures.

## Special Iterators

Iterators are often used with containers, but they have many more uses. You can define iterators for almost any sequence of objects. The standard library includes several examples of non-container iterators, most notably I/O iterators and inserters.

At the lowest level, a stream is nothing more than a sequence of characters. At a slightly higher level, you can think of a stream as a sequence of objects, which would be read with operator `>>` or written with operator `<<`. Thus, the standard library includes the following I/O iterators: `istreambuf_iterator`, `ostreambuf_iterator`, `istream_iterator`, and `ostream_iterator`. [Example 10-4](#)

shows how to use `streambuf` iterators to copy one stream to another.

#### *Example 10-4. Copying streams with streambuf iterators*

```
template<typename charT, typename traits>
void copy(std::basic_ostream<charT,traits>& out,
         std::basic_istream<charT,traits>& in)
{
    std::copy(std::istreambuf_iterator<charT>(in),
              std::istreambuf_iterator<charT>( ),
              std::ostreambuf_iterator<charT>(out));
}
```

Another kind of output iterator is an *insert iterator*, which inserts items into a sequence collection. The insert iterator requires a container and an optional iterator to specify the position where the new items should be inserted. You can insert at the back of a sequence with `back_insert_iterator`, at the front of a sequence with `front_insert_iterator`, or at a specific position in any kind of container with `insert_iterator`. Each of these iterator class templates has an associated function template that creates the iterator for you and lets the compiler deduce the container type. [Example 10-5](#) shows how to read a series of numbers from a stream, store them in reverse order in a list, and print the list, one number per line.

#### *Example 10-5. Inserting numbers in a vector*

```
#include <algorithm>
#include <iostream>
#include <iterator>
#include <list>

int main( )
{
    using namespace std;

    list<double> data;
    copy(istream_iterator<double>(cin),
         istream_iterator<double>( ),
         front_inserter(data));
}
```

```

// Use the data...
// Write the data, one number per line.
copy(data.begin( ), data.end( ),
      ostream_iterator<double>(cout, "\n"));
}

```

## Custom Iterators

The simplest way to write your own iterator is to derive from the iterator class template. Specify the iterator category as the first template parameter and the item type as the second parameter. In most cases, you can use the default template arguments for the remaining parameters. (See `<iterator>` in [Chapter 13](#) for details.) The `slist` container from [Example 10-1](#) needs an iterator and a `const_iterator`. The only difference is that a `const_iterator` returns rvalues instead of lvalues. Most of the iteration logic can be factored into a base class. [Example 10-6](#) shows iterator and `base_iterator`; `const_iterator` is almost identical to `iterator`, so it is not shown.

### *Example 10-6. Writing a custom iterator*

```

// The declaration for iterator_base is nested in slist.
class iterator_base :
    public std::iterator<std::forward_iterator_tag, T> {
    friend class slist;
public:
    bool operator==(const iterator_base& iter) const
    { return node_ == iter.node_; }
    bool operator!=(const iterator_base& iter) const
    { return ! (*this == iter); }

protected:
    iterator_base(const iterator_base& iter)
    : prev_(iter.prev_), node_(iter.node_) {}
    iterator_base(slist::link_type* prev,
                 slist::link_type* node)
    : prev_(prev), node_(node) {}
    // If node_ == 0, the iterator == end( ).
    slist::link_type* node_;
    // A pointer to the node before node_ is needed to support

```

```

erase( ). If
    // prev_ == 0, the iterator points to the head of the list.
    slist::link_type* prev_;
private:
    iterator_base( );
};

// The declaration for iterator is nested in slist.
class iterator : public iterator_base {
    friend class slist;
public:
    iterator(const iterator& iter) : iterator_base(iter) {}
    iterator& operator++( ) { // Pre-increment
        this->prev_ = this->node_;
        this->node_ = this->node_->next;
        return *this;
    }
    iterator operator++(int) { // Post-increment
        iterator tmp = *this;
        operator++( );
        return tmp;
    }
    T& operator*( ) { return this->node_->value; }
    T* operator->( ) { return &this->node_->value; }
private:
    iterator(slist::link_type* prev, slist::link_type* node)
        : iterator_base(prev, node) {}
};

```

## const\_iterators

Every container must provide an iterator type and a `const_iterator` type. Functions such as `begin` and `end` return `iterator` when called on a non-const container and return `const_iterator` when called on a const container.

Note that a `const_iterator` (with underscore) is quite different from a `const iterator` (without underscore). A `const iterator` is a constant object of type `iterator`. Being constant, it cannot change, so it cannot advance to point to a different position. A `const_iterator`, on the other hand, is a non-const object of type `const_iterator`. It is not constant, so its value can change. The key



difference between `iterator` and `const_iterator` is that `iterator` returns lvalues of type `T`, and `const_iterator` returns unmodifiable objects, either rvalues or `const` lvalues of type `T`. The standard requires that a plain `iterator` be convertible to `const_iterator`, but not vice versa.

One problem is that some members of the standard containers (most notably `erase` and `insert`) take `iterator` as a parameter, not `const_iterator`. If you have a `const_iterator`, you cannot use it as an insertion or erasure position.

Another problem is that it might be difficult to compare an `iterator` with a `const_iterator`. If the compiler reports an error when you try to compare iterators for equality or inequality, try swapping the order of the iterators, that is, if `a == b` fails to compile, try `b == a`. Most likely, the problem is that `b` is a `const_iterator` and `a` is a plain `iterator`. By swapping the order, you let the compiler convert `a` to a `const_iterator` and allow the comparison.

For a full explanation of how best to work with `const_iterator`s, see Scott Meyers's *Effective STL* (Addison-Wesley).

## Reverse Iterators

Every container that supports bidirectional or random access iterators also provides reverse iterators, that is, iterators that start with the last element and "advance" toward the first element of the container. These iterators are named `reverse_iterator` and `const_reverse_iterator`.

The standard library includes the `reverse_iterator` class template as a convenient way to implement the `reverse_iterator` type. The `reverse_iterator` class

template is an iterator adapter that runs in the reverse direction of the adapted iterator. The adapted iterator must be a bidirectional or random access iterator. You can obtain the adapted iterator, given a reverse iterator, by calling the base function.

On paper, the reverse iterator seems like a good idea. After all, a bidirectional iterator can run in two directions. There is no reason why an iterator adapter could not implement `operator++` by calling the adapted iterator's `operator--` function.

Reverse iterators share a problem with `const_iterator`s, namely that several members, such as `insert` and `erase`, do not take an iterator template parameter but require the exact iterator type, as declared in the container class. The `reverse_iterator` type is not accepted, so you must pass the adapted iterator instead, which is returned from the base function.

As an insertion point, the base iterator works fine, but for erasing, it is one off from the desired position. The solution is to increment the reverse iterator, then call `base`, as shown in [Example 10-7](#).

#### *Example 10-7. A reverse iterator*

```
int main( )
{
    std::list<int> l;
    l.push_back(10); l.push_back(42); l.push_back(99);
    print(l);
    std::list<int>::reverse_iterator ri;
    ri = std::find(l.rbegin( ), l.rend( ), 42);
    l.insert(ri.base( ), 33);
    // OK: 33 inserted before 42, from the point of view of a
    // reverse iterator,
    // that is, 33 inserted after 42

    ri = std::find(l.rbegin( ), l.rend( ), 42);
    l.erase(ri.base( ));
    // Oops! Item 33 is deleted, not item 42.
```

```
    ri = std::find(l.rbegin( ), l.rend( ), 42);  
    l.erase(++ri.base( ));  
    // That's right! In order to delete the item ri points to, you  
    must advance ri  
    // first, then delete the item.  
}
```

For a full explanation of how best to work with reverse iterators, see Scott Meyer's *Effective STL* (Addison-Wesley).

# Algorithms

The so-called algorithms in the standard library distinguish C++ from other programming languages. Every major programming language has a set of containers, but in the traditional object-oriented approach, each container defines the operations that it permits, e.g., sorting, searching, and modifying. C++ turns object-oriented programming on its head and provides a set of function templates, called *algorithms*, that work with iterators, and therefore with almost any container.

The advantage of the C++ approach is that the library can contain a rich set of algorithms, and each algorithm can be written once and work with (almost) any kind of container. And when you define a custom container, it automatically works with the standard algorithms (assuming you implemented the container's iterators correctly). The set of algorithms is easily extensible without touching the container classes. Another benefit is that the algorithms work with iterators, not containers, so even non-container iterators (such as the stream iterators) can participate.

C++ algorithms have one disadvantage, however. Remember that iterators, like pointers, can be unsafe. Algorithms use iterators, and therefore are equally unsafe. Pass the wrong iterator to an algorithm, and the algorithm cannot detect the error and produces undefined behavior. Fortunately, most uses of algorithms make it easy to avoid programming errors.

Most of the standard algorithms are declared in the `<algorithm>` header, with some numerical algorithms in `<numeric>`. Refer to the respective sections of [Chapter 13](#) for details.

## How Algorithms Work

The generic algorithms all work in a similar fashion. They are all function templates, and most have one or more iterators as template parameters. Because the algorithms are templates, you can instantiate the function with any template arguments that meet the basic requirements. For example, `for_each` is declared as follows:

```
template<typename InIter, typename Function>  
Function for_each(InIter first, InIter last, Function func);
```

The names of the template parameters tell you what is expected as template arguments: `InIter` must be an input iterator, and `Function` must be a function pointer or functor. The documentation for `for_each` further tells you that `Function` must take one argument whose type is the `value_type` of `InIter`. That's all. The `InIter` argument can be anything that meets the requirements of an input iterator. Notice that no container is mentioned in the declaration or documentation of `for_each`. For example, you can use an `istream_iterator`.

For a programmer trained in traditional object-oriented programming, the flexibility of the standard algorithms might seem strange or backwards. Thinking in terms of algorithms takes some adjustment.

For example, some object-oriented container classes define `sort` as a member function or as a function that applies only to certain kinds of objects. (For example, Java defines `sort` only on arrays and `List` objects). If you have a new kind of container, you must duplicate the implementation of `sort` or make sure the implementation of your container maps to one of the standard implementations of `sort`. In C++, you can invent any kind of crazy container, and as long as it supports a random access iterator, you can use the standard `sort` function.

Whenever you need to process the contents of a container, you should think about how the standard algorithms can help you. For example, suppose you need to read a stream of numbers into a data array. Typically, you would set up a `while` loop to read the input stream and, for each number read, append the number to the array. Now rethink the problem in terms of an algorithmic solution. What you are actually doing is copying data from an input stream to an array, so you could use the `copy` algorithm:

```
std::copy(std::istream_iterator<double>(stream),
          std::istream_iterator<double>( ),
          std::back_inserter(data));
```

The `copy` algorithm copies all the items from one range to another. The input comes from an `istream_iterator`, which is an iterator interface for reading from an `istream`. The output range is a `back_inserter` (created by the `back_inserter` function), which is an output iterator that pushes items onto a container.

At first glance, the algorithmic solution doesn't seem any simpler or clearer than a straightforward loop:

```
double x;
while (stream >> x)
    data.push_back(x);
```

More complex examples demonstrate the value of the C++ algorithms. For example, all major programming languages have a type for character strings. They typically also have a function for finding substrings. What about the more general problem of finding a subrange in any larger range? Suppose a researcher is looking for patterns in a data set and wants to see if a small data pattern occurs in a larger data set. In C++, you can use the `search` algorithm:

```
std::vector<double> data;
...
if (std::search(data.begin( ), data.end( ), pattern.begin(),
               pattern.end( )) != data.end( ))
{
```

```
    // found the pattern...
}
```

A number of algorithms take a function pointer or functor (that is, an object that overloads operator( )) as one of the arguments. The algorithms call the function and possibly use the return value. For example, `count_if` counts the number of times the function returns a true (nonzero) result when applied to each element in a range:

```
bool negative(double x)
{
    return x < 0;
}

std::vector<double>::iterator::difference_type neg_cnt;
std::vector<double> data;
...
neg_cnt = std::count_if(data.begin( ), data.end( ), negative);
```

In spite of the unwieldy declaration for `neg_cnt`, the application of `count_if` to count the number of negative items in the data vector is easy to write and read.

If you don't want to write a function to be used only with an algorithm, you might be able to use the standard functors or function objects (which are declared in the `<functional>` header). For example, the same count of negative values can be obtained with the following:

```
std::vector<double>::iterator::difference_type neg_cnt;
std::vector<double> data;
...
neg_cnt = std::count_if(data.begin( ), data.end( ),
    std::bind2nd(std::less<double>, 0.0));
```

The `std::less` class template defines operator( ), so it takes two arguments and applies operator< to those arguments. The `bind2nd` function template takes a two-argument functor and binds a constant value (in this case `0.0`) as the second argument, returning a one-argument function (which is what `count_if` requires). The use of standard function objects can make the code harder to

read, but also helps avoid writing one-off custom functions. (The Boost project expands and enhances the standard library's binders. See [Appendix B](#) for information about Boost.)

When using function objects, be very careful if those objects maintain state or have global side effects. Some algorithms copy the function objects, and you must be sure that the state is also properly copied. The numerical algorithms do not permit function objects that have side effects.

[Example 10-8](#) shows one use of a function object. It accumulates statistical data for computing the mean and variance of a data set. Pass an instance of `Statistics` to the `for_each` algorithm to accumulate the statistics. The copy that is returned from `for_each` contains the desired results.

### *Example 10-8. Computing statistics with a functor*

```
#include <algorithm>
#include <cstdint>
#include <cmath>
#include <iostream>
#include <ostream>

template<typename T>
class Statistics {
public:
    typedef T value_type;
    Statistics( ) : n_(0), sum_(0), sumsq_(0) {}
    void operator( )(double x) {
        ++n_;
        sum_ += x;
        sumsq_ += x * x;
    }
    std::size_t count( ) const { return n_; }
    T sum( ) const { return sum_; }
    T sumsq( ) const { return sumsq_; }
    T mean( ) const { return sum_ / n_; }
    T variance( ) const
        { return (sumsq_ - sum_*sum_ / n_) / (n_ - 1); }
private:
```



```

std::size_t n_;
T sum_;
T sumsq_; // Sum of squares
};

int main( )
{
    using namespace std;

    Statistics<double> stat = for_each(
        istream_iterator<double>(cin),
        istream_iterator<double>( ),
        Statistics<double>( ));

    cout << "count=" << stat.count( ) << '\n';
    cout << "mean =" << stat.mean( ) << '\n';
    cout << "var  =" << stat.variance( ) << '\n';
    cout << "stdev=" << sqrt(stat.variance( )) << '\n';
    cout << "sum  =" << stat.sum( ) << '\n';
    cout << "sumsq=" << stat.sumsq( ) << '\n';
}

```

## Standard Algorithms

[Chapter 13](#) describes all the algorithms in detail. This section presents a categorized summary of the algorithms.

It is always your responsibility to ensure that the output range is large enough to accommodate the input.

If the algorithm name ends with `_if`, the final argument must be a *predicate*, that is, a function pointer or function object that returns a Boolean result (a result that is convertible to type `bool`).

### Nonmodifying operations

The following algorithms examine every element of a sequence without modifying the order:

`count`

Returns the number of items that match a given value

`count_if`

Returns the number of items for which a predicate returns true

`for_each`

Applies a function or functor to each item

## **Comparison**

The following algorithms compare objects or sequences (without modifying the elements):

`equal`

Determines whether two ranges have equivalent contents

`lexicographical_compare`

Determines whether one range is considered less than another range

`max`

Returns the maximum of two values

`max_element`

Finds the maximum value in a range

`min`

Returns the minimum of two values

`min_element`

Finds the minimum value in a range

`mismatch`

Finds the first position where two ranges differ

## **Searching**

The following algorithms search for a value or a subsequence in a sequence (without modifying the

elements):

`adjacent_find`

Finds the first position where an item is equal to its neighbor

`find`

Finds the first occurrence of a value in a range

`find_end`

Finds the last occurrence of a subsequence in a range

`find_first_of`

Finds the first position where a value matches any one item from a range of values

`find_if`

Finds the first position where a predicate returns true

`search` , `search_n`

Finds a subsequence in a range

## **Binary search**

The following algorithms apply a binary search to a sorted sequence. The sequence typically comes from a sequence container in which you have already sorted the elements. You can use an associative containers, but they provide the last three functions as member functions, which might result in better performance.

`binary_search`

Finds an item in a sorted range using a binary search

`equal_range`

Finds the upper and lower bounds

`lower_bound`

Finds the lower bound of where an item belongs in a sorted range

`upper_bound`

Finds the upper bound of where an item belongs in a sorted range

## **Modifying sequence operations**

The following algorithms modify a sequence:

`copy`

Copies an input range to an output range

`copy_backward`

Copies an input range to an output range, starting at the end of the output range

`fill`, `fill_n`

Fills a range with a value

`generate`, `generate_n`

Fills a range with values returned from a function

`iter_swap`

Swaps the values that two iterators point to

`random_shuffle`

Shuffles a range into random order

`remove`

Reorders a range to prepare to erase all elements equal to a given value

`remove_copy`

Copies a range, removing all items equal to a given value

`remove_copy_if`

Copies a range, removing all items for which a predicate returns true

`remove_if`

Reorders a range to prepare to erase all items for which a predicate returns true

`replace`

Replaces items of a given value with a new value

`replace_copy`

Copies a range, replacing items of a given value with a new value

`replace_copy_if`

Copies a range, replacing items for which a predicate returns true with a new value

`replace_if`

Replaces items for which a predicate returns true with a new value

`reverse`

Reverses a range in place

`reverse_copy`

Copies a range in reverse order

`rotate`

Rotates items from one end of a range to the other end

`rotate_copy`

Copies a range, rotating items from one end to the other

`swap_ranges`

Swaps values in two ranges

`transform`

Modifies every value in a range by applying a transformation function

`unique`

Reorders a range to prepare to erase all adjacent, duplicate items

`unique_copy`

Copies a range, removing adjacent, duplicate items

## **Sort i n g**

The following algorithms are related to sorting and partitioning. You can supply a comparison function or functor or rely on the default, which uses the `<` operator.

`nth_element`

Finds the item that belongs at the `nth` position (if the range were sorted) and reorders the range to partition it into items less than the `nth` item and items greater than or equal to the `nth` item.

`partial_sort`

Reorders a range so the first part is sorted.

`partial_sort_copy`

Copies a range so the first part is sorted.

`partition`

Reorders a range so that all items for which a predicate is true come before all items for which the predicate is false.

`sort`

Sorts items in ascending order.

`stable_partition`

Reorders a range so that all items for which a predicate is true come before all items for which the predicate is false. The relative order of items within a partition is maintained.

`stable_sort`

Sorts items in ascending order. The relative order of equal items is maintained.

## **Merging**

The following algorithms merge two sorted sequences:

`inplace_merge`

Merges two sorted, consecutive subranges in place, so the results replace the original ranges

`merge`

Merges two sorted ranges, copying the results to a separate range

## **Set operations**

The following algorithms apply standard set operations to sorted sequences:

`includes`

Determines whether one sorted range is a subset of another

`set_difference`

Copies the set difference of two sorted ranges to an output range

`set_intersection`

Copies the intersection of two sorted ranges to an output range

`set_symmetric_difference`

Copies the symmetric difference of two sorted ranges to an output range

`set_union`

Copies the union of two sorted ranges to an output range

## **Heap operations**

The following algorithms treat a sequence as a heap data structure:

`make_heap`

Reorders a range into heap order

`pop_heap`

Reorders a range to remove the first item in the heap

`push_heap`

Reorders a range to add the last item to the heap

`sort_heap`

Reorders a range that starts in heap order into fully sorted order

## **Permutations**

The following reorder the elements of a sequence to generate permutations:

`next_permutation`

Reorders a range to form the next permutation

`prev_permutation`



Reorders a range to form the previous permutation

## Custom Algorithms

Writing your own algorithm is easy. Some care is always needed when writing function templates (as discussed in [Chapter 7](#)), but generic algorithms do not present any special or unusual challenges. Be sure you understand the requirements of the different categories of iterators and write your algorithm to use the most general category possible. You might even want to specialize your algorithm to improve its performance with some categories.

The first generic algorithm that most programmers will probably write is `copy_if`, which was inexplicably omitted from the standard. The `copy_if` function copies an input range to an output range, copying only the values for which a predicate returns true (nonzero). [Example 10-9](#) shows a simple implementation of `copy_if`.

*Example 10-9. One way to implement the `copy_if` function*

```
template<typename InIter, typename OutIter, typename Pred>
OutIter copy_if(InIter first, InIter last, OutIter result, Pred
pred)
{
    for (; first != last; ++first)
        if (pred(*first)) {
            *result = *first;
            ++result;
        }
    return result;
}
```

You can also specialize an algorithm. For example, you might be able to implement the algorithm more efficiently for a random access iterator. In this case, you can write helper functions and use the `iterator_category` trait to choose a specialized implementation. ([Chapter 8](#) has more information about traits, including an example of how to use iterator traits to optimize a function template.)

The real trick in designing and writing algorithms is being able to generalize the problem and then find an efficient solution. Before running off to write your own solution, check the standard library. Your problem might already have a solution.

For example, I recently wanted to write an algorithm to find the median value in a range. There is no median algorithm, but there is `nth_element`, which solves the more general problem of finding the element at any sorted index. Writing median became a trivial matter of making a temporary copy of the data, calling `nth_element`, and then returning an iterator that points to the median value in the original range. Because median makes two passes over the input range, a forward iterator is required, as shown in [Example 10-10](#).

*Example 10-10. Finding the median of a range*

```
template<typename FwdIter, typename Compare>
FwdIter median(FwdIter first, FwdIter last, Compare compare)
{
    typedef typename std::iterator_traits<FwdIter>::value_type
    value_type;
    std::vector<value_type> tmp(first, last);
    typename std::vector<value_type>::size_type median_pos =
    tmp.size( ) / 2;
    std::nth_element(tmp.begin( ), tmp.begin( ) + median_pos,
                    tmp.end( ), compare);
    return std::find(first, last, tmp[median_pos]);
}
```

## Chapter 11. Preprocessor Reference

The preprocessing step occurs before the main compilation step. Historically, the preprocessor has been a separate program, but compilers are not required to implement the preprocessor in that way. Because of its history, though, the preprocessor has syntax and semantics that are quite different from the rest of C++. See [Chapter 1](#) for information about all the steps in compiling a source file.

The preprocessor handles preprocessing *directives*, which can define and undefine macros, establish regions of conditional compilation, include other source files, and control the compilation process somewhat. A *macro* is a name that represents other text, called the *macro replacement text*. When the macro name is seen in the source file, the preprocessor replaces the name with the replacement text. A macro can have formal parameters, and actual arguments are substituted in the expansion.

Preprocessor directives obey different syntax rules from the rest of the language. Directives are line-oriented. Each directive starts with whitespace characters followed by # as the first non-space character on a line. After the # is more optional whitespace (no newlines are permitted) followed by the directive name. Each directive extends to the end of the line. A backslash (\) at the end of the line continues the directive onto the subsequent line.

The directive name must be one of the names listed in this chapter. Any other preprocessing token after the initial # character is an error.

## Name

# operator and directive — Stringify operator and null directive

## Synopsis

```
# // Null directive  
# identifier
```

A preprocessor directive with no directive name is called a *null directive* . It has no effect.

The # operator can also be used as a unary operator, sometimes called the *stringify operator* because it turns its operand into a string. It can be used only in the macro replacement text of a #define directive. It must be followed by a parameter name for the macro being defined. The # operator and the parameter name are replaced by a string literal whose contents are the text of the macro argument that corresponds to the macro parameter. The macro argument is not expanded before being converted to a string. Whitespace in the argument is condensed to a single space character between tokens; leading and trailing whitespace is removed.

The evaluation order of # and ## operators is undefined. If the order is important, you can control the order by using multiple macros.

## Example

The following example prints the text [now is the time]:

```
#define now then  
#define is was  
#define print(stuff)  std::cout << "[" #stuff "]" \n"  
print( now   is the   time   );
```

## See Also

[## operator](#), [#define directive](#)

# Name

`##` operator — Concatenation operator

# Synopsis

*identifier ## identifier*

The `##` operator is a binary operator, sometimes called the *concatenation operator* because it concatenates preprocessor tokens. It can be used only in the macro replacement text of a `#define` directive. It must not appear at the start or end of the macro replacement text. The operands of the `##` operator must be parameter names for the macro being defined. They are replaced by the corresponding macro arguments, which are not expanded. The tokens immediately adjacent to the `##` operator are concatenated to form a single token. If the result is not a valid token, the behavior is undefined; otherwise, the token is expanded normally.

The evaluation order of `#` and `##` operators is undefined. If the order is important, you can control the order by using multiple macros.

# Example

The following example prints `std` to `cout` because the `concat` macro assembles the token `std` from `s`, `t`, and `d`:

```
#define s this is not expanded by the concatenation operator
#define t nor is this, so the result is the token std
#define concat(part1, part2, part3) part1 ## part2 ## part3
concat(s, t, d)::cout << "std";
```

# See Also

[# operator](#), [#define directive](#)

# Name

`#define` directive — Defines a macro

## Synopsis

```
#define identifier
           definition
#define identifier( ) definition
#define identifier( ) definition
#define identifier(identifier-list) definition
```

The `#define` directive defines a macro named *identifier*. The macro's replacement text is the list of tokens shown as *definition*. The macro can be simple, with no arguments, or can have an argument list. The argument list is introduced by a left parenthesis that immediately follows the macro name. If there is a space between *identifier* and (, the ( is interpreted as the start of the definition of a simple macro. The *identifier-list* can be empty, or it can be a list of identifiers separated by commas. Whitespace is permitted in the *identifier-list* and before the closing parenthesis.

C programmers are accustomed to using macros to declare constants and simple inline functions, but C++ offers `const` declarations, true inline functions, and templates. Macros are therefore used much less often in C++ than in C. The main drawback to macros is that they do not obey scope or namespace rules. When you must use macros, a common convention is to use all uppercase letters for the macro names, and never use all uppercase letters for non-macro names.

A macro's scope is from the point of definition to the end of the source file, or until you undefine the macro with the `#undef` directive. If you try to repeat a macro definition, the

new definition must be identical to the original definition. The only way to give a macro a different definition is to undefine it first.

## Warning



If you `#include` any standard header, a macro name cannot be the same as a reserved keyword or any of the names declared in the header. Many compilers accept keywords as macro names, but your program would still be wrong and would not be portable to a compiler that is more strict about detecting this particular error. Even if you do not `#include` a standard header, using macros to redefine keywords is usually a bad idea.

Wherever a macro name appears as a distinct token after its definition, it is replaced with the replacement text. Macro names are not replaced inside string and character literals, however. The replacement text is rescanned for macro names, which are recursively expanded until no more replacements take place. During replacement, the original macro name is not expanded if it appears in any replacement text. Here is a simple example of a macro definition and use:

```
#define NAME "NAME = Tempest Software, Inc."  
char companyName[] = NAME;
```

During the macro expansion phase of compilation, the token `NAME` will be replaced by its expansion, with the following result:

```
char companyName[] = "NAME = Tempest Software, Inc.";
```

The replacement text is never interpreted as a preprocessing directive. This means, for example, you cannot `#define` a macro within a macro's replacement text.



Also, directive names are not subject to macro replacement (although directive arguments are).

You can also declare a macro with a parameter list, which is sometimes called a *function-like macro* :

```
#define DECLARE(x,y, z)    x y = z
#define PRINT(a)    (::std::cout << (a) << '\n')
```

To use a function-like macro, the macro name must be followed by a comma-separated argument list in parentheses. A single argument can contain balanced parentheses, and within those parentheses, you can have commas, which are not interpreted as argument separators. The macro invocation must have the same number of arguments as the macro definition has parameters. Newlines are permitted as ordinary whitespace characters in a macro invocation.

The following example uses the DECLARE and PRINT macros defined in the previous example:

```
int main( )
{
    DECLARE(int, x, 42);
    PRINT((x = 10, x+2));
}
```

In the macro replacement text, each occurrence of a parameter name is replaced by the corresponding argument. For example, the macro expansion for the previous example results in the following:

```
int main( )
{
    int x = 42;
    (::std::cout << (x = 10, x + 2) << '\n');
}
```

You must be extra cautious when using a template instantiation as a macro argument. The angle brackets that surround the template arguments are not treated specially for macro arguments, so commas that separate the

template arguments are interpreted as separators for the macro arguments. In the following example, the `DECL` macro attempts to declare an object named `n` with type `t`. This works fine for a simple type, such as `int`, but fails with a template instantiation. When used with `map<int, int>`, the comma separates macro arguments, so the preprocessor sees three macro arguments—`std::map<int, int>`, and `m`—and reports an error:

```
#define DECL(t, n) t n = t( )
DECL(int, zero);           // Expands to int zero = int( )
DECL(std::map<int,int>, m); // Error
```

When a macro is expanded, the macro arguments are expanded, and each parameter is replaced by its corresponding expanded argument unless the parameter is an operand to the `#` or `##` operator. After the arguments have been expanded, the `#` and `##` operators are evaluated, and the resulting text is rescanned for macros. The macro name is expanded only once, so rescanning does not expand the name of the macro being expanded.

## Predefined Macros

The following macros are predefined. Do not undefine or redefine any of the predefined macros.

`__cplusplus`

Has the value `199711L`. Future versions of the C++ standard will use a larger value. Nonconforming compilers should use a different value.



`__DATE__`

Expands to the date of compilation, as a string literal, in the form `"Mmm dd yyyy"`, in which `dd` begins with a space for numbers less than 10. An implementation is free to

substitute a different date, but the form is always the same, and the date is always valid.

`__FILE__`

Expands to the name, as a string literal, of the source file being compiled.

`__LINE__`

Expands to the line number, as a decimal constant, of the source file being compiled.



`__STDC__`

Is implementation-defined. C++ compilers might define this macro; if it is defined, the value is implementation-defined. Note that C compilers are required to define `__STDC__` as 1, and in some implementations, the same preprocessor might be used for C and C++.



`__TIME__`

Expands to the compilation time, as a string literal, in the form "hh:mm:ss". An implementation is free to substitute a different time, but the form is always the same, and the time is always valid.

An implementation is free to predefine other macros that use any of the reserved names, such as names that contain two adjacent underscores or a leading underscore followed by an uppercase letter. For example, compilers often define macros to indicate the host or target platform—e.g., `__linux__`. Consult your compiler's documentation for details.

## Examples

When writing a container class template (see [Chapter 10](#)), it is important to detect when a template parameter is an integral type. There are several ways to do this. One way is to use type traits ([Chapter 8](#)). A template declares a special tag for all integral types and a different tag for all other types. The traits template is then specialized for the integral types, which is repetitive, tedious, and error-prone. Using a macro, however, reduces the opportunity for errors because the macro body is written once. [Example 11-1](#) shows how the `DECL_IS_INTEGER` macro specializes the `is_integer` class template for each built-in integral type.

*Example 11-1. Defining type traits with a macro*

```
// Type trait to test whether a type is an integer.
struct is_integer_tag {};
struct is_not_integer_tag {};

// The default is that a type is not an integral type.
template<typename T>
struct is_integer {
    enum { value = 0 };
    typedef is_not_integer_tag tag;
};

// Explicitly override the default for all integral types.
#define DECL_IS_INTEGER(T) \
template<> \
struct is_integer<T> { \
    enum { value = 1 }; \
    typedef is_integer_tag tag; \
}
DECL_IS_INTEGER(bool);
DECL_IS_INTEGER(char);
DECL_IS_INTEGER(signed char);
DECL_IS_INTEGER(unsigned char);
DECL_IS_INTEGER(int);
DECL_IS_INTEGER(unsigned int);
DECL_IS_INTEGER(short);
DECL_IS_INTEGER(unsigned short);
DECL_IS_INTEGER(long);
DECL_IS_INTEGER(unsigned long);

#undef DECL_IS_INTEGER
```

[Example 11-2](#) shows another way that macros are used when testing the string class. The TEST macro calls a function and prints the result. The TEST macro cannot be implemented as a function because it uses the # operator.

### *Example 11-2. Testing functions*

```
#include <iostream>
#include <string>

int main( )
{
    using namespace std;

    string s("hello, world");

#define TEST(func) cout << #func "=" << s.func << '\n'

    TEST(erase(9, 1));
    TEST(erase(5));
    TEST(find_first_not_of("aeiou"));
    ...
}
```

#### **Ti p**

Most compilers have an option in which the compiler runs only the preprocessor, and you can examine the results after all macros have been expanded and all preprocessor directives have been evaluated. This mode can be helpful when debugging an incorrect macro expansion.

[Example 11-3](#) is a contrived example that illustrates how macros are expanded. Try running the example through your compiler to see if the results are correct. (Other than whitespace, the results should be the same as what is shown in the rest of this section.)

### *Example 11-3. Expanding macros*

```
#define x          x.y
#define STR(x)    #x
```

```

#define XSTR(s)      STR(s)
#define CONCAT(x, y) x ## y
#define PARENS(x)   (x)
#define APPLY(x,y)  x(y)
#define hello      HI

x                               // x.y
CONCAT(ST, R)(hello)          // "hello"
CONCAT(X,STR)(hello)          // "HI"
CONCAT(S, TR)PARENS(hello)    // STR(HI)
CONCAT(S, TR)(PARENS(hello))  // "PARENS(hello)"
APPLY(CONCAT(S, TR), hello)   // "HI"

```

The first macro expansion shows how the macro name `x` is not expanded in the replacement text. The result is simply:

```
x.y
```

The second macro expansion shows how the `CONCAT` macro forms a new token `STR` from its arguments. After the `CONCAT` macro is evaluated, the text is rescanned. The `STR` macro is then invoked with the `hello` argument. Because the `x` parameter is an operand of `#`, the argument is not expanded. Instead, `#` is applied to `hello` to produce the result:

```
"hello"
```

The third macro expansion is like the second, except it invokes `XSTR` instead of `STR`. The difference is that `XSTR` expands its argument, `s`, because the replacement text, `STR(s)`, does not use the `#` or `##` operators. Thus, `XSTR(hello)` expands to `STR(HI)`, which has the following result:

```
"HI"
```

The fourth expansion also invokes `CONCAT` to produce `STR`, but `STR` is not followed by a left parenthesis, so it is not expanded as a macro. Instead, it is followed by the `PARENS` macro. The parameter of `PARENS` is not an operand of `#` or `##`, so it is expanded, which means the argument `hello` expands to `HI`, and the final result is:

```
STR(HI)
```

The fifth expansion is just like the second, but emphasizes how the argument to `STR` is not expanded. The result is:

```
"PARENS(hello)"
```

The final macro expansion shows how to expand `hello` as an argument to `STR`, even when `STR` is the result of the `CONCAT` macro. The parameters of `APPLY` are expanded, resulting in the text `STR(HI)`, which expands to:

```
"HI"
```

## See Also

[#undef directive](#)

## Name

defined operator — Tests whether a macro is defined

## Synopsis

```
defined(identifier)  
defined identifier
```

The unary operator `defined( identifier )` (also written as `defined identifier`, without the parentheses) evaluates to 1 if *identifier* is a known macro name at the point of the defined operator, or 0 if it is not known. The operator is evaluated only in the argument to an `#if` or `#elif` directive.

The behavior is undefined if the defined operator is used in any other way, or if the token `defined` results from macro expansion.

## See Also

[#define directive](#), [#elif directive](#), [#if directive](#), [#ifdef directive](#), [#ifndef directive](#), [#undef directive](#)



## Name

`#elif` directive — Else-if for conditional compilation

## Synopsis

```
#elif constant-expression
```

The `#elif` directive marks a region of conditional compilation. Every `#elif` must be paired with an introductory directive: `#if`, `#ifdef`, or `#ifndef`. If the initial condition was false (0), and every subsequent `#elif` condition is false, and *constant-expression* is true (nonzero), subsequent statements are compiled until the next `#elif`, `#else`, or `#endif` directive is reached for this level of nesting.

## See Also

[#else directive](#), [#if directive](#)

## Name

`#else` directive — Else for conditional compilation

## Synopsis

```
#else
```

The `#else` directive marks a region of conditional compilation. Every `#else` must be paired with an introductory directive: `#if`, `#ifdef`, or `#ifndef`. There can be any number of intervening `#elif` directives. If the initial condition was false (0), and every subsequent `#elif` condition is false, statements that follow the `#else` directive are compiled until the corresponding `#endif` directive is reached.

## See Also

[`#if` directive](#)

## Name

`#endif` directive — Ends conditional compilation

## Synopsis

```
#endif
```

The `#endif` directive ends a region of conditional compilation.

## See Also

[`#if` directive](#)

## Name

`#error` directive — Issues an error message

## Synopsis

```
#error message
```

The `#error` directive tells the preprocessor to issue an error message and mark the source file as ill-formed, just as if the programmer made a programming error that the compiler detected. You can supply any sequence of preprocessor tokens as the *message*, and those tokens are echoed in the error message.

## Example

```
#if !defined(_ _cplusplus) || (_ _cplusplus < 199711L)
    #error Not a conforming C++ compiler.
#endif
```

## See Also

[#line directive](#)

# Name

`#if` directive — Tests a condition

# Synopsis

*#if constant-expression*

The `#if` directive begins a region of conditional compilation, that is, a region within a source file where preprocessor directives determine whether the code in the region is compiled. A conditional region starts with `#ifdef`, `#ifndef`, or `#if` and ends with `#endif`. Each region can have any number of `#elif` directives and an optional `#else` directive after all the `#elif` directives. The basic form to use is:

```
#if defined(__win32__ )
    const char os[] = "Microsoft Windows";
#elif defined(__linux__ ) or defined(__unix__ )
    const char os[] = "UNIX (or variant)";
#elif defined(__vms__ )
    const char os[] = "VMS";
#else
    const char os[] = "(unknown)";
#endif
```

Macros in the directive argument are expanded, except for the operands of the `defined` operator. The constant expression is evaluated, and if the result is nonzero, the `#if` condition is true, and the code in the region that immediately follows is compiled. The region ends with `#else`, `#elif`, or `#endif`. If the `#if` expression is false, the condition for the next `#elif` is evaluated, and if that expression is true, its region is compiled, and so on. If all `#elif` expressions are false, and `#else` is present, its region is compiled. Conditional processing ends with the corresponding `#endif` directive.

Conditionals can be nested. Within an inner region, the preprocessor keeps track of conditional directives even if the region is not being compiled, so conditional directives can be properly matched.

The `#if` and `#elif` directives take a single parameter, a constant expression. The expression differs slightly from non-preprocessor constant expressions:

- You can use the `defined` operator.
- Integers are long, that is, `int` values (and values that are promoted to `int`) have the same representation as `long int`, and `unsigned int` values have the same representation as `unsigned long`. All `bool` values are promoted to integers, including the keywords `true` and `false`.
- Character literals are converted to the execution character set. The numeric value of a character in a preprocessor expression is not necessarily the same as the value of the same character in a non-preprocessor expression. A character may have a negative value.
- Keywords that are alternative operators for symbolic operators (i.e., `and`, `and_eq`, `bitand`, `bitor`, `compl`, `not`, `not_eq`, `or`, `or_eq`, `xor`, and `xor_eq`) have their usual meaning, although it is ineffective to try using assignment operators (`and_eq`, `or_eq`, and `xor_eq`) in an `#if` or `#elif` condition.



Other identifiers and keywords that remain after macro expansion are replaced by `0`. (It might seem strange to convert keywords, such as `sizeof`, to the integer `0`, but that is the rule. Some compilers fail to follow this particular rule.) One consequence of this rule is that you cannot use type casts or `new`, `delete`, `sizeof`,

throw, or typeid expressions in an `#if` or `#elif` condition.

Conditional directives are most often used to guard header files from multiple inclusion. All the standard headers are guarded, so including them more than once has no harmful effects. This is important because an implementation might include one header in another header. For example, `<map>` might include `<utility>` to get the declaration for the `pair<>` template. If you explicitly `#include <map>` and `#include <utility>`, you might end up including `<utility>` more than once.

Another common use is for system- or compiler-specific code. Every compiler predefines one or more macros to identify the compiler and possibly the host operating system (such as `__linux__` or `__GNUC__`). Consult your compiler's documentation to learn which macro names are predefined.

## Examples

[Example 11-4](#) shows one way to nest conditional directives.

### *Example 11-4. Nesting conditional directives*

```
#define zero zero // Identifiers are converted to 0.
#define one true // Bool expressions are promoted to int.
#if one
// This region is compiled.
#if zero
This region can contain erroneous C++ code. The code is not
compiled, so the errors do not matter.
#else // This #else matches the inner #if.
// This region is compiled.
const int zero = 0;
#endif // This #endif matches the inner #if.
int x = zero;
#else
This #else matches the outer #if. Because the #if
condition was true, the #else region is not compiled.
#endif
```

You can guard your own headers by using conditional directives to define a guard macro and using the guard macro to ensure the file's contents are compiled only when the macro is not defined, as shown in [Example 11-5](#).

*Example 11-5. Guarding a header against multiple inclusion*

```
// In the header file employee.h #ifndef EMPLOYEE_H
    #define EMPLOYEE_H
// Thus, the entire contents of the file are compiled only when
EMPLOYEE_H is not
// defined. The first time the file is #included, the macro is
not defined, in
// which case it is immediately defined. The second and
subsequent times the same
// header is included in the same source file; the macro and
conditional
// directives ensure that the entire file is skipped.

class employee { ... };

#endif // End of employee.h
```

## See Also

[#elif directive](#), [#else directive](#), [#endif directive](#), [#ifdef directive](#), [#ifndef directive](#)



## Name

`#ifdef` directive — Tests whether a macro is defined

## Synopsis

```
#ifdef identifier
```

The `#ifdef` directive begins a region of conditional compilation. It takes a single *identifier* as an argument and is equivalent to `#if defined identifier`.

## See Also

[defined operator](#), [#if directive](#), [#ifndef directive](#)

## Name

`#ifndef` directive — Tests whether a macro is undefined

## Synopsis

```
#ifndef identifier
```

The `#ifndef` directive begins a region of conditional compilation. It takes a single *identifier* as an argument and is equivalent to `#if not defined identifier`.

## See Also

[defined operator](#), [#if directive](#), [#ifdef directive](#)

## Name

`#include` directive — Includes another source file

## Synopsis

```
#include <header>  
#include "sourcefile"
```

The `#include` directive includes the contents of a standard header or source file. The first form searches for *header* and replaces the directive with the entire contents of the header. The second form searches for *sourcefile* and replaces the directive with the entire contents of the named source file.

The basic action of `#include` is to read the named file or header as though its entire contents appeared in the source file at the position of the `#include` directive. Typically, common declarations are placed in a separate file, such as `decl.h`, and `#include "decl.h"` is used in every source file that depends on those declarations.



If a source file contains the directive `#include "filename"`, and the compiler cannot find the external file named *filename*, the compiler also tries the form `#include <filename >`. Most compilers implement these two forms of `#include` by searching in different folders or directories. For example, the quote form searches in the current directory or in the directory that contains the source file, and the angle-bracket form searches only in "system" directories. Such details are implementation-defined, and some compilers might introduce further distinctions between the two forms.

It is possible, for example, for a compiler to recognize only the standard headers in the `< header >` form and use built-in knowledge of the standard headers without referencing any external files. This hypothetical compiler might report an error for all other uses of the angle-bracket form and require all external file inclusions to use the quote form. Such a compiler would not be very popular, however, because common practice is to treat `< header >` and `"header"` as equivalent forms, except when applying the rules for locating the external file named `header`.

It is common practice to install third-party libraries in common directories and to configure compilers to look in these directories for `< header >` inclusions. For example, if you use Boost (described in [Appendix B](#)), you might use `#include <any.hpp>` to obtain the `boost::any` class template. Another common practice is to install such libraries in subdirectories. On a Unix system, for example, you might install Boost in the `boost` subdirectory of one of the standard system directories and use `#include <boost/any.hpp>`. You should be careful, however, because using system-specific filenames is not portable.

The only guarantee that the standard offers is that if `filename` consists of a sequence of letters and underscore characters followed by a period and a single letter or underscore, then the implementation must provide a unique mapping of `filename` to a source file (optionally ignoring case distinctions). The standard permits universal characters in `filename`, but you should avoid them when you need maximum portability because some compilers do not support universal characters.

The implementation defines how and where the preprocessor searches for `header` or `filename`, how `filename` maps to an external filename, whether filenames heed or ignore case distinctions, and whether different

character sequences for *filename* represent distinct external files. For example, under Windows, "foo.h" and "F00.H" are usually the same file, but under Unix, they are usually different files. If the filesystem supports links, such as Unix, two names such as "foo.h" and "bar.h" might name the same file; in other environments, you might be guaranteed that different filenames refer to distinct files.

The most common convention is that `<header>` refers only to standard headers and to vendor-supplied extensions to the standard. Compilers typically have a way for you to supply your own additional libraries and use the associated headers as `<header>` includes. The quoted form is used for all header files that are part of the application, and those are typically located in the same directory or folder as the application's source files. The most common filename convention is to end header names with `.h` (for header), although `.hpp` is also common. For example, suppose you wrote a class to represent an employee. Put the class definition in `employee.h` and the definitions of the members in `employee.cpp`. Any other file that needs to use the employee class can `#include "employee.h"` and use the class definition:

```
#include <set>
#include "employee.h"

class business_group {
private:
    std::set<employee> employees_;
    ...
};
```

You can use other preprocessor tokens in an `#include` directive, provided they expand to one of the two standard forms. Each header name or filename must be a single preprocessor token; you cannot combine tokens to form a name. To preserve portability, use macros only for the entire sequence of the `#include` argument:

```
#define HEADER "this.h"  
#include HEADER
```

## See Also

[#if directive](#)

## Name

`#line` directive — Changes the line number in error messages

## Synopsis

```
#line digits  
#line digits string
```

The `#line` directive changes the compiler's notion of the current filename and line number. The first form changes the line number (as expressed by the `__LINE__` directive and used in error messages) to *digits*. The second form changes the line number to *digits* and the filename to the contents of *string*. The new file name is used as the value of the `__FILE__` macro.

The `#line` directive is typically used by programs that generate C++ as output from some other input. The directive records the original filename and line number that produced the C++ code. Error messages and debuggers can point to the original file instead of to the intermediate C++ source file.

## See Also

[#error directive](#)

# Name

#pragma directive — Controls the compiler

# Synopsis

*#pragma tokens*



The #pragma directive is implementation-defined. An implementation can define certain pragma parameters to control the compiler. The preprocessor ignores any pragma that it does not recognize.

Because pragmas are highly compiler-dependent, you should avoid using them as much as possible. Most compilers let you control the compiler by providing command-line options, configuration files, or project files. Do your best to keep compiler-specific information out of your source files. When you must use pragmas, protect them with conditional directives for specific compilers.

# Example

```
#ifdef _BORLANDC_
    #pragma pack
#endif
#ifdef _COMO_
    #pragma instantiate
#endif
```



## Name

`#undef` directive — Undefines a macro

## Synopsis

```
#undef identifier
```

The `#undef` directive deletes the definition of the macro named *identifier*. If *identifier* is not a macro name, the directive has no effect. If you attempt to undefine the identifier defined or any predefined macro, the behavior is undefined.

## See Also

[`#define` directive](#)

## Chapter 12. Language Reference

Here begins the alphabetic reference. This chapter presents each language keyword with a syntax summary, description, and, in some cases, an example. The syntax summaries use a modified BNF (Backus Normal Form or Backus-Naur Form):

- Terminal symbols (keywords and operator symbols) are in a constant-width typeface.
- To avoid ambiguity, a terminal symbol that might be mistaken for a BNF metacharacter (e.g., a vertical bar or a square bracket) is enclosed in quotes (e.g., "|").
- Nonterminal symbols (syntax elements) are in an *italic* typeface.
- Optional elements are in square brackets ([ *like this* ]).
- Choices are separated by vertical bars (|).
- A *production* (syntax description) is introduced with := or ::= . The traditional symbol (:=) is used for a complete definition. The abbreviated symbol (:=) is used when the righthand side is incomplete. For example, here is the complete definition of *function-specifier* as it is given under *declaration*:

```
function-specifier ::= explicit | inline | virtual
```

The following is a partial production of *function-specifier*:

```
function-specifier := inline
```

The abbreviated symbol (:=) lets you see that the syntax summary is incomplete. Whenever an incomplete rule is used, a cross reference (under "See Also") leads you to the complete rule.

The starting point for parsing a C++ source file (the start symbol) is *translation-unit*, which you can find under *declaration*.

C++ syntax is complicated, and even simple statements require an understanding of many different syntax elements. To help you, this chapter duplicates some syntax rules and has plenty of cross references to help you find the parts you need.

Almost every language element is discussed at greater length in [Chapter 1-Chapter 7](#); each description in this chapter includes a reference to the relevant chapter or chapters. References to specific headers (e.g., `<new>`) are for the corresponding sections in [Chapter 13](#).

## Name

and operator — Logical and operator

## Synopsis

```
logical-and-expr := logical-and-expr && inclusive-or-expr |  
logical-and-expr and inclusive-or-expr
```

The *logical and* operator converts its operands to type `bool` and returns a `bool` result. This built-in operator is a short-circuit operator, so if the left operand is `false`, the expression yields `false` without evaluating the right operand. Note that an overloaded operator `and` cannot be short-circuited and must evaluate both operands. The keyword `and` is interchangeable with the `&&` token.

## Example

```
int* p;  
if (p != NULL and *p != 0)  
    do_stuff(*p);
```

## See Also

`bitand`, `bool`, `expression`, `not`, `or`, [Chapter 3](#), `<ciso646>`

## Name

and\_eq operator — Bitwise and assignment operator

## Synopsis

```
assignment-expr := logical-or-expr &= assignment-  
expr |  
logical-or-expr and_eq assignment-expr
```

The `and_eq` operator is an assignment operator that performs bitwise *and*. It is equivalent to `logical-or-expr = logical-or-expr & assignment-expr` except that `logical-or-expr` is evaluated only once. The keyword `and_eq` is interchangeable with the `&=` token.

## Example

```
unsigned bitmask = 0xFFFF;  
bitmask&= ~0x7E; // bitmask becomes 0xFF81.
```

## See Also

`bitand`, `expression`, `or_eq`, `xor_eq`, [Chapter 3](#), `<ciso646>`

## Name

asm definition — Inline assembler definition

## Synopsis

```
block-decl := asm-defn  
asm-defn ::= asm ( string-literal ) ;
```



The asm definition is implementation-defined. Typically, the `string-literal` contains assembler instructions. Some compilers extend the asm syntax to make it easier to write larger blocks of assembler code.

## Example

```
asm("mov 4, %eax"); // GNU on Intel IA32  
asm("mov eax, 4"); // Borland on Intel IA32
```

## See Also

declaration

# Name

auto storage class — Automatic variable specifier

# Synopsis

*storage-class-specifier* := auto

The auto storage class specifier declares a local variable to be automatic. The object is constructed when execution reaches the variable's declaration, and the object is destroyed when execution leaves the scope where it is declared.

All local variables and function parameters are auto by default, so the explicit auto specifier is rarely used.

# Example

```
int foo(auto int parm)
{
    auto int sqr = parm * parm;
    return sqr;
}
```

# See Also

declaration, register, [Chapter 2](#)

## Name

bitand operator — Bitwise and operator

## Synopsis

```
and-expr := and-expr & equality-expr | and-expr  
bitand equality-expr
```

The *bitwise and* operator requires integer or enumeration operands. It performs the usual arithmetic conversions, then does an *and* operation on pairs of bits in the operands, resulting in an integer.

The `bitand` keyword is interchangeable with the `&` token.

## Example

```
unsigned bitmask = 0xFFFF;  
bitmask = bitmask & ~0x7E; // bitmask becomes 0xFF81
```

## See Also

`and`, `and_eq`, `bitor`, `compl`, `expression`, `xor`, [Chapter 3](#), `<ciso646>`



## Name

bitor operator — Bitwise inclusive or operator

## Synopsis

```
inclusive-or-expr := inclusive-or-expr "|"
exclusive-or-expr |
inclusive-or-expr bitor exclusive-or-expr
```

The *bitwise or* operator requires integer or enumeration operands. It performs the usual arithmetic conversions, then does an *inclusive or* operation on pairs of bits in the operands, resulting in an integer. The bitor operator is interchangeable with the | token.

## Example

```
unsigned bitmask = 0xF0F0;
bitmask = bitmask| 0x0102; // bitmask becomes 0xF1F2.
```

## See Also

bitand, compl, expression, or, or\_eq, xor, [Chapter 3](#), <ciso646>

## Name

bool type — Boolean (logical) type specifier

## Synopsis

*simple-type-specifier* := bool

The `bool` type represents Boolean or logical values. The only valid values of the `bool` type are the literals `true` and `false`. A `bool` expression can be promoted to an integer: `false` becomes `0` and `true` becomes `1`. Arithmetic, enumerated, and pointer values can be converted to `bool`: `0` is `false`, a null pointer is `false`, and anything else is `true`.

## See Also

and, false, not, or, true, type, [Chapter 2](#), [Chapter 3](#)

## Name

break statement — Exits from a loop or switch statement

## Synopsis

```
statement := break ;
```

The break statement exits from the nearest enclosing loop (do, for, or while) or switch statement. Execution continues with the statement immediately following the end of the loop or switch. An error results if break is used outside of a loop or switch statement.

## Example

```
while(std::cin >> x) {  
    if (x < 0)    break;  
    data.push_back(x);  
}
```

## See Also

continue, do, for, statement, switch, while, [Chapter 4](#)

## Name

case keyword — Case label for switch statement

## Synopsis

*statement := case constant-expression : statement*

The case keyword labels a statement in a switch statement. A single statement can have multiple labels. You cannot use case outside of a switch statement.

Note that case labels have no effect on the order in which substatements are executed within a switch statement. Use the break statement to exit from a switch statement.

## Example

```
switch(c) {case '+':  
    z = add(x, y);  
    break;  
case '-':  
    z = subtract(x, y);  
    break;  
}
```

## See Also

break, default, statement, switch, [Chapter 4](#)

# Name

catch keyword — Exception handler in try statement

# Synopsis

```

                                handler ::= catch ( exception-declaration )
compound-statement
                                exception-declaration ::= type-specifier-seq
declarator |
                                type-specifier-seq abstract-declarator | type-specifier-seq |
. . .
```

The catch keyword introduces an exception handler in a try statement. A single try statement must have one or more catch blocks. The *exception-declaration* declares an exception handler object. If an exception is thrown in the try's compound-statement, the type of the exception object is compared with the type of each catch declaration. The compound-statement of the first catch block whose type matches that of the exception object is executed. A catch block (typically the last one in a try statement) can have an ellipsis ( . . . ) as the exception-declaration to match all exceptions.

# Example

```
int main( )
try {
    run_program( );
} catch(const std::exception& ex) {
    std::cerr << ex.what( ) << '\n';
    std::abort( );
} catch(...) {
    std::cerr << "Unknown exception. Program terminated.\n";
    std::abort( );
}
```

# See Also

declarator, statement, throw, try, type, [Chapter 4](#)

## Name

char type — Character type specifier

## Synopsis

*simple-type-specifier* := char



The char type represents a character, sometimes called a narrow character. The char type is distinct from signed char and unsigned char, but its representation must match one or the other. (Which one depends on the implementation.) A char can be promoted to an integer; the value depends on whether the representation is signed or unsigned.

By definition, `sizeof(char)` is 1, so a char is often used as the basic unit of memory allocation. Raw memory is often allocated as arrays of char or unsigned char.

## See Also

type, wchar\_t, [Chapter 1](#), [Chapter 2](#), <cstring>, <memory>

# Name

class keyword — Declares a class or template parameter, or elaborates a type name

## Synopsis

```
class-specifier ::= class-head { [member-
specification] }
class-head ::= class-key [identifier] [base-clause] |
class-key nested-name ::= identifier [base-clause] |
class-key [nested-name ::] template-id [base-clause] class-key
::= class |
struct | union
member-specification ::= member-decl [member-specification] |
access-specifier : [member-specification]
member-decl ::= [decl-specifier-seq] [member-declarator-list] ; |
function-definition [;] | qualified-id ; | using-decl |
template-decl
member-declarator-list ::= member-declarator |
member-declarator-list , member-declarator
member-declarator ::= declarator [= 0] | declarator
[= constant-expression] |
[identifier] : constant-expression
base-clause ::= : base-specifier-list
base-specifier-list ::= base-specifier | base-
specifier-list , base-specifier
base-specifier ::= [base-modifiers] [::] [nested-
name ::] class-name
base-modifiers ::= virtual [access-specifier] |
access-specifier [virtual]
access-specifier ::= private | protected | public
class-name ::= identifier | template-id
type-parm := class [identifier] [= type-id] |
template < template-parm-list > class [identifier] [= id-
expression]
elaborated-type-specifier := class-key [::] [nested-name ::]
identifier |
class-key [::] [nested-name ::] [template] template-id
```

The class keyword introduces a class declaration or definition, names a type template parameter (type-parm), or names a class type in an elaborated-type-specifier. In a class definition, the only difference between class and

struct is that the default access level is private for class and public for struct.

When declaring a type template parameter, typename and class are interchangeable.

## Example

```
template<class T> // Can use typename instead of class
class point {
public:
    point(T x = 0, T y = 0);
private:
    T x_, y_;
};
point<int> pt1;
class point<int> pt2; // redundant use of class
```

## See Also

declaration, declarator, expression, function, identifier, private, protected, public, struct, template, type, typename, union, virtual, [Chapter 6](#), [Chapter 7](#)



## Name

compl operator — Bitwise complement operator

## Synopsis

*unary-expr* := ~ *cast-expr* | compl *cast-expr*

The *bitwise complement* operator requires an integer or enumeration operand. It performs the usual arithmetic promotion and toggles each bit of its operand, resulting in an integer.

The compl keyword is interchangeable with the ~ token.

## Example

```
unsigned bitmask;  
bitmask = ~0xF107; // bitmask becomes 0xFFFF0EF8 (32 bits).
```

## See Also

bitand, bitor, expression, not, xor, [Chapter 3](#),  
<ciso646>

# Name

`const` qualifier — Marks objects and functions as constant

# Synopsis

```
cv-qualifier ::= const | volatile  
cv-qualifier-seq ::= const | volatile | const volatile | volatile  
const
```

The `const` keyword can be used as a qualifier when declaring objects, types, or member functions. When qualifying an object, using `const` means that the object cannot be the target of an assignment, and you cannot call any of its non-`const` member functions. When qualifying the target of a pointer, it means the destination cannot be modified. When member functions of a `const` object are called, this is a `const` pointer. When qualifying a member function, using the `const` qualifier means that within the member function, this is a `const` pointer to `const`, and the member function can be called for `const` objects. (Member functions without the `const` modifier cannot be called for a `const` object.)

When declaring pointers and references, be sure to distinguish between a `const` pointer (for which the pointer cannot be assigned, but what it points to can be modified) and a pointer to `const` (for which the pointer can be assigned, but what it points to cannot be modified).

# Example

```
struct rect {  
    rect(int x, int y) : x_(x), y_(y) {}  
    void x(int x) { x_ = x; }  
    void y(int y) { y_ = y; }  
    int x( ) const { return x_; }  
    int y( ) const { return y_; }  
    int area( ) const;  
private:
```

```

    int x_, y_;
};
int rect::area() const { return x() * y( ); }
const rect zero(0, 0);
const rect unit(1, 1);
rect p(4, 2), q(2, 4);
const rect* ptr1 = &zero;      // OK: pointer to const
ptr1->x(42);                    // Error: *ptr is const
ptr1 = &p;                      // Error: p is not const
ptr1 = &unit;                   // OK: unit is const
rect* const ptr2 = &p;         // OK: const pointer
ptr2 = &q;                       // Error: cannot set ptr2
ptr2->x(42);                    // OK: *ptr2 is not const

```

## See Also

const\_cast, declaration, mutable, type, volatile,  
[Chapter 2](#)

## Name

`const_cast` operator — Cast to modify qualifiers

## Synopsis

```
postfix-expr := const_cast < type-id > (  
expression )
```

The `const_cast` operator performs a type cast that can add or remove `const` and `volatile` qualifiers. No other modifications to the type are permitted.

If you cast away a `const` or `volatile` qualifier from an object, and then modify the object, the behavior is undefined.

## Example

```
template<typename T>  
T& unsafe(const T& x)  
{  
    return const_cast<T&>(x);  
}
```

## See Also

`const`, `dynamic_cast`, `expression`, `reinterpret_cast`, `static_cast`, `type`, `volatile`, [Chapter 3](#)

## Name

continue statement — Reiterates a loop statement

## Synopsis

```
statement := continue ;
```

The continue statement iterates a loop statement without executing the rest of the loop body. Control passes directly to the loop condition in a do or while loop or to the iterate expression in a for loop. You cannot use continue outside of a loop statement.

## Example

```
while (getline(cin, line)) {  
    if (line.empty( )) continue;  
    parse(line);  
    do_more_stuff( );  
};
```

## See Also

break, do, for, statement, while, [Chapter 4](#)

# Name

declaration — Function, namespace, object, type, template declaration

## Synopsis

```
translation-unit ::= [declaration-seq]
declaration-seq ::= declaration | declaration-seq
                    declaration
                    declaration ::= block-decl | function-decl |
template-decl | explicit-instantiation |
                    explicit-specialization | linkage-specification | namespace-
defn
                    block-decl ::= simple-decl | asm-defn | namespace-
alias-defn | using-decl |
                    using-directive
                    simple-decl ::= [decl-specifier-seq] [init-
declarator-list] ;
decl-specifier ::= storage-class-specifier | type-specifier |
function-specifier |
                    friend | typedef
decl-specifier-seq ::= decl-specifier | decl-specifier-seq
                    decl-specifier
                    storage-class-specifier ::= auto | register |
static | extern | mutable
function-specifier ::= inline | virtual | explicit
```

A source file is a sequence of zero or more declarations.  
See [Chapter 2](#) for a full discussion of declarations.

## See Also

asm, declarator, function, namespace, template, type, typedef, using, [Chapter 2](#)

# Name

declarator — Provides information about a single identifier in a declaration

## Synopsis

```
declarator ::= direct-declarator | ptr-operator
declarator
init-declarator-list ::= init-declarator | init-
declarator-list , init-declarator
init-declarator ::= declarator [initializer]
initializer ::= initializer-clause | ( expr-list )
initializer-clause ::= assignment-expr | { initializer-list [,] }
| { }
initializer-list ::= initializer-clause | initializer-list ,
initializer-clause
direct-declarator ::= declarator-id |
direct-declarator ( parm-decl-clause ) [cv-qualifier-seq]
[exception-specification] |
direct-declarator "[" constant-expr "]" | ( declarator )
ptr-operator ::= * [cv-qualifier-seq] | & | [::] nested-name :: *
[cv-qualifier-seq]
declarator-id ::= id-expression | [::] [nested-name ::] type-name
abstract-declarator ::= ptr-operator [abstract-
declarator] | direct-abstract-declarator
direct-abstract-declarator ::= [direct-abstract-
declarator] ( parm-decl-clause )
[ cv-qualifier-seq ] [ exception-specification ] |
[ direct-abstract-declarator "[" constant-expr "]" | (
abstract-declarator )
parm-decl-clause ::= [parm-decl-list] [...] | parm-decl-list ,
...
parm-decl-list ::= parm-decl | parm-decl-list , parm-decl
parm-decl ::= decl-specifier-seq
declarator [= assignment-expr] |
decl-specifier-seq [abstract-declarator] [= assignment-expr]
exception-specification ::= throw ( [type-id-list] )
type-id-list ::= type-id | type-id-list , type-id
```

A declarator provides additional information about a single identifier in a declaration. The declarator can be part of an object declaration, a class member declaration, a typedef

declaration, a parameter in a function declaration, and so on.

## **See Also**

declaration, expression, function, identifier, type, [Chapter 2](#)



## Name

default keyword — Catch-all label in switch statement

## Synopsis

*statement := default : statement*

A switch statement jumps to the default label if no case matches the switch expression. A switch statement can have at most one default label. If there is no default and the expression does not match any cases, control jumps directly to the statement that follows the switch statement.

## Example

```
switch(c) {  
  case '+': ... break;  
  case '-': ... break;default:  
    cerr << "unknown operator: " << c << '\n';  
    break;  
};
```

## See Also

break, case, statement, switch, [Chapter 4](#)

## Name

delete operator — Deletes a dynamic object or array of objects

## Synopsis

```
delete [delete-expr ::= [::] delete cast-expr | [::]  
       "cast-expr"]
```

The delete expression destroys dynamically-allocated objects and frees their memory. A scalar allocated with new must be freed with delete. An array allocated with new[] must be freed with delete[]. Do not mix scalar allocation or deallocation with array allocation or deallocation.

It is safe to delete a null pointer; nothing will happen.

You can overload operator delete and operator delete[] (as described in [Chapter 5](#)). Two global placement operator delete functions are provided by the standard library (see the <new> header); you can define additional functions if you wish.

The first argument to operator delete is the pointer to the memory that must be freed. Additional arguments can be used for placement delete operations, which cannot be used directly but are matched with placement new operations if the new expression throws an exception.

## Example

```
void operatordelete(void* p)  
{  
    debug(p);  
    std::free(p);  
}  
int* p = new int;  
int* array = new int[10];  
...
```

```
delete p;  
delete[] array;
```

## See Also

expression, new, [Chapter 3](#), [Chapter 5](#), <new>

## Name

do statement — Test-at-bottom loop statement

## Synopsis

```
statement := do statement while ( expression ) ;
```

The do statement is a loop that executes *statement*, then tests *expression*. The loop iterates while *expression* is true and ends if *expression* is false. The loop body always executes at least once.

## Example

```
do {  
    cout << "Number: ";  
    if (cin >> num)  
        data.push_back(num);  
} while(cin);
```

## See Also

break, continue, expression, for, statement, while,  
[Chapter 4](#)

## Name

double type — Double-precision, floating-point type specifier

## Synopsis

*simple-type-specifier* := double



The double type specifier represents a double-precision, floating-point number. The details are implementation-specific, but you are guaranteed that the range and precision of double can never be less than those of float.

The type long double has at least the range and precision of a plain double.

## See Also

float, long, type, [Chapter 2](#)

## Name

`dynamic_cast` operator — Polymorphic cast of class type objects

## Synopsis

```
postfix-expr := dynamic_cast < type-id > (  
expression )
```

The `dynamic_cast` operator performs a polymorphic cast on a pointer or reference to an object of class type. If *expression* is a pointer or reference to a base class, *type-id* can be a pointer or reference to a derived class. If the dynamic type of the object is not that of *type-id* or a class derived from *type-id*, the cast fails. If *expression* is a pointer, failure returns a null pointer; if *expression* is a reference, failure throws `std::bad_cast`. Casts from a derived to an unambiguous, accessible base class always succeed.

If *expression* is a null pointer, the result is a null pointer. If *type-id* is a pointer to `void`, the cast succeeds and returns a pointer to the most-derived object that *expression* represents.

## Example

```
struct base {};  
struct derived : base {};  
base* b = new derived;  
derived* d =dynamic_cast<derived*>(b);
```

## See Also

`const_cast`, `expression`, `reinterpret_cast`, `static_cast`, `type`, [Chapter 3](#), `<typeinfo>`

## Name

else keyword — Else part of if statement

## Synopsis

```
statement := if ( condition ) statement else  
statement
```

The else keyword introduces the else part of an if statement. If the condition is false, the else statement is executed.

## Example

```
template<typename T>  
T abs(T x)  
{  
    if (x < T( ))  
        return -x; else  
        return x;  
}
```

## See Also

if, statement, [Chapter 4](#)

## Name

`enum` keyword — Declares enumerated type or elaborates an enumerated type name

## Synopsis

```
enum-specifier ::= enum [identifier] {  
  [enumerator-list] }  
enumerator-list ::= enumerator-defn | enumerator-list ,  
enumerator-defn  
  enumerator-defn ::= enumerator | enumerator =  
  constant-expr  
  enumerator ::= identifier  
  elaborated-type-specifier := enum [::] [nested-  
name ::] identifier
```

The `enum` keyword declares a new enumerated type (as an `enum-specifier`) or names an existing enumerated type (in an `elaborated-type-specifier`). An enumerated type is an integral type that defines a set of named constants (the `enumerator-list`). Each enumerator is an identifier optionally followed by a value (an equal sign and a constant expression of integral or enumerated type). Without an explicit value, the value of an enumerator is one more than the value of the preceding enumerator. The implicit value of the first enumerator is 0.



Every enumerated type is stored as an integral type. The size of the integer is implementation-defined but is large enough to hold all the enumerated values. The valid values of an enumeration include the declared enumerators and every value that can be stored in the same number of bits, even if those values are not named as enumerators.

## Example



```
enum logical { no, maybe, yes };
bool logical_to_bool(enum logical x) // Redundant enum
{
    return x != no;
}
```

## See Also

expression, identifier, type, [Chapter 2](#)

# Name

explicit specifier — Explicit constructor specifier

# Synopsis

*function-specifier* := explicit

The explicit specifier can be used with a constructor to prevent implicit type conversions. It is permitted for any constructor but makes sense only for constructors that can be called with a single argument. An explicit constructor can be invoked from a declaration that uses function-like initialization or from an explicit type cast but not from a declaration that uses assignment-like initialization, nor from an implicit type cast.

# Example

```
struct point {explicit point(int x, int y = 0);
};
point p1(10); // OK
point p2 = 10; // Error: would be OK without explicit
point p3;
p3 = 20; // Error: would be OK without explicit
p3 = static_cast<point>(40); // OK
```

# See Also

class, declaration, function, static\_cast, type,  
[Chapter 2](#), [Chapter 6](#)

# Name

export specifier — Exported template specifier

# Synopsis

```
template-decl ::= [export] template < template-  
parm-list > declaration
```

The export specifier declares a template to be exported. An exported template lets you compile the definitions of a function template, or the member definitions of a class template, separately from the template's declaration.



Most compilers do not support export.

# Example

```
export template<typename T>  
class point {  
public:  
    point(T x, T y);  
};  
export template<typename T>  
point::point(T x, T y) {  
    ...  
}
```

# See Also

template, [Chapter 7](#)

# Name

expression — Any expression

## Synopsis

```
expression ::= assignment-expr | expression ,
assignment-expr
constant-expr ::= conditional-expr
expr-list ::= assignment-expr | expr-list ,
assignment-expr
assignment-expr ::= conditional-expr |
logical-or-expr
assignment-op
assignment-expr | throw-expr
assignment-op ::= = | *= | /= | %= | += | -= | >>=
| <<= | &= | and_eq | "|=" |
or_eq | ^= | xor_eq
throw-expr ::= throw [assignment-expr]
conditional-expr ::= logical-or-expr | logical-or-expr ?
expression : assignment-expr
logical-or-expr ::= logical-and-expr | logical-or-
expr "||" logical-and-expr |
logical-or-expr or logical-and-expr
logical-and-expr ::= inclusive-or-expr | logical-
and-expr && inclusive-or-expr |
logical-and-expr and inclusive-or-expr
inclusive-or-expr ::= exclusive-or-expr |
inclusive-or-expr "|" exclusive-or-expr |
inclusive-or-expr bitor exclusive-or-expr
exclusive-or-expr ::= and-expr | exclusive-or-expr
^ and-expr |
exclusive-or-expr xor and-expr
inclusive-or-expr ::= equality-expr | and-expr &
equality-expr |
and-expr bitand exclusive-or-expr
equality-expr ::= relational-expr | equality-expr
== relational-expr |
equality-expr != relational-expr | equality-expr not_eq
relational-expr
relational-expr ::= shift-expr | relational-expr <
shift-expr |
relational-expr > shift-expr | relational-expr <= shift-expr
|
relational-expr >= shift-expr
```

```

        shift-expr ::= add-expr | shift-expr << add-expr |
shift-expr >> add-expr
        add-expr ::= mul-expr | add-expr + mul-expr | add-
expr - mul-expr
        mul-expr ::= pm-expr | mul-expr * pm-expr | mul-
expr / pm-expr |
        mul-expr % pm-expr
        pm-expr ::= cast-expr | pm-expr .* cast-expr | pm-
expr ->* cast-expr
        cast-expr ::= unary-expr | ( type-id ) cast-expr
        unary-expr ::= postfix-expr | ++ cast-expr | --
cast-expr | unary-op
        cast-expr |
        sizeof unary-expr | sizeof ( type-id ) | new-expr | delete-
expr
        postfix-expr ::= primary-expr | postfix-expr "["
expression "]" |
        postfix-expr ( [expr-list] ) | simple-type-specifier ( [expr-
list] ) |
        typename [::] nested-name :: identifier ( [expr-list] ) |
        typename [::] nested-name :: [template] template-id ( [expr-
list] ) |
        postfix-expr . [template] id-expr | postfix-expr ->
[template] id-expr |
        postfix-expr . pseudo-dtor-name | postfix-expr -> pseudo-
dtor-name |
        postfix-expr ++ | postfix-expr -- | const_cast < type-id > (
expression ) |
        dynamic_cast < type-id > ( expression ) |
        reinterpret_cast < type-id > ( expression ) |
        static_cast < type-id > ( expression ) | typeid ( expression
) |
        typeid ( type-id )
pseudo-dtor-name ::= [::] [nested-name ::] type-name :: ~ type-
name |
        [::] nested-name :: template template-id :: ~ type-name
        primary-expr ::= literal | this | ( expression ) |
id-expr

```

An expression represents a value and a type. See [Chapter 3](#) for details about expressions, precedence, associativity, lvalues and rvalues, and more. Some syntax rules are presented here and in the various keyword sections. The `new-expr` and `delete-expr` rules are not duplicated here; see `new` and `delete` for details. The `id-expr` rules are

under identifier. See [Chapter 1](#) for information about literals.

## **See Also**

and, and\_eq, bitand, bitor, compl, const\_cast, delete, dynamic\_cast, identifier, new, not, not\_eq, or, or\_eq, reinterpret\_cast, sizeof, static\_cast, template, type, typeid, xor, xor\_eq, [Chapter 3](#)

# Name

extern storage class — External storage class specifier

# Synopsis

```
storage-class-specifier := extern  
linkage-specification ::= extern string-literal { [declaration-seq] } |  
    extern string-literal  
    declaration
```

The extern storage class specifier marks a declaration as having external linkage. An external object or function can be defined in one file and used in other files.



The extern keyword can also be used to specify language linkage: you provide the language name as a string literal. The "C++" language is the default. "C" linkage is also supported. Other languages are implementation-defined.

# Example

```
extern "C" int callable_from_C(int x, int y);  
extern void func(int);  
extern const double pi;
```

# See Also

declaration, static, type, [Chapter 2](#)

## Name

false literal — Boolean false literal

## Synopsis

```
literal := false
```

The false literal has type `bool` and integer value `0`.

## See Also

`bool`, `expression`, `true`, [Chapter 2](#)



## Name

float type — Single-precision, floating-point type specifier

## Synopsis

*simple-type-specifier* := float



The float type is a single-precision, floating-point number. The range and precision are implementation-defined. They can never be more than those of double or long double.

## See Also

double, long, type, [Chapter 2](#)

# Name

for statement — For loop statement

# Synopsis

```
statement := for ( for-init-stmt [condition] ;  
[expression] ) statement  
for-init-stmt ::= expression-stmt | simple-decl  
condition ::= expression | type-specifier-seq  
declarator = assignment-expr
```

The for loop is used for bounded loops and for unbounded loops that have well-defined iterations. Execution starts with *for-init-stmt*, which can be an expression statement or a declaration. (Note that the syntax for *expression-stmt* and *simple-decl* both include a terminating semicolon.)

*condition* is then evaluated. If *condition* evaluates to true, *statement* is executed. The iteration expression is then evaluated, and the *condition* is tested again. When *condition* is false, the loop ends and control passes to the statement following the end of the for statement.

Declarations in *for-int-stmt* and *condition* are in the same scope as *expression* and *statement*. A *continue* statement inside *statement* transfers control to the evaluation of the iteration expression.

# Example

```
for (int i = 0; i < 10; ++i)  
    cout << i << '\n';  
for (node* n = head; n != 0 ; n = n->next)  
    link_node(n);
```

# See Also

break, continue, do, expression, statement, while,  
[Chapter 4](#)

## Name

friend specifier — Grants access to private class members

## Synopsis

*decl-specifier* := friend

The friend specifier declares a friend function, class, or template. A friend declaration can appear only in a class or class template definition. A friend has full access to the private members of the class that contains the friend declaration. A friend function declaration can be a function definition.

If a class template contains a friend declaration, all instantiations of the template grant friendship. If the friend is a template declaration, all instantiations of the template are friends. If the friend is a template specialization or instantiation, only that specialization or instantiation is a friend.

Friendship is not transitive, nor is it inherited.

## Example

```
class list {friend bool operator==(const list& a, const list& b);  
    friend class node;  
    ...  
};
```

## See Also

class, declaration, template, [Chapter 6](#), [Chapter 7](#)

# Name

function — Function definition

## Synopsis

```
declaration := function-defn
function-defn ::= [decl-specifier-seq] declarator
[ctor-initializer] function-body |
  [decl-specifier-seq] declarator
function-try-block
function-body ::= compound-statement
function-try-block ::= try [ctor-initializer]
function-body
  handler-seq
  ctor-initializer ::= : member-initializer-list
  member-initializer-list ::= member-initializer |
    member-initializer-list , member-initializer
  member-initializer ::= member-initializer-id (
    [expr-list] )
  member-initializer-id ::= identifier | [::] [nested-name ::]
  class-name
```

See [Chapter 5](#) for information about functions and [Chapter 6](#) for information about the special member functions, such as constructors. (A `ctor-initializer` applies only to the definition of a class constructor.)

## See Also

class, declaration, declarator, statement, try, [Chapter 5](#)

# Name

goto statement — Transfers execution to a labeled statement

# Synopsis

```
statement := goto identifier ;
```

The goto statement transfers control to another statement in the same function. The identifier must match a statement label elsewhere in the function. Statement labels have the form `identifier : statement`.

Control cannot transfer into a try block. Transferring control into the middle of a block and across a declaration results in undefined behavior unless the declaration is for an uninitialized POD object.

# Example

```
while (getline(cin, line))
    for (size_t i = 0; i < line.size( ); ++i)
        if (line[i] == '.')goto exit; // Break out of nested loops.
exit:
...
```

# See Also

break, continue, statement, [Chapter 4](#)

# Name

identifier — Name of an entity

## Synopsis

```
id-expression ::= unqualified-id | qualified-id  
unqualified-id ::= identifier | operator-function-  
id | conversion-function-id |  
~ class-name | template-id  
qualified-id ::= [::] nested-name :: [template]  
unqualified-id | :: identifier |  
:: operator-function-id | :: template-id  
nested-name ::= class-or-namespace-name |  
nested-name [:: class-or-namespace-name] |  
nested-name [:: template class-name]  
class-or-namespace-name ::= class-name | namespace-name  
class-name ::= identifier | template-id  
namespace-name ::= identifier
```

An entity name can be a simple identifier, an operator name, or a qualified name. See [Chapter 1](#) for the rules that apply to valid identifiers. See [Chapter 5](#) for more information on overloaded operators. See [operator](#) for the rules that apply to *operator-function-id* and *conversion-function-id*.

## See Also

[class](#), [declarator](#), [namespace](#), [operator](#), [Chapter 1](#), [Chapter 2](#)

# Name

if statement — Selection statement

## Synopsis

```
statement := if ( condition ) statement |  
if ( condition ) statement else statement  
condition ::= expression | type-specifier-seq  
declarator = assignment-expr
```

The if statement tests condition, and if it is true, executes the first statement. If condition is false, and an else part is present, the else statement is executed; if there is no else part, execution continues with the statement immediately after the end of the if statement. If condition contains a declaration, the declared name is in the scope of both the if and else statements.

## Example

```
template<typename T>  
T abs(T x)  
{if (x < T( ))  
    return -x;  
    else  
    return x;  
}
```

## See Also

expression, else, statement, [Chapter 4](#)



# Name

inline specifier — Inline function specifier

# Synopsis

*function-specifier* := inline

The `inline` function specifier can be used with any function or member function to hint to the compiler that the function should be expanded inline at the point of call. The compiler is free to ignore the hint. The compiler is also free to expand functions inline that are declared without the `inline` specifier, so long as this does not affect the semantics of the program.

An inline function must be defined in the same source file where it is used, before it is used. An inline function can be defined in more than one file (unlike other functions); the definition in every file must be the same.

A member function that is defined within a class definition is implicitly declared `inline`.

# Example

```
struct point {inline point(int x, int y) : x_(x), y_(y) {} //  
Redundant  
    inline point( );  
private:  
    int x_, y_;  
};  
inline point::point( ) : x_(0), y_(0) {}
```

# See Also

class, declaration, function, type, [Chapter 5](#)

## Name

int type — Integer type specifier

## Synopsis

*simple-type-specifier* := int



The `int` type specifier denotes an integral type. Alone, it represents a signed integer whose size is the natural size of an integer on the host platform. With the `short` or `long` specifier, it represents an integer type whose range is possibly smaller or larger than that of plain `int`. With the `unsigned` specifier, it represents an unsigned type. The `signed` specifier is allowed, but it is the default, so it is rarely used with `int`.

## See Also

`char`, `declaration`, `declarator`, `long`, `short`, `signed`, `type`, `unsigned`, [Chapter 2](#)

## Name

long type — Long type specifier

## Synopsis

*simple-type-specifier* := long



When used alone or with `int`, the long type specifier represents a long integer whose range is at least as large as a plain `int`. It can be combined with `unsigned` to denote an unsigned integer whose range is at least as large as a plain `unsigned`. A plain `long` is signed by default, but you can use the `signed` specifier to be explicit.

When combined with `double`, `long` represents an extended-precision, floating-point number, whose range and precision are at least as great as those of a plain `double`.

## See Also

declaration, *declarator*, `double`, `float`, `int`, `short`, `signed`, `type`, `unsigned`, [Chapter 2](#)

# Name

mutable specifier — Mutable data member specifier

# Synopsis

*storage-class-specifier* := mutable

The mutable specifier can be used with the declaration of a data member. Using it means that the member can be modified even if the containing object is const.

# Example

```
// Represent a point in Cartesian coordinates. If the user
// requests polar
// coordinates, compute and cache the polar coordinates. If the
// Cartesian
// coordinates change, the polar coordinates must be recomputed.
// The polar
// coordinates are computed only if they are needed.
struct point {
    double angle( ) const {
        if (! has_angle_ ) {
            angle_ = tan2(y, x);
            has_angle_ = true;
        }
        return angle_;
    }
    void x(double x) { x_ = x; has_angle_ = has_arg_ = false; }
    ...
private:
    double x_, y_; mutable double angle_, arg_;
    mutable bool has_angle_, has_arg_;
};
int main( )
{
    const point pt(1, 2);
    std::cout << pt.angle( );
}
```

# See Also

class, const, declaration, type, [Chapter 6](#)

# Name

namespace keyword — Declares namespace

## Synopsis

```
declaration := namespace-defn
block-decl := namespace-alias-defn | using-
directive
namespace-defn ::= named-namespace-defn | unnamed-
namespace-defn
named-namespace-defn ::= namespace identifier {
namespace-body }
unnamed-namespace-defn ::= namespace { namespace-body }
namespace-body ::= [declaration-seq]
namespace-alias-defn ::= namespace identifier = namespace-
specifier ;
namespace-specifier ::= [::] [nested-name ::] namespace-name
namespace-name ::= identifier
using-directive ::= using namespace namespace-
specifier ;
```

The namespace keyword can be used in a namespace definition, a namespace alias definition, or a using directive. A namespace is a scope for declarations of classes, templates, functions, objects, and other namespaces. Outside a namespace, you can refer to a name that is declared in the namespace by qualifying the name with the scope operator (::), such as `ns::name`, or with a using directive or declaration.

Multiple namespace declarations can name the same namespace, each one adding more declarations to the namespace. The standard namespace, `std`, is built this way, with many different headers all placing their declarations in the `std` namespace. A namespace can be anonymous, which prevents the enclosed declarations from being visible in other source files.

A namespace alias defines an identifier as a synonym for an existing namespace. See `using` for information on the

using directive.

## Example

```
namespace math_version_2 {  
    const long double pi = 3.1415926535897932385L;  
};  
namespace math = math_version_2;  
using namespace math;
```

## See Also

class, declaration, identifier, using, [Chapter 2](#)

# Name

`new` operator — Allocates a dynamic object or array of objects

# Synopsis

```
new-expr ::= [::] new [placement] new-type-id
[new-initializer] |
  [::] new [placement] ( type-id ) [new-initializer]
placement ::= ( expr-list )
new-type-id ::= type-specifier-seq [new-declarator]
new-declarator ::= ptr-operator [new-declarator] | direct-new-declarator
direct-new-declarator ::= "[" expression "]" |
  direct-new-declarator "[" constant-expr "]"
new-initializer ::= ( [expr-list] )
ptr-operator ::= * [cv-qualifier-seq] | & | [::] nested-name :: *
[cv-qualifier-seq]
```

The `new` expression allocates memory and constructs an object. It has many forms, the simplest being a simple type name (e.g., `new int`). The `new-type-id` can be a sequence of type specifiers and qualifiers, with pointer operators, a reference operator, and an array size (e.g., `new int*[n][42]`, which allocates a two-dimensional array of pointers to `int` with `n` rows and 42 columns). The first dimension can be an integral expression; the second and subsequent dimensions must be constant expressions. If the type contains parentheses, such as function pointers, you should enclose it in parentheses to avoid ambiguity.

The `new` expression calls an allocator function to allocate the necessary memory, then initializes the memory. The `new-initializer` is an optional list of expressions in parentheses. If no `new-initializer` is present, the new object is initialized to its default value: POD objects are uninitialized, and other objects are initialized with their default constructors. If the `new-initializer` consists of



just empty parentheses, POD objects are initialized to 0, and other objects are initialized with their default constructors. The `new-initializer` can be the value of a scalar or a list of expressions to pass to a suitable constructor.

The allocator function is operator `new` or operator `new[]`, which can be overloaded (as described in [Chapter 5](#)). Two global placement operator `new` functions are provided by the standard library (see the `<new>` header); you can define additional functions if you wish.

The allocator function takes a `size_t` as its first parameter, which is the number of bytes of memory to allocate. It returns a pointer to the memory. The placement syntax is a list of expressions in parentheses. The expression list is passed to the allocator functions after the size argument. The compiler chooses which overloaded operator `new` according to the usual rules of overload resolution ([Chapter 5](#)).

## Example

```
void* operatornew(std::size_t nbytes)
{
    void* result = std::malloc(nbytes);
    debug(result);
    if (result == 0)
        throw std::bad_alloc("no more memory");
    return result;
}
int* p = new int;
int* array = new int[10];
int* placement = new(p) int;
...
delete p;
delete[] array;
```

## See Also

declarator, delete, expression, type, [Chapter 3](#), [Chapter 5](#), <new>

## Name

not operator — Logical negation operator

## Synopsis

*unary-expr := ! cast-expr | not cast-expr*

The not operator converts its operand to type `bool`, inverts its value, and returns a `bool` result. The not keyword is interchangeable with the `!` token.

## Example

```
template<typename C>
typename C::value_type checked_first(const C& c)
{
    if (not c.empty( ))
        return c[0];
    throw std::out_of_range("container is empty");
}
```

## See Also

and, `bool`, expression, or, [Chapter 3](#), <ciso646>

## Name

not\_eq operator — Inequality operator

## Synopsis

```
equality-expr := equality-expr != relational-expr  
|  
equality-expr not_eq relational-expr
```

The not\_eq operator compares two expressions for inequality. It returns true if the operands are different or false if they are the same. The not\_eq keyword is interchangeable with the != token. Note that there is no keyword equivalent for the == operator.

## See Also

expression, [Chapter 3](#), <ciso646>

# Name

operator keyword — Function call syntax for operators

# Synopsis

```
operator-function-id ::= operator op-symbol |
operator op-symbol < [template-arg-list] >
conversion-function-id ::= operator conversion-type-id
conversion-type-id ::= type-specifier-seq
[conversion-declarator]
conversion-declarator ::= ptr-operator [conversion-declarator]
ptr-operator ::= * [cv-qualifier-seq] | & | [::] nested-name :: *
[cv-qualifier-seq]
```

The operator keyword converts an operator symbol into function notation. You can use the operator keyword when invoking an operator or when overloading an operator.

# Example

```
int x =operator+(10, 32);
complex<double> c, d;
c.operator+=(d);
operator+=(c, d); // Same as above
bigint operator*(const bigint& a, const bigint* b);
```

[Table 12-1](#) lists the operator symbols (op-symbol) that can be overloaded. Alternative tokens (shown in parentheses) are interchangeable with their symbolic equivalents.

*Table 12-1. Operator symbols that can be overloaded*

delete	/	=	%=	<<=	++
delete [ ]	%	<	^= (xor_eq)	==	--
new	^ (xor)	>	&= (and_eq)	!= (not_eq)	,
new [ ]	& (bitand)	+=	= (or_eq)	<=	->*

+	(bitor)	-=	<<	>=	->
-	~ (compl)	*=	>>	&& (and)	( )
*	! (not)	/=	>>=	(or)	[ ]

## See Also

expression, identifier, template, type, [Chapter 5](#)

## Name

or operator — Logical or operator

## Synopsis

```
logical-or-expr := logical-or-expr "||" logical-  
and-expr |  
logical-or-expr or logical-and-expr
```

The *logical or* operator converts its operands to type `bool` and returns a `bool` result. This built-in operator is a short-circuit operator, so if the left operand is `true`, the expression yields `true` without evaluating the right operand. Note that if operator `or` is overloaded, it cannot be short-circuited and must evaluate both operands.

The keyword `or` is interchangeable with the `||` token.

## Example

```
int* p;  
if (p == NULLor *p == 0)  
    skip( );
```

## See Also

and, bitor, bool, expression, not, [Chapter 3](#), <ciso646>

## Name

or\_eq operator — Bitwise or assignment operator

## Synopsis

```
assignment-expr := logical-or-expr "|="
assignment-expr |
logical-or-expr or_eq assignment-expr
```

The or\_eq operator is an assignment operator that performs bitwise inclusive or. It is equivalent to logical-or-expr = logical-or-expr | assignment-expr except that logical-or-expr is evaluated only once.

The keyword or\_eq is interchangeable with the |= token.

## Example

```
unsigned bitmask = 0xF0F0;
bitmask|= 0x0102; // bitmask becomes 0xF1F2.
```

## See Also

and\_eq, bitor, expression, xor\_eq, [Chapter 3](#), <ciso646>



# Name

private access specifier — Restricts member access to the class

# Synopsis

*access-specifier* := private

The `private` keyword can be used within a class definition (followed by a colon) to mark subsequent member declarations as private or before a base-class name in the class header to mark the inheritance as private. The default access level of a class is private, both for members and base classes.

Private members can be used by the class itself or by friends.

Private inheritance means all public and protected members in the base class are private to the derived class.

# Example

```
class rational {public:
    rational(int n, int d) : num_(n), den_(d) { reduce( ); }
    rational( ) : num_(0), den_(1) {}
    int num( ) const { return num_; }
    int den( ) const { return den_; }
protected:
    void reduce( ); // Reduce num_ and den_ by gcd.
private:
    int num_, den_;
};
class derived : public rational {};
class myrational : private rational {
public:
    using rational::num;
    using rational::den;
};
```

## See Also

class, protected, public, struct, [Chapter 6](#)

## Name

protected access specifier — Restricts member access to the class and derived classes

## Synopsis

*access-specifier* := protected

The protected keyword can be used within a class definition (followed by a colon) to mark subsequent member declarations as protected or before a base-class name in the class header to mark the inheritance as protected.

Protected members can be used by the class itself, derived classes, or by friends. When accessing a protected member, the access must be through `this` or an instance of the same class or a derived class. It cannot be through a base class.

Protected inheritance means all public members of the base class become protected in the derived class.

See `private` for an example.

## See Also

class, private, public, struct, [Chapter 6](#)

## Name

public access specifier — No restrictions on member access

## Synopsis

*access-specifier* := public

The `public` keyword can be used within a class definition (followed by a colon) to mark subsequent member declarations as `public` or before a base-class name in the class header to mark the inheritance as `public`. The default access level of a `struct` is `public`, both for members and base classes.

Public members can be used freely by any other class or function.

Public inheritance means all members of the base class retain their accessibility levels.

See `private` for an example.

## See Also

`class`, `private`, `protected`, `struct`, [Chapter 6](#)

## Name

register storage class — Register storage class specifier

## Synopsis

*storage-class-specifier* := register

The register storage class is like auto: it can be used for local objects and function parameters, and using it means that the declared object has automatic lifetime. It also provides a hint to the compiler that the object will be used frequently, so the compiler can optimize access, perhaps by storing the object in a machine register.

Many modern compilers routinely ignore register because the compilers are better than humans at allocating registers.

## Example

```
int foo(register int parm)
{
    register int sqr = parm * parm;
    return sqr;
}
```

## See Also

auto, type, [Chapter 2](#)

## Name

`reinterpret_cast` operator — Cast for unsafe pointer conversions

## Synopsis

```
postfix-expr := reinterpret_cast < type-id > (  
expression )
```

The `reinterpret_cast` operator performs potentially unsafe type casts. It is most often used to cast a pointer to a different pointer type. Casting a pointer to a different pointer and back is usually safe and yields the original value. The limitations are that object pointers can be cast only to object pointers with similar or stricter alignment requirements, and function pointers can be cast only to function pointers. Pointers-to-members can be cast only to pointers-to-members. You can cast an integer to a pointer and vice versa. In all cases, the destination pointer must be large enough to hold the result. Casting a null pointer results in a null pointer of the destination type. Any other cast results in undefined behavior.

## Example

```
template<typename T>  
unsigned long addressof(const T& obj)  
{  
    return reinterpret_cast<unsigned long>(&obj);  
}
```

## See Also

`const_cast`, `dynamic_cast`, `expression`, `static_cast`,  
[Chapter 3](#)

## Name

return statement — Returns from a function

## Synopsis

```
statement := return [expression] ;
```

The return statement returns control from a function to its caller. If the function returns void, the expression is typically omitted, or else expression must be of type void. If the function returns a non-void type, the expression must be convertible to the return type.

A function that returns non-void must have a return statement (except for main, which has an implicit return 0; if control reaches the end of the function).

## See Also

expression, statement, [Chapter 5](#)

## Name

short type — Short integer type specifier

## Synopsis

*simple-type-specifier* := short

When used alone or with `int`, the short type specifier represents a short integer, whose range is no bigger than a plain `int`. It can be combined with `unsigned` to denote an unsigned integer whose range is no bigger than a plain `unsigned`. A plain short is signed by default, but you can use the signed specifier to be explicit.

## See Also

`int`, `long`, `signed`, `type`, `unsigned`, [Chapter 2](#)



## Name

signed specifier — Signed integer type specifier

## Synopsis

*simple-type-specifier* := signed

The signed keyword can be used alone to mean signed int or combined with other type specifiers to force the type to be signed. The int, short, and long types are implicitly signed, so it is most often used with signed char to implement a signed integral value whose range is the smallest that the compiler allows.



A signed integer can be represented using two's complement, ones' complement, or signed magnitude.

## See Also

char, int, long, short, type, unsigned, [Chapter 2](#)

# Name

sizeof operator — Size of type operator

# Synopsis

```
unary-expr := sizeof ( type-id ) | sizeof unary-expr
```

At compile time, the `sizeof` operator returns the amount of memory required to hold an object whose type is `type-id` or the type of `unary-expr`. In the latter case, `unary-expr` is not evaluated. The size of a type includes any padding that the compiler adds to it, so the size of an array of `N` elements is always equal to `N` times the size of a single element.

By definition, `sizeof(char)` is `1`, so you can think of the size of other types as multiples of the size of a character. The expression type is `std::size_t`.

# Example

```
class point { ... };
point* p = malloc(sizeof(point));
point corners[] = { { 2, 4 }, {4, 2}, ..., { 42, 10 } };
const unsigned count = sizeof(corners) / sizeof(corners[0]);
```

# See Also

expression, type, [Chapter 3](#), `<cstdlib>`

# Name

statement — Statement syntax

## Synopsis

```
statement ::= labeled-stmt | expr-stmt | compound-  
stmt | decl-stmt | try-block |  
if ( condition ) statement | if ( condition ) statement else  
statement |  
switch ( condition ) statement | while ( condition )  
statement |  
do statement while ( expression ) |  
for ( for-init-stmt [condition] ; [expression] ) statement |  
break ; |  
continue ; | goto identifier ; | return [expression] ;  
expr-stmt ::= [expression] ;  
compound-stmt ::= { [statement-seq] }  
statement-seq ::= statement | statement-seq  
labeled-stmt ::= identifier : statement | case  
constant-expr : statement |  
default : statement  
decl-stmt ::= block-decl  
for-init-stmt ::= expr-stmt | simple-decl  
condition ::= expression | type-specifier-seq  
declarator = assignment-expr
```

See [Chapter 4](#) for a complete discussion of statements.

## See Also

break, case, continue, declaration, declarator,  
default, do, else, expression, for, goto, if, return,  
switch, try, type, while, [Chapter 4](#)

# Name

static storage class — Static storage class specifier

# Synopsis

*storage-class-specifier* := `static`

The static storage class specifier can be used with objects, functions, and class members. Its purpose varies slightly depending on how you use it:

- For local objects, `static` affects the object's lifetime: instead of being constructed anew every time the function is called, a static local object is constructed at most once, and the object retains its value across function calls.
- A static class member is not tied to a specific object. Instead, it is similar to a global object or function, but in the scope of the class that declares it.
- For global declarations, using `static` means the object or function has internal linkage instead of the default external linkage, that is, the object or function is not visible in other source files. This use is deprecated; use anonymous namespaces instead.

# Example

```
struct demo {static const int size = 10;  
    static void func( ) {  
        return size;  
    }  
private:  
    int data_[size];  
};  
static int local_data = 10;  
static int local_func() { return demo::func( ); }
```

## See Also

auto, class, extern, namespace, register, type, [Chapter 2](#)

## Name

`static_cast` operator — Explicit cast operator

## Synopsis

```
postfix-expr := static_cast < type-id > (  
expression )
```

The `static_cast` operator performs type conversions from one static type to another. It cannot cast away `const` or `volatile` qualifiers. A static cast can invoke built-in type promotions or conversions or user-defined type conversions with type conversion operators or constructors (including explicit constructors).

## Example

```
char c;  
if (isalpha(static_cast<unsigned char>(c))  
    handle_alpha(c);
```

## See Also

`const_cast`, `dynamic_cast`, `explicit`, `expression`, `reinterpret_cast`, [Chapter 3](#)

## Name

struct keyword — Declares class with public members

## Synopsis

*class-key* := struct

The struct keyword declares a class that has public access by default. Inheritance is also public by default. See `class` for syntax rules.

## Example

```
struct point {
    int x, y;
};
struct point p1; // "struct" is redundant here.
point p2;
```

## See Also

class, union, [Chapter 6](#)

# Name

switch statement — Multibranch selection statement

# Synopsis

```
statement := switch ( condition ) statement  
condition ::= expression | type-specifier-seq  
declarator = assignment-expr
```

The switch statement evaluates *condition* and saves the value. It compares the value to each case label in *statement* (which is typically a compound statement). The statements that follow the matching case label are executed. If no case label matches, the statements following the default label are executed. If there is no default, execution continues with the statement following the end of the switch statement.

All case labels in *statement* must be unique. Use the break statement to exit the switch statement at the end of each case.

# Example

```
switch(c) {  
  case '+':  
    z = add(x, y);  
    break;  
  default:  
    z = noop( );  
    break;  
};
```

# See Also

break, case, declarator, default, expression, if, statement, type, [Chapter 4](#)



# Name

template keyword — Declares a template, specialization, or instantiation

## Synopsis

```

        declaration := template-decl | explicit-
instantiation | explicit-specialization
        template-decl ::= [export] template < template-
parm-list > declaration
        template-parm-list ::= template-parm | template-
parm-list , template-parm
        template-parm ::= type-parm | parm-decl
        type-parm ::= class [identifier] [= type-id] |
typename [identifier] [= type-id] |
        template < template-parm-list > class [identifier] [= id-
expr]
template-id ::= template-name < [template-arg-list] >
typename-name ::= identifier
        template-arg-list ::= template-arg | template-arg-
list , template-arg
        template-arg ::= assignment-expr | type-id | id-
expr
        explicit-instantiation ::= template declaration
        explicit-specialization ::= template < >
declaration
        elaborated-type-specifier := class-key [::]
[nested-name ::] [template]
        template-id | typename [::] nested-name :: [template]
template-id
        simple-type-specifier := [::] nested-name ::
template template-id
        postfix-expr := postfix-expr . [template] id-expr
|
        postfix-expr -> [template] id-expr
        pseudo-dtor-name := [::] nested-name :: template
template-id :: ~ class-name
        nested-name := nested-name [::] template class-
name]
qualified-id := [::] nested-name :: [template] unqualified-id
```

The template keyword declares a template, a specialization of a template, or an instance of a template. The

declaration can be a function declaration, function definition, class declaration, or class definition.

A template instance provides arguments for the template parameters, enclosing the arguments in angle brackets (<>). If the template is a member of a class, and the . or -> operator is used to access the member template, you must use the `template` keyword to tell the compiler to interpret the < symbol as the start of the template arguments. Similarly, use the `template` keyword in a qualified name when the name is used to instantiate the template in a context in which the compiler would interpret the < symbol as the less-than operator.

## Example

```
template<typename T>
T min(const T& a, const T& b) { return a < b ? a : b; }
typedef complex<float> cfloat;
template<> min<cfloat>(const cfloat& a, const cfloat& b)
{
    return cfloat(min(a.real(), b.real( )), min(a.imag(), b.imag(
    )));
}
template int min<int>(const int&, const int&);
```

## See Also

class, expression, identifier, type, [Chapter 7](#)

## Name

this keyword — Object pointer in member function

## Synopsis

*primary-expr* := this

The `this` keyword can be used only in nonstatic member functions. Its value is a pointer to the target object of a member function call. If the member function is qualified (with `const` or `volatile`), the same qualifiers apply to the type of `this`.

## Example

```
struct point {
    bool operator==(const point& that) {
        return this->x() == that.x() && this->y() == that.y( );
    }
    bool operator!=(const point& that) {
        return !(*this == that);
    }
    bool write(FILE* fp) {
        fwrite(static_cast<void*>(this), sizeof(*this), 1, fp);
    }
};
```

## See Also

class, expression, [Chapter 6](#)

## Name

throw operator — Throws an exception

## Synopsis

*throw-expr ::= throw [assignment-expr]*

The throw operator throws *assignment-expr* as an exception. The throw expression has type `void`. With no operand, throw rethrows the current pending exception. If no exception is pending, `terminate( )` is called.

## Example

```
template<typename C>
typename C::value_type checked_first(const C& c)
{
    if (not c.empty( ))
        return c[0]; throw std::out_of_range("container is empty");
}
```

## See Also

expression, try, [Chapter 3](#), <exception>, <stdexcept>

## Name

true literal — Logical true literal

## Synopsis

*literal* := true

The true literal has type `bool` and an integer value of 1.

## See Also

`bool`, `expression`, `false`, [Chapter 2](#)

# Name

try statement — Handles exceptions in statements

# Synopsis

```
statement := try-block  
try-block ::= try compound-statement  
handler-seq  
function-try-block ::= try [ctor-initializer]  
function-body  
handler-seq  
handler-seq ::= handler | handler-seq  
handler  
handler ::= catch ( exception-declaration )  
compound-statement  
exception-declaration ::= type-specifier-seq  
declarator |  
  type-specifier-seq abstract-declarator | type-specifier-seq |  
  . . .
```

The try statement executes *compound-statement*, and if an exception is thrown in any of the statements within that compound statement (and not caught and handled by another try statement), the catch handlers are tested to see if any of them can handle the exception. Each catch handler is tested in turn. The first one to match the exception type handles the exception. If no handler matches, the exception propagates up the call stack to the next try statement. If there is no further try statement, `terminate( )` is called.

# Example

```
int main( )try {  
    run_program( );  
} catch(const exception& ex) {  
    std::cerr << ex.what( ) << '\n';  
    abort( );  
} catch(...) {  
    std::cerr << "Unknown exception. Program terminated.\n";  
}
```

```
    abort( );  
}
```

## See Also

catch, declarator, function, throw, type, [Chapter 4](#),  
<exception> in [Chapter 13](#)

# Name

type — Type specifiers

## Synopsis

```
type-specifier ::= simple-type-specifier | class-  
specifier | enum-specifier |  
elaborated-type-specifier | cv-qualifier  
simple-type-specifier ::= [::] [nested-name ::]  
type-name |  
[::] nested-name :: template template-id | bool | char |  
double | float | int |  
long | short | signed | unsigned | void | wchar_t  
type-name ::= class-name | enum-name | typedef-name  
typedef-name ::= identifier  
elaborated-type-specifier ::= class-key [::]  
[nested-name ::] identifier |  
class-key [::] [nested-name ::] [template] template-id |  
enum [::] [nested-name ::] identifier | typename [::] nested-  
name :: identifier |  
typename [::] nested-name :: [template] template-id
```

Type specifiers are used throughout C++: in declarations, type casts, new expressions, and so on. Although the syntax rules shown here are quite flexible, the semantic rules impose many limits. For example, the simple syntax rules permit `short long` as a type (specifically, as a `decl-specifier-seq`; see *declaration*), but that combination of type specifiers is not permitted. See [Chapter 2](#) for details about semantics for valid type specifiers.

## See Also

class, const, declaration, declarator, enum, identifier, struct, template, typedef, typename, union, volatile, [Chapter 2](#)



## Name

`typedef` keyword — Declares a type synonym

## Synopsis

*decl-specifier* := `typedef`

A `typedef` declaration creates a synonym for an existing type. It does not create a new type (as `class` and `enum` do).

## Example

```
typedef unsigned int UINT;  
typedef map<string,string> dictionary;
```

## See Also

`class`, `declaration`, `enum`, `struct`, `typeid`, `union`, [Chapter 2](#)

# Name

typeid operator — Runtime type identification operator

# Synopsis

```
postfix-expr := typeid ( expression ) | typeid ( type-id )
```

The typeid operator returns a const reference to a type\_info object that describes type-id or the type of expression. If expression is an lvalue (not a pointer) of a polymorphic class, the type\_info of the most-derived class is returned. Otherwise, expression is not evaluated, and the type\_info of its static type is returned.

Each distinct type has its own associated type\_info object, but type synonyms (such as those created with typedef) have the same type\_info object.

# Example

```
template<typename T>
void debug(const T& obj)
{
    std::clog <<typeid(obj).name( ) << ':' << &obj << '\n';
}
```

# See Also

expression, typedef, [Chapter 3](#), <typeinfo>

# Name

typename keyword — Introduces a type name

## Synopsis

```
elaborated-type-specifier := typename [::] nested-name :: identifier |  
    typename [::] nested-name :: [template] template-id  
    using-decl := using [typename] [::] nested-name ::  
unqualified-id ;  
type-param := typename [identifier] [= id-expr]
```

The typename keyword is used in two different situations:

- When referring to a qualified member of a class template, the compiler cannot tell whether the name refers to a type, object, or function. Use `typename` before the qualified name to tell the compiler that it names a type.
- In a template declaration, use `typename` to name a type parameter. In this context, `class` means the same thing as `typename`.

## Example

```
template<typename C>  
typename C::value_type checked_first(const C& c)  
{  
    if (not c.empty( ))  
        return c[0];  
    throw std::out_of_range("container is empty");  
}
```

## See Also

class, expression, template, type, using, [Chapter 7](#)

## Name

unsigned specifier — Unsigned integer type specifier

## Synopsis

*simple-type-specifier* := unsigned

The unsigned type specifier can be used with any integral type to make the type unsigned. Unsigned integer arithmetic is always performed modulo  $2^n$ , in which  $n$  is the number of bits in the value representation of the integer. When used alone, unsigned means unsigned int.

## Example

```
char c;  
if (isalpha(static_cast<unsigned char>(c))  
    handle_alpha(c);
```

## See Also

char, int, long, short, signed, type, [Chapter 2](#)

## Name

union keyword — Declares a union

## Synopsis

*class-key* := union

The union keyword declares an aggregate type, similar to a struct, but the union object can store only one member at a time. The storage for all members overlaps. A union can have member functions (including constructors and destructors) but not virtual member functions. A union cannot be or have a base class. Union members cannot be static or references. Data members cannot have constructors, a destructor, copy-assignment operators, virtual functions, or virtual base classes. An initializer for a union can initialize only its first member.

See `class` for the syntax rules.

## Example

```
enum kind { integer, real, text };
struct data {
    kind data_kind;
    data(int i) : data_kind(integer), integerval(i) {}
    data(double d) : data_kind(real), realval(d) {}union {
        int integerval;
        double realval;
    };
};
```

## See Also

`class`, `struct`, [Chapter 2](#), [Chapter 6](#)

# Name

using keyword — Looks up names in alternate classes or namespaces

## Synopsis

```
block-decl := using-decl | using-directive  
using-decl ::= using [typename] [::] nested-name  
:: unqualified-id ; |  
    using [::] unqualified-id ;  
using-directive ::= using namespace [::] [nested-name ::]  
namespace-name ;
```

The using keyword starts a using declaration or using directive.

A using declaration imports a name from another namespace into the current namespace. It can also be used to introduce a name into a class scope; this is most often used to promote the access level of an inherited member or bring an inherited member into the derived class for overload resolution.

A using directive tells the compiler to search an additional namespace when looking up unqualified names.

## Example

```
namespace math {  
    const long double pi = 3.1415926535897932385L;  
};using math::pi;  
long double tan(long double x = pi);  
int main( )  
{  
    using namespace std;  
    cout << "pi=" << math::pi << '\n';  
}
```

## See Also

class, declaration, identifier, namespace, [Chapter 2](#),  
[Chapter 6](#)

## Name

virtual specifier — Polymorphic function specifier or shared base class

## Synopsis

```
function-specifier := virtual  
base-modifiers ::= virtual [access-specifier] | access-specifier  
[virtual]
```

The `virtual` keyword has two unrelated uses; it is used in virtual functions and virtual base classes:

- As a function specifier, `virtual` can be used only with a nonstatic member function. It makes the function and class polymorphic. A virtual function can be declared with `= 0` after the function header, which means the function is abstract. You cannot create an instance of a class with an abstract function; instead, a derived class must override the function. You can create an instance of the derived class.
- Using `virtual` as a base-class modifier means the base class subobject is shared when it is used more than once in an inheritance graph.

## Example

```
struct shape {virtual void draw(canvas&) = 0;  
    virtual void debug( );  
};  
struct square : virtual shape {  
    virtual void draw(canvas&);  
    virtual void debug( );  
};
```

## See Also

class, declaration, [Chapter 6](#)



## Name

void keyword — Absence of type or function arguments

## Synopsis

*simple-type-specifier* := void

The void keyword can be used as a type specifier to indicate the absence of a type or as a function's parameter list to indicate that the function takes no parameters.

When used as a type specifier, it is most often used as a function return type to indicate that the function does not return a value. It is also used as a generic pointer (e.g., void\*), although this usage is needed less often in C++ than in C.

C++ does not require that void be used to indicate that there are no function parameters, but it is often used in this way for compatibility with C.

## Example

```
void func(void)
{
    std::cout << "hello, world\n";
}
```

## See Also

declaration, type, [Chapter 2](#), [Chapter 5](#)

# Name

volatile qualifier — Volatile qualifier

# Synopsis

```
cv-qualifier ::= const | volatile  
cv-qualifier-seq ::= const | volatile | const volatile | volatile  
const
```

The volatile qualifier can be used with objects and member functions. The volatile qualifier tells the compiler to avoid certain optimizations because the object's value can change in unexpected ways. As a function qualifier, volatile tells the compiler to treat this as a volatile pointer in the member function body.

# Example

```
volatile sig_atomic_t interrupted = false;
```

# See Also

const, const\_cast, type, [Chapter 2](#), [Chapter 5](#)

## Name

wchar\_t type — Wide-character type specifier

## Synopsis

*simple-type-specifier* := wchar\_t



The wchar\_t type specifier denotes a wide character. The representation and interpretation of wide characters is implementation-defined.

## Example

```
wchar_t Hellas[] = L"\u0397\u03b5\u03bb\u03bb\u03b1\u03c2";
```

## See Also

char, type, [Chapter 1](#), [Chapter 2](#)

## Name

while statement — Test-at-top unbounded loop statement

## Synopsis

```
statement := while ( condition ) statement  
condition ::= expression | type-specifier-seq  
declarator = assignment-expr
```

The while loop tests *condition*, and if *condition* is true, while executes *statement*. This repeats until *condition* is false. If *condition* contains a declaration, the declaration is in the same scope as *statement*.

## Example

```
while (cin >> num)  
    data.push_back(num);
```

## See Also

break, continue, do, expression, for, statement, [Chapter 4](#)

## Name

xor operator — Bitwise exclusive or operator

## Synopsis

```
exclusive-or-expr := exclusive-or-expr ^ and-expr  
| exclusive-or-expr xor and-expr
```

The bitwise exclusive or operator requires integer or enumeration operands. It performs the usual arithmetic conversions, then does an exclusive or operation on pairs of bits in the operands, resulting in an integer.

The xor operator is interchangeable with the ^ token.

## Example

```
unsigned bitmask = 0xFFF0;  
bitmask = bitmask ^ 0xF12; // bitmask becomes 0xF0E2.
```

## See Also

bitand, bitor, expression, xor\_eq, [Chapter 3](#), <ciso646>

## Name

xor\_eq operator — Bitwise exclusive or assignment operator

## Synopsis

```
assignment-expr := logical-or-expr ^= assignment-  
expr |  
logical-or-expr xor_eq assignment-expr
```

The xor\_eq operator is an assignment operator that performs bitwise exclusive or. It is equivalent to `logical-or-expr = logical-or-expr ^ assignment-expr`, except that `logical-or-expr` is evaluated only once.

The keyword `xor_eq` is interchangeable with the `^=` token.

## Example

```
unsigned bitmask = 0xFFF0;  
bitmask ^= 0x0F12; // bitmask becomes 0xF0E2.
```

## See Also

`and_eq`, `expression`, `or_eq`, `xor`, [Chapter 3](#), <ciso646>

## Chapter 13. Library Reference

This chapter is a reference for the entire runtime library. As you can see, it is a big one. To help you find what you need, each header in this chapter is organized in alphabetical order. If you are not sure which header declares a particular type, macro, or other identifier, check the index. Once you find the right page, you can quickly see which header you must `#include` to define the identifier you need.

The subsections in each header's section describe the functions, macros, classes, and other entities declared and defined in the header. The name of the subsection tells you what kind of entity is described in the subsection—e.g., "terminate function," "basic\_string class template," and so on. Cross references in each "See Also" heading list intrasection references first, followed by references to other headers (in this chapter) and references to keywords (in [Chapter 12](#)).

The subsection for each class or class template contains descriptions of all important members. A few obvious or doing-nothing members are omitted (such as most destructors) for the sake of brevity.

The entire standard library resides in the `std` namespace, except that macros reside outside any namespace. Be sure to check the subsection name closely so you know whether an identifier is a macro or something else. To avoid cluttering the reference material, the `std::` prefix is omitted from the descriptions. Examples, however, are complete and show how each namespace prefix is properly used.

Some C++ headers are taken from the C standard. For example, the C standard `<stdio.h>` has its C++ equivalent

in `<cstdio>`. The C++ version declares all the C names (other than macros) in the `std::` namespace but reserves the same names in the global namespace, so you must not declare your own names that conflict with those of the C headers.

Each C header can be used with its C name, in which case the declarations in the header are explicitly introduced into the global namespace. For example, `<cstdio>` declares `std::printf` (and many other names), and `<stdio.h>` does the same, but adds `"using std::printf"` to bring the name `printf` into the global namespace. This use of the C headers is deprecated.

The syntax description for most macros shows the macro name as an object or function declaration. These descriptions tell you the macro's type or expected arguments. They do not reflect the macro's implementation. For macros that expand to values, read the textual description to learn whether the value is a compile-time constant.

For an overview of the standard library, see [Chapter 8](#). [Chapter 9](#) presents the I/O portions of the library, and [Chapter 10](#) discusses containers, iterators, and algorithms.

C++ permits two kinds of library implementations: *freestanding* and *hosted*. The traditional desktop computer is a hosted environment. A hosted implementation must implement the entire standard.



A freestanding implementation is free to implement a subset of the standard library. The subset must provide at least the following headers, and can optionally provide more:

`<cstdint>`



<cstdint>  
<cstdlib>  
<exception>  
<limits>  
<new>  
<typeinfo>

## <algorithm>

The <algorithm> header declares the generic algorithm function templates for operating on iterators and other objects. Refer to [Chapter 10](#) for more information about using and writing generic algorithms and about the iterators they use. See [Chapter 8](#) for a discussion of iterator traits.

### **Not e**

If you are at all confused by the removal algorithms (such as `pop_heap`, `remove`, and `unique`), be sure to read [Chapter 10](#) first.

This section uses a number of abbreviations and conventions. First, each algorithm is described using plain English. Then, a more mathematical description of the algorithm, which tends to be harder to read, is given in a "Technical Notes" section.

The names of the template parameters tell you which category of iterator is expected. The iterator category is the minimal functionality needed, so you can, for example, use a random access iterator where at least a forward iterator is needed. (See [Chapter 10](#) for more information on iterators and iterator categories.) To keep the syntax summaries short and readable, the iterator categories are abbreviated, as shown in [Table 13-1](#).

*Table 13-1. Template parameter names for iterator categories*

Parameter name	Iterator category
BidiIter	Bidirectional iterator
FwdIter	Forward iterator
InIter	Input iterator
OutIter	Output iterator
RandIter	Random access iterator

Other template parameter names are chosen to be self-explanatory. For example, any name that ends in `Predicate` is a function that returns a Boolean result (which can be type `bool` or any other type that is convertible to `bool`) or a Boolean functional object (an object that has a function call operator that returns a Boolean result).

A number of algorithms require sorted ranges or otherwise use a comparison function to test the less-than relationship. The library overloads each algorithm: the first function uses `operator<`, and the second accepts a function pointer or function object to perform the comparison. The comparison function takes two arguments and returns a Boolean result. In the "Technical Notes" sections, the `<` relationship signifies `operator<` or the caller-supplied function, depending on the version of the algorithm you are using. If you overload `operator<` or provide your own comparison function, make sure it correctly implements a less-than relationship. In particular, `a < a` must be false for any `a`.

In this section, the following conventions are used:

- Iterators are usually used in ranges that represent all the elements in the range or sequence. A range is written using standard mathematical notation: a square bracket denotes an inclusive endpoint of a range, and a parenthesis denotes an exclusive endpoint of a range. Thus,  $[x, y)$  is a range that starts at  $x$ , including  $x$ , and ends at  $y$ , excluding  $y$ . [Chapter 10](#) also discusses this aspect of iterators.
- Arithmetic expressions that involve iterators work as though each iterator were a random access iterator, even if the iterator is from a different category. For example,  $i - 1$  for an input iterator is not allowed in C++ code, but in this section, it means the input iterator that points to one position before  $i$ .
- The names used for input ranges are typically `first` and `last`, in which `first` is an iterator that points to first the element of the range, and `last` is an iterator that points to one past the end of the range. Thus, the range is written as  $[first, last)$ .
- An iterator that advances from `first` to `last` is typically called `iter`.
- Output iterators are typically named `result`. Most algorithms specify only the start of the output range. It is your responsibility to ensure that the output range has room to accommodate the entire output. The behavior is undefined if the output range overflows.

In each "Technical Notes" section, conventional mathematical notation is used with some aspects of C++ notation, such as `*`, which dereferences an iterator. Also, a single equal sign (`=`) means assignment, and a double equal sign (`==`) means comparison for equality. The following conventions are used for names:

`i`, `j`, `k`

Denote iterators, and  $*i$ ,  $*j$ , and  $*k$  denote the values that the iterators point to.

$n$ ,  $m$

Denote integers.

$a$ ,  $b$ ,  $c$

Denote values, which are usually of types that can be assigned to or from a dereferenced iterator (e.g.,  $*i = a$ ).

## Name

`adjacent_find` function template — Searches for a pair of adjacent, equal items

## Synopsis

```
template<typename FwdIter>
    FwdIter adjacent_find(FwdIter first, FwdIter last);
template<typename FwdIter, typename BinaryPredicate>
    FwdIter adjacent_find(FwdIter first, FwdIter last,
        BinaryPredicate pred);
```

The `adjacent_find` function template looks for the first set of adjacent items in the range `[first, last)` that are equal (first version) or in which `pred(*iter, *(iter+1))` is true (second version). Items are "adjacent" when their iterators differ by one position.

The return value is an iterator that points to the first of the adjacent items, or `last` if no matching items are found. See [Figure 13-1](#) for an example.



*Figure 13-1. Using `adjacent_find` to find two adjacent, equivalent items*

## Technical Notes

The `adjacent_find` function template returns `i`, in which  $i = \text{first} + n$ , and  $n$  is the smallest value such that  $*(\text{first} + n) == *(\text{first} + n + 1)$  and  $\text{first} + n + 1 < \text{last}$ , or, if there is no such  $n$ ,  $i = \text{last}$ .

Complexity is linear: the standard is muddled, but any reasonable implementation calls the predicate (operator== or `pred`) exactly  $n + 1$  times.

## See Also

[find function template](#), [find\\_if function template](#)

## Name

`binary_search` function template — Searches using a binary search

## Synopsis

```
template<typename FwdIter, typename T>
    bool binary_search(FwdIter first, FwdIter last, const T& value);
template<typename FwdIter, typename T, typename Compare>
    bool binary_search(FwdIter first, FwdIter last, const T& value,
                       Compare comp);
```

The `binary_search` function template uses a binary search to test whether `value` is in the range `[first, last)`. It returns `true` upon success and `false` if the value is not found. The contents of the range must be sorted in ascending order.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether `X < Y`.

## Technical Notes

Precondition: `!(*(i + 1) < *i)` for all `i` in `[first, last - 1)`.

The `binary_search` function template returns `true` if there is an `i` in `[first, last)` such that `!*i < value` and `!(value < *i)`. It returns `false` if there is no such `i`.

Complexity is logarithmic. The number of comparisons is at most  $\log(\text{last} - \text{first}) + 2$ . Although the iterator can be a forward iterator, the best performance is obtained with a random access iterator. With a forward or bidirectional iterator, the iterator is advanced a linear number of times, even though the number of comparisons is logarithmic.

## See Also

[equal\\_range function template](#), [find function template](#),  
[find\\_if function template](#), [lower\\_bound function template](#),  
[upper\\_bound function template](#)



## Name

copy function template — Copies every item in a range

## Synopsis

```
template<typename InIter, typename OutIter>
    OutIter copy(InIter first, InIter last, OutIter result);
```

The copy function template copies items from `[first, last)` to the output iterator starting at `result`. You must ensure that the output sequence has enough room for `last - first` items. The return value is the value of the `result` iterator after copying all the items, as shown in [Figure 13-2](#).

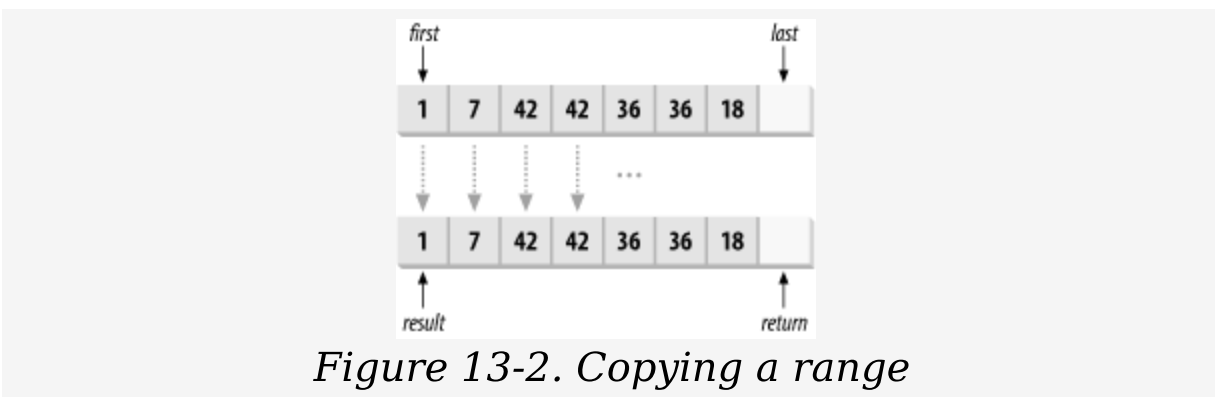


Figure 13-2. Copying a range

The `result` iterator cannot be in the source range `[first, last)`, but other parts of the destination range can overlap with the source.

See [Example 13-2](#) (under `generate`).

## Technical Notes

The copy function template assigns  $*(result + n) = *(first + n)$  for all  $n$  in the range  $[0, last - first)$ .

Complexity is linear: exactly `last - first` assignments are performed.

## See Also

[copy\\_backward function template](#), [partial\\_sort\\_copy function template](#), [replace\\_copy function template](#), [remove\\_copy function template](#), [reverse\\_copy function template](#), [rotate\\_copy function template](#), [unique\\_copy function template](#)

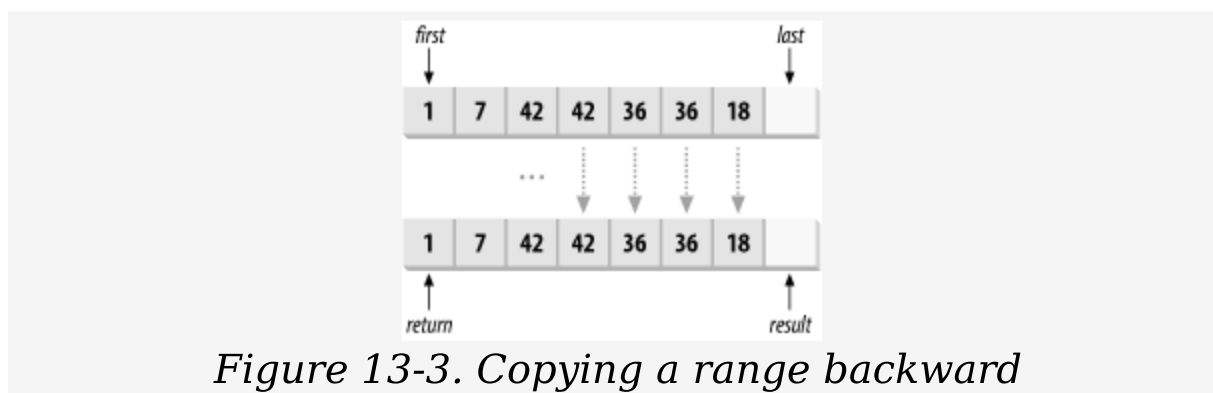
## Name

copy\_backward function template — Copies a range, starting at the end

## Synopsis

```
template<typename BidIter1, typename BidIter2>
    BidIter2 copy_backward(BidIter1 first, BidIter1 last,
                          BidIter2 result);
```

The `copy_backward` function template does the same thing as `copy`, but it works backward, starting at the element before `last` and copying elements toward `first`. The `result` iterator must point to one past the end of the destination and is decremented before copying each element. The return value is an iterator that points to the first element of the destination, as shown in [Figure 13-3](#).



*Figure 13-3. Copying a range backward*

The `result` iterator cannot be in the source range [`first`, `last`), but other parts of the destination range can overlap with the source.

## Technical Notes

The `copy_backward` function template assigns  $*(result - n) = *(last - n)$  for all  $n$  in the range  $[1, last - first]$ .

Complexity is linear: exactly last - first assignments are performed.

## **See Also**

[copy\\_function template](#), [reverse\\_copy\\_function template](#)

## Name

count function template — Counts the occurrences of a value

## Synopsis

```
template<typename InIter, typename T>
    typename iterator_traits<InIter>::difference_type
        count(InIter first, InIter last, const T& value);
```

The count function template returns the number of elements in the range [first, last) that are equal to value.

## Technical Notes

Complexity is linear: exactly last - first comparisons are performed.

## See Also

[count\\_if function template](#), [equal\\_range function template](#)

## Name

`count_if` function template — Counts the number of times a predicate returns true

## Synopsis

```
template<typename InIter, typename Predicate>
    typename iterator_traits<InIter>::difference_type
    count_if(InIter first, InIter last, Predicate pred);
```

The `count_if` function template returns the number of elements in the range `[first, last)` for which `pred(*iter)` returns true.

## Technical Notes

Complexity is linear: `pred` is called exactly `last - first` times.

## See Also

[count function template](#), [find\\_if function template](#)

## Name

equal function template — Tests whether ranges have same contents

## Synopsis

```
template<typename InIter1, typename InIter2>
    bool equal(InIter1 first1, InIter1 last1, InIter2 first2);
template<typename InIter1, typename InIter2, typename
BinaryPredicate>
    bool equal(InIter1 first1, InIter1 last1, InIter2 first2,
                BinaryPredicate pred);
```

The `equal` function template returns `true` if two ranges contain the same elements in the same order. The first range is `[first1, last1)`, and the second range has the same number of elements, starting at `first2`. The ranges can overlap.

The first form compares items using the `==` operator. The second form calls `pred(*iter1, *iter2)`.

## Technical Notes

The `equal` function template returns `true` if `*(first1 + n) == *(first2 + n)` for all `n` in the range `[0, last1 - first1)`.

Complexity is linear: at most `last1 - first1` comparisons are performed.

## See Also

[lexicographical\\_compare function template](#), [mismatch function template](#), [search function template](#)

## Name

`equal_range` function template — Finds all occurrences of a value in a sorted range using binary search

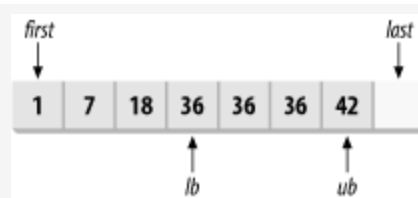
## Synopsis

```
template<typename FwdIter, typename T>
    pair<FwdIter, FwdIter>
        equal_range(FwdIter first, FwdIter last, const T& value);
template<typename FwdIter, typename T, typename Compare>
    pair<FwdIter, FwdIter>
        equal_range(FwdIter first, FwdIter last, const T& value,
                    Compare comp);
```

The `equal_range` function template determines where `value` belongs in the sorted range `[first, last)`. It returns a pair of iterators that specify the start and one past the end of the range of items that are equivalent to `value`, or both iterators in the pair point to where you can insert `value` and preserve the sorted nature of the range.

The first form compares values using the `<` operator. The second form calls `comp(*iter, value)`.

[Figure 13-4](#) shows how bounds are found with the value 36. The result of calling `equal_range` is `pair(lb, ub)`. Note that for values in the range `[19, 35]`, the upper and lower bound are both equal to `lb`, and for values in the range `[37, 41]`, the upper and lower bound are both equal to `ub`.



*Figure 13-4. Finding the limits of where the value 36 belongs in a sorted range*

## Technical Notes



Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first, last - 1)$ .

The `equal_range` function template returns the equivalent of calling the following, although the actual implementation might be different:

```
std::make_pair(std::lower_bound(first, last, value),
               std::upper_bound(first, last, value))
```

or:

```
std::make_pair(std::lower_bound(first, last, value, comp),
               std::upper_bound(first, last, value, comp))
```

Complexity is logarithmic. The number of comparisons is at most  $2 \times \log(last - first) + 1$ . Although the iterator can be a forward iterator, the best performance is obtained with a random access iterator. With a forward or bidirectional iterator, the iterator is advanced a linear number of times, even though the number of comparisons is logarithmic.

## See Also

[binary\\_search function template](#), [lower\\_bound function template](#), [upper\\_bound function template](#), `pair` in [<utility>](#)

## Name

fill function template — Fills a range with copies of a value

## Synopsis

```
template<typename FwdIter, typename T>  
    void fill(FwdIter first, FwdIter last, const T& value);
```

The fill function template fills the destination range [first, last) by assigning value to each item in the range.

## Technical Notes

The fill function template assigns `*i = value` for all `i` in the range [first, last).

Complexity is linear: exactly last - first assignments are performed.

## See Also

[fill\\_n function template](#), [generate function template](#)

## Name

`fill_n` function template — Fills a counted range with copies of a value

## Synopsis

```
template<typename OutIter, typename Size, typename T>  
    void fill_n(OutIter first, Size n, const T& value);
```

The `fill_n` function template assigns `value` to successive items in the destination range, starting at `first` and assigning exactly `n` items.

The `Size` template parameter must be convertible to an integral type.

## Technical Notes

The `fill_n` function template assigns `*(first + n) = value` for all `n` in the range `[0, n)`.

Complexity is linear: exactly `n` assignments are performed.

## See Also

[fill function template](#), [generate\\_n function template](#)

## Name

find function template — Searches for a value using linear search

## Synopsis

```
template<typename InIter, typename T>  
InIter find(InIter first, InIter last, const T& value);
```

The find function template returns an iterator that points to the first occurrence of value in [first, last). It returns last if value is not found. The == operator is used to compare items.

## Technical Notes

The find function template returns  $i = first + n$ , in which  $n$  is the smallest value such that  $*(first + n) == value$ . If there is no such  $n$ ,  $i = last$ .

Complexity is linear: at most  $last - first$  comparisons are performed.

## See Also

[find\\_end function template](#), [find\\_first\\_of function template](#), [find\\_if function template](#), [search function template](#)

## Name

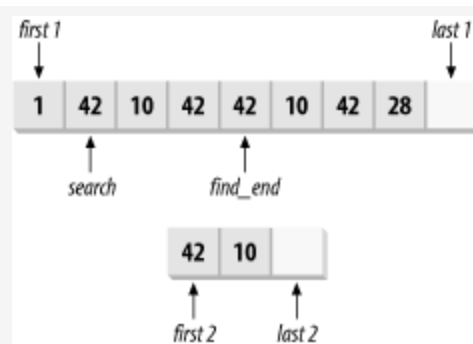
`find_end` function template — Searches for the last occurrence of a sequence

## Synopsis

```
template<typename FwdIter1, typename FwdIter2>
    FwdIter1 find_end(FwdIter1 first1, FwdIter1 last1, FwdIter2
first2,
                    FwdIter2 last2);
template<typename FwdIter1, typename FwdIter2, typename
BinaryPredicate>
    FwdIter1 find_end(FwdIter1 first1, FwdIter1 last1, FwdIter2
first2,
                    FwdIter2 last2, BinaryPredicate pred);
```

The `find_end` function template finds the last (rightmost) subsequence `[first2, last2)` within the range `[first1, last1)`, as illustrated in [Figure 13-5](#). It returns an iterator, `find_end` in [Figure 13-5](#), that points to the start of the matching subsequence or `last1` if a match cannot be found.

The first form compares items with the `==` operator. The second form calls `pred(*iter1, *iter2)`.



*Figure 13-5. Finding a subsequence with `find_end` and `search`*

## Technical Notes

Let  $\text{length1} = \text{last1} - \text{first1}$  and  $\text{length2} = \text{last2} - \text{first2}$ .

The `find_end` function template returns  $\text{first1} + n$ , in which  $n$  is the highest value in the range  $[0, \text{length1} - \text{length2})$  such that  $*(i + n + m) == (\text{first2} + m)$  for all  $i$  in the range  $[\text{first1}, \text{last1})$  and  $m$  in the range  $[0, \text{length2})$ . It returns  $\text{last1}$  if no such  $n$  can be found.

Complexity: at most  $\text{length1} \times \text{length2}$  comparisons are performed.

## See Also

[find function template](#), [search function template](#)

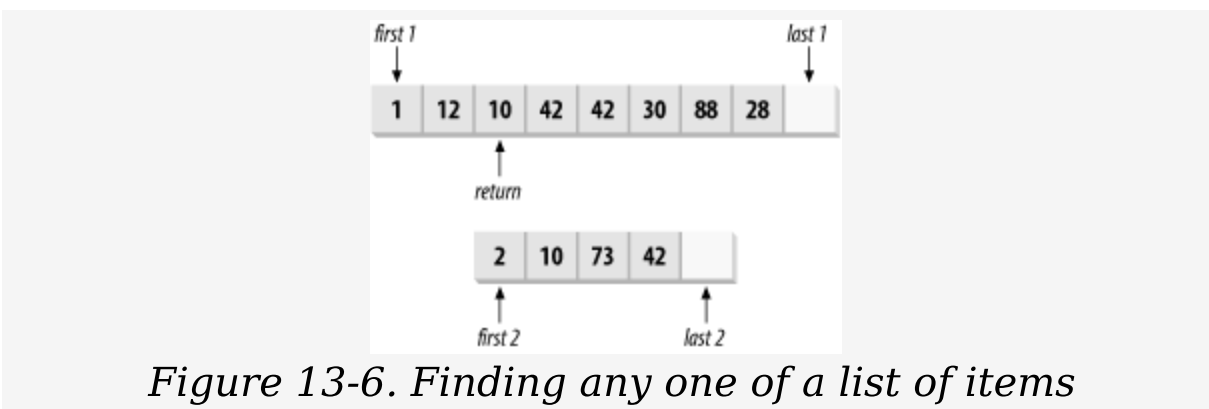
## Name

`find_first_of` function template — Searches for any one of a sequence of values

## Synopsis

```
template<typename FwdIter1, typename FwdIter2>
    FwdIter1 find_first_of(FwdIter1 first1, FwdIter1 last1,
                            FwdIter2 first2,
                            FwdIter2 last2);
template<typename FwdIter1, typename FwdIter2, typename
    BinaryPredicate>
    FwdIter1 find_first_of(FwdIter1 first1, FwdIter1 last1,
                            FwdIter2 first2,
                            FwdIter2 last2, BinaryPredicate pred);
```

The `find_first_of` function template searches the range `[first1, last1)` for any one of the items in `[first2, last2)`. If it finds a matching item, it returns an iterator that points to the matching item, in the range `[first1, last1)`. It returns `last1` if no matching item is found. [Figure 13-6](#) shows an example.



*Figure 13-6. Finding any one of a list of items*

The first form compares items with the `==` operator. The second form calls `pred(*iter1, *iter2)`.

## Technical Notes

Let  $\text{length1} = \text{last1} - \text{first1}$  and  $\text{length2} = \text{last2} - \text{first2}$ .

The `find_first_of` function template returns  $\text{first1} + n$  where  $n$  is the smallest value in the range  $[0, \text{length1})$  such that  $*(\text{first1} + n) == (\text{first2} + m)$  for some  $m$  in the range  $[0, \text{length2})$ . It returns `last1` if no such  $n$  and  $m$  can be found.

Complexity: at most  $\text{length1} \times \text{length2}$  comparisons are performed.

## See Also

[find function template](#)



## Name

`find_if` function template — Searches for when a predicate returns true

## Synopsis

```
template<typename InIter, typename Predicate>  
    InIter find_if(InIter first, InIter last, Predicate pred);
```

The `find_if` function template (similar to `find`) searches the range `[first, last)` for the first item for which `pred(*iter)` is true. It returns an iterator that points to the matching item. If no matching item is found, `last` is returned.

## Technical Notes

The `find_if` function template returns `i = first + n`, in which `n` is the smallest value such that `pred(*(first + n), value)` is true. If there is no such `n`, `i = last`.

Complexity is linear: `pred` is called at most `last - first` times.

## See Also

[find function template](#)

## Name

`for_each` function template — Calls a function for each item in a range

## Synopsis

```
template<typename InIter, typename Func>
    Function for_each(InIter first, InIter last, Func f);
```

The `for_each` function template calls `f` for each item in the range `[first, last)`, passing the item as the sole argument to `f`. It returns `f`.

## Example

[Example 13-1](#) shows how the use `for_each` to test whether a sequence is sorted. The `is_sorted` object remembers the previous item in the sequence, which it compares with the current item. The overloaded `bool` operator returns `true` if the sequence is sorted so far or `false` if the sequence is out of order. The example takes advantage of the fact that `for_each` returns the `f` parameter as its result.

*Example 13-1. Using `for_each` to test whether a list is sorted*

```
#include <iostream>
#include <algorithm>
#include <list>

template<typename T>
class is_sorted
{
public:
    is_sorted( ) : first_time(true), sorted(true) {}
    void operator( )(const T& item) {
        // for_each calls operator( ) for each item.
        if (first_time)
            first_time = false;
        else if (item < prev_item)
            sorted = false;
    }
};
```

```

    prev_item = item;
}
operator bool( ) { return sorted; }
private:
    bool first_time;
    bool sorted;
    T prev_item;
};

int main( )
{
    std::list<int> l;
    l.push_back(10);
    l.push_back(20);
    ... if (std::for_each(l.begin( ), l.end( ), is_sorted<int>(
)))
        std::cout << "list is sorted" << '\n';
}

```

## Technical Notes

Complexity is linear: f is called exactly last - first times.

## See Also

[copy function template](#), accumulate in [<numeric>](#)

## Name

generate function template — Fills a range with values returned from a function

## Synopsis

```
template<typename FwdIter, typename Generator>
    void generate(FwdIter first, FwdIter last, Generator gen);
```

The generate function template fills the sequence [first, last) by assigning the result of calling gen( ) repeatedly.

## Example

[Example 13-2](#) shows a simple way to fill a sequence with successive integers.

*Example 13-2. Using generate to fill a vector with integers*

```
#include <algorithm>
#include <iostream>
#include <iterator>
#include <vector>

// Generate a series of objects, starting with "start".
template <typename T>
class series {
public:
    series(const T& start) : next(start) {}
    T operator( )( ) { return next++; }
private:
    T next;
};

int main( )
{
    std::vector<int> v;
    v.resize(10);
    // Generate integers from 1 to 10. std::generate(v.begin( ),
v.end( ), series<int>(1));
    // Print the integers, one per line.
    std::copy(v.begin( ), v.end( ),
```

```
        std::ostream_iterator<int>(std::cout, "\\n");  
    }
```

## Technical Notes

Complexity is linear: gen is called exactly last - first times.

## See Also

[fill function template](#), [generate\\_n function template](#)

## Name

`generate_n` function template — Fills a counted range with values returned from a function

## Synopsis

```
template<typename OutIter, typename Size, typename Generator>
void generate_n(OutIter first, Size n, Generator gen);
```

The `generate_n` function template calls `gen( )` exactly `n` times, assigning the results to fill the sequence that starts at `first`. You must ensure that the sequence has room for at least `n` items. The `Size` type must be convertible to an integral type.

## Example

[Example 13-3](#) shows a simple way to print a sequence of integers.

*Example 13-3. Using `generate_n` to print a series of integers*

```
#include <algorithm>
#include <iostream>
#include <iterator>

// Use the same series template from Example 13-2.

int main( )
{
    // Print integers from 1 to 10.
    std::generate_n(std::ostream_iterator<int>(std::cout, "\n"),
                    10, series<int>(1));
}
```

## Technical Notes

Complexity is linear: `gen` is called exactly `n` times.

## See Also

[fill\\_n function template](#), [generate function template](#)

## Name

includes function template — Tests sorted ranges for subset

## Synopsis

```
template<typename InIter1, typename InIter2>
    bool includes(InIter1 first1, InIter1 last1, InIter2 first2,
                  InIter2 last2);
template<typename InIter1, typename InIter2, typename Compare>
    bool includes(InIter1 first1, InIter1 last1, InIter2 first2,
                  InIter2 last2,
                  Compare comp);
```

The `includes` function template checks for a subset relationship, that is, it returns `true` if every element in the sorted sequence `[first2, last2)` is contained in the sorted sequence `[first1, last1)`. It returns `false` otherwise.

Both sequences must be sorted. The first form uses the `<` operator to compare the elements. The second form calls `comp(*iter1, *iter2)`.

## Technical Notes

Precondition: `!(*(i + 1) < *i)` for all `i` in `[first1, last1 - 1)` and `!(*(j + 1) < *j)` for all `j` in `[first2, last2 - 1)`.

The `includes` function template returns `true` if there is an `i` in `[first1, last1)` such that `*(i + n) = *(first2 + n)` for all `n` in `[0, (last2 - first2))`. It returns `last1` if there is no such `i`.

Complexity is linear: at most,  $2 \times ((last1 - first1) + (last2 - first2)) - 1$  comparisons are performed.

## See Also

[set\\_difference function template](#), [set\\_intersection function template](#), [set\\_symmetric\\_difference function template](#),



[set\\_union function template](#)

## Name

`inplace_merge` function template — Merges sorted, adjacent ranges in place

## Synopsis

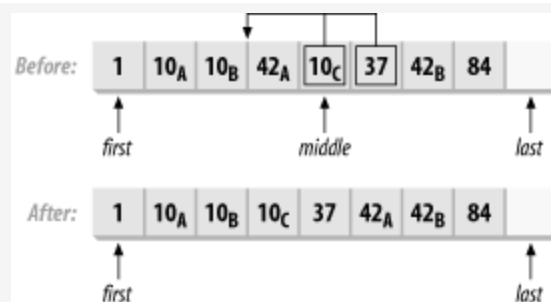
```
template<typename BidIter>
    void inplace_merge(BidIter first, BidIter middle, BidIter
last);
template<typename BidIter, typename Compare>
    void inplace_merge(BidIter first, BidIter middle, BidIter
last,
                        Compare comp);
```

The `inplace_merge` function template merges two sorted, consecutive ranges in place, creating a single sorted range. The two ranges are `[first, middle)` and `[middle, last)`. The resulting range is `[first, last)`.

The merge is stable, so elements retain their respective orders, with equivalent elements in the first range coming before elements in the second.

Both sequences must be sorted. The first form uses the `<` operator to compare elements. The second form calls `comp(*iter1, *iter2)`.

[Figure 13-7](#) shows how `inplace_merge` operates.



*Figure 13-7. Merging two sorted ranges*

## Technical Notes

Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first, middle - 1)$  and  $!(*(j + 1) < *j)$  for all  $j$  in  $[middle, last - 1)$ .

Postcondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first, last - 1)$ .

Complexity is usually linear with  $n + 1$  comparisons, but if enough temporary memory is not available, the complexity might be  $n \log n$ , in which  $n$  is  $last - first$ .

## See Also

[merge function template](#), [sort function template](#)

## Name

`iter_swap` function template — Swaps values that iterators point to

## Synopsis

```
template<typename FwdIter1, typename FwdIter2>  
    void iter_swap(FwdIter1 a, FwdIter2 b);
```

The `iter_swap` function template swaps the values pointed to by `a` and `b`. You can think of its functionality as:

```
FwdIter1::value_type tmp = *b*;  
*b = *a;  
*a = tmp;
```

## Technical Notes

Complexity is constant.

## See Also

[copy function template](#), [swap function template](#),  
[swap\\_ranges function template](#)

## Name

lexicographical\_compare function template — Compares ranges for less-than

## Synopsis

```
template<typename InIter1, typename InIter2>
    bool lexicographical_compare(InIter1 first1, InIter1 last1,
                                InIter2 first2,
                                InIter2 last2);
template<typename InIter1, typename InIter2, typename Compare>
    bool lexicographical_compare(InIter1 first1, InIter1 last1,
                                InIter2 first2,
                                InIter2 last2, Compare comp);
```

The `lexicographical_compare` function template returns true if the sequence `[first1, last1)` is less than the sequence `[first2, last2)`. If the sequences have the same length and contents, the return value is false. If the second sequence is a prefix of the first, true is returned. (The use of "lexicographical" emphasizes that the ranges are compared element-wise, like letters in words.)

The first form uses the `<` operator to compare elements. The second form calls `comp(*iter1, *iter2)`.

## Technical Notes

Let  $length1 = last1 - first1$ ,  $length2 = last2 - first2$ , and  $minlength = \min(length1, length2)$ .

The `lexicographical_compare` function template returns true if either of the following conditions is true:

- There is an  $n$  in  $[0, minlength)$  such that  $*(first1 + m) == *(first2 + m)$  for all  $m$  in  $[0, n - 1)$ , and  $*(first1 + n) < *(first2 + n)$ .

- $*(first1 + n) == *(first2 + n)$  for all  $n$  in  $[0, length2)$  and  $length2 < length1$ .

Complexity is linear: at most,  $\min(length)$  comparisons are performed.

## Example

```
#include <algorithm>
#include <iostream>
#include <ostream>

int main( )
{
    using namespace std;

    int a[] = { 1, 10, 3, 42 };
    int b[] = { 1, 10, 42, 3 };
    int c[] = { 1, 10 };

    cout << boolalpha;
    cout << lexicographical_compare(a, a+4, b, b+4); // true
    cout << lexicographical_compare(a, a+4, c, c+2); // false
    cout << lexicographical_compare(a, a+4, a, a+4); // false
    cout << lexicographical_compare(c, c+2, b, b+4); // true
}
```

## See Also

[equal function template](#)

## Name

`lower_bound` function template — Finds lower bound for a value's position in a sorted range using binary search

## Synopsis

```
template<typename FwdIter, typename T>
    FwdIter lower_bound(FwdIter first, FwdIter last, const T&
value);
template<typename FwdIter, typename T, typename Compare>
    FwdIter lower_bound(FwdIter first, FwdIter last, const T&
value, Compare comp);
```

The `lower_bound` function template determines where `value` belongs in the sorted range `[first, last)`. The return value is an iterator that points to the first (leftmost) occurrence of `value` in the range, if `value` is present. Otherwise, the iterator points to the first position where you can insert `value` and preserve the sorted nature of the range.

The first form compares values using the `<` operator. The second form calls `comp(*iter, value)`.

[Figure 13-4](#) (under `equal_range`) shows how to find the bounds for the value 36. The `lower_bound` function returns `lb` as the lower bound of 36 in the given range. Note that `lb` is the lower bound for all values in the range `[19, 36]`, and for values in the range `[37, 41]`, the lower bound is equal to `ub`.

## Technical Notes

Precondition:  $!(*(i + 1) < *i)$  for all `i` in `[first, last - 1)`.

The `lower_bound` function template returns `first + n`, in which `n` is the highest value in `[0, last - first)` such that  $*(first + m) < value$  for all `m` in `[0, n)`.

Complexity is logarithmic. The number of comparisons is at most  $\log(\text{last} - \text{first}) + 1$ . Although the iterator can be a forward iterator, the best performance is obtained with a random access iterator. With a forward or bidirectional iterator, the iterator is advanced a linear number of times, even though the number of comparisons is logarithmic.

## See Also

[binary\\_search function template](#), [equal\\_range function template](#), [upper\\_bound function template](#)



## Name

`make_heap` function template — Reorders a range to convert it into a heap

## Synopsis

```
template<typename RandIter>
    void make_heap(RandIter first, RandIter last);
template<typename RandIter, typename Compare>
    void make_heap(RandIter first, RandIter last, Compare comp);
```

The `make_heap` function template reorders the elements in the range `[first, last)` to form a heap in place.

The first form compares values using the `<` operator. The second form calls `comp(*iter, value)`.

### Heap Data Structure

A *heap* is a data structure that is ideally suited for implementing priority queues. C++ defines a range as a heap if two properties are satisfied:

- The first element of the range is the largest, which strictly means that `*first < *(first + n)` is false for all `n` in `[1, last - first)`.
- Adding or removing an element can be done in logarithmic time, and the result is still a heap.

The classic example of a heap is a binary tree, in which each node is greater than or equal to its children, but the relative order of the children is not specified. Thus, the root of the tree is the largest element in the entire tree.

## Technical Notes

Postcondition: `[first, last)` is a heap.

Complexity is linear: at most,  $3 \times (\text{last} - \text{first})$  comparisons are performed.

## See Also

[pop\\_heap function template](#), [push\\_heap function template](#),  
[sort\\_heap function template](#), [<queue>](#)

## Name

max function template — Returns the maximum of two values

## Synopsis

```
template<typename T>  
    const T& max(const T& a, const T& b);  
template<typename T, typename Compare>  
    const T& max(const T& a, const T& b, Compare comp);
```

The max function template returns the larger of a and b. If neither is larger, it returns a.

The first form compares values using the < operator. The second form calls comp(a, b).

## See Also

[max\\_element function template](#), [min function template](#)

## Name

`max_element` function template — Finds the largest element in a range

## Synopsis

```
template<typename FwdIter>
    FwdIter max_element(FwdIter first, FwdIter last);
template<typename FwdIter, typename Compare>
    FwdIter max_element(FwdIter first, FwdIter last, Compare comp);
```

The `max_element` function template returns an iterator that points to the largest element in the range `[first, last)`. If there are multiple instances of the largest element, the iterator points to the first such instance.

The first form compares values using the `<` operator. The second form calls `comp(*iter1, *iter2)`.

## Technical Notes

The `max_element` function template returns `first + n`, in which `n` is the smallest value in `[0, last - first)` such that for all `m` in `[0, last - first)`, `*(first + n) < *(first + m)` is false.

Complexity is linear: exactly `max(last - first - 1, 0)` comparisons are performed.

## See Also

[max function template](#), [min\\_element function template](#)

## Name

merge function template — Merges sorted ranges

## Synopsis

```
template<typename InIter1, typename InIter2, typename OutIter>
    OutIter merge(InIter1 first1, InIter1 last1, InIter2 first2,
                  InIter2 last2,
                  OutIter result);
template<typename InIter1, typename InIter2, typename OutIter,
        typename Compare>
    OutIter merge(InIter1 first1, InIter1 last1, InIter2 first2,
                  InIter2 last2,
                  OutIter result, Compare comp);
```

The merge function template merges the sorted ranges [first1, last1) and [first2, last2), copying the results into the sequence that starts with result. You must ensure that the destination has enough room for the entire merged sequence.

The return value is the end value of the destination iterator, that is, result + (last1 - first1) + (last2 - first2).

The destination range must not overlap either of the input ranges.

The merge is stable, so elements preserve their relative order. Equivalent elements are copied, so elements from the first range come before elements from the second range.

The first form compares values using the < operator. The second form calls comp(\*iter1, \*iter2).

## Technical Notes

Let length1 = last1 - first1 and length2 = last2 - first2.

Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first1, last1 - 1)$  and  $!(*(j + 1) < *j)$  for all  $j$  in  $[first2, last2 - 1)$ .

Postcondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[result, result + length1 + length2 - 1)$ .

Complexity is linear: at most,  $length1 + length2 - 1$  comparisons are performed.

## See Also

[inplace\\_merge function template](#), [sort function template](#)

## Name

min function template — Returns the minimum of two values

## Synopsis

```
template<typename T>  
    const T& min(const T& a, const T& b);  
template<typename T, typename Compare>  
    const T& min(const T& a, const T& b, Compare comp);
```

The `min` function template returns the smaller of `a` and `b`. If neither is smaller, it returns `a`.

The first form compares values using the `<` operator. The second form calls `comp(a, b)`.

## See Also

[max function template](#), [min\\_element function template](#)

## Name

`min_element` function template — Finds the smallest value in a range

## Synopsis

```
template<typename FwdIter>
    FwdIter min_element(FwdIter first, FwdIter last);
template<typename FwdIter, typename Compare>
    FwdIter min_element(FwdIter first, FwdIter last, Compare comp);
```

The `min_element` function template returns an iterator that points to the smallest element in the range `[first, last)`. If there are multiple instances of the smallest element, the iterator points to the first such instance.

The first form compares values using the `<` operator. The second form calls `comp(*iter1, *iter2)`.

## Technical Notes

The `min_element` function template returns `first + n`, in which `n` is the smallest value in `[0, last - first)` such that for all `m` in `[0, last - first)`, `*(first + m) < *(first + n)` is false.

Complexity is linear: exactly `max(last - first - 1, 0)` comparisons are performed.

## See Also

[max\\_element function template](#), [min function template](#)



## Name

mismatch function template — Finds first position where two ranges differ

## Synopsis

```
template<typename InIter1, typename InIter2>
    pair<InIter1, InIter2>
        mismatch(InIter1 first1, InIter1 last1, InIter2 first2);
template<typename InIter1, typename InIter2, typename
BinaryPredicate>
    pair<InIter1, InIter2>
        mismatch(InIter1 first1, InIter1 last1, InIter2 first2,
                BinaryPredicate pred);
```

The `mismatch` function template compares two sequences pairwise and returns a pair of iterators that identifies the first elements at which the sequences differ. The first sequence is `[first1, last1)`, and the second sequence starts at `first2` and has at least as many elements as the first sequence.

The return value is a pair of iterators; the first member of the pair points to an element of the first sequence, and second member of the pair points to an element of the second sequence. The two iterators have the same offset within their respective ranges. If the two sequences are equivalent, the pair returned is `last1` and an iterator that points to the second sequence at the same offset (let's call it `last2`).

The first form compares items with the `==` operator. The second form calls `pred(*iter1, *iter2)`.

[Figure 13-8](#) illustrates how the `mismatch` function template works.

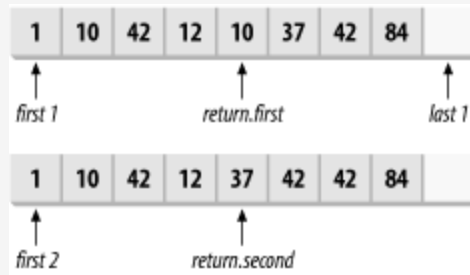


Figure 13-8. Checking two sequences for a mismatch

## Technical Notes

The mismatch function template returns the pair  $(\text{first1} + n, \text{first2} + n)$ , in which  $n$  is the smallest value in  $[0, \text{last1} - \text{first1})$  such that  $*(\text{first1} + n) == *(\text{first2} + n)$  is false and  $*(\text{first1} + m) == *(\text{first2} + m)$  is true for all  $m$  in  $[0, n)$ . If there is no such  $n$ ,  $n = \text{last1} - \text{first1}$ .

Complexity is linear: at most  $\text{last1} - \text{first1}$  comparisons are performed.

## See Also

[equal function template](#), [search function template](#), pair in [<utility>](#)

## Name

next\_permutation function template — Generates next permutation

## Synopsis

```
template<typename BidIter>
    bool next_permutation(BidIter first, BidIter last);
template<typename BidIter, typename Compare>
    bool next_permutation(BidIter first, BidIter last, Compare
        comp);
```

The next\_permutation function template rearranges the contents of the range [first, last) for the next permutation, assuming that there is a set of lexicographically ordered permutations. The return value is true if the next permutation is generated, or false if the range is at the last permutation, in which case the function cycles, and the first permutation is generated (that is, with all elements in ascending order).

[Figure 13-9](#) shows all the permutations, in order, for a sequence. The next\_permutation function swaps elements to form the next permutation. For example, if the input is permutation 2, the result is permutation 3. If the input is permutation 6, the result is permutation 1, and next\_permutation returns false.

1:	1	10	42
2:	1	42	10
3:	10	1	42
4:	10	42	1
5:	42	1	10
6:	42	10	1

*Figure 13-9. Permutations of a sequence*

## Example

[Example 13-4](#) shows a simple program that prints all the permutations of a sequence of integers. You can use this program to better understand the `next_permutation` function template.

### *Example 13-4. Generating permutations*

```
#include <algorithm>
#include <iostream>
#include <istream>
#include <iterator>
#include <ostream>
#include <vector>

void print(const std::vector<int>& v)
{
    std::copy(v.begin( ), v.end( ),
              std::ostream_iterator<int>(std::cout, " "));
    std::cout << '\n';
}

int main( )
{
    std::cout << "Enter a few integers, followed by EOF:";
    std::istream_iterator<int> start(std::cin);
    std::istream_iterator<int> end;
    std::vector<int> v(start, end);

    // Start with the first permutation (ascending order).
    std::sort(v.begin( ), v.end( ));
    print(v);

    // Print all the subsequent permutations.
    while (std::next_permutation(v.begin( ), v.end( )))
        print(v);
}
```

## Technical Notes

Complexity is linear: at most, there are  $(\text{last} - \text{first}) / 2$  swaps.

## See Also

[lexicographical\\_compare function template](#),  
[prev\\_permutation function template](#), [swap function template](#)

## Name

`nth_element` function template — Reorders a range to properly place an item at the `nth` position

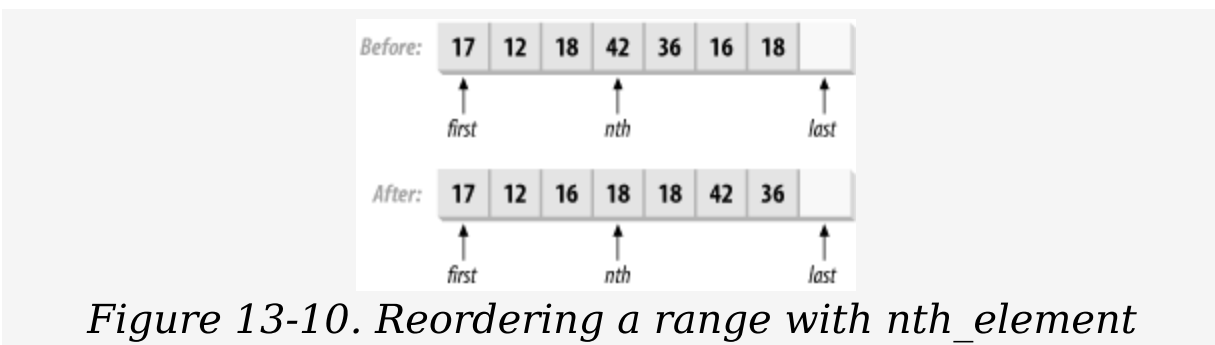
## Synopsis

```
template<typename RandIter>
    void nth_element(RandIter first, RandIter nth, RandIter last);
template<typename RandIter, typename Compare>
    void nth_element(RandIter first, RandIter nth, RandIter last,
        Compare comp);
```

The `nth_element` function template reorders the range `[first, last)` so that `*nth` is assigned the value that would be there if the entire range were sorted. It also partitions the range so that all elements in the range `[first, nth)` are less than or equal to the elements in the range `[nth, last)`.

The order is not stable—that is, if there are multiple elements that could be moved to position `nth` and preserve the sorted order, you cannot predict which element will be moved to that position.

[Figure 13-10](#) illustrates how the `nth_element` function template works.



## Technical Notes

Precondition: `nth` is in the range `[first, last)`.

Postcondition:  $*i < *nth$  for all  $i$  in  $[first, nth)$ ,  $!(*j < *nth)$  for all  $j$  in  $[nth, last)$ , and  $!(*k < *nth)$  for all  $k$  in  $[nth + 1, last)$ .

Complexity is linear for the average case but is allowed to perform worse in the worst case.

## See Also

[partition function template](#), [partial\\_sort function template](#), [sort function template](#)

## Name

`partial_sort` function template — Sorts the first part of a range

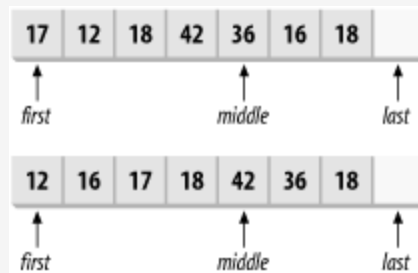
## Synopsis

```
template<typename RandIter>
    void partial_sort(RandIter first, RandIter middle, RandIter
last);
template<typename RandIter, typename Compare>
    void partial_sort(RandIter first, RandIter middle, RandIter
last,
                       Compare comp);
```

The `partial_sort` function template sorts the initial `middle - first` elements of the range `[first, last)` into the range `[first, middle)`. The remaining elements at `[middle, last)` are not in any particular order.

The first form compares values using the `<` operator. The second form calls `comp(*iter1, *iter2)`.

See [Figure 13-11](#) for an illustration of the partial-sort algorithm.



*Figure 13-11. The partial-sort algorithm*

## Technical Notes

Postcondition: for all `i` in `[first, middle - 1)`, `*(i + 1) < *i` is false, and for all `j` in `[middle, last)` and for all `i` in `[first, middle)`, `*j < *i` is false.



Complexity is logarithmic, taking about  $(\text{last} - \text{first}) \times \log(\text{middle} - \text{first})$  comparisons.

## See Also

[nth\\_element function template](#), [partial\\_sort\\_copy function template](#), [partition function template](#), [sort function template](#)

## Name

`partial_sort_copy` function template — Sorts and copies the first part of a range

## Synopsis

```
template<typename InIter, typename RandIter>
    RandIter partial_sort_copy(InIter first, InIter last, RandIter
    result_first,
                                RandIter result_last);
template<typename InIter, typename RandIter, typename Compare>
    RandIter partial_sort_copy(InIter first, InIter last, RandIter
    result_first,
                                RandIter result_last, Compare comp);
```

The `partial_sort_copy` function template copies and sorts elements from the range `[first, last)` into the range `[result_first, result_last)`. The number of items copied ( $N$ ) is the smaller of `last - first` and `result_last - result_first`. If the source range is smaller than the result range, the sorted elements are taken from the entire source range `[first, last)` and copied into the first  $N$  positions of the result range, starting at `result_first`. If the source range is larger, it is copied and sorted into the first  $N$  positions of the result range, leaving the elements in `[result_first + N, result_last)` unmodified. The return value is `result_first + N`.

The first form compares values using the `<` operator. The second form calls `comp(*iter1, *iter2)`.

## Technical Notes

Let  $n = \min(\text{last} - \text{first}, \text{result\_last} - \text{result\_first})$ .

Postcondition: for all  $i$  in `[result_first, result_first + n - 1)`, `*(i + 1) < *i` is false.

Complexity is logarithmic, taking about  $(\text{last} - \text{first}) \times \log n$  comparisons.

## See Also

[nth\\_element function template](#), [partial\\_sort\\_copy function template](#), [partition function template](#), [sort function template](#)

## Name

partition function template — Partitions a range according to a predicate

## Synopsis

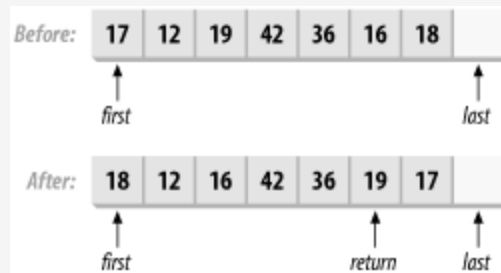
```
template<typename BidIter, typename Predicate>
    BidIter partition(BidIter first, BidIter last, Predicate
    pred);
```

The partition function template swaps elements in the range [first, last) so that all elements that satisfy pred come before those that do not. The relative order of elements is not preserved.

The return value is an iterator that points to the first element for which pred is false, or last if there is no such element.

[Figure 13-12](#) illustrates the partition function template for a predicate that tests whether a number is even:

```
function iseven(int n)
{
    return n % 2 == 0;
}
```



*Figure 13-12. Partitioning a range into even and odd numbers*

## Technical Notes

Postcondition: Let  $r$  be an iterator in the range  $[first, last]$  such that  $pred(*i)$  is true for all  $i$  in  $[first, r)$ , and  $pred(*j)$  is false for all  $j$  in  $[r, last)$ .

The partition function template returns  $r$ .

Complexity is linear:  $pred$  is called exactly  $last - first$  times, and at most  $(last - first) / 2$  swaps are performed.

## See Also

[nth\\_element function template](#), [partial\\_sort function template](#), [sort function template](#), [stable\\_partition function template](#)

## Name

pop\_heap function template — Removes largest element from a heap

## Synopsis

```
template<typename RandIter>
    void pop_heap(RandIter first, RandIter last);
template<typename RandIter, typename Compare>
    void pop_heap(RandIter first, RandIter last, Compare comp);
```

The pop\_heap function template copies the first (largest) element from the heap in [first, last) to the end of the range, that is, \*(last - 1). It then ensures that the elements remaining in [first, last - 1) form a heap.

The first form compares values using the < operator. The second form calls comp(\*iter1, \*iter2).

## Technical Notes

Precondition: [first, last) is a heap (see make\_heap for the definition of a heap).

Postcondition: [first, last - 1) is a heap, and !(\*(last - 1) < \*i) for all i in [first, last - 1).

Complexity is logarithmic: at most  $2 \times \log(\text{last} - \text{first})$  comparisons are performed.

## See Also

[make\\_heap function template](#), [push\\_heap function template](#), [sort\\_heap function template](#), [<queue>](#)

## Name

prev\_permutation function template — Generates previous permutation

## Synopsis

```
template<typename BidIter>
    bool prev_permutation(BidIter first, BidIter last);
template<typename BidIter, typename Compare>
    bool prev_permutation(BidIter first, BidIter last, Compare
    comp);
```

The `prev_permutation` function template rearranges the contents of the range `[first, last)` to the previous permutation, assuming that there is a set of lexicographically ordered permutations. The return value is `true` if the previous permutation is generated, or `false` if the range is at the first permutation, in which case the function cycles and generates the last permutation (that is, with all elements in descending order).

[Figure 13-9](#) (under `next_permutation`) shows all the permutations, in order, for a sequence. The `prev_permutation` function swaps elements to form the previous permutation. For example, if the input is permutation 3, the result is permutation 2. If the input is permutation 1, the result is permutation 6, and `prev_permutation` returns `false`.

## Technical Notes

Complexity is linear: at most  $(last - first) / 2$  swaps are performed.

## See Also

[lexicographical\\_compare function template](#),  
[next\\_permutation function template](#), [swap function  
template](#)



## Name

`push_heap` function template — Adds a value to a heap

## Synopsis

```
template<typename RandIter>
    void push_heap(RandIter first, RandIter last);
template<typename RandIter, typename Compare>
    void push_heap(RandIter first, RandIter last, Compare comp);
```

The `push_heap` function template adds the item at `last - 1` to the heap in `[first, last - 1)`, forming a new heap in the range `[first, last)`.

The first form compares values using the `<` operator. The second form calls `comp(*iter1, *iter2)`.

## Technical Notes

Precondition: `[first, last - 1)` is a heap (see `make_heap` for the definition of a heap).

Postcondition: `[first, last)` is a heap.

Complexity is logarithmic: at most  $\log(\text{last} - \text{first})$  comparisons are performed.

## See Also

[make\\_heap function template](#), [pop\\_heap function template](#), [sort\\_heap function template](#), [<queue>](#)

## Name

`random_shuffle` function template — Reorders a range into a random order

## Synopsis

```
template<typename RandIter>
    void random_shuffle(RandIter first, RandIter last);
template<typename RandIter, typename RandomNumberGenerator>
    void random_shuffle(RandIter first, RandIter last,
                        RandomNumberGenerator& rand);
```



The `random_shuffle` function template changes the order of elements in the range `[first, last)` to a random order. The first form uses an implementation-defined random number generator to produce a uniform distribution. The second form calls `rand(n)` to generate random numbers, in which `n` is a positive value of type `iterator_traits<RandIter>::difference_type`. The return value from `rand` must be convertible to the same `difference_type` and be in the range `[0, n)`.

## Technical Notes

Complexity is linear: exactly  $(last - first) + 1$  swaps are performed.

## See Also

[swap function template](#), `rand` in [<cstdlib>](#)

## Name

remove function template — Reorders a range to remove all occurrences of a value

## Synopsis

```
template<typename FwdIter, typename T>  
    FwdIter remove(FwdIter first, FwdIter last, const T& value);
```

The remove function template "removes" items that are equal to `value` from the range `[first, last)`. Nothing is actually erased from the range; instead, items to the right are copied to new positions so they overwrite the elements that are equal to `value`. The return value is one past the new end of the range. The relative order of items that are not removed is stable.

### Not e

The only way to erase an element from a container is to call one of the container's member functions.

Therefore, the remove function template does not and cannot erase items. All it can do is move items within its given range. A typical pattern, therefore, is to call `remove` to reorder the container's elements, and then call `erase` to erase the unwanted elements. To help you, the value returned from `remove` is an iterator that points to the start of the range that will be erased. For example:

```
std::vector<int> data  
...  
// Erase all values that are equal to 42.  
std::erase(std::remove(data.begin( ), data.end( ), 42),  
           data.end( ));
```

See [Figure 13-13](#) (under `remove_copy`) for an example of the removal process.

## Technical Notes

The `remove` function template assigns `*(first + n++) = *(first + m)`, in which `n` starts at 0, for all values of `m` in `[0, last - first)` in which `*(first + m) == value` is false. The return value is `first + n`.

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove\\_copy function template](#), [remove\\_copy\\_if function template](#), [remove\\_if function template](#), [replace function template](#)

## Name

`remove_copy` function template — Copies elements that are not equal to a value

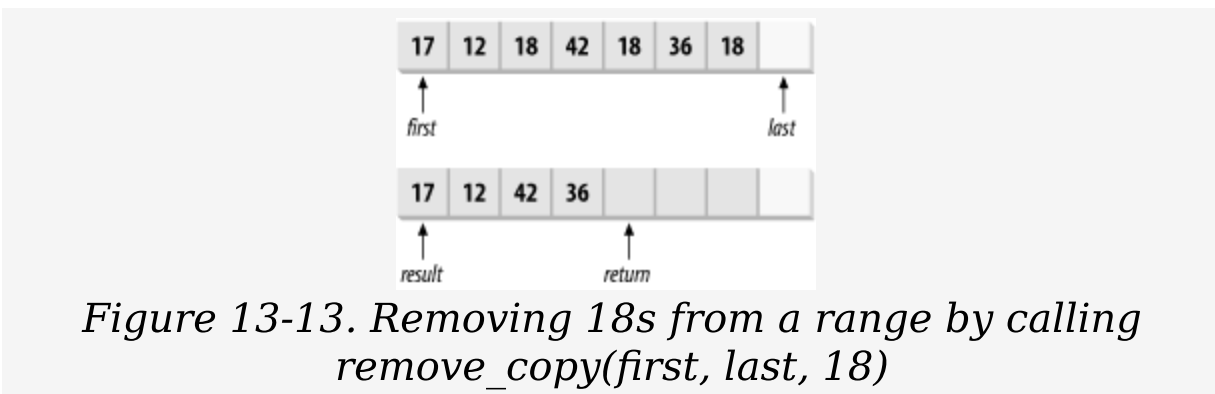
## Synopsis

```
template<typename InIter, typename OutIter, typename T>
    OutIter remove_copy(InIter first, InIter last, OutIter result,
        const T& value);
```

The `remove_copy` function template copies items from the range `[first, last)` to the range that starts at `result`. Only items that are not equal to `value` are copied, that is, cases in which `operator==` returns false.

The return value is one past the end of the result range. The relative order of items that are not removed is stable.

The source and result ranges must not overlap. [Figure 13-13](#) illustrates the removal process.



## Technical Notes

The `remove_copy` function template assigns `*(result + n++) = *(first + m)`, in which `n` starts at 0, for all values of `m` in `[0, last - first)`, in which `*(first + m) == value` is false. The return value is `result + n`.

Complexity is linear: exactly last - first comparisons are performed.

## **See Also**

[remove function template](#), [remove\\_copy\\_if function template](#), [replace\\_copy function template](#)

## Name

`remove_copy_if` function template — Copies elements for which a predicate returns false

## Synopsis

```
template<typename InIter, typename OutIter, typename Predicate>  
    OutIter remove_copy_if(InIter first, InIter last, OutIter  
    result,  
                            Predicate pred);
```

The `remove_copy_if` function template copies items from the range `[first, last)` to the range that starts at `result`. Only items for which `pred` returns false are copied.

The return value is one past the end of the result range. The relative order of items that are not removed is stable.

The source and result ranges must not overlap. See [Figure 13-13](#) (under `remove_copy`) for an example of the removal process.

## Technical Notes

The `remove_copy_if` function template assigns `*(result + n++) = *(first + m)`, in which `n` starts at 0, for all values of `m` in `[0, last - first)`, in which `pred(*(first + m))` is false. The return value is `result + n`.

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove function template](#), [remove\\_copy function template](#), [replace\\_copy\\_if function template](#)

## Name

`remove_if` function template — Reorders a range to remove elements for which a predicate returns false

## Synopsis

```
template<typename FwdIter, typename Predicate>  
    FwdIter remove_if(FwdIter first, FwdIter last, Predicate pred);
```

The `remove_if` function template "removes" items for which `pred` returns false from the range `[first, last)`. The return value is one past the new end of the range. The relative order of items that are not removed is stable.

Nothing is actually erased from the underlying container; instead, items to the right are assigned to new positions so they overwrite the elements for which `pred` returns false. See [Figure 13-13](#) (under `remove_copy`) for an example of the removal process.

## Technical Notes

The `remove_if` function template assigns `*(first + n++) = *(first + m)`, in which `n` starts at 0, for all values of `m` in `[0, last - first)`, in which `pred(*(first + m))` is false. The return value is `first + n`.

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove function template](#), [remove\\_copy\\_if function template](#), [replace\\_if function template](#)



## Name

replace function template — Replaces all occurrences of one value with another value

## Synopsis

```
template<typename FwdIter, typename T>  
    void replace(FwdIter first, FwdIter last, const T& old_value,  
                 const T& new_value);
```

The replace function template replaces all occurrences of `old_value` in `[first, last)` with `new_value`. See [Figure 13-14](#) (under `replace_copy`) for an example of the replacement process.

## Technical Notes

The replace function template assigns `*i = (*i == old_value) ? new_value : *i` for all `i` in `[first, last)`.

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove function template](#), [replace\\_copy function template](#), [replace\\_copy\\_if function template](#), [replace\\_if function template](#), [transform function template](#)

## Name

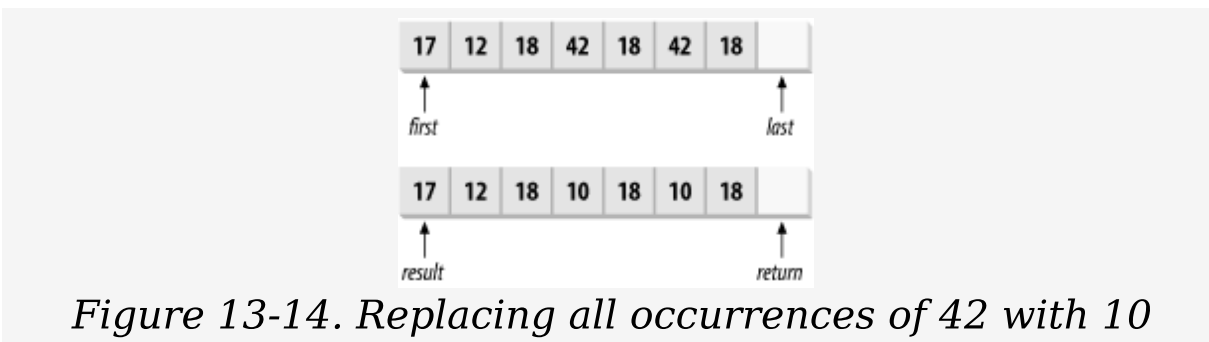
replace\_copy function template — Copies a range, replacing occurrences of one value with another value

## Synopsis

```
template<typename InIter, typename OutIter, typename T>  
    OutIter replace_copy(InIter first, InIter last, OutIter result,  
                          const T& old_value, const T& new_value);
```

The `replace_copy` function template copies values from `[first, last)` to the range that starts at `result`. Values that are equal to `old_value` are replaced with `new_value`; other values are copied without modification.

The return value is an iterator that points to one past the end of the result range. The source and result ranges must not overlap. [Figure 13-14](#) illustrates the replacement process.



## Technical Notes

The `replace_copy` function template assigns  $*(result + n) = *(first + n) == old\_value ? new\_value : *(first + n)$  for all  $n$  in  $[0, last - first)$ .

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove\\_copy function template](#), [replace function template](#),  
[replace\\_copy\\_if function template](#), [transform function  
template](#)

## Name

`replace_copy_if` function template — Copies values, replacing those that satisfy a predicate

## Synopsis

```
template<typename Iter, typename OutIter, typename Predicate,
        typename T>
    OutIter replace_copy_if(Iter first, Iter last, OutIter result,
                             Predicate pred,
                             const T& new_value);
```

The `replace_copy_if` function template copies values from `[first, last)` to the range that starts at `result`. Elements for which `pred` returns true are replaced with `new_value`; other elements are copied without modification.

The return value is an iterator that points to one past the end of the result range. The source and result ranges must not overlap. See [Figure 13-14](#) (under `replace_copy`) for an example of the replacement process.

## Technical Notes

The `replace_copy_if` function template assigns `*(result + n) = *(first + n) == pred(*(first + n)) ? new_value : *(first + n)` for all `n` in `[0, last - first)`.

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove\\_copy\\_if function template](#), [replace function template](#), [replace\\_copy function template](#), [replace\\_if function template](#), [transform function template](#)

## Name

`replace_if` function template — Replaces values that satisfy a predicate

## Synopsis

```
template<typename FwdIter, typename Predicate, typename T>
    void replace_if(FwdIter first, FwdIter last, Predicate pred,
                    const T& new_value);
```

The `replace_if` function template replaces all values in `[first, last)` for which `pred` is true with `new_value`. See [Figure 13-14](#) (under `replace_copy`) for an example of the replacement process.

## Technical Notes

The `replace_if` function template assigns `*i = pred(*i) ? new_value : *i` for all `i` in `[first, last)`.

Complexity is linear: exactly `last - first` comparisons are performed.

## See Also

[remove\\_if function template](#), [replace function template](#), [replace\\_copy function template](#), [transform function template](#)

## Name

reverse function template — Reverses the values in a range

## Synopsis

```
template<typename BidIter>
    void reverse(BidIter first, BidIter last);
```

The reverse function template reverses the order of the items in the range [first, last). See [Figure 13-15](#) (under reverse\_copy) for an example.

## Technical Notes

The reverse function template swaps  $*(first + n)$  with  $*(last - n - 1)$  for all  $n$  in  $[0, (last - first) / 2]$ .

Complexity is linear: exactly  $(last - first) / 2$  swaps are performed.

## See Also

[reverse\\_copy function template](#), [rotate function template](#), [swap function template](#)

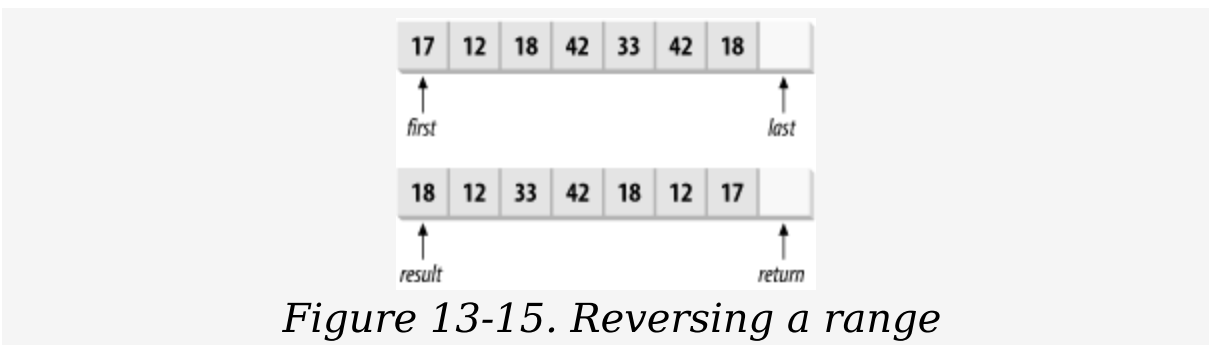
## Name

`reverse_copy` function template — Copies a range in reverse order

## Synopsis

```
template<typename BidIter, typename OutIter>
    OutIter reverse_copy(BidIter first, BidIter last, OutIter
        result);
```

The `reverse_copy` function template copies items in reverse order from the range `[first, last)` to the range that starts at `result`. In other words, `*(last - 1)` is first copied to `*result`, then `*(last - 2)` is copied to `*(result + 1)`, and so on. The return value is an iterator that points to one past the end of the result range. The source and result ranges must not overlap. [Figure 13-15](#) shows an example of reversing a range.



## Technical Notes

The `reverse_copy` function template assigns `*(result + n) = *(last - n - 1)` for all `n` in `[0, last - first)`.

Complexity is linear: exactly `last - first` assignments are performed.

## See Also

[copy\\_backward function template](#), [reverse function template](#), [rotate\\_copy function template](#)



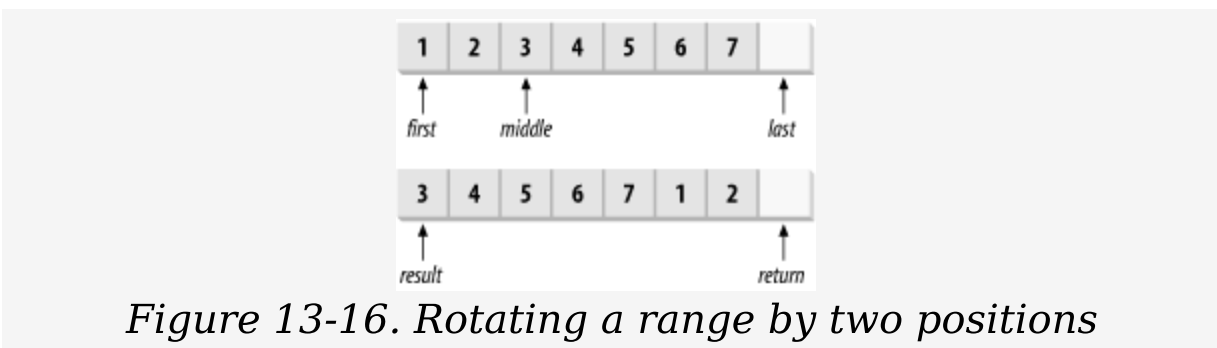
## Name

rotate function template — Rotates elements in a range

## Synopsis

```
template<typename FwdIter>  
void rotate(FwdIter first, FwdIter middle, FwdIter last);
```

The rotate function template rotates elements in the range [first, last) to the left so that the items in the range [middle, last) are moved to the start of the new sequence. Elements in the range [first, middle) are rotated to the end. See [Figure 13-16](#) for an example.



*Figure 13-16. Rotating a range by two positions*

## Technical Notes

For all  $n$  in  $[0, \text{last} - \text{first})$ , the rotate function template moves  $*(\text{first} + n)$  into position  $\text{first} + (n + (\text{last} - \text{middle})) \% (\text{last} - \text{first})$ .

Complexity is linear: at most  $\text{last} - \text{first}$  swaps are performed.

## See Also

[reverse function template](#), [rotate\\_copy function template](#)

## Name

`rotate_copy` function template — Rotates and copies items in a range

## Synopsis

```
template<typename FwdIter, typename OutIter>
    OutIter rotate_copy(FwdIter first, FwdIter middle, FwdIter
last,
                        OutIter result);
```

The `rotate_copy` function template copies elements from the range `[middle, last)` to the range that starts at `result` followed by the elements from `[first, middle)`, thereby effecting a rotation to the left. The return value is one past the end of the result range.

The source and result ranges must not overlap. [Figure 13-16](#) shows an example of rotation.

## Technical Notes

The `rotate_copy` function template assigns  $*(result + (n + (last - middle) \% (last - first)) = *(first + n)$  for all  $n$  in  $[0, last - first)$ . It returns `result + (last - first)`.

Complexity is linear: exactly `last - first` assignments are performed.

## See Also

[reverse\\_copy function template](#), [rotate function template](#)

## Name

search function template — Searches a range for a subsequence

## Synopsis

```
template<typename FwdIter1, typename FwdIter2>
    FwdIter1 search(FwdIter1 first1, FwdIter1 last1, FwdIter2
first2,
                  FwdIter2 last2);
template<typename FwdIter1, typename FwdIter2, typename
BinaryPredicate>
    FwdIter1 search(FwdIter1 first1, FwdIter1 last1, FwdIter2
first2,
                  FwdIter2 last2, BinaryPredicate pred);
```

The search function template finds the first (leftmost) subsequence [first2, last2) within the range [first1, last1). It returns an iterator that points to the start of the subsequence or last1 if the subsequence is not found.

The first form compares items with the == operator. The second form calls pred(\*iter1, \*iter2).

[Figure 13-5](#) (under find\_end) illustrates the search function template.

## Technical Notes

Let length1 = last1 - first1 and length2 = last2 - first2.

The search function template returns first1 + n, in which n is the smallest value in the range [0, length1 - length2) such that \*(i + n + m) == (first2 + m) for all m in the range [0, length2). It returns last1 if no such n can be found.

Complexity: at most length1 × length2 comparisons are performed.

## See Also

[find function template](#), [find\\_end function template](#),  
[search\\_n function template](#)

## Name

`search_n` function template — Searches a range for a repeated value

## Synopsis

```
template<typename FwdIter, typename Size, typename T>
    FwdIter search_n(FwdIter first, FwdIter last, Size count, const
T& value);
template<typename FwdIter, typename Size, typename T, typename
BinaryPredicate>
    FwdIter search_n(FwdIter first, FwdIter last, Size count, const
T& value,
                    BinaryPredicate pred);
```

The `search_n` function template finds the first (leftmost) subsequence of `count` adjacent occurrences of `value` in the range `[first, last)`. It returns an iterator that points to the start of the subsequence or `last` if the subsequence is not found.

The first form compares items with the `==` operator. The second form calls `pred(*iter, value)`.

## Technical Notes

The `search_n` function template returns `first + n`, in which `n` is the smallest value in the range `[0, last - first)` such that `*(i + n + m) == value` for all `m` in the range `[0, count)`. It returns `last` if no such `n` can be found.

Complexity: at most  $n \times (\text{last} - \text{first})$  comparisons are performed.

## See Also

[find function template](#), [search function template](#)

## Name

`set_difference` function template — Computes set difference of sorted ranges

## Synopsis

```
template<typename InIter1, typename InIter2, typename OutIter>
    OutIter set_difference(InIter1 first1, InIter1 last1, InIter2
    first2,
                            InIter2 last2, OutIter result);
template<typename InIter1, typename InIter2, typename OutIter,
    typename Compare>
    OutIter set_difference(InIter1 first1, InIter1 last1, InIter2
    first2,
                            InIter2 last2, OutIter result, Compare
    comp);
```

The `set_difference` function template copies elements from the sorted range `[first1, last1)` to the range starting at `result`. Only those elements that are not also present in the sorted range `[first2, last2)` are copied. An iterator that points to one past the end of the result range is returned.

The result range must not overlap either source range.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether  $X < Y$ .

[Figure 13-17](#) shows an example of a set difference using multisets.

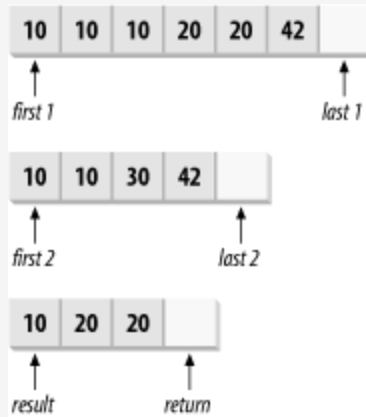


Figure 13-17. Computing the difference between sets

## Technical Notes

Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first1, last1 - 1)$  and  $!(*(j + 1) < *j)$  for all  $j$  in  $[first2, last2 - 1)$ .

Postcondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[result, return - 1)$ .

The `set_difference` function template assigns  $*(result + n++) = *(first1 + m)$  for all  $m$  in  $[first1, last1)$ , in which  $*(first1 + m)$  is not in  $[first2, last2)$ . It returns `result + n`.

Complexity is linear: at most  $2 \times ((last1 - first1) + (last2 - first2)) - 1$  comparisons are performed.

## See Also

[includes function template](#), [set\\_intersection function template](#), [set\\_symmetric\\_difference function template](#), [set\\_union function template](#)

## Name

`set_intersection` function template — Computes intersection of sorted ranges

## Synopsis

```
template<typename InIter1, typename InIter2, typename OutIter>
    OutIter set_intersection(InIter1 first1, InIter1 last1, InIter2
    first2,
                               InIter2 last2, OutIter result);
template<typename InIter1, typename InIter2, typename OutIter,
        typename Compare>
    OutIter set_intersection(InIter1 first1, InIter1 last1, InIter2
    first2,
                               InIter2 last2, OutIter result, Compare
    comp);
```

The `set_intersection` function template copies elements from the sorted range `[first1, last1)` to the range starting at `result`. Only those elements that are also present in the sorted range `[first2, last2)` are copied. An iterator that points to one past the end of the result range is returned.

The result range must not overlap either source range.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether  $X < Y$ .

[Figure 13-18](#) shows an example of intersection using multisets.



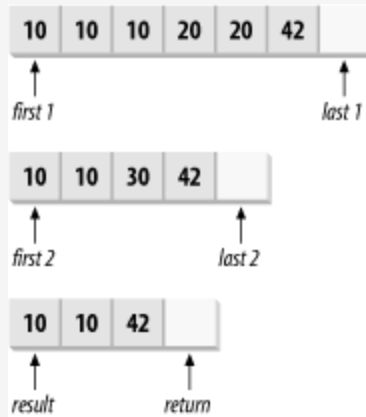


Figure 13-18. Intersecting two sets

## Technical Notes

Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first1, last1 - 1)$  and  $!(*(j + 1) < *j)$  for all  $j$  in  $[first2, last2 - 1)$ .

Postcondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[result, return - 1)$ .

The `set_intersection` function template assigns  $*(result + n++) = *(first1 + m)$  for all  $m$  in  $[first1, last1)$ , in which  $*(first1 + m)$  is in  $[first2, last2)$ . It returns `result + n`.

Complexity is linear: at most  $2 \times ((last1 - first1) + (last2 - first2)) - 1$  comparisons are performed.

## See Also

[includes function template](#), [set\\_difference function template](#), [set\\_symmetric\\_difference function template](#), [set\\_union function template](#)

## Name

`set_symmetric_difference` function template — Computes symmetric difference of sorted ranges

## Synopsis

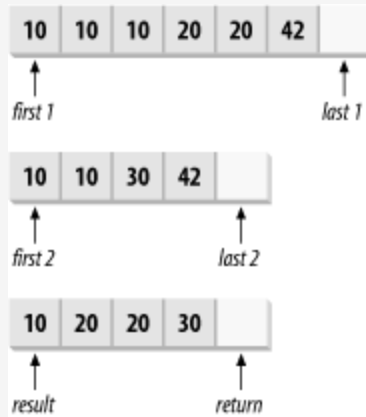
```
template<typename InIter1, typename InIter2, typename OutIter>
    OutIter set_symmetric_difference(InIter1 first1, InIter1 last1,
    InIter2 first2,
                                     InIter2 last2, OutIter
    result);
template<typename InIter1, typename InIter2, typename OutIter,
    typename Compare>
    OutIter set_symmetric_difference(InIter1 first1, InIter1 last1,
    InIter2 first2,
                                     InIter2 last2, OutIter result,
    Compare comp);
```

The `set_symmetric_difference` function template merges elements from the sorted ranges `[first1, last1)` and `[first2, last2)`, copying the sorted, merged results to the range starting at `result`. Only those elements that are not also present in the sorted range `[first2, last2)` are copied from `[first1, last1)`, and only those not present in the range `[first1, last1)` are copied from `[first2, last2)`. An iterator that points to one past the end of the result range is returned.

The result range must not overlap either source range.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether `X < Y`.

[Figure 13-19](#) shows an example of a set symmetric difference using multisets.



*Figure 13-19. Computing the symmetric difference between two sets*

## Technical Notes

Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first1, last1 - 1)$  and  $!(*(j + 1) < *j)$  for all  $j$  in  $[first2, last2 - 1)$ .

Postcondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[result, return - 1)$ .

Complexity is linear: at most  $2 \times ((last1 - first1) + (last2 - first2)) - 1$  comparisons are performed.

## See Also

[includes function template](#), [set\\_difference function template](#), [set\\_intersection function template](#), [set\\_union function template](#)

## Name

set\_union function template — Computes union of sorted ranges

## Synopsis

```
template<typename InIter1, typename InIter2, typename OutIter>
    OutIter set_union(InIter1 first1, InIter1 last1, InIter2
first2, InIter2 last2,
                    OutIter result);
template<typename InIter1, typename InIter2, typename OutIter,
typename Compare>
    OutIter set_union(InIter1 first1, InIter1 last1, InIter2
first2, InIter2 last2,
                    OutIter result, Compare comp);
```

The `set_union` function template merges elements from the sorted ranges `[first1, last1)` and `[first2, last2)`, copying the sorted, merged results to the range starting at `result`. If an element is present in both ranges, only one element is copied to the result range. If the input ranges contain duplicates, each occurrence of an element in an input range results in a copy in the result range. An iterator that points to one past the end of the result range is returned.

The result range must not overlap either source range.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether `X < Y`.

[Figure 13-20](#) shows an example of a set union using multisets.

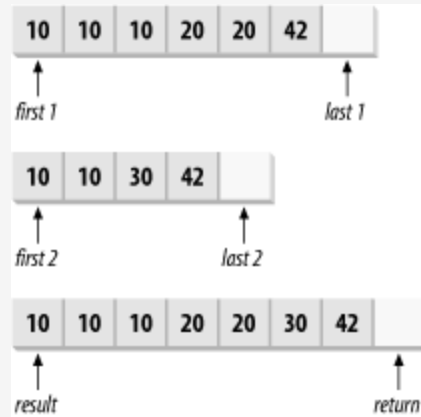


Figure 13-20. Computing the union of two sets

## Technical Notes

Precondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[first1, last1 - 1)$  and  $!(*(j + 1) < *j)$  for all  $j$  in  $[first2, last2 - 1)$ .

Postcondition:  $!(*(i + 1) < *i)$  for all  $i$  in  $[result, return - 1)$ .

Complexity is linear: at most  $2 \times ((last1 - first1) + (last2 - first2)) - 1$  comparisons are performed.

## See Also

[includes function template](#), [merge function template](#), [set\\_difference function template](#), [set\\_intersection function template](#), [set\\_symmetric\\_difference function template](#)

## Name

sort function template — Sorts a range in place

## Synopsis

```
template<typename RandIter>
    void sort(RandIter first, RandIter last);
template<typename RandIter, typename Compare>
    void sort(RandIter first, RandIter last, Compare comp);
```

The sort function template sorts the range [first, last) in place. The sort is not stable, so equivalent elements do not preserve their original order.

The first version compares items using the < operator. The second version uses comp(X, Y) to test whether X < Y.

## Technical Notes

Postcondition: !(\*i + 1) < \*i) for all i in [first, last - 1).

Complexity is  $n \log n$  comparisons in the average case, in which  $n = \text{last} - \text{first}$ . Worst case performance might be worse. Use `stable_sort` if the worst-case performance is more important than the average performance.

## See Also

[merge function template](#), [nth\\_element function template](#), [partial\\_sort function template](#), [partition function template](#), [sort\\_heap function template](#), [stable\\_sort function template](#)

## Name

sort\_heap function template — Sorts a heap in place

## Synopsis

```
template<typename RandIter>
    void sort_heap(RandIter first, RandIter last);
template<typename RandIter, typename Compare>
    void sort_heap(RandIter first, RandIter last, Compare comp);
```

The `sort_heap` function template sorts a heap in the range `[first, last)`. The sort is not stable, so equivalent elements do not preserve their original order.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether  $X < Y$ .

## Technical Notes

Precondition: `[first, last)` is a heap (see `make_heap` for the definition of a heap).

Postcondition: `!(*(i + 1) < *i)` for all `i` in `[first, last - 1)`.

Complexity is at most  $n \log n$  comparisons, in which  $n = \text{last} - \text{first}$ .

## See Also

[make\\_heap function template](#), [pop\\_heap function template](#), [push\\_heap function template](#), [sort function template](#), [<queue>](#)

## Name

`stable_partition` function template — Partitions a range in stable order

## Synopsis

```
template<typename BidIter, typename Predicate>
    BidIter stable_partition(BidIter first, BidIter last,
    Predicate pred);
```

The `stable_partition` function template swaps elements in the range `[first, last)` so that all elements that satisfy `pred` come before those that do not. The relative order of elements in each partition is preserved.

The return value is an iterator that points to the first element for which `pred` is false, or `last` if there is no such element.

[Figure 13-12](#) (under `partition`) illustrates the partition functionality.

## Technical Notes

Postcondition: Let `r` be an iterator in the range `[first, last]` such that `pred(*i)` is true for all `i` in `[first, r)`, and `pred(*j)` is false for all `j` in `[r, last)`.

The `stable_partition` function template returns `r`.

Complexity is linear if there is enough memory. Otherwise, at most  $n \log n$  swaps are performed, in which  $n = \text{last} - \text{first}$ . In all cases, *pred* is called exactly  $n$  times.

## See Also

[nth\\_element function template](#), [partial\\_sort function template](#), [partition function template](#), [stable\\_sort function template](#)



## Name

`stable_sort` function template — Sorts a range in place in stable order

## Synopsis

```
template<typename RandIter>
    void stable_sort(RandIter first, RandIter last);
template<typename RandIter, typename Compare>
    void stable_sort(RandIter first, RandIter last, Compare comp);
```

The `stable_sort` function template sorts the range `[first, last)` in place. The sort is stable, so equivalent elements preserve their original order.

The first version compares items using the `<` operator. The second version uses `comp(X, Y)` to test whether `X < Y`.

## Technical Notes

Postcondition: `!(*i + 1) < *i` for all `i` in `[first, last - 1)`.

Complexity is at most  $n \log n$  comparisons, in which  $n = \text{last} - \text{first}$ , provided that enough memory is available for temporary results. If memory is limited, the complexity is at most  $n (\log n)^2$  comparisons.

## See Also

[merge function template](#), [nth\\_element function template](#), [partial\\_sort function template](#), [sort function template](#), [sort\\_heap function template](#), [stable\\_partition function template](#)

## Name

swap function template — Swaps two values

## Synopsis

```
template<typename T> void swap(T& a, T& b);
```

The swap function template swaps the values of a and b.

The standard containers all specialize the swap template to call their swap member functions, which usually run in constant time, regardless of the number of elements in the containers.

## See Also

[iter\\_swap function template](#), [swap\\_ranges function template](#)

## Name

swap\_ranges function template — Swaps all values in two ranges

## Synopsis

```
template<typename FwdIter1, typename FwdIter2>
    FwdIter2 swap_ranges(FwdIter1 first1, FwdIter1 last1, FwdIter2
        first2);
```

The swap\_ranges function template swaps all the elements in [first1, last1) with corresponding elements in the range that starts at first2 (and has the same length as the first range). The return value is one past the end of the second range. The two ranges must not overlap.

## Technical Notes

The swap\_ranges function template performs the equivalent of `swap(*(first1 + n), *(first2 + n))` for all `n` in `[0, last1 - first1)`.

Complexity is linear: exactly `last1 - first1` swaps are performed.

## See Also

[iter\\_swap function template](#), [swap function template](#)

## Name

transform function template — Copies one or two ranges after applying an operator to each element

## Synopsis

```
template<typename InIter, typename OutIter, typename
UnaryOperation>
    OutIter transform(InIter first, InIter last, OutIter result,
                      UnaryOperation unop);
template<typename InIter1, typename InIter2, typename OutIter,
         typename BinaryOperation>
    OutIter transform(InIter1 first1, InIter1 last1, InIter2
first2,
                      OutIter result, BinaryOperation binop);
```

The transform function template assigns a new value to each element in the range that starts at `result`. In the first case, the new value is `unop(*iter)`, in which `iter` is an iterator over `[first, last)`.

In the second case, the new value is `binop(*iter1, *iter2)`, in which `iter1` ranges over `[first1, last1)` and `iter2` iterates over the range that starts at `first2`. The second input range must be at least as long as the first.

The return value is one past the end of the result range. The result range can be the same as any of the input ranges.

## Technical Notes

The first form of transform assigns  $*(result + n) = unop(*(first + n))$  for all  $n$  in  $[0, last - first)$ .

The second form assigns  $*(result + n) = binop(*(first1 + n), *(first2 + n))$  for all  $n$  in  $[0, last1 - first1)$ .

Complexity is linear: *unop* or *binop* is called exactly  $n$  times, in which  $n = last - first$  for a unary operator or

last1 - first1 for a binary operator.

## **See Also**

[copy function template, for\\_each function template](#)

## Name

unique function template — Removes adjacent, equal values from a range

## Synopsis

```
template<typename FwdIter>
    FwdIter unique(FwdIter first, FwdIter last);
template<typename FwdIter, typename BinaryPredicate>
    FwdIter unique(FwdIter first, FwdIter last, BinaryPredicate
pred);
```

The unique function template "removes" repetitions of adjacent, identical elements from the range [first, last). The return value is one past the new end of the range. For each sequence of identical elements, only the first is kept. The input range does not have to be sorted, but if it is, all duplicates are "removed," leaving only unique values (hence the function's name).

Nothing is actually erased from the underlying container; instead, items to the right are copied to new positions at lower indices (to the left) so they overwrite the elements that are duplicates. See [Figure 13-21](#) (under unique\_copy) for an example of the removal process.

The first form compares items with the == operator. The second form calls pred(a, b).

## Technical Notes

The unique function template assigns  $*(first + n++) = *(first + m)$  for all  $m$  in  $[0, last - first)$ , in which  $m == 0$  or  $*(first + m) == *(first + m - 1)$  is false. It returns  $first + n$ .

Complexity is linear: exactly  $\max(0, last - first - 1)$  comparisons are performed.

## See Also

[remove function template](#), [unique\\_copy function template](#)

## Name

unique\_copy function template — Copies unique values

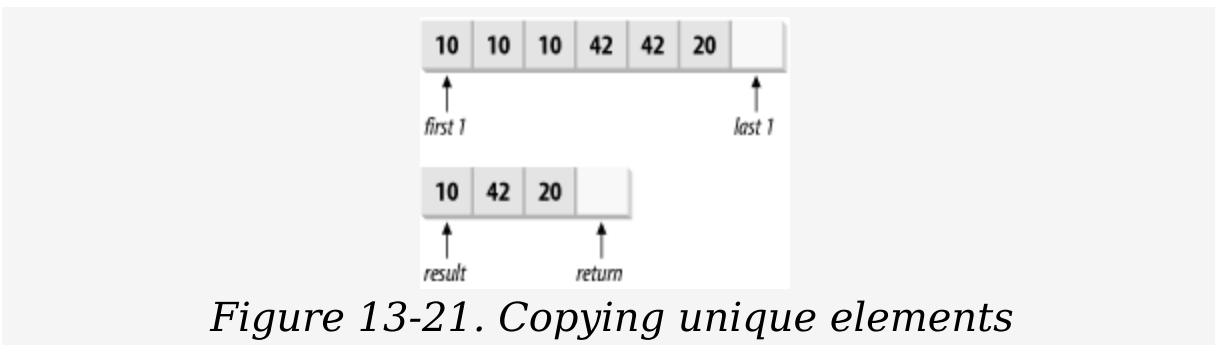
## Synopsis

```
template<typename InIter, typename OutIter>
    OutIter unique_copy(InIter first, InIter last, OutIter result);
template<typename InIter, typename OutIter, typename
BinaryPredicate>
    OutIter unique_copy(InIter first, InIter last, OutIter result,
                        BinaryPredicate pred);
```

The `unique_copy` function template copies items from `[first, last)` to the range that starts at `result`, removing duplicates. For each sequence of identical elements, only the first is kept. The return value is one past the end of the result range.

The first form compares items with the `==` operator. The second form calls `pred(a, b)`.

See [Figure 13-21](#) for an example that calls `unique_copy`.



*Figure 13-21. Copying unique elements*

## Technical Notes

The `unique_copy` function template assigns `*(result + n++) = *(first + m)` for all `m` in `[0, last - first)`, in which `m == 0` or `*(first + m) == *(first + m - 1)` is false. It returns `result + n`.



Complexity is linear: exactly last - first comparisons are performed.

## **See Also**

[remove\\_copy function template](#), [unique function template](#)

## Name

`upper_bound` function template — Finds upper bound for a value's position in a sorted range using binary search

## Synopsis

```
template<typename FwdIter, typename T>
    FwdIter upper_bound(FwdIter first, FwdIter last, const T&
value);
template<typename FwdIter, typename T, typename Compare>
    FwdIter upper_bound(FwdIter first, FwdIter last, const T&
value, Compare comp);
```

The `upper_bound` function template determines where `value` belongs in the sorted range `[first, last)`. The return value is an iterator that points to one past the last (rightmost) occurrence of `value` in the range, if `value` is present. Otherwise, the iterator points to the last position where you can insert `value` and preserve the sorted nature of the range.

The first form compares values using the `<` operator. The second form calls `comp(*iter, value)`.

[Figure 13-4](#) (under `equal_range`) shows an example of finding the bounds for the value 36. The `upper_bound` function returns `ub` as the upper bound of 36 in the given range. Note that it returns `ub` as the upper bound for all values in the range `[36, 41]`. For values in the range `[19, 35]`, the upper bound is equal to `lb`.

## Technical Notes

Precondition: `!(*(i + 1) < *i)` for all `i` in `[first, last - 1)`.

The `upper_bound` function template returns `first + n`, in which `n` is the highest value in `[0, last - first)` such that `*(first + m) < value` is false for all `m` in `[0, n)`.

Complexity is logarithmic. The number of comparisons is at most  $\log(\text{last} - \text{first}) + 1$ . Although the iterator can be a forward iterator, the best performance is obtained with a random access iterator. With a forward or bidirectional iterator, the iterator is advanced a linear number of times, even though the number of comparisons is logarithmic.

## See Also

[binary\\_search function template](#), [equal\\_range function template](#), [lower\\_bound function template](#)

## **<bitset>**

The `<bitset>` header declares a single class template, `bitset`, and some related functions. A `bitset` is a fixed-size sequence of bits. The bitwise operators (`&`, `|`, `^`, etc.) are overloaded to work with `bitsets` in the usual manner, and you can refer to individual bits by index.

The Boost project has a class template for a bit sequence that can change size at runtime. See [Appendix B](#) for information about Boost.

# Name

bitset class template — Fixed-size sequence of bits

## Synopsis

```
template<size_t N>
class bitset {
public:
    // Proxy class to simulate a bit reference
    class reference {
        friend class bitset;
        reference( );
    public:
        ~reference( );
        reference& operator=(bool x);
        reference& operator=(const reference&);
        bool operator~( ) const;
        operator bool( ) const;
        reference& flip( );
    };

    // Constructors
    bitset( );
    bitset(unsigned long val);

    template<typename charT, typename traits, typename A>
    explicit bitset(const basic_string<charT,traits,A>& s, typename
        basic_string<charT,traits,A>::size_type p=0, typename basic
        string<charT,traits,A>::size_type n =
basic_string<charT,traits,A>::npos);

    // bitset operations
    bitset<N>& operator&=(const bitset<N>& rhs);
    bitset<N>& operator|=(const bitset<N>& rhs);
    bitset<N>& operator^=(const bitset<N>& rhs);
    bitset<N>& operator<<=(size_t pos);
    bitset<N>& operator>>=(size_t pos);
    bitset<N>& set( );
    bitset<N>& set(size_t pos, int val=true);
    bitset<N>& reset( );
    bitset<N>& reset(size_t pos);
    bitset<N> operator~( ) const;
    bitset<N>& flip( );
    bitset<N>& flip(size_t pos);
```

```

// Element access
reference operator[](size_t pos);
bool operator[](size_t pos) const;

unsigned long to_ulong( ) const;
template <typename charT, typename traits, typename Alloc>
    basic_string<charT, traits, Alloc> to_string( ) const;

size_t count( ) const;
size_t size( ) const;
bool operator==(const bitset<N>& rhs) const;
bool operator!=(const bitset<N>& rhs) const;
bool test(size_t pos) const;
bool any( ) const;
bool none( ) const;
bitset<N> operator<<(size_t pos) const;
bitset<N> operator>>(size_t pos) const
};

```

The `bitset` class template offers a convenient way to manipulate a fixed-sized sequence of bits. The number of bits is specified as a template argument, so each `bitset` object can have a different size. Each bit in a `bitset` can be set (1 or `true`) or reset (0 or `false`). Bit positions are numbered from right to left, that is, 0 is the least-significant bit, and  $N - 1$  is the most-significant bit.

A `bitset` is not a standard container and does not provide iterators or support generic algorithms. For a container that holds a sequence of bit values, use `vector<int>` or `deque<bool>`. (See `<vector>` later in this chapter to learn more, including why you should not use `vector<bool>`.) In the following member function descriptions,  $N$  is the template parameter (number of bits):

### **bitset** ( )

Resets all bits.

### **bitset** (unsigned long value)

Initializes the first  $m$  bits to `value`, in which  $m == \text{CHAR\_BITS} * \text{sizeof}(\text{unsigned long})$ . If  $N > m$ , all other bits are reset to 0. If  $N < m$ , excess bits of  $m$  are ignored.

```

template<typename charT, typename traits, typename
A> explicit bitset(const
basic_string<charT,traits,A>& s, typename
basic_string<charT,traits,A>::size_type p=0,
typename basic_string<charT,traits,A>::size_type
n= basic_string<charT,traits,A>::npos)

```

Initializes the *bitset* from the character string *s*, starting at index *p* and extending for *n* characters (or to the end of the string, whichever comes first). The default is to use all characters in the string. A character equal to '0' resets a bit, '1' sets a bit, and any other character causes the constructor to throw *invalid\_argument*.

The rightmost character of the substring (that is, the character *s*[*p*+*n*-1] or the rightmost character of *s*) initializes the bit at index 0 of the *bitset*, and subsequent bits are initialized by characters at preceding indices of *s*. Bits left uninitialized by the string are reset. All of the *bitsets* in the following example are equal to **000111**:

```

bitset<6> a(string("111"));
bitset<6> b(string("000111"));
bitset<6> c(string("10110011100"), 5, 4);
bitset<6> d(string("111111"), 3, 42);

```

The unwieldy declaration is due to the *basic\_string* class template. For the common case of a plain string, you can read the declaration as:

```

bitset(const string& s, size_t p=0, size_n n=string::npos)

```

```

bitset<N>& operator&= (const bitset<N>& rhs)

```

Performs *\*this* = *\*this* & *rhs*. Returns *\*this*.

```

bitset<N>& operator|= (const bitset<N>& rhs)

```

Performs *\*this* = *\*this* | *rhs*. Returns *\*this*.

```

bitset<N>& operator^= (const bitset<N>& rhs)

```

Performs *\*this* = *\*this* ^ *rhs*. Returns *\*this*.

`bitset<N>& operator<<= (size_t pos)`

Shifts bits to the left by `pos` positions. Vacated bits are filled with `0`. Returns `*this`.

`bitset<N>& operator>>= (size_t pos)`

Shifts bits to the right by `pos` positions. Vacated bits are filled with `0`. Returns `*this`.

`bool operator== (const bitset<N> rhs)`

Returns true if every bit in `*this` has the same value as the corresponding bit in `rhs`.

`bool operator!= (const bitset<N> rhs)`

Returns true if any bit in `*this` has a different value than the corresponding bit in `rhs`.

`bitset<N> operator<< (size_t pos)`

Returns a new `bitset` with its bits shifted to the left by `pos` positions. Vacated bits are filled with `0`.

`bitset<N> operator>> (size_t pos)`

Returns a new `bitset` with its bits shifted to the right by `pos` positions. Vacated bits are filled with `0`.

`bitset<N> operator~ ( ) const`

Returns a new `bitset` with all bits flipped.

`reference operator[] (size_t pos)`

Returns a `bitset::reference` object for the bit at position `pos`. The behavior is undefined if `pos` is out of range.



`bool operator[] (size_t pos) const`

Returns the value of the bit at position `pos`. The behavior is undefined if `pos` is out of range. This member function was added to the standard as part of the technical



corrigendum (TC1), so it might not yet be supported by some compilers.

`bool any ( ) const`

Returns true if any bit is set. Returns false if all bits are 0.

`size_t count ( ) const`

Returns the number of bits set.

`bitset<N>& flip ( )`

Toggles all bits, that is, sets 0 bits to 1 and 1 bits to 0. Returns \*this.

`bitset<N>& flip (size_t pos)`

Toggles the bit at position pos. If pos is invalid, throws `out_of_range`. Returns \*this.

`bool none ( ) const`

Returns true if all bits are 0. Returns false if any bit is set.

`bitset<N>& reset ( )`

Resets all bits. Returns \*this.

`bitset<N>& reset (size_t pos)`

Resets the bit at position pos. If pos is invalid, throws `out_of_range`. Returns \*this.

`bitset<N>& set ( )`

Sets all bits. Returns \*this.

`bitset<N>& set (size_t pos, int val = true)`

Sets the bit at position pos to `val != 0`. If pos is invalid, throws `out_of_range`. Returns \*this.

`size_t size ( ) const`

Returns N.

bool **test** (size\_t pos) const

Returns the value of the bit at position pos. Throws `out_of_range` if pos is invalid.

template <class charT, class traits, class Allocator>, basic\_string<charT, traits, Allocator>, **to\_string** ( ) const

Returns a string representation of the bitset. Each bit is converted to the character '0' if reset or '1' if set. Bit position 0 is the rightmost character (position N - 1).

The compiler cannot deduce the template parameters when calling `to_string`, so you must specify them explicitly:

```
std::bitset<64> bits(std::string("101000111101010101"));
std::string str = bits.template to_string<char,
std::char_traits<char>,
std::allocator<char>
>( );
```

unsigned long **to\_ulong** ( ) const

Returns the integral value of the bitset. If N is too large for unsigned long, it throws `overflow_error`.

## See Also

[<climits>](#), [<vector>](#)

# Name

bitset::reference class — Proxy class for a bit in a bitset

# Synopsis

```
class reference {  
    friend class bitset;  
    reference( )  
public:  
    ~reference( );  
    reference& operator=(bool x);  
    reference& operator=(const reference&);  
    bool operator~( ) const;  
    operator bool( ) const;  
    reference& flip( );  
};
```

The `bitset::reference` class is a proxy that refers to a single bit in a bitset. The constructor is private, so instances can be created only by the `bitset` class, particularly by its `operator[]` function. The member functions are:

`reference& operator= (bool x), reference& operator= (const reference& x)`

Sets the referenced bit to `x` in the underlying bitset.  
Returns `*this`.

`bool operator~ ( ) const`

Returns the logical negation of the referenced bit.

`operator bool ( ) const`

Returns the value of the referenced bit.

`reference& flip ( )`

Toggles the referenced bit in the underlying bitset.  
Returns `*this`.

## See Also

[bitset class template](#)

## Name

operator& function template — Performs bitwise and of two bitsets

## Synopsis

```
template <size_t N>  
    bitset<N> operator&(const bitset<N>& a, const bitset<N>& b);
```

The & operator takes two bitsets and returns a new bitset that represents the *bitwise and* of the operands. In other words, an output bit is set only when the corresponding bit is set in both operands; otherwise, an output bit is reset.

## See Also

[bitset class template](#), [operator|\\_function template](#), [operator^\\_function template](#), [<cstdintdef>](#), [bit\\_and](#) keyword

## Name

operator| function template — Performs bitwise inclusive or of two bitsets

## Synopsis

```
template <size_t N>
    bitset<N> operator|(const bitset<N>& a, const bitset<N>& b);
```

The | operator takes two bitsets and returns a new bitset that represents the *bitwise inclusive or* of the operands. In other words, an output bit is set when the corresponding bit is set in either operand, and an output bit is reset if the corresponding bits in both operands are 0.

## See Also

[bitset class template](#), operator &, operator ^, [<cstdint>](#), bit\_or keyword

## Name

`operator^` function template — Performs bitwise exclusive or of two bitsets

## Synopsis

```
template <size_t N>  
    bitset<N> operator^(const bitset<N>& a, const bitset<N>& b);
```

The `^` operator takes two bitsets and returns a new bitset that represents the *bitwise exclusive or* of the operands. In other words, an output bit is set when the corresponding bits are not equal in either operand, and an output bit is reset if the corresponding bits in both operands are identical.

## See Also

[bitset class template](#), [operator& function template](#), [operator|\\_function template](#), [<cstdint>](#), xor keyword

## Name

operator >>function template — Reads a bitset

## Synopsis

```
template <typename charT, typename traits, size_t N>
    basic_istream<charT, traits>& operator>>(basic_istream<charT,
    traits>&
                                                    in, const bitset<N>&
    x);
```

The >> operator reads a bitset from an input stream. It extracts up to N characters and constructs a bitset object using the same format as the string constructor.

Only '0' and '1' characters are extracted. Input stops when it reaches any other character (without extracting that other character).

## See Also

[bitset class template](#), [operator<< function template, <istream>, <cstdintdef>](#)



## Name

operator<< function template — Writes a bitset

## Synopsis

```
template <typename charT, typename traits, size_t N>
    basic_ostream<charT, traits>& operator<<(basic_ostream<charT,
    traits>& in,
                                           const bitset<N>& x);
```

The << operator writes a bitset on an output stream, using the same format as the to\_string member function.

## See Also

[bitset class template](#), [operator>> function template](#), [<cstdlib>](#), [<ostream>](#)

## **<cassert>**

The `<cassert>` header (from the C standard `<assert.h>` header) declares the `assert` macro. The `<cassert>` header is unique in that you can `#include` it multiple times to obtain different effects (depending on whether the `NDEBUG` macro is defined at the time of `#include <cassert>`).

Assertions are checked at runtime. You can use templates to craft compile-time assertions. See [Appendix B](#) for information about the Boost project, which supports compile-time assertions.

Instead of assertions, consider using exceptions, which offer more flexibility and control.

## Name

assert macro — Checks an assertion at runtime

## Synopsis

```
void assert(int expression)
```



If enabled, the `assert` macro ensures the expression is true (nonzero). If so, nothing happens, and execution continues normally. If the expression evaluates to 0, `assert` prints a message to the standard error file and calls `abort`. The format of the message is implementation-defined, but it includes a textual representation of expression and the filename and line number where the `assert` call appears (that is, the values of the `__FILE__` and `__LINE__` macros).

If disabled, the `assert` macro does not evaluate expression and has no effect.

## See Also

`abort` function in [<cstdlib>](#), `throw` keyword

## Name

NDEBUG macro — Enables or disables compilation of assertions

## Synopsis

```
#define NDEBUG  
#include <cassert>
```

The NDEBUG macro is not defined by `<cassert>` or anywhere else in the standard C++ library. Instead, you can define the macro before including the `<cassert>` header to disable the `assert` macro.

In one source file, you can define and undefine NDEBUG multiple times, each time followed by `#include <cassert>`, to enable or disable the `assert` macro multiple times in the same source file.

## See Also

`assert` macro

## <cctype>

The <cctype> header (from the C standard <ctype.h> header) declares a number of functions for testing and mapping narrow character types. For working with wide characters, see <cwctype>.

All the functions take `int` parameters, but the value of the parameter must be an unsigned `char`. Most programs work with ordinary `char`, so you must cast the parameters and some of the return types:

```
char c;
if (std::isalnum(static_cast<unsigned char>(c)))
...
c = static_cast<char>(std::tolower(static_cast<unsigned char>
(c)));
```

The only other value that is permitted is EOF.



These functions get their information from the current locale, as set by calling `setlocale`. The "C" locale is the only one with behavior that is defined in the standard; all other locales can define these functions to include or exclude different characters. Even in the "C" locale, the behavior of some functions depends on the execution character set (see [Chapter 1](#)). One requirement for all locales is that `isalpha`, `iscntrl`, `isdigit`, `ispunct`, and the space character (' ') are mutually exclusive. See <locale> for information about the `setlocale` function.

See the <locale> header for a more flexible (albeit more complicated) mechanism for testing and transforming characters. Each of the functions in this section has a corresponding function in <locale> that takes an explicit locale argument. Also, the `ctype` facet supports similar functionality.

## Name

isalnum function — Tests for an alphanumeric character

## Synopsis

```
int isalnum(int c)
```

The `isalnum` function returns true (nonzero) if `c` is alphanumeric—that is, it returns `isalpha(c) || isnumeric(c)`.

## Name

isalpha function — Tests for an alphabetic character

## Synopsis

```
int isalpha(int c)
```

The `isalpha` function returns true (nonzero) if `c` is alphabetic. In the "C" locale, this includes only the characters for which `islower(c)` or `isupper(c)` is true. For other locales, other characters might be alphabetic.

## Name

isctrnl function — Tests for a control character

## Synopsis

```
int isctrnl(int c)
```

The `isctrnl` function returns true (nonzero) if `c` is a control character. The set of control characters depends on the locale and character set. In the 7-bit ASCII character set, the control characters are `'\0'`-`'\x1F'` and `'\x7F'`; other implementations might have different control characters. By definition, a control character is any character that is not a printable character (`isprint`).



## Name

isdigit function — Tests for a digit character

## Synopsis

```
int isdigit(int c)
```

The `isdigit` function returns true (nonzero) if `c` is a decimal digit character—that is, '0'-'9'—regardless of locale.

## Name

isgraph function — Tests for a graphic character

## Synopsis

```
int isgraph(int c)
```

The `isgraph` function returns true (nonzero) if `c` is any printing character (`isprint`) except space (' '). The set of printing characters varies with locale and character set.

## Name

islower function — Tests for a lowercase letter

## Synopsis

```
int islower(int c)
```

The `islower` function returns true (nonzero) if `c` is a lowercase letter. In the "C" locale, only the characters 'a'-'z' are lowercase; different locales can define other lowercase characters.

## Name

isprint function — Tests for a printable character

## Synopsis

```
int isprint(int c)
```

The `isprint` function returns true (nonzero) if `c` is a printing character, including space (' '), according to the locale and character set. Informally, a printing character occupies space on a display device.

## Name

ispunct function — Tests for a punctuation character

## Synopsis

```
int ispunct(int c)
```

The `ispunct` function returns true (nonzero) for a punctuation character, that is, any printable character (`isprint`) other than space (' ') and alphanumeric characters (`isalnum`).

## Name

isspace function — Tests for a white space character

## Synopsis

```
int isspace(int c)
```

The `isspace` function returns true (nonzero) if `c` is a whitespace character. In the "C" locale, the space (' '), form feed ('\f'), newline ('\n'), carriage return ('\r'), horizontal tab ('\t'), and vertical tab ('\v') characters are whitespace, but backspace ('\b') is not. Different locales can define other whitespace characters.

## Name

isupper function — Tests for an uppercase letter

## Synopsis

```
int isupper(int c)
```

The `isupper` function returns true (nonzero) if `c` is an uppercase letter. In the "C" locale, only the characters 'A'-'Z' are uppercase; different locales can define other uppercase characters

## Name

isxdigit function — Tests for a hexadecimal digit character

## Synopsis

```
int isxdigit(int c)
```

The `isxdigit` function returns true (nonzero) if `c` is any hexadecimal digit character—that is, '0'-'9', 'a'-'f', or 'A'-'F'—regardless of locale.



## Name

tolower function — Converts a character to lowercase

## Synopsis

```
int tolower(int c)
```

The tolower function converts uppercase characters to lowercase. If `c` is uppercase (that is, `isupper(c)` returns true), tolower returns the corresponding lowercase character (for which `islower` returns true) in the current locale, if there is such a character. Otherwise, it returns `c`.

## Name

`toupper` function — Converts a character to uppercase

## Synopsis

```
int toupper(int c)
```

The `toupper` function converts lowercase characters to uppercase. If `c` is lowercase (that is, `islower(c)` returns true), `toupper` returns the corresponding uppercase character (for which `isupper` returns true) in the current locale, if there is such a character. Otherwise, it returns `c`.

## **<cerrno>**

The `<cerrno>` header (from the C standard `<errno.h>` header) declares several macros related to error-handling in the standard library, including the `errno` object.

## Name

EDOM macro — Code for a math domain error

## Synopsis

```
int EDOM
```

Standard library functions that report domain errors set `errno` to `EDOM`. A domain error occurs when the domain of an argument is out of range, such as when asking for the square root of a negative number.

The `EDOM` macro expands to a nonzero integer constant. The value of `EDOM` is implementation-defined.

## Name

EILSEQ macro — Code for error in a multibyte character sequence

## Synopsis

```
int EILSEQ
```

Standard library functions that report errors in multibyte character sequences set `errno` to `EILSEQ`. For example, passing an invalid wide character to the `wcrtomb` function results in `EILSEQ`.

The `EILSEQ` macro expands to a nonzero integer constant. The value of `EILSEQ` is implementation-defined.

Note that `EILSEQ` is not mentioned in the C++ standard, but because C++ includes the C standard library (Amendment 1), and `EILSEQ` is part of the C standard, it is part of the C++ standard.

## Name

ERANGE macro — Code for range error

## Synopsis

int **ERANGE**

Standard library functions that report range errors set `errno` to `ERANGE`. A range error occurs when the result of a function is out of range, such as when there is overflow from the `pow` function.

The `ERANGE` macro expands to a nonzero integer constant. The value of `ERANGE` is implementation-defined.

# Name

`errno` macro — Global error code object

# Synopsis

```
int& errno
```

The `errno` macro expands into an `int` lvalue. Standard library functions can store an error code in `errno` and return an error status to the caller. You can also store a value in `errno` (e.g., resetting the error code to `0`).

When a program starts, `errno` is initially `0`. No library function resets `errno` to `0`. Any library function might set `errno` to a nonzero value, even if it is not documented to do so. Therefore, the only time it is safe to check `errno` is after a library function returns an error status and is documented to set `errno` in that case.



The C++ standard is not explicit as to whether `errno` is truly a macro (versus a variable). The intent of the standard committee is to define `errno` as a macro, so do not use `std::errno`, and if your library declares `errno` as a variable in the `std::` namespace, you can define your own macro:

```
#define errno (::std::errno)
```

In a multithreaded environment, a library implementation typically ensures that each thread gets a separate copy of `errno`. Such considerations fall outside the realm of the C++ standard. Consult your compiler and library documentation for details.

# See Also

perror in [<stdio.h>](#), strerror in [<string.h>](#)



## <float>

The <float> header is the C++ version of the C standard <float.h> header. It defines parameters that characterize floating-point types in the same way <limits> does for the integral types. The native C++ header, <limits>, defines the same information (and more) using templates instead of macros.

There are three sets of macros, each describing a different fundamental type. For each type, the corresponding set of macros has a common prefix: float (FLT\_), double (DBL\_), and long double (LDBL\_). Each set characterizes a floating-point value as a sign, a significand (sometimes called the *mantissa*), a base, and an exponent:

$$x = \text{sign} \times \text{significand} \times \text{base}^{\text{exponent}}$$

In everyday arithmetic, we are used to working with a *base* of 10. (The base is also called the *radix*.) The most common bases for computer arithmetic, however, are 16 and 2. Many modern workstations use the IEC 60559 (IEEE 754) standard for floating-point arithmetic, which uses a base of 2.

The *significand* is a string of digits in the given base. There is an implied radix point at the start of the significand so the value of the significand is always less than 1. (A *radix point* is the generalization of a decimal point for any radix.)

A floating-point value is *normalized* if the first digit of its significand is nonzero, or if the entire value is 0. A value that is not normalized is *denormalized*.

The *precision* of a floating-point type is the maximum number of places in the significand. The *range* of a floating-point type depends primarily on the minimum and maximum values for the exponent.



The value returned by each macro is implementation-defined because all the fundamental types are implementation-defined. The standard mandates minimum decimal precision and range for each floating-point type. In this section, the descriptions for the decimal characteristics of each type include the minimum value set by the standard.

Only `FLT_RADIX` expands to a constant expression. All other macros in `<float>` expand to numeric expressions, but the values might not be compile-time constants.

## Name

DBL\_DIG macro — Decimal precision

## Synopsis

```
int DBL_DIG
```

Number of significant decimal digits that can be stored in a double. The value is always at least 10.

## Name

DBL\_EPSILON macro — Limit of accuracy

## Synopsis

```
double DBL_EPSILON
```

The difference between 1 and the smallest value greater than 1 that can be stored in a double. The value is always less than or equal to  $10^{-9}$ .

## Name

DBL\_MANT\_DIG macro — Significand precision

## Synopsis

```
int DBL_MANT_DIG
```

Number of FLT\_RADIX digits in the significand of a double.

## Name

DBL\_MAX macro — Maximum finite value

## Synopsis

```
double DBL_MAX
```

Maximum finite double value. The value is always at least  $10^{37}$ .

## Name

DBL\_MAX\_10\_EXP macro — Maximum decimal exponent

## Synopsis

```
int DBL_MAX_10_EXP
```

Maximum decimal exponent for a finite double. The value is always at least 37.

## Name

DBL\_MAX\_EXP macro — Maximum exponent

## Synopsis

```
int DBL_MAX_EXP
```

Maximum exponent of a FLT\_RADIX base for a finite double.



## Name

DBL\_MIN macro — Minimum positive value

## Synopsis

double **DBL\_MIN**

Minimum normalized, positive double value. The value is always less than or equal to  $10^{-37}$ .

## Name

DBL\_MIN\_10\_EXP macro — Minimum decimal exponent

## Synopsis

```
int DBL_MIN_10_EXP
```

Minimum negative decimal exponent for a normalized double. The value is always less than or equal to -37.

## Name

DBL\_MIN\_EXP macro — Minimum exponent

## Synopsis

```
int DBL_MIN_EXP
```

Minimum negative exponent of a FLT\_RADIX base for a normalized double.

## **Name**

FLT\_DIG macro — Decimal precision

## **Synopsis**

```
int FLT_DIG
```

Number of significant decimal digits that can be stored in a float. The value is always at least 6.

## Name

FLT\_EPSILON macro — Limit of accuracy

## Synopsis

```
float FLT_EPSILON
```

The difference between 1 and the smallest value greater than 1 that can be stored in a `float`. The value is always less than or equal to  $10^{-5}$ .

## **Name**

FLT\_MANT\_DIG macro — Significant precision

## **Synopsis**

```
int FLT_MANT_DIG
```

Number of FLT\_RADIX digits in the significand of a float.

## Name

FLT\_MAX macro — Maximum finite value

## Synopsis

```
float FLT_MAX
```

Maximum finite float value. The value is always at least  $10^{37}$ .

## Name

`FLT_MAX_10_EXP` macro — Maximum decimal exponent

## Synopsis

```
int FLT_MAX_10_EXP
```

Maximum decimal exponent for a finite float. The value is always at least 37.



## Name

FLT\_MAX\_EXP macro — Maximum exponent

## Synopsis

```
int FLT_MAX_EXP
```

Maximum exponent of a FLT\_RADIX base for a finite float.

## Name

FLT\_MIN macro — Minimum positive value

## Synopsis

```
float FLT_MIN
```

Minimum normalized, positive float value. The value is always less than or equal to  $10^{-37}$ .

## Name

FLT\_MIN\_10\_EXP macro — Minimum decimal exponent

## Synopsis

```
int FLT_MIN_10_EXP
```

Minimum negative decimal exponent for a normalized float. The value is always less than or equal to -37.

## Name

FLT\_MIN\_EXP macro — Minimum exponent

## Synopsis

```
int FLT_MIN_EXP
```

Minimum negative exponent of a FLT\_RADIX base for a normalized float.

## Name

FLT\_RADIX macro — Implementation base

## Synopsis

```
int FLT_RADIX
```

The FLT\_RADIX macro is an integer constant that specifies the radix, or base, for the floating-point implementation. For example, IEC 60559 (IEEE 754) has a FLT\_RADIX of 2.

# Name

FLT\_ROUND macro — Rounding mode

# Synopsis

```
int FLT_ROUND
```

The FLT\_ROUND macro specifies how the implementation rounds floating-point numbers. [Table 13-2](#) lists the possible values as defined in the C++ standard. An implementation can define additional values with other meanings.

*Table 13-2. Floating-point rounding mode*

Value	Description
-1	Indeterminable
0	Rounds toward 0
1	Rounds to nearest
2	Rounds up (toward positive infinity)
3	Rounds down (toward negative infinity)

## Name

LDBL\_DIG macro — Decimal precision

## Synopsis

```
int LDBL_DIG
```

Number of significant decimal digits that can be stored in a long double. The value is always at least 10.

## Name

LDBL\_EPSILON macro — Limit of accuracy

## Synopsis

long double **LDBL\_EPSILON**

The difference between 1 and the smallest value greater than 1 that can be stored in a long double. The value is always less than or equal to  $10^{-9}$ .



## Name

LDBL\_MANT\_DIG macro — Significand precision

## Synopsis

```
int LDBL_MANT_DIG
```

Number of FLT\_RADIX digits in the significand of a long double.

## Name

LDBL\_MAX macro — Maximum finite value

## Synopsis

long double **LDBL\_MAX**

Maximum finite long double value. The value is always at least  $10^{37}$ .

## Name

LDBL\_MAX\_10\_EXP macro — Maximum decimal exponent

## Synopsis

```
int LDBL_MAX_10_EXP
```

Maximum decimal exponent for a finite long double. The value is always at least 37.

## Name

LDBL\_MAX\_EXP macro — Maximum exponent

## Synopsis

```
int LDBL_MAX_EXP
```

Maximum exponent of a FLT\_RADIX base for a finite double.

## Name

LDBL\_MIN macro — Minimum positive value

## Synopsis

long double **LDBL\_MIN**

Minimum normalized, positive long double value. The value is always less than or equal to  $10^{-37}$ .

## Name

LDBL\_MIN\_10\_EXP macro — Minimum decimal exponent

## Synopsis

```
int LDBL_MIN_10_EXP
```

Minimum negative decimal exponent for a normalized long double. The value is always less than or equal to -37.

## Name

LDBL\_MIN\_EXP macro — Minimum exponent

## Synopsis

```
int LDBL_MIN_EXP
```

Minimum negative exponent of a FLT\_RADIX base for a normalized long double.

## **<ciso646>**

The `<ciso646>` header (from the C standard `<iso646.h>` header) does nothing in C++. The C header defines a small number of macros, such as `and` and `for` and `&&`, but these macros are all reserved keywords in C++. The `<ciso646>` header exists only for the sake of completeness: every header in the C standard has an equivalent in C++.



Not all C++ compilers correctly implement the alternative tokens such as `and`. These compilers might use `<ciso646>` to declare these keywords as macros. For maximum portability to these nonconforming compilers, include the `<ciso646>` header when you want to use any alternative token keywords. See [Chapter 1](#) for more information about the alternative tokens.



## **<climits>**

The `<climits>` header (from the C standard `<limits.h>` header) defines parameters that characterize integral types in the same way `<cfloat>` does for the floating point types. The native C++ header, `<limits>`, defines the same information (and more) using templates instead of macros.



The types used in the descriptions of the `_MIN` and `_MAX` macros are meant as reminders and are not descriptive of the actual types of the macro expansions. The actual types are implementation-defined and can be any integral type that would be the result of normal integral promotions for the corresponding type—e.g., if unsigned char is promoted to `int`, `UCHAR_MAX` might have type `int`.

All of the macros in `<climits>` expand to constant expressions.

## Name

CHAR\_BIT macro — Bits per character

## Synopsis

```
int CHAR_BIT
```

Number of bits per character. The value is always at least 8.

## Name

CHAR\_MAX macro — Maximum char value

## Synopsis

char CHAR\_MAX

Maximum value for the char type. (Remember that the char type is the same as either signed char or unsigned char, so CHAR\_MAX has the same value as SCHAR\_MAX or UCHAR\_MAX.)

## See Also

WCHAR\_MAX in [<wchar>](#)

## Name

CHAR\_MIN macro — Minimum char value

## Synopsis

char CHAR\_MIN

Minimum value for the char type (the same value as SCHAR\_MIN or 0).

## See Also

WCHAR\_MIN in [<wchar>](#)

## Name

INT\_MAX macro — Maximum int value

## Synopsis

```
int INT_MAX
```

Maximum value for the `int` type. The value is always at least 32,767.

## Name

INT\_MIN macro — Minimum int value

## Synopsis

```
int INT_MIN
```

Minimum value for the `int` type. The value is always less than or equal to `-32,767`.

## Name

LONG\_MAX macro — Maximum long value

## Synopsis

```
long int LONG_MAX
```

Maximum value for the `long int` type. The value is always at least 2,147,483,647.

## Name

LONG\_MIN macro — Minimum long value

## Synopsis

```
long int LONG_MIN
```

Minimum value for the long int type. The value is always less than or equal to -2,147,483,647.



## Name

`MB_LEN_MAX` macro — Maximum bytes in a multibyte character

## Synopsis

```
int MB_LEN_MAX
```

Maximum number of bytes in any multibyte character, in any locale. The value is always at least 1.

## Name

SCHAR\_MAX macro — Maximum signed char value

## Synopsis

signed char **SCHAR\_MAX**

Maximum value for the signed char type. The value is always at least 127.

## Name

SCHAR\_MIN macro — Minimum signed char value

## Synopsis

signed char **SCHAR\_MIN**

Minimum value for the signed char type. The value is always less than or equal to -127.

## Name

SHRT\_MAX macro — Maximum short value

## Synopsis

short SHRT\_MAX

Maximum value for the short type. The value is always at least 32,767.

## Name

SHRT\_MIN macro — Minimum short value

## Synopsis

short **SHRT\_MIN**

Minimum value for the short type. The value is always less than or equal to -32,767.

## Name

UCHAR\_MAX macro — Maximum unsigned char value

## Synopsis

unsigned char **UCHAR\_MAX**

Maximum value for the unsigned char type. The value is always at least 255.

## Name

UINT\_MAX macro — Maximum unsigned int value

## Synopsis

```
unsigned int UINT_MAX
```

Maximum value for the unsigned int type. The value is always at least 65,535.

## Name

ULONG\_MAX macro — Maximum unsigned long value

## Synopsis

```
unsigned long ULONG_MAX
```

Maximum value for the unsigned long type. The value is always at least 4,294,967,295.



## Name

USHRT\_MAX macro — Maximum unsigned short value

## Synopsis

unsigned short **USHRT\_MAX**

Maximum value for the unsigned short type. The value is always at least 65,535.

## **<locale>**

The `<locale>` header (from the C standard `<locale.h>` header) declares types and functions to support internationalization and localization for the C standard library. C++ also offers `<locale>`, which has more flexibility and functionality, but at a cost of complexity and overhead.

The various locale settings are grouped into categories. Each category has a macro (named `LC_ category`) to identify the category in a call to `setlocale`. Ordinarily, you would use `LC_ALL` to set all the categories at once, but you can pick a category from one locale and another category from a different locale.

## **Name**

LC\_ALL macro — All locale categories

## **Synopsis**

```
int LC_ALL
```

The LC\_ALL macro expands to a constant integer, which sets all categories in a call to `setlocale`.

## Name

LC\_COLLATE macro — Collation order category

## Synopsis

```
int LC_COLLATE
```

The LC\_COLLATE macro expands to a constant integer, which sets the collation order category in a call to `setlocale`. The collation order is used by functions such as `strcoll` (in `<cstring>`).

## Name

LC\_CTYPE macro — Character type category

## Synopsis

```
int LC_CTYPE
```

The LC\_CTYPE macro expands to a constant integer, which sets the character type category in a call to `setlocale`. The `<cctype>` functions, such as `isalpha`, use character type information.

## Name

LC\_MONETARY macro — Monetary formatting category

## Synopsis

```
int LC_MONETARY
```

The LC\_MONETARY macro expands to a constant integer, which sets the monetary-formatting category in a call to `setlocale`. Call `localeconv` to retrieve this information.

## Name

LC\_NUMERIC macro — Numeric formatting category

## Synopsis

```
int LC_NUMERIC
```

The LC\_NUMERIC macro expands to a constant integer, which sets the numeric-formatting category in a call to `setlocale`. I/O functions such as `printf` (in `<cstdio>`) use this information to format numbers. Call `localeconv` to retrieve this information.

## Name

LC\_TIME macro — Time formatting category

## Synopsis

```
int LC_TIME
```

The LC\_TIME macro expands to a constant integer, which sets the time-formatting category in a call to `setlocale`. The `strftime` (in `<ctime>`) function uses this information.



# Name

lconv structure — Numeric formatting information

## Synopsis

```
struct lconv {
    char *decimal_point;
    char *thousands_sep;
    char *grouping;
    char *int_curr_symbol;
    char *currency_symbol;
    char *mon_decimal_point;
    char *mon_thousands_sep;
    char *mon_grouping;
    char *positive_sign;
    char *negative_sign;
    char int_frac_digits;
    char frac_digits;
    char p_cs_precedes;
    char p_sep_by_space;
    char n_cs_precedes;
    char n_sep_by_space;
    char p_sign_posn;
    char n_sign_posn;
};
```



The `lconv` structure stores information used to format numbers and monetary values. An implementation can add more members to the class and change the order of declaration.

The standard library is responsible for filling an `lconv` object with values that are appropriate for a locale. Do not modify the `lconv` object or its data members.

Each locale defines suitable values for the `lconv` members. The `char`-type members are all nonnegative integers, in which `CHAR_MAX` means the information is not available in the current locale. For the `char*` members, an empty string

means the information is not available. The "C" locale uses "." for `decimal_point`, an empty string ("") for all other `char*` members, and `CHAR_MAX` for all `char` members. All strings are null-terminated.

The `localeconv` function returns a pointer to the current locale's `lconv` object.

The following are descriptions of the `lconv` members:

`char*` **`currency_symbol`**

The currency symbol for the current locale (e.g., "\$").

`char*` **`decimal_point`**

The decimal point symbol for the current locale. This member is unique in that it cannot be an empty string. The default is ".".

`char` **`frac_digits`**

The number of digits to appear after the decimal point in monetary formatting.

`char*` **`grouping`**

A string that is interpreted as a series of integers, in which each integer is the number of digits in successive groups of digits for nonmonetary formatting. The value '\0' means to repeat the last grouping for the rest of the value. The value `CHAR_MAX` means not to group the remaining values. Any other value is the size of a digit group. The first character in the string specifies the size of the rightmost group of digits, the second character in the string specifies the size of the next (moving to the left) group of digits, and so forth. Digit groups are separated by `thousands_sep`.

A common value is "\3", which means to format digits in groups of three (e.g., "1,234,567").

char\* **int\_curr\_symbol**

A four-character string, in which the first three characters are the international currency symbol (according to ISO standard 4217:1987), and the fourth character is the separator between the currency symbol and the number. For example, the symbol for United States Dollars is USD. If the locale uses a space as the separator, `int_curr_symbol` would be "USD ".

char **int\_frac\_digits**

The number of digits to appear after the decimal point in an internationally-formatted monetary amount.

char\* **mon\_decimal\_point**

The monetary decimal point.

char\* **mon\_grouping**

The monetary grouping. (This works the same as `grouping`, except groups are separated by `mon_thousands_sep`.)

char\* **mon\_thousands\_sep**

The monetary thousands separator. (This works the same as `thousands_sep` in monetary groups, as specified by `mon_grouping`.)

char **n\_cs\_precedes**

Equal to 1 if the currency symbol precedes the amount when formatting a negative monetary value. Equal to 0 if the symbol follows the value.

char **n\_sep\_by\_space**

Equal to 1 if the currency symbol is separated from a negative value by a space. Equal to 0 if there is no space.

char **n\_sign\_posn**

The position of the sign for a negative monetary value. [Table 13-3](#) lists all the position values.

*Table 13-3. Position values for `n_sign_posn` and `p_sign_posn`*

Value	Position
0	Parentheses surround the value and the currency symbol.
1	The sign precedes the value and the currency symbol.
2	The sign follows the value and the currency symbol.
3	The sign appears immediately before the currency symbol.
4	The sign appears immediately after the currency symbol.

`char* negative_sign`

Marker for a negative monetary value (e.g., "-").

`char p_cs_precedes`

Equal to 1 if the currency symbol precedes the amount when formatting a nonnegative monetary value. Equal to 0 if the symbol follows the value.

`char p_sep_by_space`

Equal to 1 if the currency symbol is separated from a nonnegative value by a space. Equal to 0 if there is no space.

`char p_sign_posn`

The position of the sign for a nonnegative monetary value. [Table 13-3](#) lists all the position values.

`char* positive_sign`

Marker for a nonnegative monetary value.

char\* **thousands\_sep**

Thousands separator (e.g., ", "), which is used in digit groups, as specified by grouping.

## Name

localeconv function — Retrieves numeric-formatting information

## Synopsis

```
lconv* localeconv( );
```

The `localeconv` function returns a pointer to the current locale's `lconv` object.

### Warning

Do not modify the `lconv` object. A call to `localeconv` might overwrite the contents of the object returned from an earlier call (or simultaneous call in a multithreaded program). Calls to `setlocale` for `LC_ALL`, `LC_NUMERIC`, or `LC_MONETARY` categories might also overwrite the contents of the `lconv` object.

## Name

NULL macro — Null pointer constant

## Synopsis

```
#define NULL . . .
```

The NULL macro expands to a null pointer constant. It is defined in several C headers. See its description in `<stddef>` for details.

# Name

setlocale function — Sets or queries locale

# Synopsis

```
char* setlocale(int category, const char* locale)
```

The `setlocale` function sets the locale for a specific category, which you must specify using one of the `LC_` macros. Use `LC_ALL` to set all categories to the same locale.



The `locale` parameter is the name of the locale. The default for all categories is the "C" locale. The empty string ("") is an implementation-defined native locale. The implementation can define other possible values for `locale`.

To query the current locale, pass a null pointer as the `locale`. Note that each category might have a different locale, so, when querying for `LC_ALL`, the return value might contain multiple locale names.

The return value is a pointer to a string that contains the new locale (or current locale if you are querying with a null `locale` parameter) for the specified category. If the locale cannot be set, a null pointer is returned.

## Warning

Do not modify the string returned from `setlocale`. A call to `setlocale` might overwrite the contents of the string returned from an earlier call (or simultaneous call in a multithreaded program).



## <cmath>

The <cmath> header declares a number of mathematical functions (from the C standard <math.h>). In addition to the standard C function, most functions have overloaded versions for different parameter types; each function's syntax shows all the overloaded versions.



If an argument is out of range, a domain error occurs. The function sets `errno` to `EDOM` and returns an error value. The value is defined by the implementation, so the only portable way to test for a domain error is to check `errno`. If the function's result is an overflow, a range error occurs. The function returns `HUGE_VAL` and sets `errno` to `ERANGE`. If underflow occurs, the function returns `0` and may or may not set `errno` to `ERANGE`. (See <cerrno> for more information about `errno`.)

### **Not e**

`HUGE_VAL` is defined to be a `double`, and the C++ standard does not define a suitable value for the `float` and `long double` versions of the math functions. If you are using a system that has infinity as an explicit floating-point value (such as IEC 60559/IEEE 754, which is found on PCs, Macintoshes, and modern workstations), the overloaded versions of a function probably return infinity for overflow, so there is no problem with the `float` and `long double` versions of the functions. For maximum portability, however, use only the `double` versions of the math functions.

All the trigonometric functions use radians. The descriptions of these functions use the common

mathematical notation for ranges of values.  $[x, y)$  represents all values  $z$  such that  $x \leq z < y$ —that is, the square bracket denotes an inclusive endpoint of a range, and the parenthesis denotes an exclusive endpoint of a range.

Several other headers in the standard library declare additional mathematical functions:

`<cmath>`

Declares macros for the limits of floating-point types

`<climits>`

Declares macros for the limits of integer types

`<complex>`

Declares types and functions for working with complex numbers

`<cstdlib>`

Declares integer absolute value functions and functions that compute a quotient and remainder in a single operation

`<limits>`

Declares the `numeric_limits` class template for the limits of the numerical types—e.g., the largest float, the precision of double, and so on

`<numeric>`

Declares generic numerical algorithms

`<valarray>`

Declares types and functions for computation with arrays of numbers

## Name

abs function — Computes absolute value

## Synopsis

```
float abs(float x)
double abs(double x)
long double abs(long double x)
```

The abs function returns the absolute value of its argument: if  $x < 0$ , it returns  $-x$ ; otherwise, it returns  $x$ .

The abs function in `<cmath>` is the same as fabs. The `<cstdlib>` header declares integer versions of the abs function.

## See Also

[fabs function](#), abs function in [<cstdlib>](#)

## Name

acos function — Computes inverse cosine

## Synopsis

```
float acos(float x)  
double acos(double x)  
long double acos(long double x)
```

The `acos` function returns the inverse cosine of its argument. The parameter `x` must be in the range  $[-1, 1]$ , or a domain error occurs. The return value is in the range  $[0, \pi]$ .

## Name

asin function — Computes inverse sine

## Synopsis

```
float asin(float x)  
double asin(double x)  
long double asin(long double x)
```

The `asin` function returns the inverse sine of its argument. The parameter `x` must be in the range  $[-1, 1]$ , or a domain error occurs. The return value is in the range  $[-\pi/2, \pi/2]$ .

## Name

atan function — Computes inverse tangent

## Synopsis

```
float atan(float x)  
double atan(double x)  
long double atan(long double x)
```

The atan function returns the inverse tangent of its argument. The return value is in the range  $[-\pi/2, \pi/2]$ .

## Name

atan2 function — Computes inverse tangent

## Synopsis

```
float atan2(float y, float x)  
double atan2(double y, double x)  
long double atan2(long double y, long double x)
```



The `atan2` function returns the inverse tangent of  $y/x$  using the sign of both numbers to determine the quadrant for the return value. It correctly handles the case in which  $x$  is  $0$ . (That is, it returns  $\pi/2$  times the sign of  $y$  for nonzero  $y$ ; if  $y$  is  $0$ , the result is implementation-defined and might be a range error). The return value is in the range  $[-\pi, \pi]$ .

## Name

ceil function — Computes ceiling

## Synopsis

```
float ceil(float x)  
double ceil(double x)  
long double ceil(long double x)
```

The `ceil` function returns the smallest integer that is greater than or equal to `x`.

## See Also

[floor function](#)



## Name

cos function — Computes cosine

## Synopsis

```
float cos(float x)  
double cos(double x)  
long double cos(long double x)
```

The `cos` function returns the cosine of its argument, in radians. The return value is in the range `[-1, 1]`.

## Name

cosh function — Computes hyperbolic cosine

## Synopsis

```
float cosh(float x)  
double cosh(double x)  
long double cosh(long double x)
```

The cosh function returns the hyperbolic cosine of its argument. Note that `<cmath>` has no inverse hyperbolic trigonometric functions; the Boost project fills that gap. See [Appendix B](#) for information about Boost.

## Name

exp function — Computes exponential

## Synopsis

```
float exp(float x)  
double exp(double x)  
long double exp(long double x)
```

The exp function returns  $e^x$ . If  $x$  is too large, a range error occurs.

## See Also

[log function](#), [pow function](#)

## Name

fabs function — Computes absolute value

## Synopsis

```
float fabs(float x)  
double fabs(double x)  
long double fabs(long double x)
```

The fabs function returns the absolute value of its argument: if  $x < 0$ , it returns  $-x$ ; otherwise, it returns  $x$ .

The fabs function is the same as abs for floating-point numbers. It exists only for compatibility with C.

## See Also

[abs function](#), abs function in [<cstdlib>](#)

## Name

floor function — Computes floor

## Synopsis

```
float floor(float x)  
double floor(double x)  
long double floor(long double x)
```

The floor function returns the largest integer that is less than or equal to  $x$ .

## See Also

[ceil function](#)

# Name

fmod function — Computes modulus

## Synopsis

```
float fmod(float x, float y)  
double fmod(double x, double y)  
long double fmod(long double x, long double y)
```



The `fmod` function returns the floating-point remainder of dividing  $x$  by  $y$ . If  $y$  is  $0$ , the behavior is implementation-defined: the return value might be  $0$ , or a domain error can occur. If  $y$  is nonzero, the return value is  $x - k \times y$  for some integer  $k$ , such that the result has the same sign as  $x$  and an absolute value less than the absolute value of  $y$ .

## Name

`frexp` function — Computes binary fraction and exponent

## Synopsis

```
float frexp(float x, int* exp)
double frexp(double x, int* exp)
long double frexp(long double x, int* exp)
```

The `frexp` function separates a floating-point number into a fraction and an exponent (with a base of 2) such that  $x = frac \times 2^e$ , in which *frac* is in the range  $[1/2, 1)$  or is 0 if *x* is 0. The exponent, *e*, is stored in `*exp`. The return value is *frac*. If *x* is 0, the return value and `*exp` are 0.

## See Also

[ldexp function](#), [modf function](#)

## Name

HUGE\_VAL macro — Range error value

## Synopsis

double HUGE\_VAL



When an overflow occurs, most functions set `errno` to `ERANGE` and return `HUGE_VAL` with the correct sign of the result. The exact value of `HUGE_VAL` is implementation-defined and is not necessarily a compile-time constant. It might even be a value that can be returned as a valid result from the function. In that case, the only way to discover whether an overflow occurred is to test `errno`, as shown in [Example 13-5](#).

## Example

*Example 13-5. Computing a logarithm to any base*

```
// Return the logarithm of x to the base n.
template<typename T>
T logn(T x, T n)
{
    errno = 0;
    T logx = log(x);
    if (errno == ERANGE)
        return logx;    // Should be HUGE_VAL
    else if (errno != 0)
        return logx;    // Implementation defined
    T logn = log(n);
    if (errno == ERANGE)
        return logn;    // Should be HUGE_VAL
    else if (errno != 0)
        return logn;    // Implementation defined
    if (logn == 0) {
        errno = EDOM;
        return 0;
    }
}
```



```
}  
  return logx / logn;  
}
```

## See Also

[<cerrno>](#)

## Name

ldexp function — Makes floating point from binary fraction and exponent

## Synopsis

```
float ldexp(float frac, int exp)
double ldexp(double frac, int exp)
long double ldexp(long double frac, int exp)
```

The ldexp function returns a floating-point number that it constructs from a fractional part and an exponent (base 2). The return value is  $\text{frac} \times 2^{\text{exp}}$ .

## See Also

[frexp function](#), [modf function](#)

## Name

log function — Computes natural logarithm

## Synopsis

```
float log(float x)  
double log(double x)  
long double log(long double x)
```



The `log` function returns the natural (base  $e$ ) logarithm of its argument. A domain error occurs if  $x$  is negative. A range error might occur if  $x$  is  $0$ .

## Name

log10 function — Computes common logarithm

## Synopsis

```
float log10(float x)  
double log10(double x)  
long double log10(long double x)
```



The `log10` function returns the common (base 10) logarithm of its argument. A domain error occurs if `x` is negative. A range error might occur if `x` is 0.

## Name

modf function — Separates integer and fraction parts

## Synopsis

```
float modf(float x, float* iptr)  
double modf(double x, double* iptr)  
long double modf(long double x, long double* iptr)
```

The `modf` function splits a floating-point number into integral and fractional parts. Both parts have the same sign as `x`. The integral part is stored in `*iptr`; the return value is the fractional part.

## See Also

[frexp function](#), [ldexp function](#)

## Name

pow function — Computes power

## Synopsis

```
float pow(float x, float y)
float pow(float x, int y)
double pow(double x, double y)
double pow(double x, int y)
long double pow(long double x, long double y)
long double pow(long double x, int y)
```

The pow function raises x to the y power. If x is negative, and y is not an integral value, a domain error occurs. If x is 0, and y is less than or equal to 0, and the result cannot be represented as a real number, a domain error occurs. A range error can occur if the result is out of range.

## See Also

[exp function](#)

## Name

sin function — Computes sine

## Synopsis

```
float sin(float x)  
double sin(double x)  
long double sin(long double x)
```

The `sin` function returns the sine of its argument, in radians. The return value is in the range `[-1, 1]`.

## Name

sinh function — Computes hyperbolic sine

## Synopsis

```
float sinh(float x)  
double sinh(double x)  
long double sinh(long double x)
```

The `sinh` function returns the hyperbolic sine of its argument. Note that `<cmath>` has no inverse hyperbolic trigonometric functions; the Boost project fills that gap. See [Appendix B](#) for information about Boost.



## Name

sqrt function — Computes square root

## Synopsis

```
float sqrt(float x)  
double sqrt(double x)  
long double sqrt(long double x)
```

The `sqrt` function returns the square root of its argument. If `x` is negative, a domain error occurs. The return value is always positive or `0`.

# Name

tan function — Computes tangent

# Synopsis

```
float tan(float x)
double tan(double x)
long double tan(long double x)
```

The tan function returns the tangent of its argument. The standard does not specify the result when the tangent is undefined (that is, when  $x$  is  $k\pi + \pi/2$  for any integer  $k$ ), but a reasonable result is a range error. Due to the nature of the tangent function, the sign of the return value (HUGE\_VAL) can be positive or negative.

## Name

tanh function — Computes hyperbolic tangent

## Synopsis

```
float tanh(float x)  
double tanh(double x)  
long double tanh(long double x)
```

The tanh function returns the hyperbolic tangent of its argument. Note that `<cmath>` has no inverse hyperbolic trigonometric functions; the Boost project fills that gap. See [Appendix B](#) for information about Boost.

## **<complex>**

The `<complex>` header declares the `complex` class template and specializations for the `float`, `double`, and `long double` types. It also declares mathematical functions that work with complex values.

## Name

abs function template — Computes absolute value

## Synopsis

```
template<typename T> T abs(const complex<T>& z)
```

The `abs` function returns the absolute value (or magnitude) of `z`.

## See Also

[polar function template](#), `abs` function in [<cmath>](#)

## Name

arg function template — Computes argument (angle)

## Synopsis

```
template<typename T> T arg(const complex<T>& z)
```

The `arg` function returns the argument (angle in polar coordinates) of `z`.

## See Also

[polar function template](#)

# Name

complex class template — Complex number template

## Synopsis

```
template<typename T>
class complex {
public:
    typedef T value_type;
    complex(const T& re = T( ), const T& im = T( ));
    complex(const complex& z);
    template<typename X> complex(const complex<X>& z);
    T real( ) const;
    T imag( ) const;
    complex& operator=(const T& x);
    complex& operator+=(const T& x);
    complex& operator-=(const T& x);
    complex& operator*=(const T& x);
    complex& operator/=(const T& x);
    complex& operator=(const complex& z);
    template<typename X>
        complex& operator=(const complex<X>& z);
    template<typename X>
        complex& operator+=(const complex<X>& z);
    template<typename X>
        complex& operator-=(const complex<X>& z);
    template<typename X>
        complex& operator*=(const complex<X>& z);
    template<typename X>
        complex& operator/=(const complex<X>& z);
};
```

The `complex` class template represents a complex number. The `<complex>` header specializes the template for the `float`, `double`, and `long double` types. You can instantiate `complex<>` for any type that behaves in the manner of the fundamental numeric types.

The type definition is a straightforward representation of a complex number. Basic assignment operators are defined as member functions, and arithmetic operators are defined as global functions.

template<typename X> **complex** (const complex<X>& z)

Constructs a `complex<T>` object by copying the members from `z`. Effectively, this converts a `complex` object instantiated for one type to a `complex` object of another type.

T **real** ( ) const

Returns the real part of `*this`.

T **imag** ( ) const

Returns the imaginary part of `*this`.

complex& **operator=** (const T& x)

Assigns `x` to the real part of `*this` and `0` to the imaginary part. Returns `*this`.

complex& **operator+=** (const T& x)

Adds `x` to the real part of `*this`, leaving the imaginary part alone. Returns `*this`.

complex& **operator-=** (const T& x)

Subtracts `x` from the real part of `*this`, leaving the imaginary part alone. Returns `*this`.

complex& **operator\*=** (const T& x)

Multiplies the real and imaginary parts of `*this` by `x`. Returns `*this`.

complex& **operator/=** (const T& x)

Divides the real and imaginary parts of `*this` by `x`. Returns `*this`.

complex& **operator=** (const complex& z)

Assigns the real and imaginary parts of `z` to `*this`. Returns `*this`.



```
template<typename X> complex& operator= (const  
complex<X>& z)
```

Assigns the real and imaginary parts of z to \*this.  
Returns \*this. Note that z and \*this can have different  
template parameter types.

```
template<typename X> complex& operator+= (const  
complex<X>& z)
```

Adds z to \*this. Returns \*this. Note that z and \*this  
can have different template parameter types.

```
template<typename X> complex& operator-= (const  
complex<X>& z)
```

Subtracts z from \*this. Returns \*this. Note that z and  
\*this can have different template parameter types.

```
template<typename X> complex& operator*=( const  
complex<X>& z)
```

Multiplies \*this by z. Returns \*this. Note that z and  
\*this can have different template parameter types.

```
template<typename X> complex& operator/= (const  
complex<X>& z)
```

Divides \*this by z. Returns \*this. Note that z and  
\*this can have different template parameter types.

# Name

`complex<double>` template specialization — Double-precision complex number

# Synopsis

```
template<> class complex<double> {
public:
    typedef double value_type;
    complex(double re = 0.0, double im = 0.0);
    complex(const complex<float>&);
    explicit complex(const complex<long double>&);
    double real( ) const;
    double imag( ) const;
    complex<double>& operator=(double);
    complex<double>& operator+=(double);
    complex<double>& operator-=(double);
    complex<double>& operator*=(double);
    complex<double>& operator/=(double);
    complex<double>& operator=(const complex<double>&);
    template<typename X>
        complex<double>& operator=(const complex<X>&);
    template<typename X>
        complex<double>& operator+=(const complex<X>&);
    template<typename X>
        complex<double>& operator-=(const complex<X>&);
    template<typename X>
        complex<double>& operator*=(const complex<X>&);
    template<typename X>
        complex<double>& operator/=(const complex<X>&);
};
```

The `complex<double>` class is a straightforward specialization of the `complex` class template. It changes the operators to pass double parameters by value instead of by reference, and it adds a new constructor:

```
explicit complex (const complex<long double>& z)
```

Constructs a complex number by copying from `z`. Note that you might lose precision or overflow, so the constructor is explicit.

# Name

`complex<float>` template specialization — Single-precision complex number

# Synopsis

```
template<> class complex<float> {
public:
    typedef float value_type;
    complex(float re = 0.0f, float im = 0.0f);
    explicit complex(const complex<double>&);
    explicit complex(const complex<long double>&);
    float real( ) const;
    float imag( ) const;
    complex<float>& operator=(float);
    complex<float>& operator+=(float);
    complex<float>& operator-=(float);
    complex<float>& operator*=(float);
    complex<float>& operator/=(float);
    complex<float>& operator=(const complex<float>&);
    template<typename X>
        complex<float>& operator=(const complex<X>&);
    template<typename X>
        complex<float>& operator+=(const complex<X>&);
    template<typename X>
        complex<float>& operator-=(const complex<X>&);
    template<typename X>
        complex<float>& operator*=(const complex<X>&);
    template<typename X>
        complex<float>& operator/=(const complex<X>&);
};
```

The `complex<float>` class is a straightforward specialization of the `complex` class template. It changes the operators to pass float parameters by value instead of by reference, and it adds two new constructors:

```
explicit complex (const complex<double>& z),
explicit complex (const complex<long double>& z)
```

Constructs a complex number by copying from `z`. Note that you might lose precision or overflow, so the

constructors are explicit.

# Name

complex<long double> template specialization —  
Extended-precision complex number

# Synopsis

```
template<> class complex<long double> {
public:
    typedef long double value_type;
    complex(long double re = 0.0L, long double im = 0.0L);
    complex(const complex<float>&);
    complex(const complex<double>&);
    long double real( ) const;
    long double imag( ) const;
    complex<long double>&
        operator=(const complex<long double>&);
    complex<long double>& operator= (long double);
    complex<long double>& operator+=(long double);
    complex<long double>& operator-=(long double);
    complex<long double>& operator*=(long double);
    complex<long double>& operator/=(long double);
    template<typename X>
        complex<long double>& operator= (const complex<X>&);
    template<typename X>
        complex<long double>& operator+=(const complex<X>&);
    template<typename X>
        complex<long double>& operator-=(const complex<X>&);
    template<typename X>
        complex<long double>& operator*=(const complex<X>&);
    template<typename X>
        complex<long double>& operator/=(const complex<X>&);
};
```

The `complex<long double>` class is a straightforward specialization of the `complex` class template. It changes the operators to pass long double parameters by value instead of by reference.

## Name

conj function template — Computes conjugate

## Synopsis

```
template<typename T> complex<T> conj(const complex<T>& z)
```

The `conj` function returns the complex conjugate of `z`.

## Name

cos function template — Computes cosine

## Synopsis

```
template<typename T> complex<T> cos(const complex<T>& z)
```

The `cos` function returns the complex cosine of `z`.

## See Also

cos function in [<cmath>](#)

## Name

cosh function template — Computes hyperbolic cosine

## Synopsis

```
template<typename T> complex<T> cosh(const complex<T>& z)
```

The cosh function returns the complex hyperbolic cosine of z.

## See Also

cosh function in [<cmath>](#)



## Name

exp function template — Computes exponential

## Synopsis

```
template<typename T> complex<T> exp(const complex<T>& z)
```

The exp function returns the exponential of  $z$ , that is,  $e^z$ .

## See Also

exp function in [<cmath>](#)

## Name

imag function template — Returns imaginary part

## Synopsis

```
template<typename T> T imag(const complex<T>& z)
```

The `imag` function returns the imaginary part of `z`, that is, `z.imag( )`.

## See Also

`abs` function in [<cmath>](#)

## Name

log function template — Computes natural logarithm

## Synopsis

```
template<typename T> complex<T> log(const complex<T>& z)
```

The `log` function returns the complex natural (base  $e$ ) logarithm of  $z$ . The branch cuts are along the negative real axis, which means the imaginary part of the result is in the range  $[-\pi i, \pi i]$ .

## See Also

log function in [<cmath>](#)

## Name

log10 function template — Computes common logarithm

## Synopsis

```
template<typename T> complex<T> log10(const complex<T>& z)
```

The `log10` function returns the complex common (base 10) logarithm of  $z$ . The branch cuts are along the negative real axis, which means the imaginary part of the result is in the range  $[-\pi i, \pi i]$ .

## See Also

log10 function in [<cmath>](#)

## Name

norm function template — Computes normalized value

## Synopsis

```
template<typename T> T norm(const complex<T>& z)
```

The norm function returns the square of the absolute value of z.

## See Also

[abs function template](#)

## Name

`operator+` function template — Performs unary positive or addition

## Synopsis

```
template<typename T>
    complex<T> operator+(const complex<T>& z);

template<typename T> complex<T>
    operator+(const complex<T>& x, const complex<T>& y);
template<typename T> complex<T>
    operator+(const complex<T>& x, const T& y);
template<typename T> complex<T>
    operator+(const T& x, const complex<T>& y);
```

The unary positive operator returns  $z$ .

The binary addition operator returns the sum of its operands. If either operand is of type  $T$ , the argument is interpreted as the real part, with an imaginary part of  $T( )$  or  $0$ .

## Name

operator- function template — Performs negation or subtraction

## Synopsis

```
template<typename T>
    complex<T> operator-(const complex<T>&);

template<typename T>
    complex<T> operator-(const complex<T>&, const complex<T>&);
template<typename T>
    complex<T> operator-(const complex<T>&, const T&);
template<typename T>
    complex<T> operator-(const T&, const complex<T>&);
```

The unary negation operator returns  $-z$ .

The binary subtraction operator returns the difference of its operands. If either operand is of type  $T$ , the argument is interpreted as the real part, with an imaginary part of  $T(\ )$  or  $0$ .

## Name

`operator*` function template — Performs multiplication

## Synopsis

```
template<typename T>
    complex<T> operator*(const complex<T>&, const complex<T>&);
template<typename T>
    complex<T> operator*(const complex<T>&, const T&);
template<typename T>
    complex<T> operator*(const T&, const complex<T>&);
```

The binary `*` operator performs complex multiplication. If either operand is of type `T`, the argument is interpreted as the real part, with an imaginary part of `T( )` or `0`.



## Name

operator/ function template — Performs division

## Synopsis

```
template<typename T>
    complex<T> operator/(const complex<T>&, const complex<T>&);
template<typename T>
    complex<T> operator/(const complex<T>&, const T&);
template<typename T>
    complex<T> operator/(const T&, const complex<T>&);
```

The binary / operator performs complex division. If either operand is of type T, the argument is interpreted as the real part, with an imaginary part of T( ) or 0. Division by zero results in undefined behavior.

# Name

`operator==` function template — Checks equality

# Synopsis

```
template<typename T>
    bool operator==(const complex<T>&, const complex<T>&);
template<typename T>
    bool operator==(const complex<T>&, const T&);
template<typename T>
    bool operator==(const T&, const complex<T>&);
```

The `==` operator returns `true` if the real and imaginary parts of both values are equal. If either operand is of type `T`, the argument is interpreted as the real part, with an imaginary part of `T( )` or `0`.

## Name

`operator!=` function template — Checks inequality

## Synopsis

```
template<typename T>
    bool operator!=(const complex<T>&, const complex<T>&);
template<typename T>
    bool operator!=(const complex<T>&, const T&);
template<typename T>
    bool operator!=(const T&, const complex<T>&);
```

The `!=` operator returns `true` if the real or imaginary parts are not equal. If either operand is of type `T`, the parameter is interpreted as the real part, with an imaginary part of `T( )` or `0`.

## Name

`operator<<` function template — Writes a complex number

## Synopsis

```
template<typename T, typename charT, typename traits>
    basic_ostream<charT, traits>& operator<<(basic_ostream<charT,
    traits>&,
                                           const complex<T>& z);
```

The `<<` operator prints `z` to the output stream in the form `(x, y)`, in which `x` is the real part, and `y` is the imaginary part. [Example 13-6](#) shows how `z` is formatted. If you want more control over the formatting, you must print the value yourself.

## Example

*Example 13-6. Formatting a complex number*

```
template<class T, class charT, class traits>
std::basic_ostream<charT, traits>&
operator<<(std::basic_ostream<charT, traits>& o,
           const std::complex<T>& x)
{
    std::basic_ostringstream<charT, traits> s;
    s.flags(o.flags( ));
    s.imbue(o.getloc( ));
    s.precision(o.precision( ));
    s << "(" << x.real( ) << ", " << x.imag( ) << ")";
    return o << s.str( );
}
```

## Name

`operator>>` function template — Reads a complex number

## Synopsis

```
template<typename T, typename charT, typename traits>
    basic_istream<charT, traits>& operator>>(basic_istream<charT,
    traits>&,
                                           complex<T>& z);
```

The `>>` operator reads a complex number from an input stream into `z`. The input format can be any of the following:

`x`

The value `x` is the real part, and `T( )` or `0` is the imaginary part.

`(x)`

The value `x` is the real part, and `T( )` or `0` is the imaginary part.

`(x, y)`

The value `x` is the real part, and `y` is the imaginary part.

## Name

polar function template — Converts to polar coordinates

## Synopsis

```
template<typename T>  
complex<T> polar(const T& r, const T& theta)
```

The `polar` function returns a complex object that represents the value given in polar coordinates, in which `r` is the magnitude and `theta` is the angle (in radians). The resulting value has the following real and imaginary parts:

```
real = r * cos(theta)  
imag = r * sin(theta)
```

## See Also

[abs function template](#), [arg function template](#)

# Name

pow function template — Computes power

# Synopsis

```
template<class T>
    complex<T> pow(const complex<T>& x, int y);
template<class T>
    complex<T> pow(const complex<T>& x, const T& y);
template<class T>
    complex<T> pow(const complex<T>& x, const complex<T>& y);
template<class T>
    complex<T> pow(const T& x, const complex<T>& y);
```

The pow function returns the complex power  $xy$ . If  $x$  and  $y$  are both 0, the result is implementation-defined; otherwise, the result is  $\exp(y * \log(x))$ . The branch cuts are along the negative real axis.

# See Also

[exp function template](#), [log function template](#), pow function in [<cmath>](#)

## Name

real function template — Returns real part

## Synopsis

```
template<typename T> T real(const complex<T>& z)
```

The `real` function returns the real part of `z`, that is, `z.real( )`.



## Name

sin function template — Computes sine

## Synopsis

```
template<typename T> complex<T> sin(const complex<T>& z)
```

The `sin` function returns the complex sine of `z`.

## See Also

sin function in [<cmath>](#)

## Name

sinh function template — Computes hyperbolic sine

## Synopsis

```
template<typename T> complex<T> sinh(const complex<T>& z)
```

The sinh function returns the complex hyperbolic sine of z.

## See Also

sinh function in [<cmath>](#)

## Name

sqrt function template — Computes square root

## Synopsis

```
template<typename T> complex<T> sqrt(const complex<T>& z)
```

The `sqrt` function returns the complex square root of `z`. The branch cuts are along the negative real axis. The result always has a nonnegative real part.

## See Also

sqrt function in [<cmath>](#)

## Name

tan function template — Computes tangent

## Synopsis

```
template<typename T> complex<T> tan(const complex<T>& z)
```

The tan function returns the complex tangent of z.

## See Also

tan function in [<cmath>](#)

## Name

tanh function template — Computes hyperbolic tangent

## Synopsis

```
template<typename T> complex<T> tanh(const complex<T>& z)
```

The `tanh` function returns the complex hyperbolic tangent of `z`.

## See Also

tanh function in [<cmath>](#)

## **<csetjmp>**

The <csetjmp> header is the C++ version of the C standard <setjmp.h> header.

### **Warning**

This chapter presents only the most cursory description of this header because its use is limited in a C++ program. Use exceptions instead of the functions in <csetjmp>.

## Name

`jmp_buf` type — Jump buffer

## Synopsis

```
typedef . . . jmp_buf;
```

The `jmp_buf` type is an opaque array type that stores information for the `setjmp` and `longjmp` functions.

## Name

longjmp function — Performs nonlocal goto

## Synopsis

```
void longjmp(jmp_buf env, int val);
```

The `longjmp` function bypasses the normal function return and unwinds the call stack to the point where `setjmp` was called with the same `jmp_buf` environment. When `setjmp` returns to its caller, it returns `val`; if `val` is 0, `setjmp` returns 1.

Calling `longjmp` is similar to throwing an exception that is caught at the point of the `setjmp` call. One important difference, however, is that if any objects on the stack would have been destroyed by throwing an exception, the program's behavior is undefined if you call `longjmp`. This is why you should use exceptions instead of `longjmp`.



## Name

setjmp function — Establishes nonlocal label

## Synopsis

```
int setjmp(jmp_buf env);
```

The `setjmp` function stores the current execution environment in its argument so that the environment can be restored by a call to `longjmp`. The first time `setjmp` is called, it returns 0. When `longjmp` is called, `setjmp` returns the `val` argument that was passed to `longjmp`; that value is guaranteed to be nonzero.

## <csignal>

The <csignal> header is the C++ version of the standard C <signal.h> header. It declares functions and macros related to signal handling.

A *signal* is a condition that can arise during program execution. A signal can originate by an explicit call to `raise` or `abort`, from external sources (such as the user interrupting the program), or from internal events (such as floating-point errors or memory violations). Each signal has a *handler*, which is a function that the C++ library calls when a signal occurs (called *raising* the signal).



Signals are identified by integers. A program can establish different handlers for different signal numbers. You can choose to ignore a signal by using `SIG_IGN` as the signal handler. If the signal is raised, no handler is called, and the program continues. Each signal also has a default handler (`SIG_DFL`). When a program starts, every signal number is initialized with `SIG_IGN` or `SIG_DFL`. The details are implementation-defined.

You can set your own handler for any signal by calling the `signal` function. A handler is a function that takes one parameter: the signal number. The signal handler function is limited in what it can do. Unless a signal is raised by an explicit call to `raise` or `abort`, the only useful thing the handler can do is to set a global flag. The type of the flag must be `sig_atomic_t`.

The standard defines a basic set of signals, and an implementation is free to define additional signals. On the other hand, an implementation is not required to raise any signals. Remember that arithmetic overflow, pointer

violations, and the like result in undefined behavior. (See [Chapter 4](#).) Thus, an implementation is free to terminate a program immediately, ignore the error, raise a signal, or do anything else. Some signals are meant to reflect external events, such as the user terminating the program. How the user terminates a program is likewise implementation-defined.

Unix and Unix-like operating systems have much more extensive signal-handling facilities. The `<csignal>` handler as documented in this section is portable to all hosted C++ environments, regardless of operating system, although the portable behavior is limited. Most uses of `<csignal>` in real programs take advantage of additional, nonstandard capabilities. Consult your compiler's and library's documentation for details.

## Name

raise function — Raises a signal

## Synopsis

```
int raise(int sig);
```

The raise function sends a signal to the running program. The sig parameter is the signal number. The return value is 0 for success or nonzero for an error.

## Name

SIG\_DFL macro — Default handler

## Synopsis

```
void (*SIG_DFL)(int)
```

The SIG\_DFL macro represents the default handling of a signal. The macro expands to a constant whose value is suitable as the second argument to the signal function.

## Name

SIG\_ERR macro — Error return

## Synopsis

```
void (*SIG_ERR)(int)
```

The SIG\_ERR macro represents the value returned from `signal` in the event of an error. It expands to a constant expression.

## Name

SIG\_IGN macro — Ignore signal

## Synopsis

```
void (*SIG_IGN)(int)
```

The SIG\_IGN macro tells `signal` to ignore a signal. The macro expands to a constant whose value is suitable as the second argument to the `signal` function.

## Name

SIGABRT macro — Abort signal number

## Synopsis

```
int SIGABRT
```

The SIGABRT macro expands to a positive integer constant that represents an abnormal termination. The abort function raises SIGABRT.



## Name

SIGFPE macro — Floating-point error signal number

## Synopsis

```
int SIGFPE
```

The SIGFPE macro expands to a positive integer constant that represents a floating-point exception, such as division by zero. An implementation is not required to raise SIGFPE for a floating-point error.

## Name

SIGILL macro — Illegal instruction signal number

## Synopsis

```
int SIGILL
```

The SIGILL macro expands to a positive integer constant that represents an illegal instruction.

## Name

SIGINT macro — User interrupt signal number

## Synopsis

```
int SIGINT
```

The SIGINT macro expands to a positive integer constant that represents a user interrupt.

## Name

SIGSEGV macro — Segmentation violation signal number

## Synopsis

```
int SIGSEGV
```

The SIGSEGV macro expands to a positive integer constant that represents an addressing fault (segmentation violation).

## **Name**

SIGTERM macro — Terminate signal number

## **Synopsis**

```
int SIGTERM
```

## **Description**

The SIGTERM macro expands to a positive integer constant that represents a request to terminate the program.

## Name

`sig_atomic_t` type — Atomic type

## Synopsis

```
typedef ... sig_atomic_t;
```



The `sig_atomic_t` type is an integral type that can be accessed atomically—that is, even if a signal is delivered, the entire value is read or written. The actual type is implementation-defined.

# Name

signal function — Sets a signal handler

# Synopsis

```
void (*signal(int sig, void (*func)(int)))(int);
```

The `signal` function controls the program's behavior when a signal is delivered to the program. The first parameter (`sig`) is the signal number. The second parameter (`func`) is the function to call when signal `sig` is delivered.

The `func` parameter can also be one of the special values `SIG_DFL` or `SIG_IGN`. Use `SIG_DFL` to get the default behavior; use `SIG_IGN` to ignore a signal.



The default behavior for a signal is implementation-defined, but it usually results in the termination of the program. The signal handler must not use any C++ features (such as throwing an exception), or the results will be implementation-defined. The function must have "C" linkage.

If the `func` parameter is a function pointer, that function is called when signal `sig` is delivered. Unless the signal is delivered by calling `abort` or `raise`, the function is highly restricted in what it can do:

- The handler must not call any function in the standard library except `signal`, and the first parameter must be `sig`.
- The handler must not refer to any variable with static storage except it can assign a value to a variable of type `volatile sig_atomic_t`.

- If the signal is the result of a computational error such as SIGFPE, the signal handler must not return, but should call `abort` or `exit`. (Yes, this item contradicts the first item.)

### **Warning**

Real implementations have looser restrictions, such as allowing calls to other library functions from a signal handler. The library functions that are permitted varies, but every practical implementation allows at least `abort`. If you must use signal handlers in your program, you will probably need to rely on behavior that is dictated by your host environment, extending the limitations of the C++ standard.

If the handler returns normally, and the signal is not the result of a computational error, execution continues from the point where it was interrupted.

The return value of `signal` is the previous value of the signal handler for `sig`, or `SIG_ERR` for an error. If `SIG_ERR` is returned, `errno` is set.

[Example 13-7](#) shows a simple signal handler that sets a global flag when the user interrupts the program. Until the user interrupts it, the program reads input and counts the number of lines the user typed.

## **Example**

*Example 13-7. Reading input until the program is interrupted*

```
#include <csignal>
#include <iostream>
#include <string>

volatile std::sig_atomic_t interrupted;
```



```

// Signal handler sets a global flag
extern "C" void sigint(int sig)
{
    interrupted = 1;
}

int main( )
{
    //
    if (std::signal(SIGINT, sigint) == SIG_ERR)
        std::cerr << "Cannot set signal handler\n";
    else
    {
        unsigned long count = 0;           // Count lines.
        while(! interrupted)
        {
            std::cout << "> ";           // User prompt
            std::string s;
            if (! std::getline(std::cin, s))
                // EOF does not terminate the loop; only SIGINT does
this.
                std::cin.clear( );
            ++count;
        }
        std::cout << "I counted " << count << " line(s).\n";
    }
}

```

## <cstdarg>

The <cstdarg> header is the C++ version of the C standard <stdarg.h> header, which declares macros for accessing the arguments to a function that takes a variable number of arguments, that is, a function that is declared with an ellipsis as the last parameter.

A function that takes a variable number of arguments (called a *variadic function*) must have some way of knowing how many arguments have actually been passed to the function and what their types are. For example, the printf function (in <cstdio>) uses its format string to determine the number and type of arguments.

[Example 13-8](#) shows how a function can use the <cstdarg> macros. The max function takes at least two arguments. The first is a count of the number of remaining arguments; the count must be positive. The template parameter specifies the type of each argument that follows the count.

*Example 13-8. Finding the maximum value of any number of arguments*

```
#include <cassert>
#include <cstdarg>

// Use a trivial wrapper class to ensure that va_end is called.
class varargs {
public:
    ~varargs( ) { va_end(ap); }
    std::va_list& ap;
};

template <typename T>
T max(unsigned count, ...)
{
    assert(count > 0);
    varargs va;
    va_start(va.ap, count);
    T result = va_arg(va.ap, T); // Get first argument.
```

```
while (--count > 0) {
    T arg = va_arg(va.ap, T); // Get successive arguments.
    if (arg > result)
        result = arg; // Remember the largest.
}
return result;
}

int main( )
{
    int a, b, c, d;
    ...
    int x = max<int>(4, a, b, c, d);
    int y = max<int>(2, x, 42);
    return y;
}
```

## Name

`va_arg` macro — Gets next argument

## Synopsis

```
T va_arg(va_list ap, T)
```

The `va_arg` macro fetches the next argument, which must be of type `T`. The `ap` parameter must have been initialized by calling `va_start`. The type `T` must be a type that results from the standard type promotions ([Chapter 3](#)) or else the behavior is undefined. For example, `T` cannot be `char`, but must be `int` because the standard promotion of type `char` is to type `int` (or, in rare circumstances, `unsigned int`). The behavior is undefined if there is no next argument.

## Name

`va_end` macro — Ends getting arguments

## Synopsis

```
void va_end(va_list ap)
```

The `va_end` macro finishes fetching arguments. You must call `va_end` once for each call to `va_start`. You cannot nest calls to `va_start` and `va_end`, but you can call them multiple times sequentially in the same function.

## Name

`va_list` type — Argument list

## Synopsis

```
typedef ... va_list;
```

The `va_list` type is an opaque type that refers to the function's arguments. Declare a local variable of type `va_list` and supply the variable to the `va_start`, `va_arg`, and `va_end` macros.

## Name

`va_start` macro — Starts getting arguments

## Synopsis

```
void va_start(va_list& ap, lastNamedParm)
```

The `va_start` macro initializes `ap` and prepares to fetch function arguments with `va_arg`. You must call `va_end` to clean up and finalize `ap`. The second argument to `va_start` is the name of the function's last parameter before the ellipsis. The last named parameter must not have a function, array, or reference type.

## **<cstdint>**

The `<cstdint>` header is the C++ version of the C standard `<stdint.h>` header, which declares a few types and macros.

The C header declares the `wchar_t` type, but `wchar_t` is a reserved keyword in C++, so there is no need to `#include <cstdint>` to declare this type.



## Name

NULL macro — Null pointer

## Synopsis

```
#define NULL
```

. . .

The NULL macro expands to a null pointer constant, such as 0 or 0L.



Some C libraries declare NULL as ((void\*)0) in `stddef.h`. This definition is fine for C, but is wrong for C++. Most C++ compilers correctly declare NULL, but you should be aware of the difference.

## Name

offsetof macro — Computes member offset

## Synopsis

```
size_t offsetof(type, member-name)
```

The `offsetof` macro returns the offset, in bytes, of a member of a struct as a constant integer. The *type* must be a plain, C-style struct (Plain Old Data, or POD), and the expression `&(t.member-name)` must be an address constant, assuming `t` is an instance of *type*. In particular, this means the *member-name* must not be a bit-field, a static member, or a function member. (See [Chapter 6](#) for more information about POD types.)

## Name

`ptrdiff_t` type — Pointer difference type

## Synopsis

```
typedef ... ptrdiff_t
```



The `ptrdiff_t` type is a signed integral type that represents the difference between two pointers. The exact type is implementation-defined.

## Name

`size_t` type — `sizeof` result type

## Synopsis

```
typedef ... size_t
```



The `size_t` type is the type of the result of the `sizeof` operator. It is an unsigned integral type. The exact type is implementation-defined.

## **<cstdio>**

The `<cstdio>` header is a wrapper for the C standard `<stdio.h>` header, which declares input and output types, macros, and functions. See also `<wchar>` for wide character I/O functions.

C++ I/O streams offer more flexibility, type-safety, and clarity. On the other hand, C I/O offers simplicity and compatibility with C libraries. See [Chapter 9](#) for an overview of the C++ I/O stream classes.

## Name

`_IOFBF` macro — Full buffering

## Synopsis

```
int _IOFBF
```

When passed as the mode parameter to `setvbuf`, the `_IOFBF` macro sets an open file to full buffering. A buffer is flushed when it is full. The `_IOFBF` macro expands to a constant integer.



Support for fully-buffered streams is implementation-dependent.

## See Also

[setvbuf function](#)

## Name

`_IOLBF` macro — Line buffering

## Synopsis

```
int _IOLBF
```

When passed as the mode parameter to `setvbuf`, the `_IOLBF` macro sets an open file to line buffering. A buffer is flushed when it is full or when a newline character is read or written. The `_IOLBF` macro expands to a constant integer.



Support for line-buffered streams is implementation-dependent.

## See Also

[setvbuf function](#)

## Name

`_IONBF` macro — No buffering

## Synopsis

```
const int _IONBF
```

When passed as the mode parameter to `setvbuf`, the `_IONBF` macro disables the buffering of an open file. Characters are read or written as soon as possible, without buffering. The `_IONBF` macro expands to a constant integer.



Support for unbuffered streams is implementation-dependent. For example, a host operating system might line buffer input from a terminal, even if a program requests unbuffered input.

## See Also

[setvbuf function](#)



## Name

BUFSIZ macro — Buffer size

## Synopsis

```
int BUFSIZ
```

The BUFSIZ macro specifies the minimum buffer size for the setbuf function. The BUFSIZ macro expands to a constant integer.

## See Also

[setbuf function](#), [setvbuf function](#)

## Name

clearerr function — Clears error status

## Synopsis

```
void clearerr(FILE* stream)
```

The `clearerr` function clears the error and end-of-file indicators for `stream`.

## See Also

[feof function](#), [ferror function](#)

## Name

EOF macro — End-of-file or error

## Synopsis

```
int EOF
```

The EOF macro represents end-of-file when returned from `getchar` and other functions. Some functions return EOF to indicate an error.



The value of EOF is a negative integer constant. The precise value is implementation-defined.

## Name

fclose function — Closes a file

## Synopsis

```
int fclose(FILE* stream)
```

The `fclose` function flushes and closes an open file. It returns `0` upon success or EOF when there is an error.

## See Also

[fopen function](#)

## Name

feof function — Tests for end-of-file

## Synopsis

```
int feof(FILE* stream)
```

The feof function returns true (nonzero) if `stream` is positioned at the end-of-file, or false (0) otherwise.

## See Also

[clearerr function](#), [ferror function](#)

## Name

error function — Tests for error

## Synopsis

```
int ferror(FILE* stream)
```

The `ferror` function returns true (nonzero) if `stream` has an error condition set, or false (0) otherwise.

## See Also

[clearerr function](#), [feof function](#)

## Name

fgetc function — Reads a character

## Synopsis

```
int fgetc(FILE* stream)
```

The fgetc function reads a single character from stream. It returns the character as an unsigned char converted to int or EOF for an error or end-of-file.

## See Also

[feof function](#), [ferror function](#), [getc macro](#), [fputc function](#),  
fwgetc in [<wchar>](#)

## Name

fgetpos function — Returns file position

## Synopsis

```
int fgetpos(FILE* stream, fpos_t* pos)
```

The `fgetpos` function stores `stream`'s current position in the object that `pos` points to. The only use for the position is to save it and pass it to `fsetpos` to set the file's position. You cannot use the position arithmetically, e.g., to advance the position by one character.

The return value is 0 for success or nonzero for failure. If `fgetpos` fails, it sets `errno`.

## See Also

[fpos\\_t type](#), [fsetpos function](#), [ftell function](#)



## Name

fgets function — Reads a string

## Synopsis

```
char* fgets(char* s, int n, FILE* stream)
```

The fgets function reads a line of text from stream into the character array that s points to. It stops reading after a newline character or after n - 1 characters have been read. The newline character, if one is encountered, is copied into s.

The return value is s for success or a null pointer for an error or end-of-file. If fgets fails, the contents of the string s are undefined.

## See Also

[fgets function](#), [getc macro](#), [fputs function](#), fwgets in [<wchar>](#)

## Name

FILE type — File type

## Synopsis

```
typedef . . . FILE
```

The FILE type represents the contents of an external file. A C++ program works with FILE pointers, in which the actual FILE objects are managed by functions in the standard library. Thus, you never need to allocate or free FILE objects.

## See Also

[fclose function](#), [fopen function](#), [freopen function](#),  
[<fstream>](#)

## Name

FILENAME\_MAX macro — Maximum length of a filename

## Synopsis

```
int FILENAME_MAX
```

FILENAME\_MAX is the size you should use when declaring a character array that will store a filename. Some systems do not have a fixed maximum size for a filename, in which case FILENAME\_MAX is a recommended size, and the resulting character array might not be large enough to hold all valid filenames. The FILENAME\_MAX macro expands to a constant integer.

### **Not e**

Use `std::string` instead of a character array to avoid any problems with character arrays that are too small.

# Name

fopen function — Opens a file

# Synopsis

FILE\* **fopen**(const char\* filename, const char\* mode)

The fopen function opens a file.



The filename parameter specifies the filename in an implementation-defined manner. The mode parameter specifies how to open the file. The mode must begin with one of the strings listed in [Table 13-4](#). Additional characters can follow, and the interpretation of the extra characters is implementation-defined. Mode strings are case-sensitive.

*Table 13-4. File open modes*

Mode string	Description
a	Append: opens an existing file for appending, that is, every write is forced to the end of the file. If the file to be opened does not exist, a creates it.
r	Read: opens an existing file for reading.
w	Write: creates a new file for writing. If the file already exists, w truncates it to zero length.
ab	Append in binary mode.
rb	Read in binary mode.
wb	Write in binary mode.

Mode string	Description
a+	Append update: opens a file in append mode and allows it to be read.
r+	Read update: opens an existing file for reading, and also allows writing.
w+	Write update: creates a new file for writing, and also allows reading. If the file already exists, w+ truncates it to zero length.
ab+ or a+b	Append update in binary mode.
rb+ or r+b	Read update in binary mode.
wb+ or w+b	Write update in binary mode.

## See Also

[fclose function](#), [freopen function](#)

## Name

FOPEN\_MAX macro — Minimum limit on the number of open files

## Synopsis

```
int FOPEN_MAX
```

A typical operating system has a maximum number of files that can be open at one time. This number might be variable or fixed; FOPEN\_MAX is the guaranteed minimum value of the limit. The FOPEN\_MAX macro expands to a constant integer.

## See Also

[fopen function](#)

## Name

fpos\_t type — File position

## Synopsis

```
typedef . . . fpos_t
```

The `fpos_t` type is an opaque type that represents a position in a file. The only way to set the value of an `fpos_t` object is to call `fgetpos`, and the only things you can do with the value are assign an `fpos_t` value to it and pass it as a function argument, especially to `fsetpos`.

## See Also

[fgetpos function](#), [fsetpos function](#), `fpos` in [<ios>](#)

# Name

`fprintf` function — Writes formatted data

# Synopsis

```
int fprintf(FILE* stream, const char* format, . . . )
```

The `fprintf` function writes formatted output to `stream`. The `format` parameter contains the formatting information, and the remaining arguments are printed according to the format. The return value is the number of characters printed, or a negative value for an error.

Characters in `format` are printed verbatim except for conversion specifications, which begin with a percent sign (%). Each conversion specification is made up of the following parts (in order): flags, field width, precision, size, and conversion specifier.

The following are detailed descriptions of the parts of a conversion specification:

## Flags

The flag characters are optional and can appear in any order. [Table 13-5](#) lists the flag characters and their meanings.

*Table 13-5. Formatting flag characters*

Flag	Description
-	Left-justified (default is right-justified).
+	Signed conversions always begin with a sign (default is to use a sign only if the value is negative).



Flag	Description
Space	The output is an initial space character if a signed conversion results in an empty string or a string that does not start with a sign character (+ takes precedence over space).
#	Use an alternate form: insert a 0 for %o; insert 0x for %x or 0X for %X; always output a decimal point for floating-point conversions; do not remove trailing zeros for %g or %G; behavior is undefined for other conversions.
0	Fields are padded with leading zeros (after the sign or base indication). The - flag takes precedence over 0. For integer conversions, a precision takes precedence over the 0 flag.

## Field width

An optional number that specifies the minimum number of characters that the field will occupy. If the field is an asterisk (\*), the field width is obtained from the next argument to be processed.

## Precision

An optional number of digits for an integer, number of decimal digits for a floating-point number, or maximum size for a string. The precision is specified as a dot (.) followed by a number or an asterisk.

## Size

The character h, l, or L. h means an integer is short or unsigned short. l means an integer is long or unsigned long, a character is `wint_t`, or a string is a pointer to `wchar_t`. L means a floating-point number is long double.

## Conversion character

Specifies the type of the argument containing the data to be printed using a conversion specification. It must be

one of the following:

d, i

Signed decimal integer.

o

Unsigned octal integer.

u

Unsigned decimal integer.

x, X

Unsigned hexadecimal integer. x writes the digits a-f in lowercase, and X writes the digits A-F in uppercase.

f

Fixed-precision floating point.

e, E

Exponential floating point. The exponent is introduced with e or E, matching the conversion character.

g, G

General floating point. Use style f or e. Use style e if the exponent is less than -4 or greater than the precision; otherwise, use style f. Trailing zeros are dropped, and a trailing decimal point is dropped if it would be the last character.

c

Character. The argument must be an unsigned char promoted to int, or, if the l size modifier is used, the argument must be wchar\_t promoted to wint\_t, which is then printed as a multibyte character.

s

String. The argument is a pointer to a null-terminated array of characters, or, if the `l` size modifier is used, the argument must be a pointer to a `wchar_t` array, which is converted to a series of multibyte characters.



`p`

Pointer. The argument must be a pointer to `void`. The output format is implementation-defined.

`n`

The argument must be a pointer to an integer; `fprintf` stores in the integer the number of characters written so far. Use the `h` or `l` size modifiers if the argument is a pointer to `short int` or `long int`.

`%`

Prints a literal `%`.

### **Warning**

It is your responsibility to ensure that the argument types match the format. Any errors result in undefined behavior. Mismatches between format and the argument is a common and sometimes subtle source of error. You can avoid these problems entirely by using C++ I/O streams instead of `fprintf` and related functions.

All the `printf`-related functions interpret the format string identically.

## **Example**

The following are examples of calling `printf`:

```
long double pi = 3.141592653589792L;  
int i = 42;
```

```
const char greeting[] = "Hello, how are you?";

printf(">%d %% %Lg<\n", i, pi); // Prints >42 % 3.14159<
printf(">%4d<\n", i);          // Prints > 42<
printf(">%-16.8Le<\n", pi);    // Prints >3.14159265e+00 <
printf(">%#*. *x<\n", 8, 4, i); // Prints > 0x002a<
printf(">%.5s<\n", greeting);  // Prints >Hello<
```

## See Also

[fscanf function](#), [printf function](#), [sprintf function](#), [vfprintf function](#), [wrtomb](#) in [<wchar>](#), [fwprintf](#) in [<wchar>](#)

## Name

fputc function — Writes a character

## Synopsis

```
int fputc(int c, FILE* stream)
```

The fputc function writes a single character to stream. The character must be an unsigned char, which is automatically promoted to int, so the proper way to print a variable of type char is as follows:

```
char ch;  
fputc(static_cast<unsigned char>(ch), stream);
```

The return value is EOF for an error or c for success.

## See Also

[putc macro](#), fwputc in [<wchar>](#)

## Name

fputs function — Writes a string

## Synopsis

```
int fputs(const char* s, FILE* stream)
```

The fputs function writes the string `s` to `stream`. It returns EOF for an error or a nonnegative value for success.

## See Also

[fputc function](#), [puts function](#), fwputs in [<wchar>](#)

## Name

fread function — Reads binary data

## Synopsis

```
size_t fread(void* ptr, size_t size, size_t count, FILE* stream)
```

The fread function reads up to count elements from stream into the memory that ptr points to. The memory must have POD type (see [Chapter 6](#)). Each element is size bytes long. It returns the number of elements that were read successfully.

## See Also

[fwrite function](#)

## Name

freopen function — Opens a file with an existing stream

## Synopsis

```
FILE* freopen(const char* filename, const char* mode, FILE*  
stream)
```

The `freopen` function opens a file using an existing stream. The file previously associated with the `stream` is closed first, and the named file is opened in the same manner as if `fopen` were called. See `fopen` for a description of the mode parameter.

The main purpose of using `freopen` is to reopen one of the standard files: `stdin`, `stdout`, and `stderr`.

## See Also

[fclose function](#), [fopen function](#)



# Name

`fscanf` function — Reads formatted data

# Synopsis

```
int fscanf(FILE* stream, const char* format, . . . )
```

The `fscanf` function performs a formatted read from `stream`. The `format` parameter contains formatting information, and the remaining arguments are pointers. When `fscanf` reads items, it stores their values in the objects that successive arguments point to. The return value is the number of items read or a negative value for an error.

Items are read from `stream` and interpreted according to `format`, which contains whitespace characters, non-whitespace characters, and conversion specifications, which begin with a percent sign (%). A whitespace character directs `fscanf` to skip over whitespace in the input stream. Non-whitespace characters must match the input text. Each conversion specification is made up of the following parts (in order): assignment suppression, field width, size, and conversion specifier.

The following are descriptions of the conversion specification elements:

## Assignment suppression

An optional asterisk (\*) directs `fscanf` to read and parse the input according to the conversion specification, but not to assign the value to an argument.

## Field width

An optional number (positive decimal integer) that specifies the maximum number of characters to read.

## Size

The character `h`, `l`, or `L`. `h` means an integer is short or unsigned short. `l` means an integer is long or unsigned long; a floating-point number is double, or a string argument is a pointer to `wchar_t` for the `c`, `s`, and `[` conversion specifiers. `L` means a floating-point number is long double. The default size for an integer is `int` or unsigned `int`, `float` for a floating-point number, and `char` for any of the character conversion specifiers.

## Conversion character

Specifies the type of the argument containing the data to be printed using a conversion specification. It must be one of the following:

`d`

Signed decimal integer.

`i`

Signed integer. Reads and interprets a prefix of `0x` or `0X` for hexadecimal, `0` for octal, or anything else for decimal.

`o`

Unsigned octal integer.

`u`

Unsigned decimal integer.

`x`, `X`

Unsigned hexadecimal integer.

`e`, `E`, `f`, `g`, `G`

Floating point in fixed or exponential format.

`c`

Characters. The field width (default 1) specifies the exact number of characters to read. The corresponding argument must be a pointer to a character array large enough to hold the characters. If the `l` modifier is used, the input is read as multibyte characters, which are converted to wide characters and stored in a `wchar_t` array. In either case, no null character is appended.

s

String. Reads a sequence of non-whitespace characters. The corresponding argument must be a pointer to a character array that is large enough to hold the sequence plus a terminating null character. If the `l` modifier is used, the input is read as multibyte characters, which are converted to wide characters and stored in a `wchar_t` array, followed by a terminating null wide character.



p

Pointer. The argument must be a pointer to void. The input format is implementation-defined and matches the format that `fprintf` uses.

n

The argument must be a pointer to an integer; `fscanf` stores in the integer the number of characters read so far. The `h` or `l` size modifiers can be used if the argument is a pointer to short int or long int. Nothing is read from the input, and `%n` does not affect the count returned by `fscanf`.

[

Matches a sequence of characters. The conversion specification lists a set of characters (called the

*scanset*) in square brackets. The input string is a sequence of characters that matches any of the characters in the scanset or, if the scanset begins with a circumflex (^), any character not in the scanset. If the *l* modifier is used, the input is read as multibyte characters, which are converted to wide characters and stored in a *wchar\_t* array, followed by a terminating null wide character.

%

Matches a literal % in the input stream.

### **Warning**

It is your responsibility to ensure that the argument types match the format. Any errors result in undefined behavior. Mismatches between format and the argument are a common and sometimes subtle source of error. You can avoid these problems entirely by using C++ I/O streams instead of *fscanf* and related functions.

All the *scanf*-related functions interpret the format string identically.

## **Example**

The following is an example of calling *scanf*. The input stream is:

```
start 3.14 BeefFeed cab42.0e-01, 1234
```

and the *scanf* call is:

```
char c[10];  
double d;  
float f;  
long int l;  
unsigned short us;
```

```
scanf("start %4lf %4lx%x %9[abcdefg]%f,%hu", &d, &l, c, &f, &us);
```

which has the following result:

```
c = "cab"  
d = 3.14  
f = 4.2  
l = 48879 (0xbeef)  
us = 1234
```

## See Also

[fprintf function](#), [scanf function](#), [sscanf function](#), [vfscanf function](#), [mbrtowc](#) in [<wchar>](#), [fwscanf](#) in [<wchar>](#)

## Name

fseek function — Changes file position

## Synopsis

```
int fseek(FILE* stream, long int offset, int origin)
```

The `fseek` function seeks to a different position in `stream`. The origin must be one of `SEEK_CUR`, `SEEK_END`, or `SEEK_SET`. The offset is relative to the current position, end-of-file, or start-of-file, respectively. The end-of-file flag is cleared, and any ungetc character is also cleared.

Use `fsetpos` instead of `fseek` when using large files—that is, for which a position does not necessarily fit into a `long int`.

The return value is 0 for success or nonzero for an error.

## See Also

[fsetpos function](#), [ftell function](#), [SEEK\\_CUR macro](#), [SEEK\\_END macro](#), [SEEK\\_SET macro](#)

## Name

fsetpos function — Changes file position

## Synopsis

```
int fsetpos(FILE* stream, const fpos_t* pos)
```

The `fsetpos` function seeks to a different position in stream. The position must have been returned from an earlier successful call to `fgetpos`.

## See Also

[fpos\\_t type](#), [fseek function](#), [fgetpos function](#)

## Name

`ftell` function — Returns current file position

## Synopsis

```
long int ftell(FILE* stream)
```

The `ftell` function returns the current file position in stream. This position can be used (with an origin of `SEEK_SET`) in a subsequent call to `fseek`. If `ftell` fails, it returns `-1L` (which may be, but is not necessarily, the same value as `EOF`).

## See Also

[fgetpos function](#), [fseek function](#)



## Name

`fwrite` function — Writes binary data

## Synopsis

```
size_t fwrite(const void* ptr, size_t size, size_t count, FILE* stream)
```

The `fwrite` function writes up to `count` elements to `stream`. Each element is `size` bytes long, and `ptr` points to the first such element. The memory must have POD type (see [Chapter 6](#)).

The return value is the number of complete elements successfully written to `stream`.

## See Also

[fread function](#)

## Name

getc macro — Reads a character

## Synopsis

```
int getc(FILE* stream)
```

The `getc` macro reads a character from `stream` and returns that character as an unsigned char. The return value is EOF for end-of-file or a read error.

## See Also

[fgetc function](#), [getchar macro](#), [putc macro](#)

## Name

getchar macro — Reads a character from stdin

## Synopsis

```
int getchar( )
```

The getchar macro is equivalent to `getc(stdin)`.

## See Also

[fgetc function](#), [getc macro](#), [putchar macro](#)

## Name

gets function — Reads a string unsafely

## Synopsis

```
char* gets(char* s)
```

The gets function reads a line of text (up to and including a newline) into the string s.

### Warning

There is no way to limit the input to the size of s, so you should never call gets. Use fgets instead.

## See Also

[fgets function](#), [getchar macro](#)

## Name

L\_tmpnam macro — Length of temporary filename

## Synopsis

```
int L_tmpnam
```

L\_tmpnam is the length of a temporary filename for tmpnam. The L\_tmpnam macro expands to a constant integer.

## See Also

[tmpnam function](#)

## Name

NULL macro — Null pointer constant

## Synopsis

```
#define NULL
```

. . .

The NULL macro expands to a null pointer constant. See `<stddef>` for more information.

## See Also

NULL in [<stddef>](#)

## Name

`perror` function — Writes an error message

## Synopsis

```
void perror(const char* s)
```

The `perror` function writes an error message to `stderr`. If `s` is not null and is not an empty string, it is printed first, followed by a colon and a space. The error message is printed next and is the same text as that returned by `strerror`, passing `errno` as its argument.

## See Also

`errno` in [<cerrno>](#), `strerror` in [<cstring>](#)

## Name

printf function — Writes formatted data

## Synopsis

```
int printf(const char* format, . . . )
```

The printf function is equivalent to calling fprintf to stdout. See fprintf for information about the format string.

## See Also

[fprintf function](#), [vprintf function](#), wprintf in [<wchar>](#)



## Name

putc macro — Writes a character

## Synopsis

```
int putc(int c, FILE* stream)
```

The `putc` macro writes the character `c`, which must be cast to unsigned char, to stream.

## See Also

[fputc function](#), [putchar macro](#), `wputc` in [<wchar>](#)

## Name

putchar macro — Writes a character to stdout

## Synopsis

```
int putchar(int c)
```

The putchar macro is equivalent to `putc(c, stdout)`.

## See Also

[fputc function](#), [putc macro](#), wputchar in [<wchar>](#)

## Name

puts function — Writes a string

## Synopsis

```
int puts(const char* s)
```

The puts function writes a string to stdout.

## See Also

[fputs function](#), putchar function, wputs in [<wchar.h>](#)

## Name

remove function — Deletes a file

## Synopsis

```
int remove(const char* filename)
```

The remove function deletes the file named by `filename`. It returns 0 for success or nonzero for an error.

## See Also

[rename function](#)

## Name

rename function — Renames a file

## Synopsis

```
int rename(const char* oldname, const char* newname)
```

The rename function renames the file specified by oldname to newname. The return value is 0 for success or nonzero for failure.

## See Also

[remove function](#)

## Name

rewind function — Resets file position

## Synopsis

```
void rewind(FILE* stream)
```

The `rewind` function moves a file position to the beginning of a file and is equivalent to `fseek(stream, 0, SEEK_SET)`.

## See Also

[fseek function](#), [fsetpos function](#)

## Name

SEEK\_CUR macro — Seek from current file position

## Synopsis

```
int SEEK_CUR
```

Pass SEEK\_CUR as the last parameter to `fseek` to seek from the current file position. Positive offset values seek toward the end of the file, and negative values seek toward the beginning of the file. The SEEK\_CUR macro expands to a constant integer.

## See Also

[fseek function](#)

## Name

SEEK\_END macro — Seek from end-of-file

## Synopsis

```
int SEEK_END
```

Pass SEEK\_END as the last parameter to fseek to seek relative to the end of the file. Positive offset values seek past the end of the file, and negative values seek toward the beginning of the file. The SEEK\_END macro expands to a constant integer.

## See Also

[fseek function](#)



## Name

SEEK\_SET macro — Seek from beginning of file

## Synopsis

```
int SEEK_SET
```

Pass SEEK\_SET as the last parameter to `fseek` to seek from the start of the file. Positive offset values seek toward the end of the file. The SEEK\_SET macro expands to a constant integer.

## See Also

[fseek function](#)

## Name

setbuf function — Sets file buffer

## Synopsis

```
void setbuf(FILE* stream, char* buf)
```

The `setbuf` function sets the buffer to use when reading from or writing to `stream`. The size of `buf` must be at least `BUFSIZ` characters.

## See Also

[BUFSIZ macro](#), [setvbuf function](#)

## Name

setvbuf function — Sets file buffer

## Synopsis

```
int setvbuf(FILE* stream, char* buf, int mode, size_t size)
```

The `setvbuf` function sets the buffering for `stream`. The mode determines the buffering mode: no buffering (`_IONBF`), line buffering (`_IOLBF`), or full buffering (`_IOFBF`). You can supply a buffer in the `buf` argument, with `size` as the buffer size, or use a null pointer for the `buf` argument to let `setvbuf` allocate the buffer. (The buffer will be freed when the file is closed or `setvbuf` is called to change the buffering.)

Call `setvbuf` before performing any I/O on `stream`.

## See Also

`IOFBF` macro, [\\_IOLBF macro](#), [\\_IONBF macro](#), [setbuf function](#)

## Name

size\_t type — Size type

## Synopsis

```
typedef . . . size_t;
```



The `size_t` type is the type of the result of the `sizeof` operator. It is an unsigned integral type. The exact type is implementation-defined.

## See Also

size\_t in [<stddef>](#)

## Name

sprintf function — Writes formatted data to a string

## Synopsis

```
int sprintf(char* s, const char* format, ...)
```

The `sprintf` function is like `fprintf`, but instead of writing to an open file, it "writes" by copying characters to the string `s`. See `fprintf` for a description of the format parameter.

You must ensure that `s` is large enough for the formatted string. For some formats, this is impossible, which makes `sprintf` unsafe to use.

## See Also

[fprintf function](#), [sscanf function](#), [vsprintf function](#),  
`wsprintf` in [<wchar>](#), [<sstream>](#)

## Name

sscanf function — Reads formatted data from a string

## Synopsis

```
int sscanf(const char* s, const char* format, ...)
```

The `sscanf` function is like `fscanf`, but instead of reading from an open file, it "reads" characters from the string `s`. See `fscanf` for a description of the format parameter.

## See Also

[fscanf function](#), [sprintf function](#), `vsscanf` function, `swscanf` in [<wchar>](#), [<sstream>](#)

## Name

stderr macro — Standard error file

## Synopsis

FILE\* `stderr`

The `stderr` macro is a standard file, suitable for printing error messages. Its buffering is implementation-defined: either unbuffered or line buffered.

## See Also

`cerr` in [<iostream>](#), `clog` in [<iostream>](#)

## Name

stdin macro — Standard input file

## Synopsis

FILE\* `stdin`

The `stdin` macro is a standard file used for reading from the program's standard input device. `stdin` is fully buffered if the standard input is not an interactive device.

## See Also

cin in [<iostream>](#)



## Name

stdout macro — Standard output file

## Synopsis

FILE\* `stdout`

The `stdout` macro is a standard file used to print to the program's standard output device. `stdout` is fully buffered if the standard output is not an interactive device.

## See Also

cout in [<iostream>](#)

## Name

TMP\_MAX macro — Size of temporary filename

## Synopsis

```
int TMP_MAX
```

The TMP\_MAX macro is the number of unique names the tmpnam function generates.

## See Also

[tmpnam function](#)

## Name

tmpfile function — Opens a temporary file

## Synopsis

```
FILE* tmpfile( )
```

The tmpfile function generates a unique filename and opens a file with that name using mode "wb+". When the program terminates normally, the file is automatically deleted. If the temporary file cannot be created, a null pointer is returned.

## See Also

[tmpnam function](#)

# Name

tmpnam function — Returns a temporary filename

# Synopsis

```
char* tmpnam(char* s)
```

The tmpnam function generates a unique filename and returns a pointer to that name. If the s parameter is not null, it must point to an array of at least L\_tmpnam characters, the new filename is copied into that array, and s is returned. If s is null, a static character array is returned; the contents of the array are overwritten each time tmpnam is called.

## Warning

The tmpnam function has a race condition: after it generates a "unique" filename, but before the file is actually created, another program might generate the same "unique" filename. Use tmpfile instead of tmpnam to ensure that the file is created, thereby ensuring that its name is unique.

# See Also

[L\\_tmpnam macro](#), [tmpfile function](#)

## Name

`ungetc` function — Pushes a character back for reading later

## Synopsis

```
int ungetc(int c, FILE* stream)
```

The `ungetc` function pushes back the character `c` (which must be an unsigned char), so the next read from `stream` will return `c`. The standard guarantees that you can push back just one character, though in some situations, you may be able to push back more.

The return value is `c` for success or `E0F` for an error.

## See Also

[fgetc function](#), [getc macro](#), [getchar macro](#), `ungetwc` in [<wchar>](#)

# Name

vfprintf function — Writes formatted data

# Synopsis

```
#include <cstdarg>
int vfprintf(FILE* stream, const char* format, va_list arg)
```

The `vfprintf` function is like `fprintf`, but the values to print are taken from successive arguments in `arg` (obtained by calling `va_start(arg, param)`). Use `vfprintf` to write your own `fprintf`-like function.

# Example

The following shows how you can implement `fprintf` in terms of `vfprintf`:

```
int fprintf(std::FILE* stream, const char* format, . . . )
{
    std::va_list ap;
    va_start(ap, format);
    int result = vfprintf(stream, format, ap);
    va_end(ap);
    return result;
}
```

# See Also

[fprintf function](#), `vfwprintf` in [<wchar>](#), [<cstdarg>](#)

## Name

vprintf function — Writes formatted data

## Synopsis

```
#include <cstdarg>
int vprintf(const char* format, va_list arg)
```

The vprintf function is like printf, but the values to print are taken from successive arguments in arg (obtained by calling va\_start(arg, *param*)). Use vprintf to write your own printf-like function.

## See Also

[printf function](#), vwprintf in [<wchar>](#), [<cstdarg>](#)

## Name

vsprintf function — Writes formatted data to a string

## Synopsis

```
#include <cstdarg>
int vsprintf(char* s, const char* format, va_list arg)
```

The vsprintf function is like sprintf, but the values to print are taken from successive arguments in arg (obtained by calling va\_start(arg, param)). Use vsprintf to write your own sprintf-like function.

## See Also

[sprintf function](#), vswprintf in [<wchar>](#), [<cstdarg>](#)



## **<cstdlib>**

The `<cstdlib>` header is the C++ version of the C standard `<stdlib.h>` header, which declares macros, types, and functions of general utility. Several functions in `<cstdlib>` convert character arrays to numbers. You can also use a string stream (see `<sstream>`) if you want more control over the conversion, or if you need to convert a number to a string.

The multibyte functions (`mblen`, etc.) have counterparts in `<wchar>` that are reentrant, that is, the `<wchar>` functions take an explicit `mbstate_t*` parameter to avoid using an internal, static shift state.

## Name

abort function — Terminates the program abnormally

## Synopsis

```
void abort( )
```

The abort function raises the SIGABRT signal. Unless the program has registered a handler for SIGABRT, the default action is to terminate the program without destroying any automatic or static objects and without calling any atexit functions.

## See Also

[atexit function](#), [exit function](#), raise in [<csignal>](#), SIGABRT in [<csignal>](#)

## Name

abs function — Computes an absolute value

## Synopsis

```
int abs(int x)
long abs(long x)
```

The `abs` function returns the absolute value of `x`. The `<cmath>` header declares floating-point versions of the `abs` function.

## See Also

[labs function](#), `abs` function in [<cmath>](#)

## Name

atexit function — Calls a function at program termination

## Synopsis

```
extern "C"    int atexit(void (*func)( ))  
extern "C++" int atexit(void (*func)( ))
```

The `atexit` function registers a parameterless function, `func`, which is called when the program exits normally.

Multiple functions can be registered, and registered functions are called in the opposite order of registration. A function can be registered more than once, in which case it will be called as many times as it was registered.

The `atexit` functions are not called if the program exits due to a signal, such as `SIGABRT`.

## See Also

[abort function](#), [exit function](#)

## Name

atof function — Converts a string to floating point

## Synopsis

```
double atof(const char* str)
```

The `atof` function reads a floating-point number from the character array `str` and returns the value of the floating-point number. The conversion is similar to calling `strtod(str, NULL)`.

## See Also

[atoi function](#), [atol function](#), [strtod function](#)

## Name

atoi function — Converts a string to an integer

## Synopsis

```
int atoi(const char* str)
```

The `atoi` function reads an integer from the character array `str` and returns the value of the number. The `atoi` function is equivalent to calling `static_cast<int>(strtol(str, NULL, 10))`.

## See Also

[atof function](#), [atol function](#), [strtod function](#)

## Name

atol function — Converts a string to a long integer

## Synopsis

```
long atol(const char* str)
```

The `atol` function reads an integer from the character array `str`, and returns the value of the number. The conversion is similar to calling `strtol(str, NULL, 10)`.

## See Also

[atof function](#), [atoi function](#), [strtol function](#), [strtoul function](#)

# Name

bsearch function — Performs a binary search

# Synopsis

```
extern "C"  
    void* bsearch(const void* key, const void* base, size_t count,  
size_t size, int  
                (*compare)(const void*, const void*))  
extern "C++"  
    void* bsearch(const void* key, const void* base, size_t count,  
size_t size, int  
                (*compare)(const void*, const void*))
```

The `bsearch` function uses a binary search to search for `key` in the array `base`, in which each element takes up `size` bytes. There are `count` elements in the array.

The `compare` function is called with `key` as the first argument and a pointer into the array as the second. The function should return an integer less than zero if the key is less than the array element, greater than zero if the key is larger, or `0` if the key is the same as the array element.

The `bsearch` function returns a pointer to the array element that matches `key` or a null pointer if no element matches.

Two versions of `bsearch` are declared so the `compare` function can have "C" linkage or "C++" linkage.

# See Also

[qsort function](#), `binary_search` in [<algorithm>](#)



## Name

calloc function — Allocates memory

## Synopsis

```
void* calloc(size_t count, size_t size)
```

The `calloc` function allocates `count` elements, each of `size` bytes, and initializes the allocated memory to all zeros. It returns a pointer to the start of the newly allocated memory or a null pointer if there is insufficient memory to fulfill the request. The pointer is suitably aligned for any type.

C++ programs should use the `new` operator instead of calling `calloc`. Call `fill` or `memset` to fill the allocated memory with zeros.

## See Also

[free function](#), [malloc function](#), [realloc function](#), `new` keyword, `fill` in [<algorithm>](#), `memset` in [<cstring>](#)

## Name

div function — Computes quotient and remainder

## Synopsis

```
div_t div(int numerator, int denominator)  
ldiv_t ldiv(long numerator, long denominator)
```

The div function divides numerator by denominator and returns the quotient and the remainder in a structure.

## See Also

[div\\_t type](#), [ldiv function](#)

## Name

`div_t` type — Quotient and remainder type

## Synopsis

```
struct div_t { int quot, rem; }
```

The `div_t` type is used only by the `div` function to return the quotient and remainder of an integer division. The order of the `quot` and `rem` members in the structure may vary between implementations.

## See Also

[div function](#), [ldiv\\_t type](#)

## Name

exit function — Terminates the program normally

## Synopsis

```
void exit(int code)
```



The `exit` function terminates the program normally. Automatic objects are not destroyed, but static objects are. Then, all functions registered with `atexit` are called in the opposite order of registration. The code is returned to the operating system. An exit code of `0` or `EXIT_SUCCESS` means successful completion. If code is `EXIT_FAILURE`, an indication of program failure is returned to the operating system. Other values of code are implementation-defined.

## See Also

[abort function](#), [atexit function](#)

## Name

EXIT\_FAILURE macro — Exit status for unsuccessful termination

## Synopsis

```
int EXIT_FAILURE
```

Pass EXIT\_FAILURE to exit to terminate the program normally and informs the operating system that the program was unsuccessful. Returning EXIT\_FAILURE from main is the same as calling exit. The EXIT\_FAILURE macro expands to an integer constant.

## See Also

[exit function](#)

## Name

EXIT\_SUCCESS macro — Exit status for successful termination

## Synopsis

```
int EXIT_SUCCESS
```

Pass EXIT\_SUCCESS to exit to terminate the program normally and inform the operating system that the program was successful. Returning EXIT\_SUCCESS from main is the same as calling exit.

The value 0 also means success, but EXIT\_SUCCESS is not necessarily equal to 0. The EXIT\_SUCCESS macro expands to an integer constant.

## See Also

[exit function](#)

## Name

free function — Releases allocated memory

## Synopsis

```
void free(void* ptr)
```

The free function releases the memory that ptr points to. The pointer must have been returned by a call to malloc or calloc. After freeing the pointer, do not refer to the memory again.

## See Also

[calloc function](#), [malloc function](#), [realloc function](#), delete keyword

## Name

getenv function — Gets environment variable

## Synopsis

```
char* getenv(const char* name)
```



The `getenv` function obtains the value of an environment variable. The environment is a system-defined list of name/value pairs. The `getenv` function searches the environment for `name` and returns the associated value. The `getenv` function returns a null pointer if the specified name is not found.



## Name

labs function — Computes absolute value

## Synopsis

```
long labs(long x)
```

The labs function returns the absolute value of x.

## See Also

[abs function](#), abs function in [<cmath>](#)

## Name

ldiv function — Computes quotient and remainder

## Synopsis

```
ldiv_t ldiv(long numerator, long denominator)
```

The `ldiv` function divides numerator by denominator and returns the quotient and remainder in a structure.

## See Also

[div function](#), [ldiv\\_t type](#)

## Name

ldiv\_t type — Quotient and remainder type

## Synopsis

```
struct ldiv_t { long quot, rem; }
```

The `ldiv_t` type is used only as the return type for the `ldiv` function and the overloaded `div` function. It stores the quotient and remainder of a long integer division.

## See Also

[div\\_t type](#), [ldiv function](#)

## Name

malloc function — Allocates memory

## Synopsis

```
void* malloc(size_t size)
```

The malloc function allocates size bytes of memory. It returns a pointer to the start of the newly allocated memory or a null pointer if there is insufficient memory to fulfill the request. The pointer is suitably aligned for any type.

C++ programs should use the new operator instead of calling malloc.

## See Also

[calloc function](#), [free function](#), [realloc function](#), new keyword

## Name

MB\_CUR\_MAX function — Maximum size of a multibyte character

## Synopsis

```
int MB_CUR_MAX
```

The MB\_CUR\_MAX macro is the maximum number of bytes required to represent a multibyte character in the extended character set, in any locale.

## See Also

[mblen function](#), [mbtowl function](#), [wctomb function](#),  
MB\_LEN\_MAX in [<climits>](#), [<locale>](#)

## Name

`mblen` function — Returns number of bytes in a multibyte character

## Synopsis

```
int mblen(const char* s, size_t n)
```

The `mblen` function returns the length of the multibyte character pointed to by `s`. The character array `s` can be null, it can point to an empty string, or it must have at least `n` bytes, which must form a valid multibyte character.

If `s` is null, the return value depends on whether multibyte characters have state-dependent encodings. (See [Chapter 8](#) for a discussion of shift state.) The `mblen` function returns a nonzero value if encodings are state-dependent or 0 if encodings are not state-dependent.

If `s` points to an empty string, 0 is returned.

If `s` points to a valid multibyte character, the number of bytes that make up that character is returned. If `s` points to an invalid multibyte character, -1 is returned.

## See Also

`MB_CUR_MAX` macro, [mbtowc function](#), `mbrlen` in [<wchar>](#)

## Name

`mbstowcs` function — Converts multibyte string to wide string

## Synopsis

```
size_t mbstowcs(wchar_t* dst, const char* src, size_t n)
```

The `mbstowcs` function converts a multibyte string to a wide character string. The `src` parameter points to the null-terminated multibyte string. Up to `n` wide characters are stored in `dst`. If fewer than `n` characters are stored, a null wide character is appended to the wide character array.

The return value is the number of wide characters stored in `dst`. If any multibyte character in `src` is not valid, the return value is `static_cast<size_t>(-1)`.

## See Also

[mbtowc function](#), `wcstombc` function, `mbcrtowcs` in [<wchar>](#)

## Name

`mbtowc` function — Converts multibyte character to wide character

## Synopsis

```
int mbtowc(wchar_t* pwc, const char* src, size_t n)
```

The `mbtowc` function converts a multibyte character sequence to a single wide character. It starts by counting the number of bytes in `src` that make up the first multibyte character. It examines only the first `n` bytes.

If `src` is null, the return value depends on whether multibyte characters have state-dependent encodings. (See [Chapter 8](#) for a discussion of shift state.) The `mbtowc` function returns a nonzero value if encodings are state-dependent or `0` if encodings are not state-dependent.

If `src` points to an empty string, `0` is returned.

If `src` points to a valid multibyte character, the number of bytes that make up that character is returned. If `dst` is not null, the multibyte character is converted to its equivalent wide character, and the wide character is stored in `*dst`.

If `src` points to an invalid multibyte character, `-1` is returned.

## See Also

[mblen function](#), [mbstowcs function](#), [wctomb function](#), `mbrtowc` in [<wchar>](#)



## Name

NULL macro — NULL pointer constant

## Synopsis

```
#define NULL . . .
```

The NULL macro expands to a null pointer constant. See `<stddef>` for more information.

## See Also

NULL in [<stddef>](#)

# Name

qsort function — Sorts an array

## Synopsis

```
extern "C"  
    void qsort(void* base, size_t count, size_t size, int  
                (*compare)(const void*, const void*))  
extern "C++"  
    void qsort(void* base, size_t count, size_t size, int  
                (*compare)(const void*, const void*))
```

The `qsort` function sorts in ascending order an array of `count` elements, each of size `size` bytes, in which `base` is the pointer to the first element. The array must have POD type. The sort is not stable, that is, the relative order of identical elements is not necessarily preserved.

The compare function takes two pointers into the array and compares the elements. It returns an integer: negative if the first element is less than the second, positive if the first is greater than the second, or 0 if the two elements are equal.

The name `qsort` derives from the original implementation, which used the Quick Sort algorithm. The current standard does not specify which sort algorithm is used, nor does it specify any performance characteristics of the sort algorithm.

Two versions of `qsort` are declared so the compare function can have "C" linkage or "C++" linkage.

## See Also

[bsearch function](#), sort in [<algorithm>](#)

## Name

rand function — Generates a pseudo-random number

## Synopsis

```
int rand( )
```

The rand function returns a pseudo-random integer in the range 0 to RAND\_MAX, inclusive.

## See Also

[RAND\\_MAX macro](#), [srand function](#)

## Name

RAND\_MAX macro — Maximum value returned by rand

## Synopsis

```
int RAND_MAX
```

RAND\_MAX is the maximum value that rand can return. The RAND\_MAX macro expands to an integer constant.

## See Also

[rand function](#)

## Name

realloc function — Reallocates memory

## Synopsis

```
void* realloc(void* ptr, size_t size)
```

The `realloc` function changes the size of the allocated memory that `ptr` points to. The new size is `size` bytes. The return value is a pointer to the newly resized memory block, which might be at a different address than the original block. The pointer is suitably aligned for any type.

The contents of the original memory are preserved, up to the smaller of the new and old sizes. If the new size is larger than the old size, the extra memory above the old size is uninitialized.

The memory might be copied, so you can store only POD values in the memory that you reallocate with `realloc`.

If `ptr` is null, `realloc` behaves just like `malloc(size)`. If `size` is 0, `realloc` is like `free` and frees `ptr`.

If there is insufficient memory to fulfill the request, the original memory is untouched, and a null pointer is returned.

## See Also

[calloc function](#), [free function](#), [malloc function](#), [Chapter 6](#)

## Name

size\_t type — Size type

## Synopsis

```
typedef . . . size_t
```



The `size_t` type is the type of the result of the `sizeof` operator. It is an unsigned integral type. The exact type is implementation-defined.

## See Also

size\_t in [<stddef>](#)

## Name

srand function — Sets seed for pseudo-random number generator

## Synopsis

```
void srand(unsigned int seed)
```

The srand function saves seed as the seed for a new sequence of pseudo-random numbers to be returned by successive calls to rand. The default seed is 1.

## See Also

[rand function](#)

## Name

strtod function — Converts a string to double

## Synopsis

```
double strtod(const char* str, char** end)
```

The `strtod` function converts a character array to a floating-point number. The string `str` is divided into three parts: optional whitespace, the text of the floating-point value, and a trailing part, which starts with the first character that cannot be part of a floating-point number. The first part is skipped, and the second part is converted to a floating-point value. If the second part is empty, `0` is returned. If `end` is not null, `*end` is assigned a pointer to the start of the third part of `str`. If the third part is empty, `*end` points to the terminating null character.

If the result would cause overflow, positive or negative `HUGE_VAL` is returned, and `errno` is set to `ERANGE`. If the result causes underflow, `0` is returned, and `errno` is set to `ERANGE`.

## See Also

[atoi function](#), [strtol function](#), [strtoul function](#), `wcstod` in [<wchar>](#)



## Name

strtol function — Converts a string to a long integer

## Synopsis

```
long int strtol(const char* str, char** end, int base)
```

The `strtol` function converts a character array to a long integer. The string `str` is divided into three parts: optional whitespace, the text of the integer value, and a trailing part, which starts with the first character that cannot be part of an integer. The first part is skipped, and the second part is converted to a long integer. If the second part is empty, 0 is returned. If `end` is not null, `*end` is assigned a pointer to the start of the third part of `str`. If the third part is empty, `*end` points to the terminating null character.

If `base` is 0, the base is determined from the prefix of the integer text: a leading `0x` or `0X` means hexadecimal, a leading `0` means octal, and anything else is decimal.

Otherwise, `base` must be between 2 and 36, in which the letters `a-z` (of either case) represent digits with values of 10-35. Only letters that are appropriate for the base are permitted, that is, the corresponding digit value must be less than the base.

If the resulting value is too large or too small to fit in a `long int`, the value `LONG_MAX` or `LONG_MIN` is returned, and `errno` is set to `ERANGE`.

## See Also

[atol function](#), [strtod function](#), [strtoul function](#), `wcstol` in [<wchar>](#)

## Name

strtoul function — Converts a string to unsigned long

## Synopsis

```
unsigned long strtoul(const char* str, char** end, int base)
```

The `strtoul` function converts a character array to an unsigned long integer. The string `str` is divided into three parts: optional whitespace, the text of the integer value, and a trailing part, which starts with the first character that cannot be part of an integer. The first part is skipped, and the second part is converted to an unsigned long integer. If the second part is empty, `0` is returned. If `end` is not null, `*end` is assigned a pointer to the start of the third part of `str`. If the third part is empty, `*end`` points to the terminating null character.

If `base` is `0`, the base is determined from the prefix of the integer text: a leading `0x` or `0X` means hexadecimal, a leading `0` means octal, and anything else is decimal. Otherwise, `base` must be between 2 and 36, in which the letters a-z (of either case) represent digits with values of 10-35. Only letters that are appropriate for the base are permitted, that is, the corresponding digit value must be less than the base.

If the resulting value is too large to fit in an unsigned long int, the value `ULONG_MAX` is returned, and `errno` is set to `ERANGE`.

## See Also

[atol function](#), [strtod function](#), [strtol function](#), `wcstoul` in [<wchar>](#)

## Name

system function — Runs a program

## Synopsis

```
int system(const char* command)
```



The `system` function passes `command` to the host operating system to run as an external command. The use and interpretation of the command string is implementation-defined.

The return value is implementation-defined.

If `command` is null, the return value is true (nonzero) if a command processor is available; it is false (0) if no command processor is available.

## Name

wctomb function — Converts a wide character to a multibyte character

## Synopsis

```
int wctomb(char* s, wchar_t wc)
```

The `wctomb` function converts a wide character to a multibyte character. It first determines the number of bytes needed to represent `wc` as a multibyte character. If `s` is not null, the sequence of multibyte characters is stored there. At most, `MB_CUR_MAX` bytes are stored, and the return value is the actual number of bytes written to `s`. If `wc` does not have a valid multibyte encoding, `-1` is returned.

If `s` is null, the return value is true (nonzero) if multibyte characters have state-dependent encodings, or false (0) if they do not.

## See Also

[mbtowc function](#), [wcstombs function](#), `wcrtomb` in [<wchar>](#)

## Name

wcstombs function — Converts a wide string to a multibyte string

## Synopsis

```
size_t wcstombs(char* dst, const wchar_t* src, size_t n)
```

The `wcstombs` function converts a wide string `src` to a string `dst` of multibyte characters. At most, `n` bytes of `dst` are written to. If the conversion of `src` requires fewer than `n` bytes, a trailing null byte is appended to `dst`.

If any wide characters cannot be represented as a multibyte character, `static_cast<size_t>(-1)` is returned. Otherwise, the return value is the number of bytes written to `dst` (not counting a trailing null byte).

## See Also

[mbstowcs function](#), [wctomb function](#), `wcsrtoombs` in [<wchar>](#)

## <cstring>

The <cstring> header is for the C++ version of the C standard <string.h> header, which declares string-handling functions.

The functions in this section fall into two categories, identified by the first three letters of the function name:

mem . . .

The mem functions operate on arbitrary chunks of memory, treating the memory as arrays of unsigned char. The caller must specify the size of each memory chunk.

str . . .

The str functions operate on null-terminated character arrays. Even though the function parameters are declared as type char, they are always interpreted as unsigned char when comparing two characters.

See also <wchar> for wide character string functions.

### **Not e**

Instead of using C-style, null-terminated character arrays, C++ code should use the string and wstring classes that are declared in the <string> header. C++ strings offer high performance, more flexibility, more safety, and greater ease of use. The char\_traits class template in <string> also provides member functions for working with narrow and wide character arrays.

## Name

memchr function — Searches for a byte

## Synopsis

```
const void* memchr(const void* mem, int c, size_t n)
void* memchr(void* mem, int c, size_t n)
```

The memchr function searches the memory that mem points to, of size n bytes, for the byte whose value is c (converted to unsigned char). The return value is a pointer into the mem array to the first occurrence of c, or a null pointer if c is not present in the first n bytes of mem.

## See Also

[strchr function](#) , find in  [<algorithm>](#) , wmemchr in  [<wchar>](#)

## Name

memcmp function — Compares memory

## Synopsis

```
int memcmp(const void* s1, const void* s2, size_t n)
```

The memcmp function compares the first n bytes of s1 and s2 as arrays of unsigned char. If all n bytes are equal, the return value is 0. Otherwise, the return value is positive if s1 is greater than s2 or negative if s1 is less than s2.

## See Also

[strcmp function](#), [strncmp function](#), equal in [<algorithm>](#), lexicographical\_compare in [<algorithm>](#), mismatch in [<algorithm>](#), wmemcmp in [<cwchar>](#)



## Name

memcpy function — Copies memory

## Synopsis

```
void* memcpy(void* dst, const void* src, size_t n)
```

The memcpy function copies `n` bytes from `src` to `dst`. If `src` and `dst` overlap, the results are undefined. The return value is `dst`.

If you copy the memory that contains any non-POD objects, the results are undefined. See [Chapter 6](#) for more information about POD objects.

## See Also

[memmove function](#), [strcpy function](#), [strncpy function](#), copy in [<algorithm>](#), wmemcpy in [<wchar>](#), [Chapter 6](#)

## Name

memmove function — Copies possibly overlapping memory

## Synopsis

```
void* memmove(void* dst, const void* src, size_t n)
```

The memmove function copies n bytes from src to dst. The memory regions can overlap. The return value is dst.

If you copy the memory that contains any non-POD objects, the results are undefined. See [Chapter 6](#) for more information about POD objects.

## See Also

[memcpy function](#), [strcpy function](#), [strncpy function](#), copy in [<algorithm>](#), copy\_backward in [<algorithm>](#), wmemmove in [<wchar>](#), [Chapter 6](#)

## Name

memset function — Fills memory with a byte

## Synopsis

```
void* memset(void* s, int c, size_t n)
```

The memset function fills the array s with n copies of c (converted to unsigned char). The return value is s.

## See Also

[memcpy function](#), fill\_n in [<algorithm>](#), wmemset in [<wchar>](#)

## Name

NULL macro — NULL pointer constant

## Synopsis

```
#define NULL . . .
```

The NULL macro expands to a null pointer constant. See `<stddef>` for more information.

## See Also

NULL in [<stddef>](#)

## Name

size\_t type — Size type

## Synopsis

```
typedef . . . size_t
```



The `size_t` type is the type of the result of the `sizeof` operator. It is an unsigned integral type. The exact type is implementation-defined.

## See Also

size\_t in [<stddef>](#)

## Name

strcat function — Concatenates strings

## Synopsis

```
char* strcat(char* dst, const char* src)
```

The `strcat` function concatenates `src` onto the end of `dst`, overwriting the null byte that ends `dst`. The `src` and `dst` arrays cannot overlap. The caller must ensure that `dst` points to a region of memory that is large enough to hold the concatenated result, including its null terminator.

## See Also

[strcpy function](#), [strncat function](#), `wcscat` in [<wchar>](#)

## Name

strchr function — Searches a string for a character

## Synopsis

```
const char* strchr(const char* s, int c)
char* strchr(      char* s, int c)
```

The strchr function returns a pointer to the first occurrence of c (converted to unsigned char) in the null-terminated string s. If c does not appear in s, a null pointer is returned.

## See Also

[memchr function](#), [strcspn function](#), [strpbrk function](#), [strrchr function](#), [strspn function](#), wcschr in [<wchar>](#)

## Name

strcmp function — Compares strings

## Synopsis

```
int strcmp(const char* s1, const char* s2)
```

The strcmp function compares two null-terminated strings as arrays of unsigned char. If the strings are equal, the return value is 0. Otherwise, the return value is positive if s1 is greater than s2 or negative if s1 is less than s2. If one string is a prefix of the other, the longer string is greater than the shorter string.

## See Also

[memcmp function](#), [strncmp function](#), wcscmp in [<wchar>](#)



## Name

strcoll function — Compares strings using locale's collation order

## Synopsis

```
int strcoll(const char* s1, const char* s2)
```

The `strcoll` function compares two null-terminated strings, interpreting the strings according to the `LC_COLLATE` (defined in `<locale>`) category of the current C locale. The return value is the same as that of `strcmp`.

## See Also

[strcmp function](#), `wscoll` in [<wchar>](#), `collate` in [<locale>](#), [<locale>](#)

## Name

strcpy function — Copies a string

## Synopsis

```
char* strcpy(char* dst, const char* src)
```

The `strcpy` function copies the null-terminated string `src` to `dst`. The caller must ensure that `dst` points to a region of memory that is large enough to hold the entire `src` string plus its null terminator. The return value is `dst`.

## See Also

[memcpy function](#), [strncpy function](#), `wcscpy` in [<wchar>](#)

## Name

strcspn function — Counts initial characters not in a span set

## Synopsis

```
size_t strcspn(const char* str, const char* spanset)
```

The strcspn function returns the number of characters at the start of str that are not in the string spanset. Thus, the c in its name means complement, that is, strcspn counts characters that are in the complement of the span set.

## See Also

[strchr function](#), [strpbrk function](#), [strspn function](#), [strstr function](#), wcscspn in [<wchar>](#)

## Name

strerror function — Retrieves error message text

## Synopsis

```
char* strerror(int errnum)
```

The `strerror` function returns a pointer to an error message string that corresponds to the error number `errnum`. The message is the same as that printed by the `perror` function.

A program must not modify the array returned by `strerror`, and subsequent calls to `strerror` can overwrite the array.

## See Also

`perror` in [<cstdio>](#), [<cerrno>](#)

## Name

strlen function — Computes length of a string

## Synopsis

```
size_t strlen(const char* s)
```

The `strlen` function returns the length of the null-terminated string `s`, that is, the number of bytes that come before the null byte at the end of the string.

## See Also

wcslen in [<wchar>](#)

## Name

strncat function — Concatenates strings

## Synopsis

```
char* strncat(char* dst, const char* src, size_t n)
```

The `strncat` function concatenates `src` onto the end of `dst`, overwriting the null byte at the end of `dst`. At most, `n` characters are copied from `src`. A terminating null character is always appended to the end of `dst`. The caller must ensure that `dst` points to a region of memory that is large enough to hold the concatenated result, including the null terminator. The return value is `dst`.

## See Also

[strcat function](#), `wcsncat` in [<wchar>](#)

## Name

strncmp function — Compares strings

## Synopsis

```
int strncmp(const char* s1, const char* s2, size_t n)
```

The `strncmp` function compares at most `n` characters of two null-terminated strings as arrays of unsigned `char`. If the strings are equal, the return value is `0`. Otherwise, the return value is positive if `s1` is greater than `s2` or negative if `s1` is less than `s2`. If one string is a prefix of the other, the longer string is greater than the shorter string.

## See Also

[memcmp function](#), [strcmp function](#), `wcsncmp` in [<wchar>](#)

## Name

strncpy function — Copies a string

## Synopsis

```
char* strncpy(char* dst, const char* src, size_t n)
```

The `strncpy` function copies at most `n` characters from the null-terminated string `src` to `dst`. If `src` is shorter than `dst`, null characters are appended to the end so that exactly `n` characters are always written to `dst`.

The return value is `dst`.

## See Also

[memcpy function](#), [strcpy function](#), `wcsncpy` in [<wchar>](#)



## Name

strpbrk function — Locates a character in a span set

## Synopsis

```
const char* strpbrk(const char* str, const char* spanset)
char* strpbrk(      char* str, const char* spanset)
```

The `strpbrk` function searches `str` for any of the characters in `spanset` and returns a pointer to the first occurrence of such a character. If none of the characters in `spanset` appear in `str`, `strpbrk` returns a null pointer.

## See Also

[strchr function](#), [strcspn function](#), [strspn function](#), `wcspbrk` in [<wchar>](#)

## Name

strrchr function — Locates rightmost occurrence of a character

## Synopsis

```
const char* strrchr(const char* str, int c)
char* strrchr(char* str, int c)
```

The `strrchr` function returns a pointer to the last (rightmost) occurrence of `c` (converted to unsigned char) in the null-terminated string `s`. If `c` does not appear in `s`, `NULL` is returned.

## See Also

[memchr function](#), [strchr function](#), `wcsrchr` in [<wchar>](#)

## Name

strspn function — Counts characters in a span set

## Synopsis

```
size_t strspn(const char* str, const char* spanset)
```

The `strspn` function returns the number of characters at the start of `str` that are in the string spanset.

## See Also

[strchr function](#), [strcspn function](#), [strpbrk function](#), `wcsspn`  
in [<wchar>](#)

## Name

strstr function — Finds a substring

## Synopsis

```
const char* strstr(const char* str, const char* substr)
char* strstr(char* str, const char* substr)
```

The `strstr` function returns the address in `str` of the first occurrence of `substr` or a null pointer if `substr` does not appear in `str`.

## See Also

[strchr function](#), `wcsstr` in [<wchar>](#)

# Name

strtok function — Tokenizes a string

# Synopsis

```
char* strtok(char* str, const char* delimset)
```

The `strtok` function splits `str` into separate tokens, separated by one or more characters from `delimset`. The contents of `str` are modified when each token is found.

To parse a string `str`, you must call `strtok` multiple times. The first time, pass `str` as the first parameter to `strtok`; for the second and subsequent calls, pass a null pointer. Because `strtok` saves `str`, only one series of `strtok` calls can be active at a time. Each call to `strtok` can use a different `delimset`.

The `strtok` function skips over initial delimiters, searching `str` for the first character that is not in `delimset`. If it reaches the end of the string without finding any token characters, it returns a null pointer. Otherwise, it saves a pointer to the first non-delimiter character as the start of the token. It then searches for the next delimiter character, which ends the token. It changes the delimiter character to a null character and returns a pointer to the start of the token. When `strtok` is called with a null pointer as the first parameter, it starts searching for the next token at the point where the previous search ended.

# See Also

[strcspn function](#), [strpbrk function](#), [strspn function](#), `wcstok` in [<wchar>](#)

## Name

strxfrm function — Transforms a string for collation

## Synopsis

```
size_t strxfrm(char* dst, const char* src, size_t n)
```

The `strxfrm` function transforms the `src` string by converting each character to its collation order equivalent. The equivalent is copied into `dst`. Thus, after transforming two different strings with `strxfrm`, the transformed strings can be compared by calling `strcmp` to obtain the same result as calling `strcoll` on the original strings.

No more than `n` bytes are stored in `dst`, including the trailing null character. If `n` is 0, `dst` can be null.

The return value is the number of transformed characters written to `dst`.

## See Also

[strcmp function](#), [strcoll function](#), `wcsxfrm` in [<wchar>](#),  
`collate` in [<locale>](#), [<locale>](#)

## <ctime>

The <ctime> header is the C++ version of the C standard <time.h> header, which declares types and functions for working with dates and times.



The `time_t` type is the fundamental representation of a date and time. The details of this type and how it encodes a date and time are implementation-defined. Unix programmers recognize this type as the number of seconds since January 1, 1970, but that is only one possible implementation.

A date and time have a secondary representation as a `tm` object, which breaks down a date and time into constituent parts. The parts facilitate formatting dates and times for output, or you can read a date and time, parse the parts, and build a `tm` object, which you can then convert to a `time_t` object.

This section describes the types and functions for working with dates and times. The Boost project has additional date and time classes. See [Appendix B](#) for information about Boost.

## Name

asctime function — Converts a time to a string

## Synopsis

```
char* asctime(const tm* tmptr)
```

The `asctime` function formats the date and time pointed to by `tmptr` as a character string. It returns a pointer to a static buffer that is overwritten with each call. (The static buffer can be shared with `ctime`.)

The returned value has the format "Ddd Mmm DD HH:MM:SS YYYY\n" followed by a terminating null character. Thus, the result always has a length of 25. The day of the week (Ddd) and month name (Mmm) are English abbreviations and are not localized—that is, Monday is represented by "Mon" regardless of locale. The day of the month (DD) always takes up the same number of characters, using a leading space if necessary. The hours (HH), minutes (MM), and seconds (SS) use a leading zero if necessary.

## See Also

[ctime function](#), [gmtime function](#), [localtime function](#),  
time\_put in [<locale>](#)



## Name

clock function — Gets the processor time

## Synopsis

```
clock_t clock( )
```



The `clock` function returns the amount of processor time used since an implementation-defined start time. The time is returned in implementation-defined units. There are `CLOCKS_PER_SEC` units per second.

The value returned by the `clock` function is not useful by itself but is intended to be used by comparing it with the value returned from an earlier call to `clock`. For example, the first statement in the `main` function might be a call to `clock`; the difference between subsequent calls and the original call tell you how much time has elapsed since the program started.

If the environment cannot provide the processor time, `static_cast<clock_t>(-1)` is returned.

## See Also

[clock\\_t type](#), [CLOCKS\\_PER\\_SEC macro](#), [time function](#)

## Name

clock\_t type — Represents processor time

## Synopsis

```
typedef . . . clock_t
```

The clock\_t type is an arithmetic type returned by the clock function.

## See Also

[clock function](#)

## Name

CLOCKS\_PER\_SEC macro — Processor time resolution

## Synopsis

```
int CLOCKS_PER_SEC
```

The CLOCKS\_PER\_SEC macro returns the number of clock ticks per second. It is not necessarily a compile-time constant.

## See Also

[clock function](#)

## Name

ctime function — Converts a time to a string

## Synopsis

```
char* ctime(const time_t* timeptr)
```

The ctime function converts the date and time pointed to by timeptr to local time and formats the time as a string. It is equivalent to:

```
std::asctime(std::localtime(timeptr));
```

The text is written to a static buffer, and a pointer to that buffer is returned. Subsequent calls can overwrite the buffer contents. The static buffer can be shared with asctime.

## See Also

[asctime function](#), [localtime function](#)

## Name

difftime function — Computes the difference between two times

## Synopsis

```
double difftime(time_t t1, time_t t2);
```

The difftime function computes the difference between two times: t1 - t2. The return value is in seconds.

## See Also

[time function](#), [time\\_t type](#)

## Name

gmtime function — Extracts parts of a UTC time

## Synopsis

```
tm* gmtime(const time_t* timeptr)
```

The `gmtime` function expands the calendar time pointed to by `timeptr` into a static `tm` object using Coordinated Universal Time (UTC). It returns a pointer to the static object. Subsequent calls to `gmtime` overwrite the object. The static object can be shared with `localtime`.

If UTC is not available (for example, the host operating system does not provide the time zone offset), a null pointer is returned.

## See Also

[localtime function](#), [tm struct](#)

## Name

localtime function — Extracts parts of a local time

## Synopsis

```
tm* localtime(const time_t* timeptr)
```

The `localtime` function expands the calendar time pointed to by `timeptr` into a static `tm` object using local time. It returns a pointer to the static object. Subsequent calls to `localtime` overwrite the object. The static object can be shared with `gmtime`.

## See Also

[gmtime function](#), [tm struct](#)

## Name

mktime function — Makes a time from parts

## Synopsis

```
time_t mktime(tm* tmptr)
```

The mktime function makes a `time_t` time by assembling the parts in a `tm` object, interpreted as local time. The `tm_wday` and `tm_yday` members are ignored, and other fields are permitted to be outside their normal ranges.

If the conversion is successful, the corresponding time is returned, the `tm_wday` and `tm_yday` members are set, and the other fields are changed if necessary to reflect their normal ranges.

If the time cannot be represented as a `time_t` value, `static_cast<time_t>(-1)` is returned.

## See Also

[localtime function](#), [time\\_t type](#), [tm struct](#), `time_get` in [<locale>](#)



## Name

NULL macro — NULL pointer constant

## Synopsis

```
#define NULL . . .
```

The NULL macro expands to a null pointer constant. See `<stddef>` for more information.

## See Also

NULL in [<stddef>](#)

## Name

size\_t type — Size type

## Synopsis

```
typedef . . . size_t
```



The `size_t` type is the type of the result of the `sizeof` operator. It is an unsigned integral type. The exact type is implementation-defined.

## See Also

size\_t in [<stddef>](#)

# Name

strptime function — Formats a time as a string

# Synopsis

```
size_t strptime(char* str, size_t n, const char* fmt, const tm*  
tmpr)
```

The `strptime` function formats a `tm` object as a string. Up to `n` bytes are stored in `str`, including a terminating null character. The return value is the number of characters actually stored, not counting the final null character. If the formatted result requires more than `n` characters, the return value is `0`.

Characters from `fmt` are copied to `str`, except conversion specifiers, which start with a percent sign (%) and are followed by one of the letters shown in [Table 13-6](#). The `LC_TIME` category in the current C locale controls the text that is copied to `str` for each conversion specifier.

*Table 13-6. Conversion specifiers for strptime*

Specified	Description
a	Abbreviated weekday name
A	Full weekday name
b	Abbreviated month name
B	Full month name
C	Complete date and time
D	Day of the month (01-31)

Specified	Description
H	Hour (00-23); 24-hour clock
I	Hour (01-12); 12-hour clock
j	Day of the year (001-366)
m	Month (01-12)
M	Minutes (00-59)
P	A.M./P.M. designation for use with a 12-hour clock
S	Second (00-61); up to two leap seconds
U	Week number (00-53); week 1 starts with the first Sunday
w	Weekday (0-6); Sunday is day 0
W	Week number (00-53); week 1 starts with first Monday
x	Date
X	Time
y	Year in century (00-99)
Y	Year
Z	Time zone name or abbreviation, or empty string if time zone is unknown
%	Literal %

**See Also**

[asctime function](#), [ctime function](#), [tm struct](#), time\_put in  
[<locale>](#), [<ctype>](#)

## Name

time function — Gets the current date and time

## Synopsis

```
time_t time(time_t *timeptr)
```



The time function returns the current date and time in an implementation-defined format. If the host environment cannot provide the date and time, `static_cast<time_t>(-1)` is returned.

If `timeptr` is not null, the return value is also stored in `*timeptr`.

## See Also

[clock function](#), [time\\_t type](#)

## Name

`time_t` type — Represents a date and time

## Synopsis

```
typedef . . . time_t
```



The `time_t` type is an arithmetic type that represents a date and time. The actual type and the encoding of the date and time are implementation-defined.

## See Also

[clock\\_t type](#), [time function](#), [tm struct](#)

# Name

tm struct — Represents the parts of a date and time

# Synopsis

```
struct tm {
    int tm_sec;    /* Seconds: 0-61 */
    int tm_min;    /* Minutes: 0-59 */
    int tm_hour;   /* Hours: 0-23 */
    int tm_mday;   /* Day of month: 1-31 */
    int tm_mon;    /* Month: 0-11 */
    int tm_year;   /* Years since 1900 */
    int tm_wday;   /* Days since Sunday: 0-6 */
    int tm_yday;   /* Days since January 1: 0-365 */
    int tm_isdst;  /* Daylight Savings Time */
}
```

The tm structure stores parts of a date and time. The values returned by `localtime` and `gmtime` will always be in the ranges shown above. (Note that two extra leap seconds are allowed for `tm_sec`.)

The `tm_isdst` member is positive when Daylight Savings Time is in effect, 0 when it is not in effect, or negative if it is unknown.



The order of the members is implementation-defined. An implementation can have additional members.

# See Also

[gmtime function](#), [localtime function](#), [mktime function](#),  
`time_t` type



## <wchar>

The <wchar> header is the C++ version of the C standard <wchar.h> header, which declares types and functions for working with wide characters. Many of these functions are wide versions of functions found in <stdio> and <string> plus improved versions of the multibyte functions from <stdlib>.

You can use narrow (byte-oriented) I/O functions, as declared in <stdio>, or wide I/O functions, as declared in <wchar>, but you cannot mix wide and narrow functions on a single stream without explicitly changing the stream's orientation (see the `fwide` function in this section for details).

Wide I/O treats a file as a sequence of multibyte characters. When reading, multibyte characters are converted to wide characters, and when writing, wide characters are converted to multibyte characters. The conversion depends on the C locale (set with `setlocale` in <locale>).

See [Chapter 1](#) for information about character sets, [Chapter 8](#) for information about multibyte characters and shift states, and [Chapter 9](#) for information about wide characters and I/O.

When working with wide characters, consider using the C++ I/O streams and `wstring` class instead of the C functions. (See <string> later in this chapter for the `wstring` class and the `char_traits` class template.) The <locale> header provides additional support for converting between narrow and wide characters (the `codecvt` and related facets).

## Name

`btowc` function — Converts a multibyte character to a wide character

## Synopsis

```
wint_t btowc(int c)
```

The `btowc` function returns a wide character representation of `c`, which is a multibyte character that can be represented in a single byte (as an unsigned `char`). If `c` is not a valid one-byte, multibyte character, or if `c` is EOF, WEOF is returned.

## See Also

[mbrtowc function](#), [wctob function](#), [WEOF macro](#), `codecv`t in [<locale>](#)

## Name

fgetwc function — Reads a wide character

## Synopsis

```
wint_t fgetwc(FILE* stream)
```

The fgetwc function reads the next wide character from stream. It returns the character read, or WEOF for end-of-file or an error.

## See Also

[fgetws function](#), [fputwc function](#), [getwc macro](#), fgetc in [<cstdio>](#)

## Name

fgetws function — Reads a wide string

## Synopsis

```
wchar_t* fgetwc(wchar_t* str, int n, FILE* stream)
```

The `fgetws` function reads a line of wide characters from `stream` and stores them in `str`. The newline character is also stored. At most, `n` wide characters are stored in `str`, including a terminating null wide character.

The return value is `str` for success or a null pointer for end-of-file or an error.

## See Also

[fgetwc function](#), [fputws function](#), `fgets` in [<stdio>](#)

## Name

fputc function — Writes a wide character

## Synopsis

```
wint_t fputc(wchar_t wc, FILE* stream)
```

The fputc function writes a wide character, `wc`, to `stream`. It returns `wc`, or `WEOF` for an error.

## See Also

[fgetc function](#), [putc macro](#), `fputc` in [<stdio>](#)

## Name

fputws function — Writes a wide string

## Synopsis

```
int fputws(const wchar_t* str, FILE* stream)
```

The fputws function writes the wide string `str` to `stream`. It returns EOF (not WEOF) for an error, or a nonnegative value for success.

## See Also

[fgetws function](#), [fputwc function](#), fputs in [<stdio>](#)

## Name

`fwide` function — Gets or sets stream orientation

## Synopsis

```
int fwide(FILE* stream, int mode);
```

The `fwide` function gets or sets the orientation of `stream`. The orientation is wide or narrow (byte). When a file is opened, it starts without orientation. Calling any wide I/O function on the stream gives it wide orientation. Calling any narrow I/O function on the stream gives it narrow orientation. Mixing narrow and wide functions on a stream results in an error—that is, calling a narrow function on a stream with wide orientation or calling a wide function on a stream with narrow orientation results in an error.

Before performing any I/O on a newly opened stream, you can force the stream's orientation by calling `fwide`. Once the orientation is set, it cannot be changed except by closing and reopening the stream (for example, by calling `freopen` in `<cstdio>`).

If `mode` is positive, the orientation is set to wide. If `mode` is negative, the orientation is set to narrow. If the orientation has already been set, it is not changed, and the stream's true orientation is returned. If `mode` is `0`, the orientation is queried without being changed.

The return value indicates the new orientation: positive for wide, negative for narrow, or `0` if the stream has no orientation.

## See Also

`fopen` in [<cstdio>](#), `freopen` in [<cstdio>](#)

## Name

fwprintf function — Writes formatted data

## Synopsis

```
int fwprintf(FILE* stream, const wchar_t* format, . . . )
```

The `fwprintf` function writes wide output to `stream`, formatted according to the conversion specifiers in `format`. See `fprintf` in `<cstdio>` for more information.

## See Also

`fprintf` in [<cstdio>](#)



## Name

fwscanf function — Reads formatted data

## Synopsis

```
int fwscanf(FILE* stream, const wchar_t* format, . . . )
```

The fwscanf function reads wide input from `stream` and interprets it according to the conversion specifiers in `format`. See `fscanf` in `<cstdio>` for more information.

## See Also

`fscanf` in [<cstdio>](#)

## Name

getwc macro — Reads a wide character

## Synopsis

```
wint_t getwc(FILE* stream)
```

The `getwc` macro reads a wide character from `stream`. It returns the character converted to `wint_t`, or `WEOF` for end-of-file or an error.

## See Also

[fgetwc function](#), [getwchar macro](#), `getc` in [<stdio>](#)

## Name

getwchar macro — Reads a wide character

## Synopsis

```
wint_t getwchar( )
```

The getwchar macro is equivalent to getwc(stdin).

## See Also

[getwc macro](#), getchar in [<cstdio>](#)

## Name

`mbrlen` function — Gets number of bytes in a multibyte character

## Synopsis

```
size_t mbrlen(const char* str, size_t n, mbstate_t* ps)
```

The `mbrlen` function counts the number of bytes needed to complete the next multibyte character that `str` points to. At most, `n` bytes of `str` are examined.

The `ps` parameter points to the shift state, which keeps track of the conversion state between calls to `mbrlen`. If `ps` is a null pointer, an internal shift state is used (which is similar to calling `mblen` in `<cstdlib>`).

The return value is one of the following:

0

If the multibyte character represents the null wide character

```
static_cast<size_t>(-1)
```

If `str` does not point to a valid multibyte character

```
static_cast<size_t>(-2)
```

If `n` is too small

Anything else

If the multibyte character is valid, in which case the value returned is the number of bytes in the multibyte character

## See Also

[mbrtowc function](#), [mbstate\\_t type](#), `mblen` in [<cstdlib>](#)

## Name

`mbrtowc` function — Converts a multibyte character to a wide character

## Synopsis

```
size_t mbrtowc(wchar_t* pwc, const char* str, size_t n,  
mbstate_t* ps)
```

The `mbrtowc` function converts a multibyte character to a wide character. First, it counts the number of bytes needed to complete the next multibyte character that `str` points to. At most, `n` bytes of `str` are examined. If `str` points to a valid multibyte character, that character is converted to a wide character, which is stored in `*pwc`.

The `ps` parameter points to the shift state, which keeps track of the conversion state between calls to `mbrtowc`. If `ps` is a null pointer, an internal shift state is used (which is similar to calling `mbtowc` in `<stdlib.h>`).

The return value is one of the following:

0

If the multibyte character represents the null wide character

```
static_cast<size_t>(-1)
```

If `str` does not point to a valid multibyte character

```
static_cast<size_t>(-2)
```

If `n` is too small

Anything else

If the multibyte character is valid, in which case the value returned is the number of bytes in the multibyte character

## See Also

[mbstate\\_t type](#), mbtowc in [<cstdlib>](#), codecvr in [<locale>](#)

## Name

`mbsinit` function — Determines whether a state is the initial shift state

## Synopsis

```
int mbsinit(const mbstate_t* ps)
```

The `mbsinit` function returns true (nonzero) if `ps` is a null pointer or it points to an `mbstate_t` object that is in the initial shift state; otherwise, it returns false (0).

## See Also

[`mbstate\_t` type](#)

## Name

`mbsrtowcs` function — Converts a multibyte string to a wide string

## Synopsis

```
size_t mbsrtowcs(wchar_t* dst, const char** src, size_t len,
                 mbstate_t* ps)
```

The `mbsrtowcs` converts a multibyte string to a wide character string. The `src` parameter indirectly points to the null-terminated multibyte string, that is, `*src` points to the start of the multibyte string.

If `dst` is not null, up to `len` wide characters are stored in `dst`. If fewer than `len` characters are stored, a trailing null character is appended to the wide character array. If conversion stops upon reaching a null character in the `src` string, a null pointer is assigned to `*src`; otherwise, `*src` is assigned a pointer to the byte immediately past the end of the last multibyte character converted.

The `dst` parameter can be a null pointer, in which case no wide characters are stored and `*src` is not altered, but `ps` is updated and the return value is the same as it would be if `dst` were large enough to hold the entire converted string.

The `ps` parameter points to the shift state, which keeps track of the conversion state between calls to `mbsrtowcs`. If `ps` is a null pointer, an internal shift state is used (which is similar to calling `mbstowcs` in `<cstdlib>`). If the conversion ends without a terminating null character in `*src`, the shift state is reset to an `mbstate_t` initial state.

The return value is the number of wide characters successfully converted. If any multibyte character is not valid, the return value is `static_cast<size_t>(-1)`.



## See Also

[mbrtowc function](#), [mbstate\\_t type](#), mbstowcs in [<cstdlib>](#),  
codecvt in [<locale>](#)

## Name

`mbstate_t` type — Represents a multibyte shift state

## Synopsis

```
typedef . . . mbstate_t
```

The `mbstate_t` type is an opaque, POD type that stores the conversion state used to convert between multibyte and wide characters. The type is implementation-defined, but it is not an array type, so an `mbstate_t` object can be returned from a function.

A value of 0 for an `mbstate_t` object corresponds to the initial shift state, although other values might also represent the initial state. Thus, to initialize an `mbstate_t` object, use a default constructor:

```
std::mbstate_t mbs = std::mbstate_t( );
```

If two `mbstate_t` objects are identical, they represent the same shift state, but the reverse is not necessarily true.

There is no way to compare two `mbstate_t` objects to determine whether they represent the same state, but you can call `mbsinit` to determine whether a state is the initial state.

Typically, you would use an `mbstate_t` object by initializing it to the initial shift state, then passing it to any of the multibyte functions (such as `mbrtowc`) repeatedly. Each call to the multibyte function reads the shift state and uses that information for the conversion, updating the shift state depending on which multibyte characters were provided as input. You should not alter the `mbstate_t` object between calls to the multibyte function.

## See Also

[mbrlen function](#), [mbrtowc function](#), [mbsinit function](#),  
[mbsrtowcs function](#), [wctomb function](#), [wcsrtombs function](#)

## Name

NULL macro — NULL pointer constant

## Synopsis

```
#define NULL . . .
```

The NULL macro expands to a null pointer constant. See `<stddef>` for more information.

## See Also

NULL in [<stddef>](#)

## Name

putwc macro — Writes a wide character

## Synopsis

```
wint_t putwc(wchar_t wc, FILE* stream)
```

The `putwc` macro writes the wide character `wc`. The return value is the character converted to `wint_t`, or `WEOF` for an error.

## See Also

[fputwc function](#), [putwchar macro](#), `putc` in [<cstdio>](#)

## Name

putwchar macro — Writes a wide character to stdout

## Synopsis

```
wint_t putwchar(wchar_t wc)
```

The putwchar macro is equivalent to `puwc(wc, stdout)`.

## See Also

[putwc macro](#), putchar in [<stdio>](#)

## Name

size\_t type — Size type

## Synopsis

```
typedef . . . size_t
```



The `size_t` type is the type of the result of the `sizeof` operator. It is an unsigned integral type. The exact type is implementation-defined.

## See Also

size\_t in [<stddef>](#)

## Name

swprintf function — Writes formatted data to a wide string

## Synopsis

```
int swprintf(wchar_t* dst, size_t n, const wchar_t* format, . . . )
```

The `swprintf` function is similar to `sprintf`, except it stores the formatted output in a wide string, `dst`, and the format is a wide string. Another difference is that `n` is the maximum number of wide characters (including a terminating null wide character) that can be written to `dst`.

The return value is the number of wide characters actually stored in `dst` (not counting the terminating null wide character) or a negative value if the formatted output requires `n` or more characters (not including the terminating null character).

## See Also

[fwprintf function](#), [vswprintf function](#), `sprintf` in [<stdio>](#)



## Name

swscanf function — Reads formatted data from a wide string

## Synopsis

```
int swscanf(const wchar_t* str, const wchar_t* format, . . . )
```

The swscanf function is similar to sscanf, except it reads from a wide string, `str`, and the format string is also wide. Like sscanf, the return value is the number of items converted.

## See Also

[fwscanf function](#), sscanf in [<cstdio>](#)

# Name

tm struct — Represents the parts of a date and time

# Synopsis

```
struct tm {
    int tm_sec;    /* Seconds: 0-61 */
    int tm_min;    /* Minutes: 0-60 */
    int tm_hour;   /* Hours: 0-24 */
    int tm_mday;   /* Day of month: 1-31 */
    int tm_mon;    /* Month: 1-12 */
    int tm_year;   /* Years since 1900 */
    int tm_wday;   /* Days since Sunday: 0-6 */
    int tm_yday;   /* Days since January 1: 0-365 */
    int tm_isdst;  /* Daylight Savings Time */
}
```

The tm structure stores parts of a date and time. It is the same structure definition as that found in `<ctime>`. See `<ctime>` for details.

# See Also

tm struct in [<ctime>](#)

## Name

ungetwc function — Pushes back a wide character

## Synopsis

```
wint_t ungetwc(wint_t wc, FILE* stream)
```

The `ungetwc` function pushes back the wide character `wc`, so the next read from `stream` will return `wc`. The standard guarantees that you can push back just one character, though in some situations you may be able to push back more.

The return value is `wc` if the pushback was successful, or `WEOF` if the pushback was not successful.

## See Also

[fgetwc function](#), `getwc` function, `ungetc` in [<cstdio>](#)

## Name

vfwprintf function — Writes formatted data

## Synopsis

```
int vfwprintf(FILE* stream, const wchar_t* format, va_list arg)
```

The vfwprintf function is similar to vfprintf in `<stdio.h>`, except it prints wide characters to stream, and the format parameter is a wide string.

## See Also

vfprintf in [<stdio.h>](#), [<stdio.h>](#)

## Name

vswprintf function — Writes formatted data to a wide string

## Synopsis

```
int vswprintf(wchar_t* dst, size_t n, const wchar_t* format,  
va_list arg)
```

The vswprintf function is similar to vsprintf in `<cstdio>`, except it stores its output in a wide string, `dst`, and the format parameter is a wide string. Another difference is that no more than `n` wide characters are written to `dst`, including a terminating null character.

## See Also

[swprintf function](#), vsprintf in [<cstdio>](#), [<cstdarg>](#)

## Name

vwprintf function — Writes formatted data

## Synopsis

```
int vwprintf(const wchar_t* format, va_list arg)
```

The `vwprintf` function is similar to `vprintf` in `<cstdio>`, except it prints wide characters to `stdout`, and the format parameter is a wide string.

## See Also

[wprintf function](#), `vprintf` in [<cstdio>](#), [<cstdarg>](#)

## Name

WCHAR\_MAX macro — Largest value of a wide character

## Synopsis

```
wchar_t WCHAR_MAX
```

The WCHAR\_MAX macro is the largest value that can be represented by the wchar\_t type. It is not necessarily a valid character in the extended character set.

## See Also

[WCHAR\\_MIN macro](#), CHAR\_MAX in [<climits>](#), [<limits>](#)

## Name

WCHAR\_MIN macro — Smallest value of a wide character

## Synopsis

```
wchar_t WCHAR_MIN
```

The WCHAR\_MIN macro is the smallest value that can be represented by the wchar\_t type. It is not necessarily a valid character in the extended character set.

## See Also

[WCHAR\\_MAX macro](#), CHAR\_MIN in [<climits>](#), <limits>



## Name

wcrtomb function — Converts a wide character to a multibyte character

## Synopsis

```
size_t wcrtomb(char* dst, wchar_t wc, mbstate_t* ps)
```

The `wcrtomb` function converts a wide character to a multibyte character. It first determines the number of bytes needed to represent `wc` as a multibyte character. If `dst` is not null, the sequence of multibyte characters is stored there. At most, `MB_CUR_MAX` (defined in `<cstdlib>`) bytes are stored, and the return value is the actual number of bytes written to `dst`. If `wc` does not have a valid multibyte encoding, `static_cast<size_t>(-1)` is returned.

If `dst` is null, `wcrtomb` ignores `wc` and converts the null wide character using a private, internal buffer (e.g., `wcrtomb( buffer, L'\0', ps)`).

The `ps` parameter points to the shift state, which keeps track of the conversion state between calls to `wcrtomb`. If `ps` is null, an internal shift state is used (which is similar to calling `wctomb` in `<cstdlib>`).

## See Also

[mbrtowc function](#), [mbstate\\_t type](#), `MB_CUR_MAX` in [<cstdlib>](#), `wctomb` in [<cstdlib>](#), `codecv`t in [<locale>](#)

## Name

wcscat function — Concatenates wide strings

## Synopsis

```
wchar_t* wcscat(wchar_t* dst, const wchar_t* src)
```

The `wcscat` function concatenates `src` onto the end of `dst`, overwriting the null character at the end of `dst`. The caller must ensure that `dst` points to a region of memory that is large enough to hold the entire string plus its null terminator. The return value is `dst`.

## See Also

[wcscpy function](#), [wcsncat function](#), `strcat` in [<cstring>](#)

## Name

wcschr function — Searches for a wide character in a wide string

## Synopsis

```
const wchar_t* wcschr(const wchar_t* str, wchar_t wc)
wchar_t* wcschr(      wchar_t* str, wchar_t wc)
```

The `wcschr` function returns a pointer to the first occurrence of `wc` in the null-terminated wide string `str`. If `wc` does not appear in `str`, a null pointer is returned.

## See Also

[wmemchr function](#), [wcsncpy function](#), [wcpbrk function](#), [wchrchr function](#), [wcsspn function](#), `strchr` in [<cstring>](#)

## Name

wcscmp function — Compares wide strings

## Synopsis

```
int wcscmp(const wchar_t* s1, const wchar_t* s2)
```

The `wcscmp` function compares two null-terminated wide strings. If the strings are equal, the return value is 0. Otherwise, the return value is positive if `s1` is greater than `s2` or negative if `s1` is less than `s2`. If one string is a prefix of the other, the longer string is greater than the shorter string.

## See Also

[wmemcmp function](#), [wcsncmp function](#), `strcmp` in [<cstring>](#)

## Name

wscoll function — Compares wide strings using locale's collation order

## Synopsis

```
int wscoll(const wchar_t* s1, const wchar_t* s2)
```

The `wscoll` function compares two null-terminated wide strings, interpreting the strings according to the `LC_COLLATE` (defined in `<locale>`) category of the current C locale. The return value is the same as that of `wscmp`.

## See Also

[wscmp function](#), `strcoll` in [<cstring>](#), [<locale>](#),  
`collate` in [<locale>](#)

## Name

wcscpy function — Copies wide strings

## Synopsis

```
wchar_t* wcscpy(wchar_t* dst, const wchar_t* src)
```

The `wcscpy` function copies the null-terminated wide string `src` to `dst`. The caller must ensure that `dst` points to a region of memory that is large enough to hold the entire `src` string plus its null terminator. The return value is `dst`.

## See Also

[wmemcpy function](#), [wcsncpy function](#), `strcpy` in [<cstring>](#)

## Name

wcscspn function — Counts initial characters that do not match a span set

## Synopsis

```
size_t wcscspn(const wchar_t* str, const wchar_t* spanset)
```

The wcscspn function returns the number of wide characters at the start of `str` that are not in the wide string spanset. Thus, the `c` in its name means complement, that is, wcscspn counts characters that are in the complement of the span set.

## See Also

[wcschr function](#), [wcpbrk function](#), [wcspn function](#), [wcsstr function](#), `strspn` in [<cstring>](#)

## Name

wcsftime function — Formats a time as a wide string

## Synopsis

```
size_t wcsftime(wchar_t* str, size_t n, const wchar_t* format,  
const tm* tmptr)
```

The `wcsftime` function is similar to `strftime` in `<ctime>`, except it formats the result as a wide string, `str`, and the format parameter is a wide string.

## See Also

`strftime` in [<ctime>](#)



## Name

wcslen function — Gets length of a wide string

## Synopsis

```
size_t wcslen(const wchar_t* str)
```

The wcslen function returns the number of wide characters (not including the terminating null wide character) in str.

## See Also

strlen in [<cstring>](#)

## Name

wcsncat function — Concatenates wide strings

## Synopsis

```
wchar_t* wcsncat(wchar_t* dst, const wchar_t* src, size_t n)
```

The `wcsncat` function concatenates `src` onto the end of `dst`. At most, `n` wide characters are copied from `src`. A terminating null wide character is always appended to the end of `dst`. You must ensure that `dst` points to a region of memory that is large enough to hold the concatenated result plus the null terminator. The return value is `dst`.

## See Also

[wcsncat function](#), `strncat` in [<cstring>](#)

## Name

wcsncmp function — Compares wide strings

## Synopsis

```
int wcsncmp(const wchar_t* s1, const wchar_t* s2, size_t n)
```

The `wcsncmp` function compares at most `n` wide characters of two null-terminated wide strings. If the strings are equal, the return value is `0`. Otherwise, the return value is positive if `s1` is greater than `s2` or negative if `s1` is less than `s2`. If one string is a prefix of the other, the longer string is greater than the shorter string.

## See Also

[wcscmp function](#), `strncmp` in [<cstring>](#)

## Name

wcsncpy function — Copies wide strings

## Synopsis

```
wchar_t* wcsncpy(wchar_t* dst, const wchar_t* src, size_t n)
```

The `wcsncpy` function copies at most `n` wide characters from the null-terminated wide string `src` to `dst`. If `src` is shorter than `dst`, null wide characters are appended to the end so that exactly `n` characters are always written to `dst`.

The return value is `dst`.

## See Also

[wcscopy function](#), `strncpy` in [<cstring>](#)

## Name

wcspbrk function — Locates a span set member in a wide string

## Synopsis

```
const wchar_t* wcspbrk(const wchar_t* str, const wchar_t*  
spanset)  
wchar_t* wcspbrk(      wchar_t* str, const wchar_t*  
spanset)
```

The `wcspbrk` function searches `str` for any of the wide characters in `spanset` and returns a pointer to the first occurrence of such a character. If none of the characters in `spanset` appears in `str`, `wcspbrk` returns a null pointer.

## See Also

[wchr function](#), [wcscspn function](#), [wcssp function](#),  
`strpbrk` in [<cstring>](#)

## Name

wcsrchr function — Locates rightmost occurrence of a wide character

## Synopsis

```
const wchar_t* wcsrchr(const wchar_t* str, wchar_t wc)
wchar_t* wcsrchr(wchar_t* str, wchar_t wc)
```

The wcsrchr function returns a pointer to the last (rightmost) occurrence of wc in the null-terminated wide string str. If wc does not appear in str, the function returns a null pointer .

## See Also

[wmemchr function](#), [wcschr function](#), strchr in [<cstring>](#)

## Name

wcsrtombs function — Converts a wide string to a multibyte string

## Synopsis

```
size_t wcsrtombs(char* dst, const wchar_t** src, size_t len,
mbstate_t* ps)
```

The `wcsrtombs` function converts a wide string to a string of multibyte characters. The `src` parameter points indirectly to the source wide string, that is, `*src` points to the start of the wide string.

If `dst` is not null, up to `len` bytes are stored in `dst`. If fewer than `len` bytes are stored, a trailing null character is appended to the narrow character array. If conversion stops upon reaching a null wide character in the `src` string, a null pointer is assigned to `*src`; otherwise, `*src` is assigned a pointer to the character immediately past the end of the last wide character converted.

The `dst` parameter can be null, in which case no narrow characters are stored and `*src` is not altered, but `ps` is updated and the return value is the same as it would be if `dst` were large enough to hold the entire converted string.

The `ps` parameter points to the shift state, which keeps track of the conversion state between calls to `wcsrtombs`. If `ps` is null, an internal shift state is used (which is similar to calling `wcstombs` in `<cstdliblib>`). If the conversion ends without a terminating null character in `*src`, the shift state is reset to an `mbstate_t` initial state.

If any of the wide characters cannot be represented as a multibyte character, `static_cast<size_t>(-1)` is returned. Otherwise, the return value is the number of

bytes successfully converted from wide characters (not counting the trailing null byte).

## See Also

[mbstate\\_t type](#), `wmbsrtowcs` function, [wrtomb function](#), `wcstombs` in [<cstdlib>](#), `codecv`t in [<locale>](#)



## Name

wcsspn function — Counts characters that match a span set

## Synopsis

```
size_t wcsspn(const wchar_t* str, const wchar_t* spanset)
```

The wcsspn function returns the number of wide characters at the start of `str` that are in the string `spanset`.

## See Also

[wcschr function](#), [wcscspn function](#), [wcpbrk function](#),  
strspn in [<cstring>](#)

## Name

wcsstr function — Finds a wide substring

## Synopsis

```
const wchar_t* wcsstr(const wchar_t* str, const wchar_t* substr)
wchar_t* wcsstr(wchar_t* str, const wchar_t* substr)
```

The `wcsstr` function returns the index in `str` of the first occurrence of `substr`, or a null pointer if `substr` does not appear in `str`.

## See Also

[wcschr function](#), `strstr` in [<cstring>](#)

## Name

wcstod function — Converts a wide string to double

## Synopsis

```
double wcstod(const wchar_t* str, wchar_t** end)
```

The `wcstod` function converts a wide string to double. It is similar to the `strtod` function.

## See Also

[wcstol function](#), [wcstoul function](#), `strtod` in [<cstdliblib>](#)

# Name

wcstok function — Tokenizes a wide string

# Synopsis

```
wchar_t* wcstok(wchar_t* str, const wchar_t* delimset, wchar_t**  
ptr)
```

The `wcstok` function is similar to `strtok` in `<cstring>`, except it works with wide strings. Another difference is that it is reentrant, taking a third parameter, `ptr`, which is the address of a wide string. The `wcstok` function uses `ptr` for storing working information, which it uses when `str` is null.

To parse a string `str`, you must call `wcstok` multiple times. The first time, pass `str` as the first parameter to `wcstok`; for the second and subsequent calls, pass a null pointer. For the final argument, `ptr`, pass the address of a `wchar_t*` object. For subsequent calls to `wcstok` (when `str` is null), pass the address of the same `ptr` object. Do not modify `ptr` between successive calls to `wcstok` when parsing a single wide string.

Each call to `wcstok` can use a different `delimset`.

# See Also

[wcscspn function](#), [wcspbrk function](#), [wcspn function](#), `strtok` in [<cstring>](#)

## Name

wcstol function — Converts a wide string to a long integer

## Synopsis

```
long int wcstol(const wchar_t* str, wchar_t** end)
```

The wcstol function converts a wide string to long int. It is similar to the strtol function in <cstdlib>.

## See Also

[wcstod function](#), [wcstoul function](#), strtol in [<cstdlib>](#)

## Name

wcstoul function — Converts a wide string to an unsigned long integer

## Synopsis

```
unsigned long int wcstoul(const wchar_t* str, wchar_t** end)
```

The `wcstoul` function converts a wide string to unsigned long int. It is similar to the `strtoul` function in `<cstdlib>`.

## See Also

[wcstod function](#), [wcstol function](#), `strtoul` in [<cstdlib>](#)

## Name

wcsxfrm function — Transforms a wide string for collation

## Synopsis

```
size_t wcsxfrm(wchar_t* dst, const wchar_t* src, size_t n)
```

The `wcsxfrm` function transforms the `src` wide string by converting each wide character to its collation order equivalent. The functionality and return value are similar to `strxfrm` in `<cstring>`, except `wcsxfrm` works with wide strings.

## See Also

[wcscmp function](#), [wcscoll function](#), `strxfrm` in [<cstring>](#), `collate` in [<locale>](#), [<clocale>](#)

## Name

wctob function — Converts a wide character to a single byte

## Synopsis

```
int wctob(wint_t wc)
```

If the wide character `wc` has a single-byte representation as a multibyte character, `wctob` returns that byte; otherwise, it returns EOF.

## See Also

[btowc function](#), EOF in [<stdio>](#), `codecv`t in [<locale>](#)



## Name

WEOF macro — End-of-file or error

## Synopsis

```
wint_t WEOF
```

The WEOF macro expands to a constant integer value that does not correspond to any valid wide character value. Unlike EOF, WEOF is not guaranteed to be negative.

## See Also

[wint\\_t type](#), EOF in [<stdio.h>](#)

## Name

wint\_t type — Integer representation of a wide character

## Synopsis

```
typedef . . . wint_t
```

The wint\_t type is an integral type that represents wide characters. It can hold the value for any character in the extended character set plus the value WEOF.

## See Also

[WEOF macro](#)

## Name

wmemchr function — Searches for a wide character

## Synopsis

```
const wchar_t* wmemchr(const wchar_t* mem, wchar_t c, size_t n)
wchar_t* wmemchr(wchar_t* mem, wchar_t c, size_t n)
```

The `wmemchr` function searches the memory that `mem` points to, of size `n` wide characters, for the wide character whose value is `c`. The return value is a pointer in the `mem` array that points to the first occurrence of `c`, or a null pointer if `c` is not present in the first `n` wide characters of `mem`.

## See Also

[wcschr function](#), `find` in [<algorithm>](#), `memchr` in [<cstring>](#)

## Name

wmemcmp function — Compares wide strings

## Synopsis

```
int wmemcmp(const wchar_t* s1, const wchar_t* s2, size_t n)
```

The wmemcmp function compares the first n wide characters of s1 and s2. If all n wide characters are equal, the return value is 0. Otherwise, the return value is positive if s1 is greater than s2 or negative if s1 is less than s2.

## See Also

[wcscmp function](#), [wcsncmp function](#), equal in [<algorithm>](#), lexicographical\_compare in [<algorithm>](#), mismatch in [<algorithm>](#), memcmp in [<cstring>](#)

## Name

wmemcpy function — Copies wide strings

## Synopsis

```
wchar_t* wmemcpy(wchar_t* dst, const wchar_t* src, size_t n)
```

The `wmemcpy` function copies `n` wide characters from `src` to `dst`. If `src` and `dst` overlap, the results are undefined. The return value is `dst`.

## See Also

[wcscpy function](#), [wcsncpy function](#), [wmemmove function](#),  
copy in [<algorithm>](#), memcpy in [<cstring>](#)

## Name

wmemmove function — Copies overlapping wide strings

## Synopsis

```
wchar_t* memmove(wchar_t* dst, const wchar_t* src, size_t n)
```

The wmemmove function copies n wide characters from src to dst. The memory regions can overlap. The return value is dst.

## See Also

[wcscpy function](#), [wcsncpy function](#), [wmemcpy function](#),  
copy in [<algorithm>](#), copy\_backward in [<algorithm>](#),  
memmove in [<cstring>](#)

## Name

wmemset function — Fills a wide string with an integer

## Synopsis

```
wchar_t* wmemset(wchar_t* str, wchar_t wc, size_t n)
```

The wmemset function fills the array `str` with `n` copies of the wide character `wc`. The return value is `str`.

## See Also

[wmemcpy function](#), `fill_n` in [<algorithm>](#), `memset` in [<cstring>](#)

## Name

wprintf function — Writes formatted wide data

## Synopsis

```
int wprintf(const wchar_t* format, . . . )
```

The `wprintf` function is similar to `printf` in `<stdio>`, except it prints wide characters, and the format parameter is a wide string.

## See Also

`wfprintf` function, `wsprintf` function, `wvprintf` function, `printf` in [<stdio>](#)



## Name

wscanf function — Reads formatted wide data

## Synopsis

```
int wscanf(const wchar_t* format, . . . )
```

The wscanf function is similar to scanf in <stdio>, except it reads wide characters, and the format parameter is a wide string.

## See Also

wfscanf function, wsscanf function, scanf in [<stdio>](#)

## `<cwctype>`

The `<cwctype>` header is the C++ version of the C standard `<wctype.h>` header, which declares types and functions for classifying and converting wide characters.

Most of the functions in this header are wide equivalents of functions found in `<cctype>`. For example, `iswalnum` determines whether a wide character is alphanumeric, just as `isalnum` determines whether a narrow (byte) character is alphanumeric. The behavior of the wide functions is similar to their narrow equivalents. In particular, for any narrow character `c`, its wide character equivalent `wc`, and classification functions `isxyz` and `iswxyz`, if `isxyz(c)` is true, then `iswxyz(wc)` is true and vice versa. The only exception is that `iswgraph` and `iswpunct` behave slightly differently than `isgraph` and `ispunct` for whitespace characters other than ' '.

The behavior of the `<cwctype>` functions depend on the C locale, as set with the `setlocale` function in `<locale>`. For more flexibility in dealing with multiple locales, you can use C++ locales, in particular the `cctype` facet in `<locale>`.

## Name

iswalnum function — Determines whether a wide character is alphanumeric

## Synopsis

```
int iswalnum(wint_t wc)
```

The `iswalnum` function returns true (nonzero) if either `iswalpha(wc)` or `iswdigit(wc)` is true.

## See Also

[iswalpha function](#), [iswdigit function](#), `isalnum` in [<cctype>](#)

## Name

`iswalpha` function — Determines whether a wide character is alphabetic

## Synopsis

```
int iswalpha(wint_t wc)
```

The `iswalpha` function returns true (nonzero) if `wc` is an alphabetic character, that is, a wide character for which `iswcntrl`, `iswdigit`, `iswpunct`, and `iswspace` all return false (0).

## See Also

[iswcntrl function](#), [iswdigit function](#), [iswpunct function](#), [iswspace function](#), `isalpha` in [<cctype>](#)

## Name

`iswcntrl` function — Determines whether a wide character is a control character

## Synopsis

```
int iswcntrl(wint_t wc)
```

The `iswcntrl` function returns true (nonzero) if `wc` is a wide control character, that is, a wide character for which `iswalnum`, `iswpunct`, and `iswspace` all return false (0).

## See Also

[iswalnum function](#), [iswpunct function](#), [iswspace function](#), `iscntrl` in [<cctype>](#)

## Name

`iswctype` function — Tests any category of a wide character

## Synopsis

```
int iswctype(wint_t wc, wctype_t desc)
```

The `iswctype` function tests any category of the wide character `wc`. The category to test is specified by `desc`, which must be obtained by calling `wctype`. The setting of the `LC_CTYPE` category must be the same for the call to `iswctype` and the call to `wctype` that returned `desc`.

Using `iswctype`, you can implement all the `isw...` functions. For example, you can implement the `iswalnum` function as follows:

```
int iswalnum(wint_t wc)
{
    return std::iswctype(wc, std::wctype("alnum"));
}
```

## See Also

[wctype function](#), [wctype\\_t type](#)

## Name

iswdigit function — Determines whether a wide character is a digit

## Synopsis

```
int iswdigit(wint_t wc)
```

The `iswdigit` function returns true (nonzero) if `wc` is a decimal digit character—that is, '0'-'9'—regardless of locale.

## See Also

[iswxdigit function](#), `isdigit` in [<cctype>](#)

## Name

`iswgraph` function — Determines whether a wide character is graphic

## Synopsis

```
int iswgraph(wint_t wc)
```

The `iswgraph` function returns true (nonzero) if `iswprint(wc)` is true and `iswspace(wc)` is false. Note that `iswgraph` is slightly different from `isgraph` in that it returns true for whitespace characters other than ' '.

## See Also

[iswprint function](#), [iswspace function](#)



## Name

`iswlower` function — Determines whether a wide character is lowercase

## Synopsis

```
int iswlower(wint_t wc)
```

The `iswlower` function returns true (nonzero) if `wc` is a lowercase character.

## See Also

[iswalph](#) function, [iswupper](#) function, [towlower](#) function, `islower` in [<cctype>](#)

## Name

iswprint function — Determines whether a wide character is printable

## Synopsis

```
int iswprint(wint_t wc)
```

The `iswprint` function returns true (nonzero) if `wc` is a printable wide character.

## See Also

[iswgraph function](#), `isprint` in [<cctype>](#)

## Name

iswpunct function — Determines whether a wide character is punctuation

## Synopsis

```
int iswpunct(wint_t wc)
```

The iswpunct function returns true (nonzero) if iswalnum(wc), iswcntrl(wc), and iswspace(wc) are false.

## See Also

[iswalnum function](#), [iswcntrl function](#), [iswspace function](#),  
ispunct in [<cctype>](#)

## Name

iswspace function — Determines whether a wide character is whitespace

## Synopsis

```
int iswspace(wint_t wc)
```

The `iswspace` function returns true (nonzero) if `wc` is a whitespace character.

## See Also

[iswgraph function](#), [iswprint function](#), `isspace` in [<cctype>](#)

## Name

`iswupper` function — Determines whether a wide character is uppercase

## Synopsis

```
int iswupper(wint_t wc)
```

The `iswupper` function returns true (nonzero) if `wc` is an uppercase character.

## See Also

[iswalph](#) function, [iswlower](#) function, [towupper](#) function, `isupper` in [<cctype>](#)

## Name

`iswxdigit` function — Determines whether a wide character is a hexadecimal digit

## Synopsis

```
int iswxdigit(wint_t wc)
```

The `iswxdigit` function returns true (nonzero) if `wc` is a hexadecimal digit character—that is, '0'-'9', 'a'-'f', or 'A'-'F'—regardless of locale.

## See Also

[iswdigit function](#), `isxdigit` in [<cctype>](#)

## Name

towctrans function — Translates a wide character's case

## Synopsis

```
wint_t towctrans(wint_t wc, wctrans_t desc)
```

The towctrans function translates the wide character `wc` according to the description `desc`, which was returned from the `wctrans` function. For example, the `tolower` function can be implemented using `towctrans`:

```
wint_t tolower(wint_t wc)
{
    return std::towctrans(wc, std::wctrans("tolower"));
}
```

## See Also

[wctrans function](#), [wctrans\\_t type](#)

## Name

`tolower` function — Converts a wide character to lowercase

## Synopsis

```
wint_t tolower(wint_t wc)
```

The `tolower` function maps the wide character `wc` to lowercase. If `iswupper(wc)` is false, or if `wc` has no lowercase mapping, `wc` is returned unchanged.

## See Also

[towctrans function](#), [towupper function](#), `tolower` in [<cctype>](#)



## Name

towupper function — Converts a wide character to uppercase

## Synopsis

```
wint_t towupper(wint_t wc)
```

The towupper function maps the wide character `wc` to uppercase. If `iswlower(wc)` is false, or if `wc` has no uppercase mapping, `wc` is returned unchanged.

## See Also

[towctrans function](#), [tolower function](#), toupper in [<cctype>](#)

## Name

wctrans function — Construct a wctrans\_t object

## Synopsis

```
wctrans_t wctrans(const char* property)
```

The wctrans function constructs a wctrans\_t object according to the given property. [Table 13-7](#) lists the properties defined in the standard.

*Table 13-7. Character translation properties*

Property	Description
"tolower"	Maps from uppercase to lowercase
"toupper"	Maps from lowercase to uppercase

## See Also

[towctrans function](#), [wctrans\\_t type](#)

## Name

wctrans\_t type — Represents a wide character translation

## Synopsis

```
typedef . . . wctrans_t
```

The wctrans\_t type is a scalar type used to represent character mappings for the towctrans function.

## See Also

[towctrans function](#), [wctrans function](#), [<locale>](#)

# Name

wctype function — Constructs a wctype\_t object

## Synopsis

```
wctype_t wctype(const char* property)
```

The wctype function constructs a wctype\_t object that describes wide characters that have the given property. [Table 13-8](#) lists the properties that are supported by this standard.

*Table 13-8. Character classification properties*

Property	Description
"alnum"	Alphanumeric
"alpha"	Alphabetic
"cntrl"	Control
"digit"	Digit
"graph"	Non-space printable
"lower"	Lowercase
"print"	Printable or whitespace
"punct"	Punctuation
"space"	Whitespace
"upper"	Uppercase
"xdigit"	Hexadecimal digit

## See Also

[iswctype function](#), [wctype\\_t type](#), [<locale>](#)

## Name

wctype\_t type — Represents a character classification

## Synopsis

```
typedef . . . wctype_t
```

The wctype\_t type is a scalar type used to represent character classifications for the iswctype function.

## See Also

[iswctype function](#), [wctype function](#), [<locale>](#)

## Name

WEOF macro — End-of-file or error

## Synopsis

```
wint_t WEOF
```

The WEOF macro expands to a constant integer value that does not correspond to any valid wide character value. Unlike EOF, WEOF is not guaranteed to be negative.

## See Also

WEOF in [<wchar>](#), EOF in [<stdio>](#)

## Name

wint\_t type — Integer representation of a wide character

## Synopsis

```
typedef . . . wint_t
```

The wint\_t type is an integral type that represents wide characters. It can hold the value for any character in the extended character set plus the value WEOF.

## See Also

wint\_t in [<wchar>](#)



## <deque>

The <deque> header is one of the standard container template headers. It declares the deque class template and a few global functions that operate on deque objects.

A *deque*, short for double-ended queue, is similar to a vector, but the performance is constant when adding to or removing from the collection at the beginning and at the end.

If you need a vector of `bool` that behaves as a normal C++ container, you should use `deque<bool>` instead of `vector<bool>`. See <vector> later in this chapter for an explanation.

See [Chapter 10](#) for information about containers in general.

# Name

deque class template — Double-ended queue

## Synopsis

```
template <class T, class Alloc = allocator<T> >
class deque {
public:
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef T value_type;
    typedef Alloc allocator_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;

    explicit deque(const Alloc& = Alloc( ));
    explicit deque(size_type n, const T& value = T( ), const
Alloc& = Alloc( ));
    template <class InputIterator>
    deque(InputIterator first, InputIterator last, const Alloc& =
Alloc( ));
    deque(const deque<T,Alloc>& x);

    ~deque( );

    deque<T,Alloc>& operator=(const deque<T,Alloc>& x);
    template <class InputIterator>
    void assign(InputIterator first, InputIterator last);
    void assign(size_type n, const T& t);
    allocator_type get_allocator( ) const;

    iterator begin( );
    const_iterator begin( ) const;
    iterator end( );
    const_iterator end( ) const;
    reverse_iterator rbegin( );
    const_reverse_iterator rbegin( ) const;
```

```

reverse_iterator rend( );
const_reverse_iterator rend( ) const;

size_type size( ) const;
size_type max_size( ) const;
void resize(size_type sz, T c = T( ));
bool empty( ) const;

reference operator[](size_type n);
const_reference operator[](size_type n) const;
reference at(size_type n);
const_reference at(size_type n) const;
reference front( );
const_reference front( ) const;
reference back( );
const_reference back( ) const;

void push_front(const T& x);
void push_back(const T& x);
iterator insert(iterator position, const T& x);
void insert(iterator position, size_type n, const T& x);
template <class InputIterator>
void insert (iterator position, InputIterator first,
InputIterator last);
void pop_front( );
void pop_back( );
iterator erase(iterator position);
iterator erase(iterator first, iterator last);
void swap(deque<T,Alloc>&);
void clear( );
};

```

The deque class template represents a double-ended queue. It is one of the standard container types, like list and vector. Like a list, a deque yields amortized, constant performance when adding and removing items from the beginning and end of the container. Like a vector, performance is constant when accessing items at any index in the deque. Performance for inserting or removing items not at the start or end is linear with respect to the size of the container.

After inserting items at the beginning or end of the deque, all iterators become invalid. All references and pointers to

items in the deque remain valid. After inserting in the middle of the deque, all iterators, references, and pointers to items in the deque become invalid.

After erasing an element from the beginning or end of the deque, all iterators and references remain valid, except those pointing to the erased element. After erasing an element from the middle of the deque, all iterators, references, and pointers to items in the deque become invalid.

```
explicit deque (const Alloc& = Alloc( ))
```

Constructs an empty deque.

```
explicit deque (size_type n, const T& value = T( ), const Alloc& = Alloc( ))
```

Constructs a deque with n copies of value.

```
template <class InputIterator>, deque  
(InputIterator first, InputIterator last, const  
Alloc& alloc = Alloc( ))
```

Constructs a deque with copies of the elements in [first, last), unless InputIterator is an integral type, in which case the deque is constructed as though the arguments were cast as follows:

```
deque(static_cast<size_type>(first), static_cast<value_type>  
(last),  
alloc);
```

```
template <class InputIterator>, void assign  
(InputIterator first, InputIterator last)
```

Erases the current contents of the deque and inserts the elements in [first, last), unless InputIterator is an integral type, in which case the arguments are interpreted as though they were cast as follows:

```
assign(static_cast<size_type>(first), static_cast<value_type>  
(last));
```

void **assign** (size\_type n, const T& t)

Erases the current contents of the deque and inserts n copies of t.

allocator\_type **get\_allocator** ( ) const

Returns the allocator object.

reference **operator[]** (size\_type n), const\_reference **operator[]** (size\_type n) const

Returns the element at index n. If n >= size( ), the behavior is undefined.

reference **at** (size\_type n), const\_reference **at** (size\_type n) const

Returns the element at index n. If n >= size( ), it throws out\_of\_range.

reference **back** ( ), const\_reference **back** ( ) const

Returns the last element in the deque. The behavior is undefined if the deque is empty.

iterator **begin** ( ), const\_iterator **begin** ( ) const

Returns an iterator that points to the first element of the deque.

void **clear** ( )

Erases all elements from the deque.

bool **empty** ( ) const

Returns size( ) == 0.

iterator **end** ( ), const\_iterator **end** ( ) const

Returns an iterator that points to the last element of the deque.

iterator **erase** (iterator position)

Erases the element at position.

iterator **erase** (iterator first, iterator last)

Erases all the elements in the range [first, last).

reference **front** ( ), const\_reference **front** ( ) const

Returns the first element of the deque. The behavior is undefined if the deque is empty.

iterator **insert** (iterator position, const T& x)

Inserts x at position. If position is begin( ) or end( ), the performance is constant; at any other position, the performance is linear.

void **insert** (iterator pos, size\_type n, const T& x)

Inserts n copies of x at pos.

template <class InputIterator>, void **insert**  
(iterator position, InputIterator first,  
InputIterator last)

Inserts the elements in the range [first, last) starting at position, unless InputIterator is an integral type, in which case the arguments are interpreted as though they were cast:

```
insert(position, static_cast<size_type>(first),  
        static_cast<value_type>(last));
```

If an exception is thrown, such as `bad_alloc` when there is insufficient memory for a new element, the deque is unchanged, and all iterators and references remain valid. If the exception is thrown from an element's copy constructor or assignment operator, however, the behavior is unspecified.

size\_type **max\_size** ( ) const

Returns the size of the largest possible deque.

void **pop\_front** ( )

Erases the first element of the deque. The behavior is undefined if the deque is empty.

**void pop\_back ( )**

Erases the last element of the deque. The behavior is undefined if the deque is empty.

**void push\_front (const T& x)**

Inserts x as the new first element of the deque.

**void push\_back (const T& x)**

Inserts x as the new last element of the deque.

**reverse\_iterator rbegin ( ), const\_reverse\_iterator rbegin ( ) const**

Returns a reverse iterator that points to the last element of the deque.

**reverse\_iterator rend ( ), const\_reverse\_iterator rend ( ) const**

Returns a reverse iterator that points to one position before the first element of the deque.

**size\_type size ( ) const**

Returns the number of elements in the deque.

**void resize (size\_type n, T c = T( ))**

Changes the size of the deque to n. If  $n > \text{size}()$ , one or more copies of c are added to the end of the deque to reach the desired size. If the new size is smaller than the current size, the first n elements are unchanged, and elements are erased from the end to reach the new size.

**void swap (deque<T, Alloc>& that)**

Exchanges all the elements in the deque with all the elements in that.

## See Also

<list>, [<vector>](#)



## Name

`operator==` function template — Compares two deques for equality

## Synopsis

```
template<typename T, typename A>  
bool operator==(const deque<T,A>& x, const deque<T,A>& y)
```

The `==` operator returns true if `x` and `y` are the same size and their elements are equal, that is, `x.size( ) == y.size( ) && equals(x.begin( ), x.end( ), y.begin( ))`.

## See Also

`equals` in [<algorithm>](#)

## Name

`operator!=` function template — Compares two deques for inequality

## Synopsis

```
template<typename T, typename A>  
bool operator!=(const deque<T,A>& x, const deque<T,A>& y)
```

The `!=` operator is equivalent to `!(x == y)`.

## Name

`operator<` function template — Compares two deques for less-than

## Synopsis

```
template<typename T, typename A>  
bool operator<(const deque<T,A>& x, const deque<T,A>& y)
```

The `<` operator determines whether `x` is less than `y` using the same algorithm as

```
lexicographical_compare(x.begin( ), x.end( ),  
y.begin( ), y.end( )).
```

## See Also

`lexicographical_compare` in [<algorithm>](#)

## Name

`operator<=` function template — Compares two deques for less-than-or-equal

## Synopsis

```
template<typename T, typename A>  
bool operator<=(const deque<T,A>& x, const deque<T,A>& y)
```

The `<=` operator is equivalent to `!(y < x)`.

## Name

`operator>` function template — Compares two deques for greater-than

## Synopsis

```
template<typename T, typename A>  
bool operator>(const deque<T,A>& x, const deque<T,A>& y)
```

The `>` operator is equivalent to `(y < x)`.

## Name

`operator>=` function template — Compares two deques for greater-than-or-equal

## Synopsis

```
template<typename T, typename A>  
bool operator>=(const deque<T,A>& x, const deque<T,A>& y)
```

The `>=` operator is equivalent to `!(x < y)`.

## Name

swap function template specialization — Swaps the contents of two deques

## Synopsis

```
template<typename T, typename Alloc>  
void swap(deque<T, Alloc>& x, deque<T, Alloc>& y)
```

The swap function template specialization is equivalent to calling `x.swap(y)`.

## See Also

swap in [<algorithm>](#)

## <exception>

The <exception> header declares classes, types, and functions related to fundamental exception handling. See <stdexcept> for additional exception classes.

### **Not e**

This section describes a complicated system of exceptions, function calls, and handlers. The end result is simpler than it might seem at first:

*A function can throw only the exception types listed in its exception specification, or else the program terminates immediately.*

If a function does not have an exception specification, it can throw any exception. (Virtual functions work a little differently; see [Chapter 6](#) for details.)



# Name

`bad_exception` class — Wrong exception type

# Synopsis

```
class bad_exception : public exception {
public:
    bad_exception( ) throw( );
    bad_exception(const bad_exception&) throw( );
    bad_exception& operator=(const bad_exception&) throw( );
    virtual const char* what( ) const throw( );
};
```

A `bad_exception` object is thrown from the unexpected function when unexpected throws an exception that is not listed in the exception specification that caused unexpected to be called. Most programs do not throw or catch `bad_exception`. You can list `bad_exception` in an exception specification if you want to handle this unusual situation differently.

See [Chapter 5](#) and unexpected (in this section) for more details.

# See Also

terminated function, [unexpected function](#), throw keyword

## Name

exception class — Base class for all standard exceptions

## Synopsis

```
class exception {  
public:  
    exception( ) throw( );  
    exception(const exception&) throw( );  
    exception& operator=(const exception&) throw( );  
    virtual ~exception( ) throw( );  
    virtual const char* what( ) const throw( );  
};
```

The `exception` class is the base class for all exception objects thrown by the standard library or by code generated by the compiler. By convention, user-defined exception classes also derive from `exception` or from one of its derived classes.



```
virtual const char* what ( ) const throw( );
```

Returns a message that describes the nature of the exception. The exact contents of the string are implementation-defined; it might be a multibyte string, which can be converted to a wstring.

## See Also

[bad\\_exception\\_class](#), `bad_alloc` in [<new>](#), `bad_cast` in [<typeinfo>](#), `bad_typeid` in [<typeinfo>](#), `ios_base::failure` in [<ios>](#), `logic_error` in [<stdexcept>](#), `runtime_error` in [<stdexcept>](#)

## Name

set\_terminate function — Changes the terminate( ) handler

## Synopsis

```
typedef void (*terminate_handler)( );  
terminate_handler set_terminate(terminate_handler f) throw( );
```

The set\_terminate function saves f to be used by calls to terminate. The previous value of the terminate handler is returned.

## See Also

[terminate function](#)

## Name

set\_unexpected function — Changes the unexpected( ) handler

## Synopsis

```
typedef void (*unexpected_handler)( );  
unexpected_handler set_unexpected(unexpected_handler f)  
    throw( );
```

The set\_unexpected function saves f to be used by calls to unexpected. The previous value of the unexpected handler is returned.

## See Also

[unexpected function](#)

## Name

terminate function — Terminates program when exception handling fails

## Synopsis

```
void terminate( )
```

The terminate function is called when normal exception handling cannot handle an exception for any reason—for example, when there is no matching catch block for an exception, or when an exception is thrown and, while the stack is unwinding, another exception is thrown by a destructor.

A program might also call terminate explicitly.

You can change the behavior of the terminate function by calling `set_terminate`. The default behavior is to call `abort`.

The terminate function is a last resort because normal exception handling failed. For this reason, you cannot rely on the usual destruction of static objects and objects on the stack.

## See Also

[set\\_terminate function](#), [unexpected function](#), `abort` function in [<cstdlib>](#), `catch` keyword

## Name

`uncaught_exception` function — Determines whether exception handling is currently underway

## Synopsis

```
bool uncaught_exception( )
```

The `uncaught_exception` function returns `true` while an exception is processed: after evaluating the argument of a `throw` expression but before a matching exception declaration is initialized in a `catch` block. It also returns `true` after `terminate` is called (but not for an explicit call to `terminate`). It returns `false` at other times.

Call `uncaught_exception` to learn whether exception handling is currently underway. If it returns `true`, throwing a new exception results in a call to `terminate`.

## See Also

[terminate function](#), `catch` keyword, `throw` keyword

## Name

unexpected function — Handles an unexpected exception type

## Synopsis

```
void unexpected( )
```

If a function has an exception specification and throws an exception that is not listed in the exception specification, the unexpected function is called to handle the unexpected exception.

You can implement your own unexpected function. If you do so, you must ensure that unexpected does not return normally. It can terminate the program—for example, by calling terminate—or it can throw an exception.

If your unexpected function throws an exception that is not listed in the function's exception specification, a new exception of type `bad_exception` is created and thrown. If the function's exception specification does not list `bad_exception`, `terminate` is called automatically.

The default implementation of unexpected calls `terminate`.

In other words, if a function has an exception specification, it is guaranteed that only the specified exceptions can be thrown out of the function, or else the application will be terminated.

## See Also

[bad\\_exception class](#), [set\\_unexpected function](#), `abort` in [<cstdlib>](#), `throw` keyword

## **<fstream>**

The `<fstream>` header declares classes and other types for performing I/O with external files. A file in C++ is a sequence of bytes. A narrow (byte) stream or buffer simply reads or writes those bytes. A wide stream or buffer reads multibyte characters and converts them to wide characters (according to the stream's locale) or converts wide characters to their multibyte equivalents for writing.

See [Chapter 9](#) for a general discussion of I/O and related topics (stream buffers, locales, and facets), [Chapter 1](#) for more information about character sets, and the `<iostream>` section in this chapter for information about the base-class templates required by the `fstream` class templates. Refer to [Chapter 8](#) for information about traits in general and to the `<string>` section in this chapter for detailed information about the `char_traits` class template. Refer to the `<streambuf>` section in this chapter for information about the `basic_streambuf` class template.

To open a file for reading, use `ifstream`; for writing, use `ofstream`; for reading and writing, use `fstream`; for wide character I/O, use `wifstream`, `wofstream`, or `wfstream`.



# Name

`basic_filebuf` class template — Class template for file buffers

# Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_filebuf : public basic_streambuf<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    basic_filebuf( );
    virtual ~basic_filebuf( );

    bool is_open( ) const;
    basic_filebuf<charT,traits>*
        open(const char* filename, ios_base::openmode mode);
    basic_filebuf<charT,traits>*
        open(const char* filename, ios_base::open_mode mode);
    basic_filebuf<charT,traits>* close( );
protected:
    virtual streamsize showmanyc( );
    virtual int_type underflow( );
    virtual int_type uflow( );
    virtual int_type pbackfail(int_type c = traits::eof( ));
    virtual int_type overflow(int_type c = traits::eof( ));
    virtual basic_streambuf<charT,traits>*
        setbuf(char_type* s, streamsize n);
    virtual pos_type seekoff(off_type off, ios_base::seekdir way,
        ios_base::openmode = ios_base::in | ios_base::out);
    virtual pos_type seekpos(pos_type newpos,
        ios_base::openmode which = ios_base::in | ios_base::out);
    virtual int sync( );
    virtual void imbue(const locale& loc);
};
```

The `basic_filebuf` class template implements a stream buffer that is associated with an external file. The

connection to the external file is equivalent to calling C I/O functions (declared in `<cstdio>`) but may or may not actually call the C functions. For example, the `open` function is equivalent to calling `fopen` and having the `filebuf` object store the returned `FILE` pointer.

The external file is treated as a series of bytes, which might form multibyte characters. When converting the multibyte character sequences to characters in the file buffer, which are of type `charT`, the file buffer uses a code conversion facet from the buffer's locale. This `codecvt` facet is equivalent to `a_codecvt`, which is declared and initialized as follows:

```
std::codecvt<charT, char, typename traits::state_type>
    a_codecvt = use_facet<codecvt<charT, char,
        typename traits::state_type> >(getloc( ));
```

Some of the function descriptions in this section refer to `a_codecvt` and describe functionality, assuming that an actual `codecvt` facet object exists. An implementation does not have to create a `codecvt` facet object as long as the file stream acts as though it did create and use an explicit `codecvt` facet. (See `<locale>` for more information about `codecvt`.)

Remember that `codecvt<char, char, mbstate_t>` is essentially a no-op, mapping a character to itself, so the `codecvt` facet is most important when `charT` is `wchar_t`.

See `<streambuf>` for a description of the required behavior of a stream buffer, especially for the virtual functions that `basic_filebuf` overrides.

The following are the member functions of `basic_filebuf`:

**`basic_filebuf ( )`**

Initializes the file buffer in a closed state.

**`virtual ~basic_filebuf ( )`**

Calls `close( )` and finalizes the file buffer.

`basic_filebuf<charT, traits>*` **close** ( )

Closes the file, severing the connection between the external file and the file buffer. If the file is already closed, it returns a null pointer. Otherwise, it calls `overflow(EOF)` to flush the output buffer. If the buffer recently called `overflow`, the external character stream might have an incomplete shift sequence of a multibyte character. The `close` function therefore calls `a_codecvt.unshift( )` as often as needed to complete the shift sequence, and then calls `overflow(EOF)` again to flush the buffer. Finally, the external file is closed by calling `fclose` or its equivalent.

The return value is `this` for success or a null pointer for failure.

virtual void **imbue** (const locale& loc)

Changes the file buffer's locale, in particular the `codecvt` facet. It is safe to change the locale when the file is positioned at its beginning, when the character encoding (`a_codecvt.encoding( )`) is not state-dependent, or the old and new locales have the same `codecvt` facet.

bool **is\_open** ( ) const

Returns `true` if the file is open or `false` if the file is closed.

`basic_filebuf<charT, traits>*`, **open** (const char\* filename, ios\_base::openmode mode),

`basic_filebuf<charT, traits>*`, **open** (const char\* filename, ios\_base::open\_mode mode)

Opens the file `filename`. If the file is already open (`is_open( )` returns `true`), a null pointer is returned immediately; otherwise, the file buffer is initialized and the named file is opened by calling the equivalent of

`fopen(filename, modestr )`. The *modestr* is determined from the mode (without the `ios_base::ate` bit), as shown in [Table 13-9](#). No other mode combinations are allowed. If the mode includes `ios_base::ate`, the opened file is positioned at its end by calling the equivalent of `fseek(file, 0, SEEK_END)`.

The second form is deprecated. It has the same functionality as the first form. See `ios_base::openmode` in `<ios>` for details.

If the file is opened successfully, this is returned; otherwise, the return value is a null pointer.

*Table 13-9. File open modes*

<code>ios_base</code> mode bits	<code>fopen</code> equivalent <code>modestr</code>
out	"w"
out   app	"a"
out   trunc	"w"
in	"r"
in   out	"r+"
in   out   trunc	"w+"
binary   out	"wb"
binary   out   app	"ab"
binary   out   trunc	"wb"
binary   in	"rb"

<code>ios_base::mode</code> bits	file open equivalent mode string
binary   in   out	"r+b"
binary   in   out   trunc	"w+b"

```
virtual int_type overflow (int_type c = traits::eof(
))
```

Converts its output using `a_codecvt.out` and writes the converted characters to the external file. The return value is `traits::eof( )` for failure, which also occurs if the file is not open. For success, the return value is `traits::not_eof(c)`.

```
virtual int_type pbackfail (int_type c =
traits::eof( ))
```

Tries to push back the character `c` so it will be the next character read from the input buffer. If a push-back position is not available (see `<streambuf>` for a definition of "push-back position"), the file buffer attempts to make one available (e.g., by moving the file position).

If `c` is `traits::eof( )` or the same character as `gptr( )[-1]`, the file buffer decrements the `gptr( )` pointer; otherwise, if the input array is assignable, `c` is assigned to `gptr( )[-1]` and `gptr( )` is decremented.

The return value is `traits::eof( )` for failure or `traits::not_eof(c)` for success.

```
virtual pos_type seekoff (off_type off,
ios_base::seekdir way, ios_base::openmode =
ios_base::in | ios_base::out)
```

Tries to seek to a new position in the file as an offset from a defined position. If the file is not open, the attempt to seek fails. If the character set encoding uses shift states or otherwise does not have a fixed size per


character, `off` must be 0. Otherwise, if the destination is not the current position (`off != 0` or `way != basic_ios::cur`), the output buffer is flushed and unshift sequences are written as needed (`a_codecvt.unshift`). The new file position is set by calling the equivalent of `fseek(file, width * off, origin)`, in which `width` is `a_codecvt.encoding()` and `origin` is determined as shown in [Table 13-10](#). If `width < 0`, `off` must be 0, so the call is `fseek(file, 0, origin)`. The return value is the new file position or -1 for an error or if the new position is unknown. Note that `seekoff` does not use its final parameter.

*Table 13-10. Seek origins*

<code>ios_base::seekdir</code> options	<code>fseek</code> equivalents
<code>basic_ios::beg</code>	<code>SEEK_SET</code>
<code>basic_ios::cur</code>	<code>SEEK_CUR</code>
<code>basic_ios::end</code>	<code>SEEK_END</code>

```
virtual pos_type seekpos (pos_type newpos,
ios_base::openmode
which=ios_base::in|ios_base::out)
```


Attempts to set the file position to `newpos`, which must be the result of calling `seekoff` or `seekpos` on the same file. If `which` includes the `ios_base::in` bit, the input sequence is updated; if `which` includes the `ios_base::out` bit, the output sequence is updated, and any necessary unshift characters are written prior to setting the file position. If neither bit is set in `which`, an error results. The return value is -1 for an error or `newpos` for success.

 virtual basic\_streambuf<charT, traits>\* **setbuf**  
(char\_type\* s, streamsize n)

Sets the buffer. If you call `setbuf(0, 0)` before any I/O operations are performed on a file, the file is set to unbuffered. The behavior is implementation-defined for any other argument values. The return value is this.

virtual streamsize **showmanyc** ( )

Returns an estimate of the number of characters immediately available for input. In other words, it does the same thing as the base class `showmanyc` function.

 virtual int **sync** ( )

Flushes output to the external file. The behavior for input is implementation-defined.

virtual int\_type **uflow** ( )

Fills the input buffer in the same manner as `underflow`.

virtual int\_type **underflow** ( )

Fills the input buffer. Multibyte characters are read from the external file and converted to `charT` characters by calling the equivalent of `a_codecvt.in`.

## See Also

[filebuf class](#), [wfilebuf class](#), `basic_streambuf` in [<streambuf>](#)

# Name

`basic_fstream` class template — Class template for file input and output streams

# Synopsis

```
template <class charT, class traits=char_traits<charT> >
class basic_fstream : public basic_iostream<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    basic_fstream( );
    explicit basic_fstream(const char* filename,
                           ios_base::openmode mode =
ios_base::in|ios_base::out);

    basic_filebuf<charT,traits>* rdbuf( ) const;
    bool is_open( );
    void open(const char* filename,
              ios_base::openmode mode =
ios_base::in|ios_base::out);
    void close( );
};
```

The `basic_fstream` class template supports reading and writing to and from named files using a `basic_filebuf<charT, traits>` object. (See `<istream>` for a description of the base-class template, `basic_iostream`.) In the following member function descriptions, the file buffer object is assumed to be a private data member with the name *buf*.

## **basic\_fstream** ( )

Constructor initializes the base class with `basic_iostream(& buf )` and initializes *buf* with its



default constructor.

```
explicit basic_fstream (const char* filename,  
ios_base::openmode mode = ios_base::in |  
ios_base::out)
```

Initializes the base class and *buf*, then calls `open(filename, mode)`. If `open` returns a null pointer, the constructor calls `setstate(failbit)`.

```
basic_filebuf<charT, traits>* rdbuf ( ) const
```

Returns & *buf*.

```
bool is_open ( )
```

Returns `rdbuf( )->is_open( )`.

```
void open (const char* filename, ios_base::openmode  
mode = ios_base::in | ios_base::out)
```

Calls `rdbuf( )->open(filename, mode)`. If that function returns a null pointer, `open` calls `setstate(failbit)`.

```
void close ( )
```

Calls `rdbuf( )->close( )`. If that function fails, `close` calls `setstate(failbit)`.

## See Also

[basic\\_filebuf class template](#), `basic_ios` in [<ios>](#),  
`basic_iostream` in [<istream>](#)

# Name

`basic_ifstream` class template — Class template for file input streams

# Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_ifstream : public basic_istream<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;
    basic_ifstream( );
    explicit basic_ifstream(const char* s, ios_base::openmode mode
= ios_base::in);

    basic_filebuf<charT,traits>* rdbuf( ) const;
    bool is_open( );
    void open(const char* s, ios_base::openmode mode =
ios_base::in);
    void open(const char* s, ios_base::open_mode mode);
    void close( );
};
```

The `basic_ifstream` class template supports reading from named files using a `basic_filebuf<charT, traits>` object. (See `<istream>` for a description of the base-class template, `basic_istream`.) In the following member function descriptions, the file buffer object is assumed to be a private data member with the name *buf*:

## **basic\_ifstream** ( )

Initializes the base class with `basic_istream(&buf)` and initializes *buf* with its default constructor.

```
explicit basic_ifstream (const char* filename,
ios_base::openmode mode = ios_base::in)
```

Initializes the base class and *buf*, then calls `open(filename, mode)`. If `open` returns a null pointer, the constructor calls `setstate(failbit)`.

`basic_filebuf<charT, traits>* rdbuf ( ) const`

Returns `&buf`.

`bool is_open ( )`

Returns `rdbuf( )->is_open( )`.

`void open (const char* filename, ios_base::openmode mode = ios_base::in), void open (const char* filename, ios_base::open_mode mode)`

Calls `rdbuf( )->open(filename, mode)`. If that function returns a null pointer, `open` calls `setstate(failbit)`. The second form is deprecated. It has the same functionality as the first form. See `ios_base::openmode` in `<ios>` for details.

`void close ( )`

Calls `rdbuf( )->close( )`. If that function fails, `close` calls `setstate(failbit)`.

## See Also

[basic\\_filebuf class template](#), `basic_ios` in [<ios>](#),  
`basic_istream` in [<istream>](#)

# Name

`basic_ofstream` class template — Class template for file output streams

# Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_ofstream : public basic_ostream<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    basic_ofstream( );
    explicit basic_ofstream(const char* s, ios_base::openmode mode
=
                                ios_base::out);
    basic_filebuf<charT,traits>* rdbuf( ) const;
    bool is_open( );
    void open(const char* s, ios_base::openmode mode =
ios_base::out);
    void open(const char* s, ios_base::open_mode mode);
    void close( );
};
```

The `basic_ofstream` class template supports writing to named files using a `basic_filebuf<charT, traits>` object. (See `<ostream>` for a description of the base-class template, `basic_ostream`.) In the following member function descriptions, the file buffer object is assumed to be a private data member with the name *buf*:

## **basic\_ofstream** ( )

Initializes the base class with `basic_ostream(&buf)` and initializes *buf* with its default constructor.

explicit **basic\_ofstream** (const char\* filename,  
ios\_base::openmode mode = ios\_base::out)

Initializes the base class and *buf*, then calls  
open(filename, mode). If open returns a null pointer, the  
constructor calls setstate(failbit).

basic\_filebuf<charT, traits>\* **rdbuf** ( ) const

Returns & *buf*.

bool **is\_open** ( )

Returns rdbuf( )->is\_open( ).

void **open** (const char\* filename, ios\_base::openmode  
mode = ios\_base::out), void **open** (const char\*  
filename, ios\_base::open\_mode mode)

Calls rdbuf( )->open(filename, mode). If that function  
returns a null pointer, open calls setstate(failbit).

The second form is deprecated. It has the same  
functionality as the first form. See ios\_base::openmode  
in <ios> for details.

void **close** ( )

Calls rdbuf( )->close( ). If that function fails, close  
calls setstate(failbit).

## See Also

[basic\\_filebuf class template](#), basic\_ios in [<ios>](#),  
basic\_ostream in [<ostream>](#)

## Name

filebuf class — File buffer

## Synopsis

```
typedef basic_filebuf<char> filebuf;
```

The filebuf class is a specialization of the basic\_filebuf template for char characters.

## See Also

[basic\\_filebuf class template](#), [wfilebuf class](#)

## Name

`fstream` class — Input and output file stream

## Synopsis

```
typedef basic_fstream<char> fstream;
```

The `fstream` class is a specialization of the `basic_fstream` template for `char` characters.

## See Also

[basic\\_fstream class template](#), [wfstream class](#), `iostream` in [<istream>](#)

# Name

ifstream class — Input file stream

# Synopsis

```
typedef basic_ifstream<char> ifstream;
```

The `ifstream` class is a specialization of the `basic_ifstream` template for `char` characters. [Example 13-9](#) shows a simple use of the `ifstream` and `ofstream` classes.

# Example

*Example 13-9. Copying a file using `ifstream` and `ofstream`*

```
#include <cstdlib>
#include <fstream>
#include <cstdio>    // For perror( )
#include <iostream> // For cerr

int main(int argc, char** argv)
{
    if (argc != 3) {
        std::cerr << "usage: copy FROM TO\n";
        return EXIT_FAILURE;
    }

    // Open the input file.
    std::ifstream in(argv[1]);
    if (! in) {
        std::perror(argv[1]);
        return EXIT_FAILURE;
    }

    // Open the output file.
    std::ofstream out(argv[2]);
    if (! out) {
        std::perror(argv[2]);
        return EXIT_FAILURE;
    }

    // Copy the input to the output, one character at a time.
```



```
char c;
while (in.get(c))
    out.put(c);
out.close( );
// Make sure the output was written.
if (! out) {
    std::perror(argv[2]);
    return EXIT_FAILURE;
}
}
```

## See Also

[basic\\_ifstream class template](#), [wofstream class](#), `istream` in [<istream>](#)

## Name

ofstream class — Output file stream

## Synopsis

```
typedef basic_ofstream<char> ofstream;
```

The `ofstream` class is a specialization of the `basic_ofstream` template for `char` characters.

## See Also

[basic\\_ofstream class template](#), [wofstream class](#), `ostream` in [<ostream>](#)

## Name

wfilebuf class — Wide character file buffer

## Synopsis

```
typedef basic_filebuf<wchar_t> wfilebuf;
```

The wfilebuf class is a specialization of the basic\_filebuf template for wchar\_t characters.

## See Also

[basic\\_filebuf class template](#), [filebuf class](#)

## Name

wfstream class — Wide character input and output stream

## Synopsis

```
typedef basic_fstream<wchar_t> wfstream;
```

The wfstream class is a specialization of the basic\_fstream template for wchar\_t characters.

## See Also

[basic\\_fstream class template](#), [fstream class](#), wiostream in [<istream>](#)

## Name

wfstream class — Wide character input stream

## Synopsis

```
typedef basic_ifstream<wchar_t> wfstream;
```

The wfstream class is a specialization of the basic\_ifstream template for wchar\_t characters.

## See Also

[basic\\_ifstream class template](#), [ifstream class](#), wistream in [<istream>](#)

## Name

wofstream class — Wide character output stream

## Synopsis

```
typedef basic_ofstream<wchar_t> wofstream;
```

The `wofstream` class is a specialization of the `basic_ofstream` template for `wchar_t` characters.

## See Also

[basic\\_ofstream class template](#), [ofstream class](#), `wostream` in [<ostream>](#)

## <functional>

The <functional> header defines several *functionals*, or *function objects*. A function object is an object that has an operator( ), so it can be called using the same syntax as a function. Function objects are most often used with the standard algorithms.

For example, to copy a sequence of integers, adding a fixed amount (42) to each value, you could use the following expression:

```
std::transform(src.begin( ), src.end( ), dst.begin( ),
               std::bind2nd(std::plus<int>( ), 42))
```

The result of combining `bind2nd` and `plus<int>` is a function object that adds the value 42 when it is applied to any integer. The `transform` algorithm copies all the elements from `src` to `dst`, applying the functional argument to each element. See the detailed description of `bind2nd` and `plus` in this section for details.

The standard function objects are defined for C++ operators; for binding function arguments; and for adapting functions, member functions, etc., as function objects.

Boost defines functionals that extend and improve on those in the standard library. See [Appendix B](#) for information about Boost.

# Name

binary\_function class template — Base class for binary functionals

# Synopsis

```
template <typename Arg1, typename Arg2, typename Result>
struct binary_function {
    typedef Arg1 first_argument_type;
    typedef Arg2 second_argument_type;
    typedef Result result_type;
};
```

The `binary_function` template is a base-class template for all the function classes that represent binary operations. It provides standard names for the argument and result types.

The base template has separate template parameters for each of the argument types and the return type. Many of the predefined function objects in this section use the same type for all three parameters, but you can use different types when defining your own function object, as shown in [Example 13-10](#).

# Example

*Example 13-10. Functional to round off a floating-point number*

```
// Functional for a binary function that rounds off a floating-
// point number (of
// type FltT) to a certain number of decimal places (supplied as
// an unsigned)
template<typename FltT>
struct roundoff : std::binary_function<FltT,unsigned,FltT> {
    FltT operator( )(FltT x, unsigned digits) const {
        FltT y = std::pow(10.0, static_cast<FltT>(digits));
        FltT z = x * y;
        return (z < 0 ? std::ceil(z - 0.5) : std::floor(z + 0.5)) /
        y;
```



```
    }  
};  
...  
// Copy seq to seq2, rounding off to two decimal places.  
std::transform(seq.begin( ), seq.end( ), seq2.begin( ),  
    std::bind2nd(roundoff<double>( ), 2));
```

## See Also

[binary\\_negate class template](#), [const\\_mem\\_fun1\\_ref\\_t class template](#), [const\\_mem\\_fun1\\_t class template](#), [mem\\_fun1\\_ref\\_t class template](#), [mem\\_fun1\\_t class template](#), [pointer\\_to\\_binary\\_function class template](#), [unary\\_function class template](#)

## Name

`binary_negate` class template — Logical negation of a binary predicate

## Synopsis

```
template <typename P>
class binary_negate : public
    binary_function<typename P::first_argument_type,
                  typename P::second_argument_type, bool>
{
public:
    explicit binary_negate(const P& predicate);
    bool operator( )(const typename P::first_argument_type& x,
                   const typename P::second_argument_type& y)
    const;
};
```

The `binary_negate` class template is a binary functional that returns the logical negation of another binary functional—that is, `operator( )` returns `!predicate(x, y)`. The simplest way to use `binary_negate` is to use the `not2` function template.

## See Also

[unary\\_negate class template](#), [not2 function template](#)

## Name

`bind1st` function template — Creates a `binder1st` function object

## Synopsis

```
template <typename Operation, typename T>
binder1st<Operation> bind1st(const Operation& op, const T& x);
```

The `bind1st` function is a convenient way to construct a `binder1st` object. Use `bind1st` when you have a binary function and always want to supply the same value as the first argument to the function.

## Example

Suppose you have a container of data points, and you want to count the number of points that exceed a threshold—in other words, where the threshold is less than or equal to the data point. Here is one way to do this:

```
std::cout
  << std::count_if(data.begin( ), data.end( ),
                 std::bind1st(std::less_equal<double>( ), threshold))
  << '\n';
```

## See Also

[bind2nd function template](#), [binder1st class template](#)

## Name

`bind2nd` function template — Creates a `binder2nd` function object

## Synopsis

```
template <typename Operation, typename T>  
binder2nd<Operation> bind2nd(const Operation& op, const T& x);
```

The `bind2nd` function is a convenient way to construct a `binder2nd` object. Use `bind2nd` when you have a binary function and always want to supply the same value as the first argument to the function.

## Example

Suppose you have a container of data points, and you want to count the number of points that exceed a threshold. Here is one way to do this:

```
std::cout << std::count_if(data.begin( ), data.end( ),  
                          std::bind2nd(std::greater<double>( ), threshold))  
<< '\n';
```

## See Also

[bind1st function template](#), [binder2nd class template](#)

## Name

binder1st class template — Binds a value to the first argument of a binary function

## Synopsis

```
template <typename Operation>
class binder1st : public unary_function<
    typename Operation::second_argument_type,
    typename Operation::result_type>
{
protected:
    Operation op;
    typename Operation::first_argument_type value;
public:
    binder1st(const Operation& x,
              const typename Operation::first_argument_type& y);
    typename Operation::result_type operator( )
        (const typename Operation::second_argument_type& x) const;
};
```

The `binder1st` class template is a unary functional that binds a fixed value as the first argument to a binary function object. The constructor initializes the `op` and `value` data members with the `x` and `y` arguments. The `operator( )` member function returns `op(value, x)`.

See the `bind1st` function template for an easier way to construct and use the `binder1st` class template.

## See Also

[bind1st function template](#), [binder2nd class template](#)

## Name

binder2nd class template — Binds a value to the second argument of a binary function

## Synopsis

```
template <typename Operation>
class binder2nd : public unary_function<
    typename Operation::first_argument_type,
    typename Operation::result_type>
{
protected:
    Operation op;
    typename Operation::second_argument_type value;
public:
    binder2nd(const Operation& x,
              const typename Operation::second_argument_type& y);
    typename Operation::result_type operator( )
        (const typename Operation::first_argument_type& x) const;
};
```

The binder2nd class template is a unary functional that binds a fixed value as the second argument to a binary function object. The constructor initializes the op and value data members with the x and y arguments. The operator( ) member function returns op(x, value).

See the bind2nd function template for an easier way to construct and use the binder2nd class template.

## See Also

[bind2nd function template](#), [binder1st class template](#)

## Name

`const_mem_fun_ref_t` class template — Calls a member function of a constant reference object

## Synopsis

```
template <typename Rtn, typename T>
class const_mem_fun_ref_t : public unary_function<T, Rtn>
{
public:
    explicit const_mem_fun_ref_t(Rtn (T::*p)( ) const);
    Rtn operator( )(const T& p) const;
};
```

The `const_mem_fun_ref_t` class template is a unary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, and the `T` template parameter is the class that declares the member function. The argument to the constructor is a pointer to the member function, which takes no arguments. The member function is called from `operator( )` using a reference to the `const` object.

See the `mem_fun_ref` function template for an easier way to construct and use the `const_mem_fun_ref_t` class template.

## See Also

[const\\_mem\\_fun\\_t class template](#), [const\\_mem\\_fun1\\_ref\\_t class template](#), [mem\\_fun\\_ref function template](#), [mem\\_fun\\_ref\\_t class template](#)

## Name

`const_mem_fun_t` class template — Calls a member function of a constant object

## Synopsis

```
template <class Rtn, class T>
class const_mem_fun_t : public unary_function<T*, Rtn>
{
public:
    explicit const_mem_fun_t(Rtn (T::*p)( ) const);
    Rtn operator( )(const T* p) const;
};
```

The `const_mem_fun_t` class template is a unary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, and the `T` template parameter is the class that declares the member function. The argument to the constructor is a pointer to the member function, which takes no arguments. The member function is called from `operator( )` using a pointer to the `const` object.

See the `mem_fun` function template for an easier way to construct and use the `const_mem_fun_t` class template.

## See Also

[const\\_mem\\_fun\\_ref\\_t class template](#), [const\\_mem\\_fun1\\_t class template](#), [mem\\_fun function template](#), [mem\\_fun\\_t class template](#)



## Name

`const_mem_fun1_ref_t` class template — Calls a member function of a constant reference object with an argument

## Synopsis

```
template <typename Rtn, typename T, typename Arg>
class const_mem_fun1_ref_t :
    public binary_function<T, Arg, Rtn>
{
public:
    explicit const_mem_fun1_ref_t(Rtn (T::*p)(Arg) const);
    Rtn operator( )(const T& p, Arg x) const;
};
```

The `const_mem_fun1_ref_t` class template is a binary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, the `T` template parameter is the class that declares the member function, and the `Arg` template parameter is the type of the member function's sole argument.

The argument to the constructor is a pointer to the member function. The member function is called from `operator( )` using a reference to the `const` object.

See the `mem_fun_ref` function template for an easier way to construct and use the `const_mem_fun1_ref_t` class template.

## See Also

[const\\_mem\\_fun\\_ref\\_t class template](#), [const\\_mem\\_fun1\\_t class template](#), [mem\\_fun\\_ref function template](#), [mem\\_fun1\\_ref\\_t class template](#)

## Name

`const_mem_fun1_t` class template — Calls a member function of a constant object with an argument

## Synopsis

```
template <typename Rtn, typename T, typename Arg>
class const_mem_fun1_t: public binary_function<T*, Arg, Rtn>
{
public:
    explicit const_mem_fun1_t(Rtn (T::*p)(Arg) const);
    Rtn operator( )(const T* p, Arg x) const;
};
```

The `const_mem_fun1_t` class template is a binary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, the `T` template parameter is the class that declares the member function, and the `Arg` template parameter is the type of the member function's sole argument.

The argument to the constructor is a pointer to the member function. The member function is called from `operator( )` using a pointer to the `const` object.

See the `mem_fun` function template for an easier way to construct and use the `const_mem_fun1_t` class template.

## See Also

[const\\_mem\\_fun\\_t class template](#), [const\\_mem\\_fun1\\_ref\\_t class template](#), [mem\\_fun function template](#), [mem\\_fun1\\_t class template](#)

## Name

divides class template — Binary functional to divide

## Synopsis

```
template <typename T>
struct divides : binary_function<T, T, T> {
    T operator( )(const T& x, const T& y) const;
};
```

The divides class template is a binary functional in which operator( ) returns  $x / y$ .

## See Also

[binary\\_function class template](#), [minus class template](#), [modulus class template](#), [multiplies class template](#), [negate class template](#), [plus class template](#)

## Name

`equal_to` class template — Binary functional to compare for equality

## Synopsis

```
template <typename T>
struct equal_to : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The `equal_to` class template is a binary functional in which `operator( )` returns `x == y`.

## See Also

[binary\\_function class template](#), [greater class template](#),  
[greater\\_equal class template](#), [less class template](#),  
[less\\_equal class template](#), [not\\_equal\\_to class template](#)

## Name

greater class template — Binary functional to compare for greater-than

## Synopsis

```
template <typename T>
struct greater : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The greater class template is a binary functional in which `operator( )` returns  $x > y$ .

## See Also

[binary\\_function class template](#), [equal\\_to class template](#), [greater\\_equal class template](#), [less class template](#), [less\\_equal class template](#), [not\\_equal\\_to class template](#)

## Name

greater\_equal class template — Binary functional to compare for greater-than-or-equal

## Synopsis

```
template <typename T>
struct greater_equal : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The greater\_equal class template is a binary functional in which operator( ) returns  $x \geq y$ .

## See Also

[binary\\_function class template](#), [equal\\_to class template](#), [greater class template](#), [less class template](#), [less\\_equal class template](#), [not\\_equal\\_to class template](#)

## Name

less class template — Binary functional to compare for less-than

## Synopsis

```
template <typename T>
struct less : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The less class template is a binary functional in which operator( ) returns  $x < y$ .

## See Also

[binary\\_function class template](#), [equal\\_to class template](#), [greater class template](#), [greater\\_equal class template](#), [less\\_equal class template](#), [not\\_equal\\_to class template](#)

## Name

less\_equal class template — Binary functional to compare for less-than-or-equal

## Synopsis

```
template <typename T>
struct less_equal : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The less\_equal class template is a binary functional in which operator( ) returns  $x \leq y$ .

## See Also

[binary\\_function class template](#), [equal\\_to class template](#), [greater class template](#), [greater\\_equal class template](#), [less class template](#), [not\\_equal\\_to class template](#)



## Name

`logical_and` class template — Binary functional for logical conjunction

## Synopsis

```
template <typename T>
struct logical_and : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The `logical_and` class template is a binary functional in which `operator( )` returns `x && y`. Note that no short-circuiting occurs because both arguments must be evaluated before `operator( )` can be called.

## See Also

[logical\\_not class template](#), [logical\\_or class template](#)

## Name

`logical_not` class template — Binary functional for logical negation

## Synopsis

```
template <typename T>
struct logical_not : unary_function<T, bool> {
    bool operator( )(const T& x) const;
};
```

The `logical_not` class template is a unary functional in which `operator( )` returns `!x`.

## See Also

[logical\\_and class template](#), [logical\\_or class template](#), [not1 function template](#), [not2 function template](#)

## Name

`logical_or` class template — Binary functional for logical disjunction

## Synopsis

```
template <typename T>
struct logical_or : binary_function<T, T, bool> {
    bool operator( )(const T& x, const T& y) const;
};
```

The `logical_or` class template is a binary functional in which `operator( )` returns `x || y`. Note that no short-circuiting occurs because both arguments must be evaluated before `operator( )` can be called.

## See Also

[logical\\_and class template](#), [logical\\_not class template](#)

## Name

`mem_fun` function template — Creates a function object to call a member function via a pointer

## Synopsis

```
template<typename Rtn, typename T>
    const_mem_fun_t<Rtn,T> mem_fun(Rtn (T::*f)( ) const);
template<typename Rtn, typename T, typename Arg>
    const_mem_fun1_t<Rtn,T,Arg> mem_fun(Rtn (T::*f)(Arg) const);
template<typename Rtn, typename T>
    mem_fun_t<Rtn,T> mem_fun(Rtn (T::*f)( ));
template<typename Rtn, typename T, typename Arg>
    mem_fun1_t<Rtn,T,Arg> mem_fun(Rtn (T::*f)(Arg));
```

The `mem_fun` function template takes a pointer to a member function as an argument and returns a function object that can call the member function. The function object must be applied to a pointer to `T` (or a derived class). The `Rtn` template parameter is the return type of the member function, and the `T` template parameter is the object that has the member function. The optional `Arg` template parameter is the type of the argument to the member function.

The `mem_fun` function is usually the simplest way to create a function object that wraps a member function. In normal use, the compiler deduces the template parameters.

Suppose you have an `Employee` class and a container of `Employee` pointers. One of the member functions of `Employee` is `gets_bonus`, which returns a `bool`: `true` if the employee is lucky and gets a bonus this year, and `false` if the employee is unlucky. [Example 13-11](#) shows how to remove all the unlucky employees from the container.

## Example

*Example 13-11. Wrapping a member function called via a pointer as a function object*

```
class Employee {
public:
    int      sales( )      const { return sales_; }
    std::string name( )    const { return name_; }
    bool     gets_bonus( ) const { return sales( ) > bonus; }
    ...
};

std::list<Employee*> empptrs;
// Fill empptrs with pointers to Employee objects.
...
// Remove the employees who will NOT receive bonuses.
std::list<Employee*>::iterator last =
    std::remove_if(empptrs.begin( ), empptrs.end( ),
                  std::not1(std::mem_fun(&Employee::gets_bonus)));
```

## See Also

[const mem\\_fun t class template](#), [const mem\\_fun1 t class template](#), [mem\\_fun\\_ref function template](#), [mem\\_fun t class template](#), [mem\\_fun1 t class template](#), [ptr\\_fun function template](#)

# Name

`mem_fun_ref` function template — Creates a function object to call a member function via a reference

# Synopsis

```
template<typename Rtn, typename T>
    const_mem_fun_ref_t<Rtn,T>
        mem_fun_ref(Rtn (T::*f)( ) const);
template<typename Rtn, typename T, typename Arg>
    const_mem_fun1_ref_t<Rtn,T,Arg>
        mem_fun_ref(Rtn (T::*f)(Arg) const);
template<typename Rtn, typename T>
    mem_fun_ref_t<Rtn,T> mem_fun_ref(Rtn (T::*f)( ));
template<typename Rtn, typename T, typename Arg>
    mem_fun1_ref_t<Rtn,T,A> mem_fun_ref(Rtn (T::*f)(Arg));
```

The `mem_fun_ref` function template takes a pointer to a member function as an argument and returns a function object that can call the member function. The function object must be applied to an object of type `T` (or a derived class). The object is passed by reference to the functional. The `Rtn` template parameter is the return type of the member function; the `T` template parameter is the object that has the member function. The optional `Arg` template parameter is the type of the argument to the member function.

The `mem_fun_ref` function is usually the simplest way to create a function object that wraps a member function. In normal use, the compiler deduces the template parameters.

Suppose you have an `Employee` class and a container of `Employee` objects. As in [Example 13-11](#), one of the member functions of `Employee` is `gets_bonus`, which returns a `bool`: `true` if the employee is lucky and gets a bonus this year, or `false` if the employee is unlucky. [Example 13-12](#) shows

how to remove all the unlucky employees from the container.

## Example

*Example 13-12. Wrapping a member function called via a reference as a function object*

```
class Employee {
public:
    int      sales( )      const { return sales_; }
    std::string name( )    const { return name_; }
    bool     gets_bonus( ) const { return sales( ) > bonus; }
    ...
};

std::list<Employee> emps;
// Fill emps with Employee objects.
...
// Erase the employees who will NOT receive bonuses. The call to
remove_if
// rearranges emps; the call to erase removes the unlucky
employees from the
// list.
emps.erase(
    std::remove_if(emps.begin( ), emps.end( ),
        std::not1(std::mem_fun_ref(&Employee::gets_bonus))),
    emps.end( ));
```

## See Also

[const\\_mem\\_fun\\_ref\\_t class template](#), [const\\_mem\\_fun1\\_ref\\_t class template](#), [mem\\_fun function template](#), [mem\\_fun\\_ref\\_t class template](#), [mem\\_fun1\\_ref\\_t class template](#), [ptr\\_fun function template](#)

## Name

`mem_fun_ref_t` class template — Calls a member function of a reference object

## Synopsis

```
template <typename Rtn, typename T>
class mem_fun_ref_t : public unary_function<T, Rtn>
{
public:
    explicit mem_fun_ref_t(Rtn (T::*p)( ));
    Rtn operator( )(T& p) const;
};
```

The `mem_fun_ref_t` class template is a unary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, and the `T` template parameter is the class that declares the member function. The argument to the constructor is a pointer to the member function, which takes no arguments. The member function is called from `operator( )` using a reference to the object.

See the `mem_fun_ref` function template for an easier way to construct and use the `mem_fun_ref_t` class template.

## See Also

[const mem\\_fun\\_ref\\_t class template](#), [mem\\_fun\\_ref function template](#), [mem\\_fun\\_t class template](#), [mem\\_fun1\\_ref\\_t class template](#)



## Name

`mem_fun_t` class template — Calls a member function of a constant object

## Synopsis

```
template <class Rtn, class T>
class mem_fun_t : public unary_function<T*, Rtn>
{
public:
    explicit mem_fun_t(Rtn (T::*p)( ) const);
    Rtn operator( )(const T* p) const;
};
```

The `mem_fun_t` class template is a unary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, and the `T` template parameter is the class that declares the member function. The argument to the constructor is a pointer to the member function, which takes no arguments. The member function is called from `operator( )` using a pointer to the object.

See the `mem_fun` function template for an easier way to construct and use the `mem_fun_t` class template.

## See Also

[const\\_mem\\_fun\\_t class template](#), [mem\\_fun function template](#), [mem\\_fun\\_ref\\_t class template](#), [mem\\_fun1\\_t class template](#)

## Name

`mem_fun1_ref_t` class template — Calls a member function of a constant reference object with an argument

## Synopsis

```
template <typename Rtn, typename T, typename Arg>
class mem_fun1_ref_t :
    public binary_function<T, Arg, Rtn>
{
public:
    explicit mem_fun1_ref_t(Rtn (T::*p)(Arg) const);
    Rtn operator( )(const T& p, Arg x) const;
};
```

The `mem_fun1_ref_t` class template is a binary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, the `T` template parameter is the class that declares the member function, and the `Arg` template parameter is the type of the member function's sole argument.

The argument to the constructor is a pointer to the member function. The member function is called from `operator( )` using a `const` reference to the object.

See the `mem_fun_ref` function template for an easier way to construct and use the `mem_fun1_ref_t` class template.

## See Also

[const\\_mem\\_fun1\\_ref\\_t class template](#), [mem\\_fun\\_ref function template](#), [mem\\_fun\\_ref\\_t class template](#), [mem\\_fun1\\_t class template](#)

## Name

`mem_fun1_t` class template — Calls a member function of an object with an argument

## Synopsis

```
template <typename Rtn, typename T, typename Arg>
class mem_fun1_t: public binary_function<T*, Arg, Rtn>
{
public:
    explicit mem_fun1_t(Rtn (T::*p)(Arg));
    Rtn operator( )(T* p, Arg x) const;
};
```

The `mem_fun1_t` class template is a binary functional that wraps a member function pointer. The `Rtn` template parameter is the member function's return type, the `T` template parameter is the class that declares the member function, and the `Arg` template parameter is the type of the member function's sole argument.

The argument to the constructor is a pointer to the member function. The member function is called from `operator( )` using a pointer to the object.

See the `mem_fun` function template for an easier way to construct and use the `mem_fun1_t` class template.

## See Also

[const\\_mem\\_fun1\\_t class template](#), [mem\\_fun function template](#), [mem\\_fun\\_t class template](#), [mem\\_fun1\\_ref\\_t class template](#)

## Name

minus class template — Binary functional for subtraction

## Synopsis

```
template <typename T>
struct minus : binary_function<T, T, T> {
    T operator( )(const T& x, const T& y) const;
};
```

The minus class template is a binary functional in which operator( ) returns  $x - y$ .

## See Also

[binary\\_function class template](#), [divides class template](#), [modulus class template](#), [multiplies class template](#), [negate class template](#), [plus class template](#)

## Name

modulus class template — Binary functional for modulus (remainder)

## Synopsis

```
template <typename T>
struct modulus : binary_function<T, T, T> {
    T operator( )(const T& x, const T& y) const;
};
```

The modulus class template is a binary functional in which `operator( )` returns `x % y`.

## See Also

[binary\\_function class template](#), [divides class template](#), [minus class template](#), [multiplies class template](#), [negate class template](#), [plus class template](#)

## Name

multiplies class template — Binary functional for multiplication

## Synopsis

```
template <typename T>
struct multiplies : binary_function<T, T, T> {
    T operator( )(const T& x, const T& y) const;
};
```

The multiplies class template is a binary functional in which operator( ) returns  $x * y$ .

## See Also

[binary\\_function class template](#), [divides class template](#), [minus class template](#), [modulus class template](#), [negate class template](#), [plus class template](#)

## Name

negate class template — Unary functional for arithmetic negation

## Synopsis

```
template <typename T>
struct negate : unary_function<T,T> {
    T operator( )(const T& x) const;
};
```

The negate class template is a unary functional that performs arithmetic negation, that is, operator( ) returns  $-x$ .

## See Also

[divides class template](#), [minus class template](#), [modulus class template](#), [multiplies class template](#), [plus class template](#), [unary\\_function class template](#)

## Name

not1 function template — Returns a unary\_negate object

## Synopsis

```
template <typename Predicate>  
unary_negate<Predicate> not1(const Predicate& pred);
```

The not1 function template is a convenient way to construct a unary\_negate function object that performs the logical negation of pred. See [Example 13-11](#) earlier in this section.

## See Also

[logical\\_not class template](#), [not2 function template](#),  
[unary\\_negate class template](#)



## Name

not2 function template — Returns a binary\_negate object

## Synopsis

```
template <typename Predicate>  
binary_negate<Predicate> not2(const Predicate& pred);
```

The not2 function template is a convenient way to construct a binary\_negate function object that performs the logical negation of pred.

## See Also

[binary\\_negate class template](#), [logical\\_not class template](#), [not1 function template](#)

## Name

`not_equal_to` class template — Binary functional for inequality

## Synopsis

```
template <typename T>
struct not_equal_to : binary_function<T, T, bool>
{
    bool operator( )(const T& x, const T& y) const;
};
```

The `not_equal_to` class template is a binary functional in which `operator( )` returns `x != y`.

## See Also

[binary\\_function class template](#), [equal\\_to class template](#), [greater class template](#), [greater\\_equal class template](#), [less class template](#), [less\\_equal class template](#)

## Name

plus class template — Binary functional for addition

## Synopsis

```
template <typename T>
struct plus : binary_function<T, T, T>
{
    T operator( )(const T& x, const T& y) const;
};
```

The plus class template is a binary functional in which `operator( )` returns `x + y`.

## See Also

[binary\\_function class template](#), [divides class template](#), [minus class template](#), [modulus class template](#), [multiplies class template](#), [negate class template](#)

## Name

`pointer_to_binary_function` class template — Functional for a pointer to a binary function

## Synopsis

```
template <class Arg1, class Arg2, class Rtn>
class pointer_to_binary_function :
    public binary_function<Arg1,Arg2,Rtn>
{
public:
    explicit pointer_to_binary_function(Rtn (*f)(Arg1, Arg2));
    Rtn operator( )(Arg1 x, Arg2 y) const;
};
```

The `pointer_to_binary_function` class template is a function object that wraps a pointer to a function, in which the function is an ordinary (nonmember) function that takes two arguments. The `ptr_fun` function template is the most convenient way to create a `pointer_to_binary_function` object.

## See Also

[pointer\\_to\\_unary\\_function class template](#), [ptr\\_fun function template](#)

## Name

`pointer_to_unary_function` class template — Functional for a pointer to a unary function

## Synopsis

```
template <typename Arg, typename Rtn>
class pointer_to_unary_function :
    public unary_function<Arg, Rtn>
{
public:
    explicit pointer_to_unary_function(Result (*f)(Arg));
    Rtn operator( )(Arg x) const;
};
```

The `pointer_to_unary_function` class template is a function object that wraps a pointer to a function, in which the function is an ordinary (nonmember) function that takes one argument. The `ptr_fun` function template is the most convenient way to create a `pointer_to_unary_function` object.

## See Also

[pointer\\_to\\_binary\\_function class template](#), [ptr\\_fun function template](#)

## Name

`ptr_fun` function template — Creates a pointer to a function object

## Synopsis

```
template <typename Arg1, typename Arg2, typename Rtn>
    pointer_to_binary_function<Arg1,Arg2,Rtn>
        ptr_fun(Rtn (*f)(Arg1, Arg2));
template <typename Arg, typename Rtn>
    pointer_to_unary_function<Arg, Rtn>
        ptr_fun(Rtn (*f)(Arg));
```

The `ptr_fun` function template creates a function object from a pointer to a function. The resulting function object has an `operator()` that calls the function. Functions of one and two arguments are supported.

For example, suppose you have two numeric vectors, `a` and `b`, and you want to raise each element of `a` to the power of the corresponding element in `b`, saving the result in a third vector, `c`. There is no predefined power function object, so you can use `ptr_fun` and your own power function instead, as shown in [Example 13-13](#).

## Example

*Example 13-13. Wrapping the `pow` function in a function object*

```
std::vector<double> a, b, c;
double power(double x, double y)
{
    return std::pow(x, y);
}
...
std::transform(a.begin(), a.end(), b.begin(), c.begin(),
    ptr_fun(power));
```

## See Also

[mem\\_fun function template](#), [mem\\_fun\\_ref function template](#), [pointer to binary\\_function class template](#), [pointer to unary\\_function class template](#)

## Name

`unary_negate` class template — Logical negation of a unary predicate

## Synopsis

```
template <typename P>
class unary_negate :
    public unary_function<typename P::argument_type, bool>
{
public:
    explicit unary_negate(const P& predicate);
    bool operator( )(const typename P::argument_type& x) const;
};
```

The `unary_negate` class template is a binary functional that returns the logical negation of another unary functional— that is, `operator( )` returns `!predicate(x)`. The simplest way to use `unary_negate` is to use the `not1` function template.

## See Also

[binary\\_negate class template](#), [not1 function template](#)



## Name

`unary_function` class template — Base class for unary functionals

## Synopsis

```
template <typename Arg, typename Result>
struct unary_function {
    typedef Arg argument_type;
    typedef Result result_type;
};
```

The `unary_function` template is a base class for all the function classes that represent unary operations. It provides standard names for the argument and result types.

## See Also

[binary\\_function class template](#), [binder1st class template](#), [binder2nd class template](#), [const\\_mem\\_fun\\_ref\\_t class template](#), [const\\_mem\\_fun\\_t class template](#), [mem\\_fun\\_ref\\_t class template](#), [mem\\_fun\\_t class template](#), [negate class template](#), [pointer\\_to\\_unary\\_function class template](#), [unary\\_negate class template](#)

## <iomanip>

The <iomanip> header declares several I/O manipulators. An I/O manipulator is a function object that can be used in a sequence of input or output operators to manipulate the I/O stream. The manipulators are simple wrappers for functionality that is available as member functions of the `ios_base` class, but manipulators are simpler to use in some situations.

For example, to print formatted output, the following two code fragments are equivalent:

```
// Using manipulators
std::cout << std::setw(16) << std::setprecision(12) << x;

// Without manipulators
std::cout.width(16);
std::cout.precision(12);
std::cout << x;
```

At a basic level, manipulators are easy to use. If you want to understand exactly how they work, perhaps to write your own, see [Chapter 9](#) for a thorough discussion of I/O, including manipulators.



The return type of each manipulator is implementation-defined. For the following function descriptions, this type is shown as *manip\_t*.

Use a manipulator by applying it to a stream—that is, *out* << *manip*, in which *out* is an instance of `basic_ostream`, or *in* >> *manip*, in which *in* is an instance of `basic_istream`. In the following function descriptions, *stream* refers to the input or output stream to which the manipulator is being applied.

## Name

resetiosflags function — Clears specified flags

## Synopsis

```
manip_t  
resetiosflags(ios_base::fmtflags mask)
```

The `resetiosflags` function clears the flag bits in `mask` for a stream. In other words, it performs the equivalent of `stream.setf(ios_base::fmtflags(0), mask)`.

## See Also

[setiosflags function](#), `ios_base` in [<ios>](#)

## Name

setbase function — Sets conversion radix

## Synopsis

```
manip_t  
setbase(int base)
```

The setbase function sets the conversion radix for a stream. In other words, it performs the equivalent of *stream*.setf(newbase, ios\_base::basefield), in which *newbase* depends on base, as shown in [Table 13-11](#). Notice that any value of a base other than 8, 10, or 16 is treated the same as 0.

Table 13-11. Conversion radix

set base argument	mt flags equivalent
8	ios_base::oct
10	ios_base::dec
16	ios_base::hex
Anything else	0

## See Also

ios\_base in [<ios>](#)

## Name

setfill function template — Sets pad character

## Synopsis

```
template <typename charT>  
manip_t  
                setfill(charT c)
```

The `setfill` function template sets the fill character for a stream to `c`. In other words, it performs the equivalent of `stream.fill(c)`.

## See Also

`ios_base` in [<ios>](#)

## Name

setiosflags function — Sets specified flags

## Synopsis

```
manip_t  
setiosflags(ios_base::fmtflags mask)
```

The `setiosflags` function sets the flag bits in `mask` for a stream. In other words, it performs the equivalent of `stream .setf(mask)`.

## See Also

[resetiosflags function](#), `ios_base` in [<ios>](#)

## Name

setprecision function — Sets precision

## Synopsis

```
manip_t  
setprecision(int n)
```

The `setprecision` function template sets the output precision for a stream to `n`. In other words, it performs the equivalent of `stream.precision(n)`.

## See Also

`ios_base` in [<ios>](#)

## Name

setw function — Sets field width

## Synopsis

```
manip_t  
setw(int n)
```

The `setw` function template sets the output field width for a stream to `n`. In other words, it performs the equivalent of `stream.width(n)`.

## See Also

`ios_base` in [<ios>](#)



## <ios>

The <ios> header declares the classes, types, and manipulator functions that form the foundation of the C++ I/O library (which is often called *I/O streams*). The class `ios_base` is the base class for all I/O stream classes. The class template `basic_ios` derives from `ios_base` and declares the behavior that is common to all I/O streams (e.g., establishing a stream buffer and defining the I/O state).

Refer to [Chapter 9](#) for more information about input and output, including the use of manipulators, formatting flags, streams, and stream buffers.

The <ios> header `#includes` <iosfwd>.

# Name

`basic_ios` class template — Base class template for all I/O streams

# Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_ios : public ios_base
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;
    // Status
    operator void*( );
    const bool operator!( );
    const iostate rdstate( ) const;
    void clear(iostate state = goodbit);
    void clear(io_state state);
    void setstate(iostate state);
    void setstate(io_state state);
    bool good( ) const;
    bool eof( ) const;
    bool fail( ) const;
    bool bad( ) const;
    iostate exceptions( ) const;
    void exceptions(iostate except);
    void exceptions(io_state except);
    explicit basic_ios(basic_streambuf<charT,traits>* sb);
    virtual ~basic_ios( );
    basic_ostream<charT,traits>* tie( ) const;
    basic_ostream<charT,traits>* tie(basic_ostream<charT,traits>*
tiestr);
    basic_streambuf<charT,traits>* rdbuf( ) const;
    basic_streambuf<charT,traits>*
rdbuf(basic_streambuf<charT,traits>* sb);
    basic_ios& copyfmt(const basic_ios& rhs);
    char_type fill( ) const;
    char_type fill(char_type ch);
    locale imbue(const locale& loc);
    char narrow(char_type c, char deflt) const;
    char_type widen(char c) const;
```

```

protected:
    basic_ios( );
    void init(basic_streambuf<charT,traits>* buf);
private:
    basic_ios(const basic_ios& ); // Not defined
    basic_ios& operator=(const basic_ios&); // Not defined
};

```

The `basic_ios` class template is the root of all I/O stream class templates. It provides a common functionality for all derived stream classes; in particular, it manages a stream buffer. In the following descriptions, the name *buf* refers to a private data member that points to the stream buffer. An implementation can use any name.

The following are the member functions of `basic_ios`:

### **basic\_ios** ( )

The default constructor leaves the data members uninitialized. In a derived class, the default constructor must call `init` to initialize the members.

### **basic\_ios** (const **basic\_ios**&)

The copy constructor is declared private and is not defined, which prevents the copying of any I/O stream objects.

### explicit **basic\_ios** (basic\_streambuf<charT,traits>\* sb)

Calls `init(sb)` to initialize the members.

### **operator void\*** ( )

If `fail( )` returns true, the `void*` operator returns a null pointer; otherwise, it returns a non-null pointer to indicate success. `operator void*` is most often used as an implicit conversion in a conditional (e.g., `while (cin) cin >> data[i++]`).

### const **bool operator!** ( )

Returns `fail( )`. `operator !` is most often used in a conditional (e.g., `if (!cout) cerr << "output error\n"`).

`basic_ios& operator= (const basic_ios&)`

The assignment operator, like the copy constructor, is private, so it cannot be used and is not defined.

Assigning or copying an I/O stream would corrupt the stream buffer.

`bool bad ( ) const`

Returns true if `badbit` is set in `rdstate( )`, or false otherwise.

`void clear (iostate state = goodbit), void clear (io_state state)`

Sets the I/O state to `state`. If `rdbuf( )` is a null pointer, `badbit` is also set (to `state | ios_base::badbit`). After setting the state, if any state bit is an exception bit (`(rdstate( ) & exceptions( )) != 0`), `basic_ios::failure` is thrown.

The second form is deprecated. See `ios_base::iostate` later in this section for details.

`basic_ios& copyfmt (const basic_ios& rhs)`

Copies formatting information from `rhs`. In particular, the format flags, fill character, locale, and the contents of the `isword( )` and `isword( )` arrays are copied. The I/O state and stream buffer are not copied. Before copying any callback functions, each one is called with `erase_event`. The callbacks are then replaced with those copied from `rhs`, and each one is called with `copyfmt_event`. (See `ios_base` for information about callbacks.) The `exceptions( )` mask is copied last. The return value is `*this`.

`bool eof ( ) const`

Returns true if eofbit is set in rdbuf( ), or false otherwise.

ios\_base **exceptions** ( ) const, void **exceptions** (ios\_base except), void **exceptions** (ios\_base except)

Returns or sets the exception mask. (See the clear function for how and when an exception is thrown.) The third form is deprecated. See ios\_base::exceptions later in this section for details.

bool **fail** ( ) const

Returns true if badbit is set or if failbit is set in rdbuf( ), or false if neither bit is set.

char\_type **fill** ( ) const, char\_type **fill** (char\_type ch)

Returns or changes the fill character (also called the *pad character*). When setting the fill character, the old fill character is returned.

bool **good** ( ) const

Returns true if the I/O state is clean—that is, it returns rdbuf( ) == 0.

locale **imbue** (const locale& loc)

Calls ios\_base::imbue(loc) and rdbuf( )->pubimbue(loc) (if rdbuf( ) is not null). The return value is the previous value of ios\_base::imbue( ).

void **init** (basic\_streambuf<charT,traits>\* buf)

Initializes the basic\_ios object. [Table 13-12](#) lists the observable effects of initialization. Also, the arrays for iword( ) and pword( ) are initially null pointers.

*Table 13-12. Effects of calling basic\_ios::init*

Member function	Return value
<code>exceptions( )</code>	<code>goodbit</code>
<code>fill( )</code>	<code>widen(' ')</code>
<code>flags( )</code>	<code>skipws   dec</code>
<code>getloc( )</code>	Current global locale, that is, <code>std::locale( )</code>
<code>precision( )</code>	<code>6</code>
<code>rdbuf( )</code>	<code>buf</code>
<code>rdstate( )</code>	<code>buf !=0 ? goodbit : badbit</code>
<code>tie( )</code>	Null pointer
<code>width( )</code>	<code>0</code>

`char narrow (char_type c, char deflt) const`

Narrows the character `c` by returning the following:

```
std::use_facet<ctype<char_type>>(getloc( )).narrow(c, deflt)
```

`basic_streambuf<charT,traits>* rdbuf ( ) const,`  
`basic_streambuf<charT,traits>*, rdbuf`  
`(basic_streambuf<charT,traits>* sb)`

Returns or changes the stream buffer, `buf`. After changing the stream buffer, the `rdbuf` function calls `clear( )`. The function returns the previous value of `rdbuf( )`.

`const iostate rdstate ( ) const`

Returns the current I/O state bitmask. See the `bad`, `eof`, `fail`, and `good` functions for convenient ways to test different bits in the state mask.

```
void setstate (iostate state), void setstate  
(io_state state)
```

Sets the specified bits in the I/O state bitmask—that is, it calls `clear(rdstate( ) | state)`. The second form is deprecated. See `ios_base::iostate` later in this section for details.

```
basic_ostream<charT,traits>* tie ( ) const,  
basic_ostream<charT,traits>* tie  
(basic_ostream<charT,traits>* tiestr)
```

Ties a stream (typically an input stream) to an output stream, `tiestr`. Any input operation on this stream is prefaced by flushing `tiestr`. Tying streams can be used to ensure that prompts appear at the proper time. With no arguments, the `tie` function returns the currently tied stream, or `0` if no stream is tied.

```
char_type widen (char c) const
```

Widens the character `c` by returning the following:

```
std::use_facet<ctype<char_type> >(getloc( )).widen(c)
```

## See Also

[ios\\_base class](#), `ctype` in [<locale>](#), `basic_streambuf` in [<streambuf>](#)

## Name

boolalpha function — Manipulator for reading and writing bool as text

## Synopsis

```
ios_base& boolalpha(ios_base& stream)
```

The boolalpha function is a manipulator that sets the boolalpha flag, which tells the stream to read or write a bool value as text, according to the stream's locale. Specifically, the function calls `stream.setf(ios_base::boolalpha)` and returns stream.

## See Also

[ios\\_base::fmtflags type](#), [nboolalpha function](#), num\_get in [<locale>](#), num\_put in [<locale>](#)



## Name

dec function — Manipulator for decimal integers

## Synopsis

```
ios_base& dec(ios_base& stream)
```

The dec function is a manipulator that sets the conversion radix to base 10. The function calls `stream.setf(ios_base::dec, ios_base::basefield)` and returns stream.

## See Also

[hex function](#), [ios\\_base::fmtflags type](#), [noshowbase function](#), [oct function](#), [showbase function](#), `num_get` in [<locale>](#), `num_put` in [<locale>](#)

## Name

fixed function — Manipulator for fixed-point output

## Synopsis

```
ios_base& fixed(ios_base& stream)
```

The fixed function is a manipulator that sets the floating-point output style to fixed-point. The function calls `stream.setf(ios_base::fixed, ios_base::floatfield)` and returns stream.

## See Also

[ios\\_base::fmtflags type](#), [noshowpoint function](#), [scientific function](#), [showpoint function](#), `num_get` in [<locale>](#), `num_put` in [<locale>](#)

# Name

fpos class template — Represents a file position

## Synopsis

```
template <typename stateT>
class fpos{
public:
    stateT state( ) const;
    void state(stateT);

    // The following functionality is required, although not
    // necessarily as member
    // functions.
    fpos(int i);
    fpos(streamoff offset);
    operator streamoff( ) const;
    bool operator==(const fpos& rhs) const;
    bool operator!=(const fpos& rhs) const;
    fpos operator+(streamoff offset) const;
    fpos operator-(streamoff offset) const;
    fpos& operator+=(streamoff offset);
    fpos& operator-=(streamoff offset);
    streamoff operator-(const fpos& rhs) const;
};
```

The `fpos` class template represents a position in a stream. The `stateT` template parameter is a multibyte shift state, such as `mbstate_t`. Objects of type `fpos` can be compared for equality or inequality, they can be subtracted to yield a stream offset, or a stream offset can be added to an `fpos` position to produce a new `fpos`. Also, stream offsets can be converted to and from `fpos` values. Although the declaration in this section shows these functions as member functions, they might be global functions or be provided in some other fashion.

**The Shift State Problem**

The C++ standard requires a library implementation to implement `fpos<>` and the `streamoff` types so that you can convert an object whose type is any `fpos<>` instance into a `streamoff` object and back again, recovering the original value.

Most library implementations do not work this way.

Instead, `streamoff` is usually implemented as an integral type, such as `int` or `long`. The problem is that `fpos<>` must store a stream position and a shift state, and the stream position is usually implemented with the same type as `streamoff`. In other words, converting from `fpos<>` to `streamoff` discards the shift state. The library cannot convert `streamoff` back to the original `fpos<>` because the original shift state is gone.

Note that `char_traits` has the same problem when specialized for `char` and `wchar_t` because its `pos_type` is defined to be `streampos`, which is defined in `<iosfwd>` to be `fpos<mbstate_t>`, and `off_type` is defined as `streamoff`.

The solution is not to convert a stream position to a stream offset. If you must do this, remember that you might be sacrificing the shift state. Save the shift state separately so you can restore it when converting a stream offset back to a stream position.

## See Also

[streamoff type](#), `mbstate_t` in [<cwchar>](#)

## Name

hex function — Manipulator for hexadecimal integers

## Synopsis

```
ios_base& hex(ios_base& stream)
```

The hex function is a manipulator that sets the conversion radix to base 16. The function calls `stream.setf(ios_base::hex, ios_base::basefield)` and returns stream.

## See Also

[dec function](#), [ios\\_base::fmtflags type](#), [noshowbase function](#), [oct function](#), [showbase function](#), `num_get` in [<locale>](#), `num_put` in [<locale>](#)

## Name

internal function — Manipulator to align output on an internal point

## Synopsis

```
ios_base& internal(ios_base& stream)
```

The `internal` function is a manipulator that sets the stream's alignment to internal. The function calls `stream.setf(ios_base::internal, ios_base::adjustfield)` and returns `stream`. Internal padding works as follows:

- If the formatted number begins with a sign, insert the padding after the sign.
- If the formatted number begins with `0x` or `0X`, insert the padding after the `x` or `X`.
- Otherwise, insert the padding before the number (like `ios_base::right`).

## See Also

[ios\\_base::fmtflags type](#), [left function](#), [right function](#)

# Name

ios\_base class — Root class for I/O declarations

# Synopsis

```
class ios_base{
public:
    class failure;
    typedef . . .   fmtflags;
    typedef . . .   iostate;
    typedef . . .   io_state;
    typedef . . .   openmode;
    typedef . . .   open_mode;
    typedef . . .   seekdir;
    typedef . . .   seek_dir;
    typedef . . .   streamoff;
    typedef . . .   streampos;
    class Init;

    // Destructor
    virtual ~ios_base( );
    // Formatting
    fmtflags flags( ) const;
    fmtflags flags(fmtflags fmtfl);
    fmtflags setf(fmtflags fmtfl);
    fmtflags setf(fmtflags fmtfl, fmtflags mask);
    void unsetf(fmtflags mask);
    streamsize precision( ) const;
    streamsize precision(streamsize prec);
    streamsize width( ) const;
    streamsize width(streamsize wide);
    // Locales
    locale imbue(const locale& loc);
    locale getloc( ) const;
    // Storage
    static int xalloc( );
    long& iword(int index);
    void*& pword(int index);
    // Callbacks
    enum event { erase_event, imbue_event, copyfmt_event };
    typedef void (*event_callback)(event, ios_base&, int index);
    void register_callback(event_callback fn, int index);
    static bool sync_with_stdio(bool sync = true);
protected:
```

```

    ios_base( );
private:
    ios_base(const ios_base&);
    ios_base& operator=(const ios_base&);
};

```

The `ios_base` class is the root class for all the I/O stream classes. It declares fundamental types that are used throughout the I/O library. It also has members to keep track of formatting for input and output, storing arbitrary information for derived classes, and registering functions to be called when something interesting happens to the stream object.

The `io_state`, `open_mode`, `seek_dir`, `streamoff`, and `streampos` types are deprecated and might not be included in a future revision of the C++ standard. The first three are integer types that are equivalent to `iostate`, `openmode`, and `seekdir`. (See their respective subsections later in this section for details.) The `streamoff` and `streampos` types have equivalent types at namespace scope. See `streamoff` later in this section and `streampos` in `<iosfwd>` for details.

The following are the member functions of `ios_base`:

### **ios\_base**

The default constructor is protected so you cannot accidentally declare an object of type `ios_base`. It does not initialize its members. That is left to the `basic_ios::init` function.

### **ios\_base**

The copy constructor is private and not defined so you cannot copy objects of type `ios_base` or its derived classes.

### **~ios\_base**

Calls every registered callback with the `erase_event` if the `ios_base` object has been properly initialized. See



`basic_ios::init.`

`ios_base& operator= (const ios_base&)`

The assignment operator is private and not defined to prevent the assignment of `ios_base` objects or its derivatives.

`fmtflags flags ( ) const, fmtflags flags (fmtflags fmtfl)`

Returns the current format flags or sets the flags. When setting the flags, the previous flags are returned.

`locale getloc ( ) const`

Returns the stream's currently imbued locale.

`locale imbue (const locale& loc)`

Saves `loc` as the new locale and calls all registered callbacks with `imbue_event`. The new locale is stored before calling any callbacks, so if a callback function calls `getloc`, it gets the new locale.

`long& iword (int index)`

Returns a reference to a long integer that is stored in a private array, at index `index`. If `iword` has been called before with the same index, a reference to the array element is returned. Otherwise, the array is extended as needed so that `index` is a valid index, and the new entry is initialized to 0. A reference to the new element is returned.

The structure of the internal array is implementation-defined, and it might be a sparse array. The internal array grows as needed, and all prior values stored in the array are preserved, although the references might become invalid in any of the following situations:

- After a call to `iword` with a different index

- After calling `basic_ios::copyfmt` for this object
- When the object is destroyed

If `word` fails (perhaps because the internal array cannot grow), it returns a reference to a valid `long&` with a value that is initially `0`. If the member function is called for an object whose class derives from `basic_ios<>`, `badbit` is set (which might throw `ios_base::failure`).

See the `xalloc` member function to learn how to obtain a suitable index.

`streamsize` **precision** ( ) `const`, `streamsize` **precision** (`streamsize` `prec`)

Returns or sets the precision (places after the decimal point) used to format floating-point numbers for output. When setting a new precision, the previous precision is returned.

`void*&` **pword** (`int` `index`)

Returns a reference to a `void*` that is stored in a private array, at index `index`. If `pword` has been called before with the same index, a reference to the array element is returned. Otherwise, the array is extended as needed so that `index` is a valid index, and the new entry is initialized to a null pointer. A reference to the new element is returned.

The structure of the internal array is implementation-defined, and it might be a sparse array. The internal array grows as needed, and all prior values stored in the array are preserved, although the references might become invalid in any of the following situations:

- After a call to `pword` with a different index
- After calling `basic_ios::copyfmt` for this object
- When the object is destroyed

If `pwd` fails (perhaps because the internal array cannot grow), it returns a reference to a valid `void*&` with a value that is initially `0`. If the object derives from `basic_ios<>`, `badbit` is set (which might throw `ios_base::failure`).

See the `xalloc` member function to learn how to obtain a suitable index.

```
void register_callback (event_callback fn, int index)
```

Registers a function `fn` to be called when one of three events occurs for the `ios_base` object:

- The object is destroyed (`erase_event`)
- `copyfmt` is called (`erase_event` followed by `copyfmt_event`)
- `imbue` is called (`imbue_event`)

Each callback function is registered with an integer index. The index is passed to the callback function. Functions are called in the opposite order of registration. The callback function must not throw exceptions.

For example, suppose a program stores some debugging information with each stream. It allocates a struct and stores a pointer to the struct in the stream's `pwd` array. When `copyfmt` is called, the debugging information should also be copied. [Example 13-14](#) shows how to use callbacks to make sure the memory is managed properly.

*Example 13-14. Copying information associated with streams*

```
void manage_info(std::ios_base::event event,
                std::ios_base& stream, int index)
{
    infostruct* ip;
```

```

switch(event) {
case std::ios_base::erase_event:
    ip = static_cast<infostruct*>(stream.pword(index));
    stream.pword(index) = 0;
    delete ip;
    break;
case std::ios_base::copyfmt_event:
    stream.pword(index) = new infostruct;
    break;
default:
    break; // imbue_event does not affect storage.
}
}

void openread(std::ifstream& f, const char* name)
{
    f.open(name);
    int index = f.xalloc( );
    f.pword(index) = new infostruct;
    f.register_callback(manage_info, index);
}

```

`fmtflags` **setf** (fmtflags addflags)

Sets the addflags bits of the formatting flags. It is equivalent to calling `flags( ) | addflags`.

`fmtflags` **setf** (fmtflags newflags, fmtflags mask)

Clears the mask bits from the formatting flags and then sets the newflags & mask bits. It is equivalent to calling `flags((flags( ) & ~mask) | (newflags & mask))`. The two-argument version of `setf` is most often used with multiple-choice flags (e.g., `setf(ios_base::dec, ios_base::basefield)`).

`static bool` **sync\_with\_stdio** (bool sync = true)

Determines whether the standard C++ I/O objects are synchronized with the C I/O functions. Initially, they are synchronized.



If you call `sync_with_stdio(false)` after any I/O has been performed, the behavior is implementation-defined.

`void unsetf (fmtflags mask)`

Clears the mask bits from the formatting flags. It is equivalent to calling `flags( ) & ~mask`.

`streamsize width ( ) const, streamsize width  
(streamsize wide)`

Returns or sets the minimum field width. When setting the width, the previous width is returned.

`static int xalloc ( )`

Returns a unique integer, suitable for use as an index to the `isword` or `isword` functions. You can think of `ios_base` as having a static integer data member, `xalloc_index`, and `xalloc` is implemented so it returns `xalloc_index++`.

## See Also

[basic\\_ios class template](#)

## Name

`ios_base::event` type — Callback event type

## Synopsis

```
enum event { erase_event, imbue_event, copyfmt_event };
```

The `ios_base::event` type denotes an interesting event in the lifetime of an I/O stream object. See the `register_callback` function in the `ios_base` class, earlier in this section, to learn how to register a function that is called when one of these events occurs.

## See Also

[ios\\_base class](#), [ios\\_base::event\\_callback type](#)

## Name

`ios_base::event_callback` type — Callback function type

## Synopsis

```
typedef void (*event_callback)(event, ios_base&, int index);
```

The `ios_base::event_callback` type denotes a callback function. See the `register_callback` function in the `ios_base` class, earlier in this section, to learn how to register a callback function, which a stream object calls when an interesting event occurs.

## See Also

[ios\\_base class](#), [ios\\_base::event type](#)

## Name

`ios_base::failure` class — Exception class for I/O failure

## Synopsis

```
class ios_base::failure : public exception
{
public:
    explicit failure(const string& msg);
    virtual ~failure( );
    virtual const char* what( ) const throw( );
};
```

The `ios_base::failure` class is the base class for I/O-related exceptions. Its use of the constructor's `msg` parameter and `what( )` member function are consistent with the conventions of the exception class.

## See Also

`basic_ios::clear` function, exception in [<exception>](#)



# Name

ios\_base::fmtflags type — Formatting flags

## Synopsis

```
typedef . . . fmtflags;  
static const fmtflags boolalpha;  
static const fmtflags dec;  
static const fmtflags fixed;  
static const fmtflags hex;  
static const fmtflags internal;  
static const fmtflags left;  
static const fmtflags oct;  
static const fmtflags right;  
static const fmtflags scientific;  
static const fmtflags showbase;  
static const fmtflags showpoint;  
static const fmtflags showpos;  
static const fmtflags skipws;  
static const fmtflags unitbuf;  
static const fmtflags uppercase;  
static const fmtflags adjustfield;  
static const fmtflags basefield;  
static const fmtflags floatfield;
```



The `fmtflags` type is an integer, enum, or bitmask type (the exact type is implementation-defined) that represents formatting flags for input and output. In the `ios_base` class, several static constants are also defined, which can be implemented as enumerated literals or as explicit constants. [Table 13-13](#) lists the flag literals.

*Table 13-13. `fmtflags` literals*

Literal name	Description
<code>boolalpha</code>	Reads and writes bool values as text, according to the locale

Literal name	Description
dec	Reads and writes decimal integers
fixed	Writes floating-point values in fixed notation
hex	Reads and writes hexadecimal integers
internal	Aligns output to internal point (e.g., after sign or 0x)
left	Left-aligns output
oct	Reads and writes octal integers
right	Right-aligns output
scientific	Writes floating-point values in scientific notation
showbase	Writes a prefix for an integer radix ( e.g., 0x for hexadecimal)
showpoint	Writes decimal point, even if not needed
showpos	Writes plus sign (+), even if not needed
skipws	Skips whitespace before input
unitbuf	Flushes output after each operation
uppercase	Uses uppercase in generated output (e.g., 0X prefix)

Some formatting items are Boolean: a flag is set or cleared. For example, the uppercase flag can be set to perform output in uppercase (that is, the 0X hexadecimal prefix or E in scientific notation), or the flag can be cleared for lowercase output. Other flags are set in *fields*. You can set a field to one of a number of values. [Table 13-14](#) lists the field names, definitions, and the default behavior if the field

value is 0. Each field name is used as a mask for the two-argument form of the `ios_base::setf` function.

*Table 13-14. fmtflags constants*

Constant name	Value	Default
adjustfield	left   internal   right	right
basefield	dec   hex   oct	Output: decInput: leading 0x or 0X is hex, 0 is oct, anything else is dec
floatfield	fixed   scientific	scientific if exponent is $< -4$ or $\geq$ precision, else fixed; strip trailing zeros and unneeded decimal point

## See Also

[ios\\_base class](#), `ctype` in [<locale>](#), `num_get` in [<locale>](#), `num_put` in [<locale>](#)

# Name

`ios_base::Init` class — Initialization class

# Synopsis

```
class ios_base::Init {
public:
    Init( );
    ~Init( );
};
```

The `Init` class is used to ensure that the construction of the standard I/O stream objects occurs. The first time an `ios_base::Init` object is constructed, it constructs and initializes `cin`, `cout`, `cerr`, `clog`, `wcin`, `wcout`, `wcerr`, and `wclog`. A static counter keeps track of the number of times `ios_base::Init` is constructed and destroyed. When the last instance is destroyed, `flush( )` is called for `cout`, `cerr`, `clog`, `wcout`, `wcerr`, and `wclog`.

For example, suppose a program constructs a static object, and the constructor prints a warning to `cerr` if certain conditions hold. To ensure that `cerr` is properly initialized and ready to receive output, declare an `ios_base::Init` object before your static object, as shown in [Example 13-15](#).

# Example

*Example 13-15. Ensuring proper initialization of standard I/O streams*

```
class myclass {
public:
    myclass( ) {
        if (! okay( ))
            std::cerr << "Oops: not okay!\n";
    }
};
```

```
static std::ios_base::Init init;  
static myclass myobject;
```

## See Also

[<iostream>](#)

# Name

`ios_base::iostate` type — I/O status

# Synopsis

```
typedef . . . iostate
typedef . . . io_state
static const iostate badbit static const iostate eofbit
static const iostate failbit
static const iostate goodbit = iostate(0);
```



The `ios_base::iostate` type is an integer, enum, or bitset type (the exact type is implementation-defined) that represents the status of an I/O stream. The `io_state` type is an integral type that represents the same information. Some functions that take an `iostate` parameter have an overloaded version that accepts an `io_state` parameter and has the same functionality as its `iostate` counterpart. The `io_state` type and related functions are deprecated, so you should use the `iostate` versions.

[Table 13-15](#) lists the `iostate` literals and their meanings. The `basic_ios` class template has several member functions for setting, testing, and clearing `iostate` bits.

*Table 13-15. iostate literals*

<b>Literal</b>	<b>Description</b>
<code>badbit</code>	Irrecoverable error, such as a null <code>streambuf</code> pointer or a write failure
<code>eofbit</code>	End-of-file when reading
<code>failbit</code>	Failure to read or write expected characters ( e.g., trying to read an integer from nonnumeric input)

Litera	Description
goodbit	No problems; value is 0

## See Also

[basic\\_ios class template](#), <bitset>

# Name

`ios_base::openmode` type — Open mode bits

## Synopsis

```
typedef . . . openmode
typedef . . . open_mode
static const openmode app
static const openmode ate
static const openmode binary
static const openmode in
static const openmode out
static const openmode trunc
```



The `ios_base::openmode` type is an integer, enum, or bitset type (the exact type is implementation-defined) that defines the mode for opening a file. The `open_mode` type is an integral type that represents the same information. Some functions that take an `openmode` parameter have an overloaded version that accepts an `open_mode` parameter and has the same functionality as its `openmode` counterpart. The `open_mode` type and related functions are deprecated, so you should use the `openmode` versions.

[Table 13-16](#) lists the `openmode` literals and their meanings. (Refer to the `<fstream>` section of this chapter for the most common use of `ios_base::openmode` and the permitted combinations of `openmode` literals.) The `openmode` type is also used for the `basic_streambuf::pubseekoff` and `pubseekpos` functions, and for related functions.

*Table 13-16. openmode literals*

Literal	Description
---------	-------------



Literal	Description
app	Seeks to end-of-file before each write
ate	Seeks to end-of-file immediately after opening
binary	Reads and writes in binary mode (default is text)
in	Opens for input (reading)
out	Opens for output (writing)
trunc	Truncates file to zero length

## See Also

<bitset>, basic\_filebuf in [<fstream>](#), basic\_streambuf in [<streambuf>](#)

# Name

`ios_base::seekdir` type — Seek direction

## Synopsis

```
typedef . . . seekdir  
typedef . . . seek_dir  
static const seekdir beg  
static const seekdir cur  
static const seekdir end
```



The `ios_base::seekdir` type is an implementation-defined enumerated type that specifies the origin for seeking to a new file position. The `seek_dir` type is an integral type that represents the same information. Some functions that take a `seekdir` parameter have overloaded versions that accept a `seek_dir` parameter and have the same functionality as their `seekdir` counterparts. The `seek_dir` type and related functions are deprecated, so you should use the `seekdir` versions.

[Table 13-17](#) lists the `seekdir` literals. Note that the order and integer values of the literals are implementation-defined.

*Table 13-17. seekdir literals*

Literal	Description
beg	Seeks from the beginning of the stream (e.g., offset is absolute position)
cur	Seeks from the current position; positive is toward end of stream, and negative offsets are toward the beginning of the stream

Literal	Description
end	Seeks relative to the end of the stream; negative offsets are towards the beginning

## See Also

basic\_istream in [<istream>](#), basic\_ostream in [<ostream>](#), basic\_streambuf in [<streambuf>](#)

## Name

left function — Manipulator to left-align output

## Synopsis

```
ios_base& left(ios_base& stream)
```

The `left` function is a manipulator that selects left-alignment for output to `stream`. The function calls `stream.setf(ios_base::left, ios_base::adjustfield)` and returns `stream`.

## See Also

[internal function](#), [ios\\_base::fmtflags type](#), [right function](#)

## Name

noboolalpha function — Manipulator to disable reading and writing bool as text

## Synopsis

```
ios_base& noboolalpha(ios_base& stream)
```

The noboolalpha function is a manipulator that clears the boolalpha flag, causing the stream to read or write bool values as integers. Specifically, the function calls `stream.unsetf(ios_base::boolalpha)` and returns `stream`.

## See Also

[boolalpha function](#), [ios\\_base::fmtflags type](#), ctype in [<locale>](#)

## Name

noshowbase function — Manipulator to disable showing output radix

## Synopsis

```
ios_base& noshowbase(ios_base& stream)
```

The noshowbase function is a manipulator that clears the showbase flag, which tells an output stream to write a prefix for integer output: 0x for hexadecimal or 0 for octal. Specifically, the function calls stream.unsetf(ios\_base::showbase) and returns stream.

## See Also

[hex function](#), [ios\\_base::fmtflags type](#), [oct function](#), [nouppercase function](#), [showbase function](#), [uppercase function](#), num\_put in [<locale>](#)

## Name

noshowpoint function — Manipulator to suppress unnecessary decimal points

## Synopsis

```
ios_base& noshowpoint(ios_base& stream)
```

The `noshowpoint` function is a manipulator that clears the `showpoint` flag, causing an output stream to write a decimal point for floating-point output, even if the point is unnecessary (only zeros appear after the decimal point). Specifically, the function calls `stream.unsetf(ios_base::showpoint)` and returns `stream`.

## See Also

[fixed function](#), [ios\\_base::fmtflags type](#), [scientific function](#), [showpoint function](#), `num_put` in [<locale>](#)

## Name

noshowpos function — Manipulator to suppress plus sign in nonnegative output

## Synopsis

```
ios_base& noshowpos(ios_base& stream)
```

The noshowpos function is a manipulator that clears the showpos flag, which causes an output stream to always write a plus sign (+) in front of a number even if the sign is unnecessary (the value is 0 or positive). Specifically, the function calls `stream.unsetf(ios_base::showpos)` and returns `stream`.

## See Also

[ios\\_base::fmtflags type](#), [showpos function](#), `num_put` in [<locale>](#)



## Name

`noskipws` function — Manipulator to disable skipping whitespace before reading

## Synopsis

```
ios_base& noskipws(ios_base& stream)
```

The `noskipws` function is a manipulator that clears the `skipws` flag, which tells an input stream to skip whitespace before reading most fields. Specifically, the function calls `stream.unsetf(ios_base::skipws)` and returns `stream`.

## See Also

[ios\\_base::fmtflags type](#), [skipws function](#), `num_get` in [<locale>](#)

## Name

nounitbuf function — Manipulator to use unit buffering

## Synopsis

```
ios_base& nounitbuf(ios_base& stream)
```

The `nounitbuf` function is a manipulator that clears the `unitbuf` flag, so the stream is not flushed after each output operation. Specifically, the function calls `stream.unsetf(ios_base::unitbuf)` and returns `stream`.

## See Also

[ios\\_base::fmtflags](#) type, [unitbuf](#) function

## Name

nouppercase function — Manipulator to use lowercase in generated output

## Synopsis

```
ios_base& nouppercase(ios_base& stream)
```

The nouppercase function is a manipulator that clears the uppercase flag, which tells an output stream to use uppercase letters for generated output (e.g., 0X for hexadecimal prefix or E for exponents). Specifically, the function calls `stream.unsetf(ios_base::uppercase)` and returns `stream`.

## See Also

[hex function](#), [ios\\_base::fmtflags type](#), [scientific function](#), [uppercase function](#), `num_put` in [<locale>](#)

## Name

oct function — Manipulator for octal integers

## Synopsis

```
ios_base& oct(ios_base& stream)
```

The oct function is a manipulator that sets the conversion radix to base 8. The function calls `stream.setf(ios_base::oct, ios_base::basefield)` and returns stream.

## See Also

[dec function](#), [hex function](#), [ios\\_base::fmtflags type](#), [noshowbase function](#), [showbase function](#), `num_get` in [<locale>](#), `num_put` in [<locale>](#)

## Name

right function — Manipulator to right-align output

## Synopsis

```
ios_base& right(ios_base& stream)
```

The `right` function is a manipulator that selects right-alignment for output to `stream`. The function calls `stream.setf(ios_base::right, ios_base::adjustfield)` and returns `stream`.

## See Also

[internal function](#), [ios\\_base::fmtflags type](#), [left function](#)

## Name

scientific function — Manipulator to use scientific notation for output

## Synopsis

```
ios_base& scientific(ios_base&)
```

The fixed function is a manipulator that sets the floating-point output style to scientific or exponential notation. The function calls `stream.setf(ios_base::scientific, ios_base::floatfield)` and returns `stream`.

## See Also

[fixed function](#), [ios\\_base::fmtflags type](#), `num_get` in [<locale>](#), `num_put` in [<locale>](#)

## Name

showbase function — Manipulator to show output radix

## Synopsis

```
ios_base& showbase(ios_base& stream)
```

The showbase function is a manipulator that sets the showbase flag, which tells an output stream to write a prefix for integer output: 0x for hexadecimal or 0 for octal. Specifically, the function calls `stream.setf(ios_base::showbase)` and returns stream.

## See Also

[hex function](#), [ios\\_base::fmtflags type](#), [oct function](#), [noshowbase function](#), [nouppercase function](#), [uppercase function](#), `num_put` in [<locale>](#)

## Name

showpoint function — Manipulator to show decimal point even when unnecessary

## Synopsis

```
ios_base& showpoint(ios_base& stream)
```

The showpoint function is a manipulator that sets the showpoint flag, which tells an output stream to write a decimal point for floating-point output, even if the point is unnecessary (only zeros appear after the decimal point). Specifically, the function calls `stream.setf(ios_base::showpoint)` and returns `stream`.

## See Also

[fixed function](#), [ios\\_base::fmtflags type](#), [noshowpoint function](#), [scientific function](#), `num_put` in [<locale>](#)



## Name

showpos function — Manipulator to show plus sign for nonnegative numbers

## Synopsis

```
ios_base& showpos(ios_base& stream)
```

The showpos function is a manipulator that sets the showpos flag, which tells an output stream to write a plus (+) sign, even if the sign is unnecessary (the value is 0 or positive). Specifically, the function calls `stream.setf(ios_base::showpos)` and returns stream.

## See Also

[ios\\_base::fmtflags type](#), [noshowpos function](#), num\_put in [<locale>](#)

## Name

skipws function — Manipulator to skip whitespace before reading

## Synopsis

```
ios_base& skipws(ios_base& stream)
```

The skipws function is a manipulator that sets the skipws flag, which tells an input stream to skip whitespace before reading most fields. Specifically, the function calls `stream.setf(ios_base::skipws)` and returns stream.

## See Also

[ios\\_base::fmtflags type](#), [noskipws function](#), `num_get` in [<locale>](#)

## Name

streamoff type — Stream offset type

## Synopsis

```
typedef . . . streamoff
```



The `streamoff` type is an implementation-defined type that represents a signed offset in a stream. See the `fpos` type for more information about working with file positions and offsets.

## See Also

`fpos` type, `char_traits` in `<string>`

## Name

streamsize type — Stream size type

## Synopsis

```
typedef . . . streamsize
```



The `streamsize` type is an implementation-defined type that is used to represent the size of various stream entities, such as number of characters to read or write. It is a synonym for one of the signed integral types. You can convert a `streamsize` to a `streamoff` without loss of information. You can also convert a `streamoff` back to a `streamsize`. If the `streamoff` is the result of converting a `streamsize`, converting the `streamoff` back to `streamsize` yields the original `streamsize`.

## See Also

`fpos` type, `char_traits` in [<string>](#)

## Name

unitbuf function — Manipulator to use unit buffering

## Synopsis

```
ios_base& unitbuf(ios_base& stream)
```

The `unitbuf` function is a manipulator that sets the `unitbuf` flag, which causes the stream to be flushed after each output operation. Specifically, the function calls `stream.setf(ios_base::unitbuf)` and returns `stream`.

## See Also

[ios\\_base::fmtflags type](#), [nunitbuf function](#)

## Name

uppercase function — Manipulator to use uppercase for generated output

## Synopsis

```
ios_base& uppercase(ios_base& stream)
```

The uppercase function is a manipulator that sets the uppercase flag, which tells an output stream to use uppercase letters for generated output (e.g., 0X for hexadecimal prefix or E for exponents). Specifically, the function calls `stream.setf(ios_base::uppercase)` and returns `stream`.

## See Also

[hex function](#), [ios\\_base::fmtflags type](#), [nouppercase function](#), [scientific function](#), `num_put` in [<locale>](#)

## <iosfwd>

The <iosfwd> header provides forward declarations of the various I/O-related classes and templates. The forward declarations are incomplete type declarations. (Read about complete and incomplete types in [Chapter 6](#).) By using <iosfwd> you can sometimes avoid including the complete definitions of the I/O classes, and thereby save some compilation time.

It is always safe to include <iosfwd> and any of the other I/O headers, even if they both declare the same type.

Because <iosfwd> does not provide any complete type declarations, this section contains only [Table 13-18](#), which lists each type name with a reference to the header (if any) that is required for the complete type.

*Table 13-18. Forward type declarations in <iosfwd>*

Type declaration	Header
template<typename T> class <b>allocator</b>	<memory>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_filebuf</b>	<fstream>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_fstream</b>	<fstream>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_ifstream</b>	<fstream>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_ios</b>	<ios>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_istream</b>	<istream>

Type decl ar at i on	Header
template<typename charT, typename traits=char_traits<charT> > class <b>basic_istream</b>	<istream>
template<typename charT, typename traits=char_traits<charT>, typename Alloc = allocator<charT> > class <b>basic_istreamream</b>	<sstream>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_ofstream</b>	<fstream>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_ostream</b>	<ostream>
template<typename charT, typename traits=char_traits<charT>, typename Alloc = allocator<charT> > class <b>basic_ostreamream</b>	<sstream>
template<typename charT, typename traits=char_traits<charT> > class <b>basic_streambuf</b>	<streambuf>
template<typename charT, typename traits=char_traits<charT>, typename Alloc = allocator<charT> > class <b>basic_stringbuf</b>	<sstream>
template<typename charT, typename traits=char_traits<charT>, typename Alloc = allocator<charT> > class <b>basic_stringstream</b>	<sstream>
template<typename charT> class <b>char_traits</b>	<string>
template<> class <b>char_traits&lt;char&gt;</b>	<string>
template<> class <b>char_traits&lt;wchar_t&gt;</b>	<string>
typedef basic_filebuf<char> <b>filebuf</b>	<fstream>
template <typename state> class <b>fpos</b>	<ios>
typedef basic_fstream<char> <b>fstream</b>	<fstream>
typedef basic_ifstream<char> <b>ifstream</b>	<fstream>
typedef basic_ios<char> <b>ios</b>	<ios>



Type decl ar at i on	Header
typedef basic_iostream<char> <b>iostream</b>	<iostream>
typedef basic_istream<char> <b>istream</b>	<istream>
template<typename charT, typename traits=char_traits<charT> > class <b>istreambuf_iterator</b>	<iterator>
typedef basic_istringstream<char> <b>istringstream</b>	<sstream>
typedef basic_ofstream<char> <b>ofstream</b>	<fstream>
typedef basic_ostream<char> <b>ostream</b>	<ostream>
template<typename charT, typename traits=char_traits<charT> > class <b>ostreambuf_iterator</b>	<iterator>
typedef basic_ostringstream<char> <b>ostringstream</b>	<sstream>
typedef basic_streambuf<char> <b>streambuf</b>	<streambuf>
typedef fpos<char_traits<char>::state_type> <b>streampos</b>	N/A
typedef basic_stringbuf<char> <b>stringbuf</b>	<sstream>
typedef basic_stringstream<char> <b>stringstream</b>	<sstream>
typedef basic_filebuf<wchar_t> <b>wfilebuf</b>	<fstream>
typedef basic_fstream<wchar_t> <b>wfstream</b>	<fstream>
typedef basic_ifstream<wchar_t> <b>wifstream</b>	<fstream>
typedef basic_ios<wchar_t> <b>wios</b>	<ios>
typedef basic_iostream<wchar_t> <b>wiostream</b>	<iostream>

Type decl ar at i on	Header
typedef basic_istream<wchar_t> <b>wistream</b>	<istream>
typedef basic_istringstream<wchar_t> <b>wistringstream</b>	<sstream>
typedef basic_ofstream<wchar_t> <b>wofstream</b>	<fstream>
typedef basic_ostream<wchar_t> <b>wostream</b>	<ostream>
typedef basic_ostringstream<wchar_t> <b>wstringstream</b>	<sstream>
typedef basic_streambuf<wchar_t> <b>wstreambuf</b>	<streambuf>
typedef fpos<char_traits<wchar_t>::state_type> <b>wstreampos</b>	N/A
typedef basic_stringbuf<wchar_t> <b>wstringbuf</b>	<sstream>
typedef basic_stringstream<wchar_t> <b>wstringstream</b>	<sstream>

## <iostream>

The <iostream> header declares the eight standard stream objects: cerr, cin, clog, cout, wcerr, wcin, wclog, and wcout. These objects are initialized when the first instance of ios\_base::Init is constructed (or earlier), or before the main program starts. They are not destroyed during normal program execution, so any static object's destructor or other function can use the standard I/O objects.

Each of the standard I/O objects is associated with a standard C FILE pointer (see <cstdio>). (You can sever the connection by calling ios\_base::sync\_with\_stdio(false), as described in <ios>.) You can use narrow or wide I/O objects, but once you have performed any I/O on an underlying C stream, you cannot switch from narrow to wide or wide to narrow. For example, after writing to cerr, you cannot write to wclog because both objects use the same C stream, stderr. See [Chapter 9](#) for more information about I/O.

Many C++ programmers assume that <iostream> automatically #includes <istream> and <ostream>, but the standard does not guarantee that behavior. Always #include every header you need, for example:

```
#include <cstdlib>
#include <iostream>
#include <istream>
#include <ostream>

// Copy standard input to standard output.
int main( )
{
    std::cout << std::cin.rdbuf( );
    return std::cout ? EXIT_SUCCESS : EXIT_FAILURE;
}
```

## Name

cerr variable — Error message stream

## Synopsis

```
extern ostream cerr
```

The cerr object is a standard output stream associated with the C stderr file. It is typically used for error messages. When the cerr object is initialized, it sets the unitbuf flag, causing cerr to flush the output buffer after every operation.

## See Also

[clog variable](#), [wcerr variable](#), ostream in [<ostream>](#)

## Name

cin variable — Normal input stream

## Synopsis

```
extern istream cin
```

The cin object is a standard input stream associated with the C stdin file. It is used for normal program input. It is tied to cout (see the tie member function of basic\_ios in <ios>).

## See Also

[wcin variable](#), basic\_ios in [<ios>](#), istream in [<istream>](#)

## Name

clog variable — Log output stream

## Synopsis

```
extern ostream clog
```

The `clog` object is a standard output stream associated with the C `stderr` file. Unlike `cerr`, the `clog` object does not set the `unitbuf` flag. This makes it more suitable for writing less critical messages that are not considered normal program output, such as debugging or logging messages.

## See Also

[cerr variable](#), [wclog variable](#), ostream in [<ostream>](#)

## Name

cout variable — Normal output stream

## Synopsis

```
extern ostream cout
```

The cout object is a standard output stream associated with the C stdout file. It is used for normal program output.

## See Also

[wcout variable](#), ostream in [<ostream>](#)

## Name

wcerr variable — Wide error message stream

## Synopsis

```
extern wostream wcerr
```

The wcerr object is a wide character standard output stream associated with the C stderr file. It is typically used for error messages. When the wcerr object is initialized, it sets the unitbuf flag, causing wcerr to flush the output buffer after every operation.

## See Also

[cerr variable](#), [wclog variable](#), wostream in [<ostream>](#)



## Name

wcin variable — Wide input stream

## Synopsis

```
extern wistream wcin
```

The wcin object is a wide character standard input stream associated with the C stdin file. It is used for wide program input. It is tied to wcout (see the tie( ) member function of basic\_ios in <ios>).

## See Also

[cin variable](#), basic\_ios in [<ios>](#), wistream in [<istream>](#)

## Name

wclog variable — Wide log output stream

## Synopsis

```
extern wostream wclog
```

The `wclog` object is a wide character standard output stream associated with the C `stderr` file. Unlike `wcerr`, the `wclog` object does not set the `unitbuf` flag. This makes it more suitable for writing less critical messages that are not considered normal program output, such as debugging or logging messages.

## See Also

[clog variable](#), [wcerr variable](#), `wostream` in [<ostream>](#)

## Name

wcout variable — Wide output stream

## Synopsis

```
extern wostream wcout
```

The wcout object is a wide character standard input stream associated with the C stdout file. It is used for wide program output.

## See Also

[cout variable](#), wostream in [<ostream>](#)

## <istream>

The <istream> header declares the input stream classes, templates, and an input manipulator.

### **Not e**

istream is declared in <istream>, and ostream is declared in <ostream>, but iostream is declared in <istream>, not <iostream>.

See <fstream> for derived-class templates that read from files and <sstream> for derived-class templates that read from strings. See <ios> for the base-class declarations. Refer to [Chapter 9](#) for general information about I/O.

## Name

`basic_iostream` class template — Base class for input and output stream

## Synopsis

```
template <typename charT, typename traits = char_traits<charT> >
class basic_iostream : public basic_istream<charT,traits>,
                      public basic_ostream<charT,traits>
{
public:
    explicit basic_iostream(basic_streambuf<charT,traits>* sb);
    virtual ~basic_iostream( );
};
```

The `basic_iostream` class template represents a stream that can perform input and output. (For details, see its base class templates, `basic_istream` in this section and `basic_ostream` in `<ostream>`.) Note that the two base-class templates share a common stream buffer, `sb`. Note also that `basic_istream` and `basic_ostream` inherit virtually from `basic_ios`; `basic_iostream` inherits a single instance of the formatting flags, `iostate`, and so on.

## See Also

[basic\\_istream class template](#), [iostream class](#), [wiostream class](#), `basic_iostream` in [<ios>](#), `basic_ostream` in [<ostream>](#)

# Name

basic\_istream class template — Base class for input stream

## Synopsis

```
template <typename chT, typename traits = char_traits<chT> >
class basic_istream : virtual public basic_ios<chT,traits>
{
public:
    // Types
    typedef chT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    explicit basic_istream(basic_streambuf<chT,traits>* sb);
    virtual ~basic_istream( );
    class sentry;
    // Formatted input
    basic_istream<chT,traits>& operator>>
(basic_istream<chT,traits>&
    (*pf)(basic_istream<chT,traits>&));
    basic_istream<chT,traits>& operator>> (basic_ios<chT,traits>& |
    (*pf)(basic_ios<chT,traits>&));
    basic_istream<chT,traits>& operator>>(ios_base& (*pf)
(ios_base&));
    basic_istream<chT,traits>& operator>>(bool& n);
    basic_istream<chT,traits>& operator>>(short& n);
    basic_istream<chT,traits>& operator>>(unsigned short& n);
    basic_istream<chT,traits>& operator>>(int& n);
    basic_istream<chT,traits>& operator>>(unsigned int& n);
    basic_istream<chT,traits>& operator>>(long& n);
    basic_istream<chT,traits>& operator>>(unsigned long& n);
    basic_istream<chT,traits>& operator>>(float& f);
    basic_istream<chT,traits>& operator>>(double& f);
    basic_istream<chT,traits>& operator>>(long double& f);
    basic_istream<chT,traits>& operator>>(void*& p);
    basic_istream<chT,traits>& operator>>
(basic_streambuf<char_type,traits>* sb);
    // Unformatted input
    streamsize gcount( ) const;
    int_type get( );
    basic_istream<chT,traits>& get(char_type& c);
```

```

    basic_istream<chT,traits>& get(char_type* s, streamsize n);
    basic_istream<chT,traits>& get(char_type* s, streamsize n,
char_type delim);
    basic_istream<chT,traits>&
get(basic_streambuf<char_type,traits>& sb);
    basic_istream<chT,traits>&
get(basic_streambuf<char_type,traits>& sb,
        char_type delim);
    basic_istream<chT,traits>& getline(char_type* s, streamsize n);
    basic_istream<chT,traits>& getline(char_type* s, streamsize n,
        char_type delim);
    basic_istream<chT,traits>& ignore (streamsize n = 1, int_type
        delim=traits::eof( ));

    int_type peek( );
    basic_istream<chT,traits>& read(char_type* s, streamsize n);
    streamsize readsome(char_type* s, streamsize n);
    basic_istream<chT,traits>& putback(char_type c);
    basic_istream<chT,traits>& unget( );
    int sync( );
    pos_type tellg( );
    basic_istream<chT,traits>& seekg(pos_type);
    basic_istream<chT,traits>& seekg(off_type, ios_base::seekdir);
};

```

The `basic_istream` class template is the base for all input streams. It declares members for reading from streams and for managing streams. The act of reading from a stream is also known as *extracting* from the stream.

All reads go through a stream buffer, which provides the low-level access to the stream data. (See the `sbumpc` and `sgetc` functions for `basic_streambuf` in the `<streambuf>` header.) If the stream buffer returns `traits::eof( )` for an input operation, the stream sets `failbit | eofbit` in the stream's I/O state. If the buffer object throws an exception, the stream sets `badbit`.

Before performing an input operation (e.g., `operator>>` or a `get( )` function), the stream constructs a sentry object. If the sentry evaluates to `true`, the read operation continues. If the read throws an exception, `badbit` is set. The second (`noskipws`) argument to the sentry constructor is `false` for the formatted functions (`operator>>`) and `true`

for all other input operations. The sentry object is destroyed before the input function returns. See `basic_istream::sentry` later in this section for more information.

When an input operation throws an exception, the stream sets `badbit`. If `badbit` is set in the `exceptions( )` mask, the stream does not throw `ios_base::failure`, but instead rethrows the original exception. If any input operation sets `eofbit` or `failbit`, and that bit is set in the `exceptions( )` mask, the usual `ios_base::failure` exception is thrown.

The following are the `basic_istream` member functions:

explicit **`basic_istream`**

`(basic_streambuf<chT,traits>* sb)`

Calls the `basic_ios(sb)` constructor and initializes `gcount( )` to 0.

**`basic_istream<chT,traits>& operator>>`**

`(basic_istream<chT,traits>& (*manip)`  
`(basic_istream<chT,traits>&))`

Returns `manip(*this)`. See the `ws` function for an example of such a manipulator.

**`basic_istream<chT,traits>& operator>>`**

`(basic_ios<chT,traits>&(*manip)`  
`(basic_ios<chT,traits>&)),`  
**`basic_istream<chT,traits>& operator>>` `(ios_base&`  
`(*manip)(ios_base&))`**

Calls `manip(*this)` and returns `*this`. See the `dec` function in `<ios>` for an example of such a manipulator.

**`basic_istream<chT,traits>& operator>>`** `(bool& n),`  
**`basic_istream<chT,traits>& operator>>`** `(short& n),`  
**`basic_istream<chT,traits>& operator>>`** `(unsigned`  
`short& n),` **`basic_istream<chT,traits>& operator>>`**



```
(int& n), basic_istream<chT,traits>& operator>>
(unsigned int& n), basic_istream<chT,traits>&
operator>> (long& n), basic_istream<chT,traits>&
operator>> (unsigned long& n),
basic_istream<chT,traits>& operator>> (float& f),
basic_istream<chT,traits>& operator>> (double& f),
basic_istream<chT,traits>& operator>> (long
double& f), basic_istream<chT,traits>& operator>>
(void*& p)
```

Reads a formatted item from the stream. These formatted input functions use the `num_get` facet of the stream's imbued locale as shown in [Example 13-16](#). See the `<locale>` header for information about `num_get` and locales.

*Example 13-16. Using the `num_get` facet to format input*

```
typedef
    std::num_get<char_type,
        std::istreambuf_iterator<char_type, traits_type> >
    numget;
std::ios_base::iostate err = 0;
std::use_facet<numget>(getloc( )).(*this, 0, *this, err,
val);
this->setstate(err);
```

```
basic_istream<chT,traits>& operator>>
(basic_streambuf<char_type,traits>* sb)
```

Copies input to the stream buffer `sb`. If `sb` is a null pointer, the stream sets `failbit` and returns immediately. Otherwise, the function copies characters from the input stream to `sb` until it reaches the end of the input stream, writing to `sb` fails, or an exception is caught. If no characters are copied to `sb`, `failbit` is set. The return value is `*this`.

```
streamsize gcount ( ) const
```

Returns the number of characters extracted by the most recent call to an unformatted member function (`get`,

getline, ignore, peek, putback, read, readsome, and unget).

int\_type **get** ( )

Reads a single character and returns it. If no character is available, the function sets failbit and returns traits::eof( ).

basic\_istream<chT,traits>& **get** (char\_type& c)

Reads a single character and stores it in c. If no character is available, the function sets failbit and does not modify c. The return value is \*this.

basic\_istream<chT,traits>& **get** (char\_type\* s, streamsize n)

Returns get(s, n, widen('\n')).

basic\_istream<chT,traits>& **get** (char\_type\* s, streamsize n, char\_type delim)

Reads up to n - 1 characters into the character array that s points to. Reading stops at the end-of-file (which sets eofbit) or when the next character would be delim (but the delimiter character is not read from the buffer). If no characters are stored in s, failbit is set.

A null character is always appended to the end of s. The return value is \*this.

basic\_istream<chT,traits>& **get**  
(basic\_streambuf<char\_type,traits>& sb)

Returns get(sb, widen('\n')).

basic\_istream<chT,traits>& **get**  
(basic\_streambuf<char\_type,traits>& sb, char\_type delim)

Copies characters to the output buffer sb. Copying stops when an end-of-file is reached, when writing to sb fails,

or when the next character to read would be `delim` (the delimiter is not read from the input buffer). If no characters are copied to `sb`, `failbit` is set.

```
basic_istream<chT,traits>& getline (char_type* s,  
streamsize n)
```

Returns `getline(s, n, widen('\n'))`.

```
basic_istream<chT,traits>& getline (char_type* s,  
streamsize n, char_type delim)
```

Reads up to `n - 1` characters into the character array that `s` points to. Reading stops at the end-of-file (which sets `eofbit`) or when `delim` is read (the delimiter character is read from the buffer but not stored in `s`).

If no characters are read from the stream, `failbit` is set. If exactly `n - 1` characters are read before reaching the end-of-file, `eofbit` is set; otherwise, if the limit of `n - 1` characters is reached before reading a delimiter, `failbit` is set.

A null character is always appended to the end of `s`. The return value is `*this`.

```
basic_istream<chT,traits>& ignore (streamsize n=1,  
int_type delim=traits::eof( ))
```

Reads characters from the input buffer and discards them. The operation stops when one of the following occurs:

- `n` characters have been read and discarded if `n != numeric_limits<streamsize>::max`.
- End-of-file is reached, in which case `eofbit` is set.
- `delim` is read if `delim != traits::eof( )`.

The return value is `*this`.

```
int_type peek ( )
```

Looks ahead to the next character in the input stream. If `good( )` returns true, `rdbuf( )->sgetc( )` is returned; otherwise, `traits::eof( )` is returned. No characters are extracted from the stream, so a subsequent call to `gcount( )` returns 0.

`basic_istream<chT,traits>& read (char_type* s, streamsize n)`

Reads up to `n` characters and stores them in the array that `s` points to. Before reading anything, if `good( )` is false, `failbit` is set, and nothing is read. Otherwise, if end-of-file is reached before reading `n` characters, `eofbit` and `failbit` are set. A null character is not appended to `s`. The return value is `*this`.

`streamsize readsome (char_type* s, streamsize n)`

Tries to read as many immediately available characters as possible into the array that `s` points to. If `good( )` is false before reading, `failbit` is set, and `readsome` returns immediately. Otherwise, `readsome` calls `rdbuf( )->in_avail( )` and does one of the following:

- `in_avail( ) == -1`: `eofbit` is set and no characters are read
- `in_avail( ) == 0`: no characters are read
- `in_avail( ) > 0`: `min(in_avail( ), n)` characters are read into `s`

A null character is not appended to `s`. The return value is the number of characters stored in `s`.

`basic_istream<chT,traits>& putback (char_type c)`

Tries to push back the character `c` so it will be the next character read from the input stream. If `good( )` is false, `failbit` is set and `putback` returns. Otherwise, if `rdbuf( )` is not null, `putback` calls `rdbuf( )-`

>sputbackc(c). If rdbuf( ) is null or sputbackc returns traits::eof( ), badbit is set. The return value is \*this. A subsequent call to gcount( ) returns 0.

```
basic_istream<chT,traits>& seekg (pos_type pos),  
basic_istream<chT,traits>& seekg (off_type off,  
ios_base::seekdir dir)
```

Tries to seek to a new position in the stream. The first form specifies the position explicitly; the second form specifies the new position as an offset from a known position (beginning of file, current position, or end-of-file). If fail( ) is false, seekg calls rdbuf( )->pubseekoff(pos) or rdbuf( )->pubseekoff(off, dir). The return value is \*this.

```
int sync ( )
```

Synchronizes the stream buffer. If rdbuf( ) is null, sync returns -1; otherwise, sync calls rdbuf( )->pubsync( ). If pubsync returns -1, badbit is set, and sync returns -1. Otherwise, sync returns 0.

```
pos_type tellg ( )
```

Returns the current position in the stream. If fail( ) is true, the return value is pos\_type(-1); otherwise, the return value is rdbuf( )->pubseekoff(0, ios\_base::cur, ios\_base::in).

```
basic_istream<chT,traits>& unget ( )
```

Tries to push back the last input character. If good( ) is false, failbit is set and unget returns. Otherwise, if rdbuf( ) is not null, putback calls rdbuf( )->sungetc( ). If rdbuf( ) is null or sungetc returns traits::eof( ), badbit is set. The return value is \*this. A subsequent call to gcount( ) returns 0.

## See Also

[istream class](#), [wistream class](#)

# Name

`basic_istream::sentry` class — Sentry class for input streams

# Synopsis

```
template <typename chT, typename traits=char_traits<chT> >
class basic_istream<chT,traits>::sentry{
    public: explicit sentry(basic_istream<chT,traits>& stream,
                           bool noskipws = false);

    ~sentry( );
    operator bool( ) const;
private:
    sentry(const sentry&); // Not defined
    sentry& operator=(const sentry&); // Not defined
};
```

A `basic_istream` object constructs a temporary sentry object prior to each input operation. The sentry object is destroyed when the input operation finishes and the function returns. The sentry manages tied streams and is responsible for skipping whitespace prior to a formatted read.

The stream passes itself and a flag to the sentry's constructor. Formatted reads (`operator>>`) use `false` for the second argument; unformatted reads (`get`, `getline`, etc.) use `true`.

If `stream.good( )` is true, the sentry first flushes any tied stream. That is, if `stream.tie( )` is not null, the sentry calls `stream.tie( )->flush( )`. This ensures that prompts and similar output appears before the input is requested. Then, if the `noskipws` argument is false, and if the `skipws` bit is not set in the stream's formatting flags (`((ios_base::skipws & stream.formatflags( )) == 0)`), the sentry's constructor reads and discards whitespace

characters. The sentry uses code similar to that shown in [Example 13-17](#).

## Example

*Example 13-17. Skipping whitespace in an input sentry*

```
const std::ctype<char_type>& ctype =
    std::use_facet<ctype<char_type> >(stream.getloc( ));
int_type c;
while ((c = stream.rdbuf( )->snextc( )) != traits::eof( ))
    if (ctype.is(ctype.space,c) == 0) {
        // Put back the non-space character.
        stream.rdbuf( )->sputbackc(c);
        break;
    }
```

If anything goes wrong, the sentry calls `stream.setstate(failbit)`. The sentry's operator `bool( )` returns true if `stream.good( )` is true, and false otherwise.

## See Also

[basic\\_istream class template](#), `basic_ostream::sentry` in [<ostream>](#)



## Name

`iostream` class — Narrow input and output stream

## Synopsis

```
typedef basic_istream<char> istream;
```

The `istream` class specializes `basic_istream` for the `char` type.

## See Also

[basic\\_istream class template](#), [istream class](#), [wistream class](#), `ostream` in [<ostream>](#)

## Name

istream class — Input stream

## Synopsis

```
typedef basic_istream<char> istream;
```

The `istream` class specializes `basic_istream` for the `char` type.

## See Also

[basic\\_istream class template](#), [iostream class](#), [wistream class](#), `ostream` in [<ostream>](#)

## Name

`operator>>` function template — Input operator for single characters

## Synopsis

```
template<typename charT, typename traits>
    basic_istream<charT,traits>& operator>>
        (basic_istream<charT,traits>& stream, charT& c);

template<typename traits> basic_istream<char,traits>&
    operator>>(basic_istream<char,traits>& stream, unsigned char&
c);
template<typename traits> basic_istream<char,traits>&
    operator>>(basic_istream<char,traits>& stream, signed char& c);
```

The `operator>>` function reads a character from an input stream using the rules for a formatted read (that is, a sentry object is constructed and initial whitespace is skipped). The character is stored in `c`. If no character is available, `failbit` is set. The return value is `stream`.

Note that the first form of `operator>>` reads the stream's character type. If the stream's character type is `char`, the second and third forms read `signed char` and `unsigned char`.

## See Also

[basic\\_istream class template](#)

## Name

`operator>>` function template — Input operator for character arrays

## Synopsis

```
template<typename charT, typename traits>
    basic_istream<charT,traits>& operator>>
        (basic_istream<charT,traits>& stream, charT* str);

template<typename traits> basic_istream<char,traits>&
    operator>>(basic_istream<char,traits>& stream, unsigned char*
str);
template<typename traits> basic_istream<char,traits>&
    operator>>(basic_istream<char,traits>& stream, signed char*
str);
```

The `operator>>` function reads characters from `stream` into `str`. As with any formatted input, a sentry object is created and initial whitespace is skipped. Characters are then read from the stream into `str` until an end-of-file is reached or the next character read would be a whitespace character (the space is not read from the stream).

To limit the number of characters read (always a good idea), set the stream's width to  $n$ . At most,  $n - 1$  characters will be read into `s`. The function resets the width to 0 after reading the stream. A null character is always appended to the end of the string.

If no characters are read from the stream, `failbit` is set.

Note that the first form of `operator>>` reads with the stream's character type. If the stream's character type is `char`, the second and third forms read signed `char` and unsigned `char`.

## See Also

[basic\\_istream class template](#)

## Name

wiostream class — Wide input and output stream

## Synopsis

```
typedef basic_istream<wchar_t> wiostream;
```

The `wiostream` class specializes `basic_istream` for the `wchar_t` type.

## See Also

[basic\\_istream class template](#), [istream class](#), [wistream class](#), `wostream` in [<ostream>](#)

## Name

wistream class — Wide input stream

## Synopsis

```
typedef basic_istream<wchar_t> wistream;
```

The wistream class specializes basic\_istream for the wchar\_t type.

## See Also

[basic\\_istream class template](#), [istream class](#), [wostream class](#), wostream in [<ostream>](#)

## Name

ws function — Skips whitespace manipulator

## Synopsis

```
template <class charT, class traits>
    basic_istream<charT,traits>& ws(basic_istream<charT,traits>&
    stream);
```

The `ws` function is an input stream manipulator that skips whitespace characters in `stream` using the same technique as `basic_istream::sentry`. If the manipulator reaches the end-of-file, it sets `eofbit`, but not `failbit`.

## Example

Suppose you want to read a line of text by calling `getline`, but you also want to skip whitespace at the beginning of the line. (Because `getline` is an unformatted input function, it does not automatically skip whitespace.) The following example shows one way to skip whitespace and read the line using `ws`:

```
char buffer[BUFSIZ];
. . .
in >> ws;
in.getline(buffer, sizeof(buffer));
```

## See Also

[basic\\_istream class template](#), [basic\\_istream::sentry class](#)



## <iterator>

The <iterator> header declares classes and templates for defining and using iterators. (See [Chapter 10](#).) Iterators are especially important when using the standard algorithms in <algorithm>.

An iterator gives a program access to the contents of a container or other sequence, such as an I/O stream. You can think of an iterator as an abstraction of a pointer; the syntax for using iterators resembles that of pointers. Conceptually, an iterator points to a single element in a container or sequence and can be advanced to the next element with the ++ (increment) operator. The unary \* (dereference) operator returns the element that the iterator points to. Iterators (except for output iterators) can be compared: two iterators are equal if they point to the same position in the same sequence, or if they both point to one position past the end of the same sequence.

There are five categories of iterators:

### *Input iterators*

Permit one pass to read a sequence. The increment operator advances to the next element, but there is no decrement operator. The dereference operator does not return an lvalue, so you can read elements but not modify them.

### *Output iterators*

Permit one pass to write a sequence. The increment operator advances to the next element, but there is no decrement operator. You can dereference an element only to assign a value to it. You cannot compare output iterators.

### *Forward iterators*

Are like a combination of an input and an output iterator. You can use a forward iterator anywhere an input iterator is required or where an output iterator is required. A forward iterator, as its name implies, permits unidirectional access to a sequence. You can refer to a single element and modify it multiple times before advancing the iterator.

### *Bidirectional iterators*

Are like forward iterators but also support the decrement operator ( - - ) to move the iterator backward by one position.

### *Random access iterators*

Are like bidirectional iterators but also support the [ ] (subscript) operator to access any index in the sequence. Also, you can add or subtract an integer to move a random access iterator by more than one position at a time. Subtracting two random access iterators yields an integer distance between them. Thus, a random access iterator is most like a conventional pointer, and, in fact, a pointer can be used as a random access iterator.

An input, forward, bidirectional, or random access iterator can be a `const_iterator`. Dereferencing a `const_iterator` yields a constant (an rvalue or a const lvalue). [Chapter 10](#) discusses `const_iterator`s in more depth.

An iterator can point to any element in a sequence or to one position past the end. You cannot dereference an iterator that points to one past the end, but you can compare it with other iterators (provided it is not an output iterator) or, if it is a bidirectional or random access iterator,

decrease its position so it points to a valid element of the sequence.

Iterators are often used in *ranges*. A range has two iterators: a starting point and an ending point. The end iterator typically points to one position past the last element in the range. Thus, a range is often written using the mathematical notation of `[first, last)`, in which the square bracket means `first` is included in the range, and the parenthesis means `last` is excluded from the range.

Like ordinary pointers, iterators can be uninitialized or otherwise have invalid values, such as `x.begin() - 1`, in which `x` is any container. It is your responsibility to ensure that you dereference only valid iterators, that you don't let iterators run past their valid end-points, and so on.

To create your own iterator class, see the `iterator class template`.

## Name

advance function template — Moves iterator forward or backward

## Synopsis

```
template <typename InputIterator, typename Distance>  
    void advance(InputIterator& i, Distance n);
```

The advance function template advances an input iterator *i* by a distance *n*. If the iterator is bidirectional or random access, *n* can be negative. If the iterator is random access, the advance function is specialized as *i* + *n*; other iterators apply the ++ operator *n* times (or -- for a bidirectional iterator when *n* is negative).

## See Also

[distance function template](#)

## Name

`back_insert_iterator` class template — Output iterator to push items back onto a container

## Synopsis

```
template <typename Container>
    class back_insert_iterator :
        public iterator<output_iterator_tag,void,void,void,void>
    {
protected:
    Container* container;
public:
    typedef Container container_type;
    explicit back_insert_iterator(Container& x);
    back_insert_iterator<Container>&
        operator=(typename Container::const_reference value);
    back_insert_iterator<Container>& operator*( );
    back_insert_iterator<Container>& operator++( );
    back_insert_iterator<Container> operator++(int);
    };
```

The `back_insert_iterator` class template implements an output iterator that stores elements in a container by calling the container's `push_back` function. The most convenient way to create a `back_insert_iterator` object is to use the `back_inserter` function template.

The way `back_insert_iterator` works seems slightly unconventional, although it is perfectly reasonable for an output iterator: the `*` operator returns the iterator, not an element of the container. Thus, the expression `*iter = value` assigns `value` to the iterator itself. The iterator's assignment operator appends `value` to the underlying container by calling the container's `push_back` function. Thus, the iterator does not maintain any notion of a position, and the increment operator is a no-op.

The following are the member functions of `back_insert_iterator`:

explicit **back\_insert\_iterator** (Container& x)

Initializes the container member with &x.

**back\_insert\_iterator**<Container>& **operator=**  
(typename Container::const\_reference value)

Calls container->push\_back(value). The return value is \*this.

**back\_insert\_iterator**<Container>& **operator\*** ( )

Returns \*this.

**back\_insert\_iterator**<Container>& **operator++** ( ),  
**back\_insert\_iterator**<Container>& **operator++** (int)

Returns \*this with no side effects.

## See Also

[back\\_inserter function template](#), [front\\_insert\\_iterator class template](#), [insert\\_iterator class template](#)

## Name

`back_inserter` function template — Creates a `back_insert_iterator`

## Synopsis

```
template <typename Container>
    back_insert_iterator<Container> back_inserter(Container& x);
```

The `back_inserter` function template constructs a `back_insert_iterator` object for the container `x`. [Example 13-18](#) shows how to use `back_inserter` to read integers from a file into a vector.

## Example

*Example 13-18. Using `back_inserter` to add numbers to a vector*

```
std::ifstream in("experiment.dat");
std::vector<int> data;
std::copy(std::istream_iterator<int>(in),
          std::istream_iterator<int>( ),
          std::back_inserter(data));
```

## See Also

[back\\_insert\\_iterator class template](#), [front\\_inserter function template](#), [inserter function template](#)

## Name

`bidirectional_iterator_tag` class — Tag for a bidirectional iterator

## Synopsis

```
struct bidirectional_iterator_tag :  
    public forward_iterator_tag {};
```

Use the `bidirectional_iterator_tag` class as the iterator category when declaring a new bidirectional iterator class. When writing a generic algorithm or similar function, you can use the iterator's category to write specialized implementations for different kinds of iterators. See [Example 13-19](#) (under the `distance` function template).

## See Also

[bidirectional\\_iterator\\_tag class](#), [forward\\_iterator\\_tag class](#), [input\\_iterator\\_tag class](#), [iterator class template](#), [output\\_iterator\\_tag class](#), [random\\_access\\_iterator\\_tag class](#)



## Name

distance function template — Counts elements between two iterators

## Synopsis

```
template<typename InputIterator>
    typename iterator_traits<InputIterator>::difference_type
    distance(InputIterator first, InputIterator last);
```

The distance function returns the number of elements between `first` and `last`. The function is specialized for random access iterators to use the `-` operator; for other input iterators, the function applies the `++` operator to `first` until `first == last`. The behavior is undefined if `first` and `last` refer to different containers or if `last` points to an element earlier than `first`.

[Example 13-19](#) shows a simple implementation of the distance function. The first specialized implementation works for any iterator (except output iterators, which do not support comparison with the `!=` operator). The second one works only with random access iterators and uses the subtraction operator to compute the distance in constant time. The compiler picks the more specialized function when it can, so random access iterators can compute distances in constant time, compared to linear time for other iterators.

## Example

*Example 13-19. A simple implementation of distance*

```
namespace std {
    template<typename InputIter>
        typename iterator_traits<InputIter>::difference_type
        specialize_distance(InputIter first, InputIter last, ...)
        {
            typename iterator_traits<InputIter>::difference_type n;
```

```

    for (n = 0; first != last; ++first)
        ++n;
    return n;
}

template<typename InputIter>
typename iterator_traits<InputIter>::difference_type
specialize_distance(InputIter first, InputIter last,
                    random_access_iterator_tag)
{
    return last - first;
}

template<typename InputIter>
typename
iterator_traits<InputIter>::difference_type distance(InputIter
first, InputIter last)
{
    return specialize_distance(first, last,
        iterator_traits<InputIter>::iterator_category( ));
}
}

```

## See Also

[advance function template](#)

## Name

`forward_iterator_tag` class — Tag for a forward iterator

## Synopsis

```
struct forward_iterator_tag : public input_iterator_tag {};
```

Use the `forward_iterator_tag` class as the iterator category when declaring a new forward iterator class. When writing a generic algorithm or similar function, you can use the iterator's category to write specialized implementations for different kinds of iterators. See [Example 13-19](#) (under the distance function template).

## See Also

[bidirectional\\_iterator\\_tag class](#), [forward\\_iterator\\_tag class](#), [input\\_iterator\\_tag class](#), [iterator class template](#), [output\\_iterator\\_tag class](#), [random\\_access\\_iterator\\_tag class](#)

## Name

`front_insert_iterator` class template — Iterator to insert elements at the front of a container

## Synopsis

```
template <typename Container>
class front_insert_iterator :
    public iterator<output_iterator_tag,void,void,void,void>
{
protected:
    Container* container;
public:
    typedef Container container_type;
    explicit front_insert_iterator(Container& x);
    front_insert_iterator<Container>&
        operator=(typename Container::const_reference value);
    front_insert_iterator<Container>& operator*( );
    front_insert_iterator<Container>& operator++( );
    front_insert_iterator<Container> operator++(int);
};
```

The `front_insert_iterator` class template implements an output iterator that stores elements in a container by calling the container's `push_front` function. The most convenient way to create a `front_insert_iterator` object is to use the `front_inserter` function template.

The way `front_insert_iterator` works seems slightly unconventional, although it is perfectly reasonable for an output iterator: the `*` operator returns the iterator, not an element of the container. Thus, the expression `*iter = value` assigns `value` to the iterator itself. The iterator's assignment operator adds `value` to the underlying container by calling the container's `push_front` function. Thus, the iterator does not maintain any notion of a position, and the increment operator is a no-op.

The following are the member functions of `front_insert_iterator`:

explicit **front\_insert\_iterator** (Container& x)

Initializes the container member with &x.

**front\_insert\_iterator**<Container>& **operator=**  
(typename Container::const\_reference value)

Calls container->push\_front(value). The return value is \*this.

**front\_insert\_iterator**<Container>& **operator\*** ( )

Returns \*this.

**front\_insert\_iterator**<Container>& **operator++** ( ),  
**front\_insert\_iterator**<Container> **operator++** (int)

Returns \*this with no side effects.

## See Also

[back\\_insert\\_iterator class template](#), [front\\_inserter function template](#), [insert\\_iterator class template](#)

## Name

`front_inserter` function template — Creates a `front_insert_iterator`

## Synopsis

```
template <typename Container>
    front_insert_iterator<Container>
        front_inserter(Container& x);
```

The `front_inserter` function template constructs a `front_insert_iterator` object for the container `x`. [Example 13-20](#) shows how to use `front_inserter` to read integers from a file into a list in reverse order.

## Example

*Example 13-20. Using a `front_inserter` to add numbers to a list*

```
std::ifstream in("experiment.dat");
std::list<int> data;
std::copy(std::istream_iterator<int>(in),
          std::istream_iterator<int>( ),
          std::front_inserter(data));
```

## See Also

[back\\_inserter function template](#), [front\\_insert\\_iterator class template](#), [inserter function template](#)

## Name

`input_iterator_tag` class — Tag for an input iterator

## Synopsis

```
struct input_iterator_tag {};
```

Use the `input_iterator_tag` class as the iterator category when declaring a new input iterator class. When writing a generic algorithm or similar function, you can use the iterator's category to write specialized implementations for different kinds of iterators. See [Example 13-19](#) (under the distance function template).

## See Also

[bidirectional\\_iterator\\_tag class](#), [forward\\_iterator\\_tag class](#), [iterator class template](#), [output\\_iterator\\_tag class](#), [random\\_access\\_iterator\\_tag class](#)

# Name

`insert_iterator` class template — Iterator to insert elements in a container

# Synopsis

```
class insert_iterator :
    public iterator<output_iterator_tag,void,void,void,void>
{
protected:
    Container* container;
    typename Container::iterator iter;
public:
    typedef Container container_type;
    insert_iterator(Container& cont, typename Container::iterator
iter);
    insert_iterator<Container>&
        operator=(typename Container::const_reference value);
    insert_iterator<Container>& operator*( );
    insert_iterator<Container>& operator++( );
    insert_iterator<Container>& operator++(int);
};
```

The `insert_iterator` class template implements an output iterator that stores elements in a container by calling the container's `insert` function. The most convenient way to create a `insert_iterator` object is to use the `inserter` function template.

The way `insert_iterator` works seems slightly unconventional, although it is perfectly reasonable for an output iterator: the `*` operator returns the iterator, not an element of the container. Thus, the expression `*iter = value` assigns `value` to the iterator itself. The iterator's assignment operator adds `value` to the underlying container by calling the container's `insert` function. Thus, the iterator does not maintain any notion of a position, and the increment operator is a no-op.



The following are the member functions of `insert_iterator`:

**`insert_iterator`** (Container& x, typename Container::iterator i)

Initializes the container member with &x and iter with i. Thus, the elements to be inserted in the container will be inserted at position i.

`insert_iterator`<Container>& **operator=** (typename Container::const\_reference value)

Assigns a value to an element in the iterator's container, performing the equivalent of the following:

```
iter = container->insert(iter, value);  
++iter;  
return *this;
```

`insert_iterator`<Container>& **operator\*** ( )

Returns \*this.

`insert_iterator`<Container>& **operator++** ( ),  
`insert_iterator`<Container> **operator++** (int)

Returns \*this with no side effects.

## See Also

[back\\_insert\\_iterator class template](#), [front\\_insert\\_iterator class template](#), [inserter function template](#)

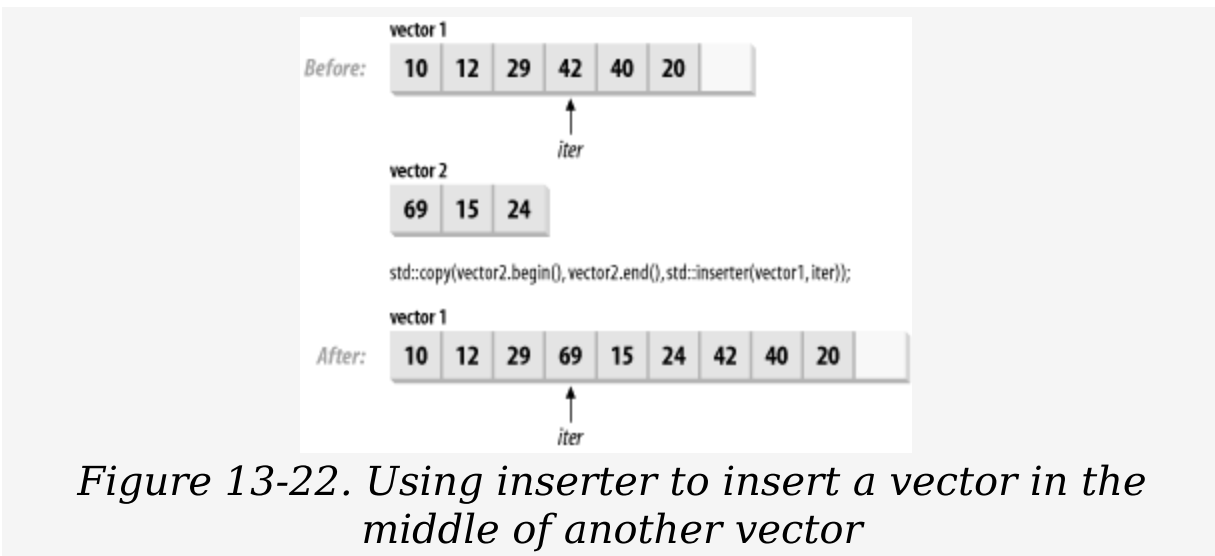
# Name

inserter function template — Creates an insert\_iterator

# Synopsis

```
template <typename Container, typename Iterator>
    insert_iterator<Container>
        inserter(Container& x, Iterator i);
```

The inserter function template constructs an insert\_iterator object for the container x to insert items starting at position i. [Figure 13-22](#) illustrates a simple use of the inserter function.



*Figure 13-22. Using inserter to insert a vector in the middle of another vector*

# See Also

[back\\_inserter function template](#), [front\\_inserter function template](#), [insert\\_iterator class template](#)

# Name

`istream_iterator` class template — Input iterator to read items from an `istream`

# Synopsis

```
template <typename T, typename charT = char,
          typename traits = char_traits<charT>, typename Distance
          = ptrdiff_t>
class istream_iterator :
    public iterator<input_iterator_tag,
                  T, Distance, const T*, const T>
{
public:
    typedef charT char_type;
    typedef traits traits_type;
    typedef basic_istream<charT,traits> istream_type;
    istream_iterator( );
    istream_iterator(istream_type& stream);
    istream_iterator(const
istream_iterator<T,charT,traits,Distance>& x);
    ~istream_iterator( );
    const T& operator*( ) const;
    const T* operator->( ) const;
    istream_iterator<T,charT,traits,Distance>& operator++( );
    istream_iterator<T,charT,traits,Distance> operator++(int);
};
```

The `istream_iterator` class template wraps an input iterator around an input stream (an instance of `basic_istream`), making the stream appear to be a sequence of items, each of type `T`.

[Example 13-20](#) (under the `front_inserter` function template) shows how an `istream_iterator` can be used to read a series of integers from a file.

The following are the member functions of `istream_iterator`:

**istream\_iterator** ( )

Constructs an `istream_iterator` that denotes the end of the stream. End-of-stream iterators are equal to each other and are not equal to any other `istream_iterator`.

**`istream_iterator`** (`istream_type& stream`)

Constructs an `istream_iterator` to read from a stream. The constructor might read the first item from the stream. An `istream_iterator` that wraps a stream is equal to an end-of-stream iterator when `stream.eof( )` is true.

**`istream_iterator`** (`const istream_iterator<T, charT, traits, Distance>& iter`)

Constructs a copy of `iter`. Note that two `istream_iterator` objects are the same (`operator==` is true) if they point to the same stream object.

`const T& operator* ( ) const`

Returns the item that was read most recently from the stream.

**`operator->`**

Returns a pointer to the item that was read most recently from the stream.

`istream_iterator<T, charT, traits, Distance>& operator++ ( )`

Reads the next item from the stream using `operator>>`. The return value is a copy of `*this`, made prior to reading from the stream.

## See Also

[istreambuf\\_iterator class template](#), [ostream\\_iterator class template](#), `basic_istream` in [<istream>](#)

## Name

`istreambuf_iterator` class template — Input iterator to read characters from a `streambuf`

## Synopsis

```
template<typename charT, typename traits=char_traits<charT> >
class istreambuf_iterator :
    public iterator<input_iterator_tag, charT, typename
traits::off_type, charT*,
                charT&>
{
public:
    typedef charT char_type;
    typedef traits traits_type;
    typedef typename traits::int_type int_type;
    typedef basic_streambuf<charT,traits> streambuf_type;
    typedef basic_istream<charT,traits> istream_type;
    class proxy; // Exposition only

    istreambuf_iterator( ) throw( );
    istreambuf_iterator(istream_type& s) throw( );
    istreambuf_iterator(streambuf_type* s) throw( );
    istreambuf_iterator(const proxy& p) throw( );
    charT operator*( ) const;
    istreambuf_iterator<charT,traits>& operator++( );
    proxy
        operator++(int);
    bool equal(istreambuf_iterator& b) const;
};
```

The `istreambuf_iterator` class template wraps a stream buffer object (instance of `basic_streambuf`) as an input iterator to read characters from the stream buffer. [Example 13-21](#) shows how to use `streambuf_iterator`s to copy files.

## Examples

*Example 13-21. Copying files using streambuf iterators*

```

void copyfile(const char* from, const char* to)
{
    std::ifstream in(from);
    std::ofstream out(to);

    std::copy(std::istreambuf_iterator<char>(in),
              std::istreambuf_iterator<char>( ),
              std::ostreambuf_iterator<char>(out));
}

```

The post-increment operator (++) returns a *proxy object*, which is an object that stands in for the `istreambuf_iterator` object. Its use is largely transparent, and you rarely need to think about it. The definition and name of the proxy class are implementation-defined, but the class has at least the capability to return the input character and the underlying stream buffer. This section assumes that the class name is *proxy*. [Example 13-22](#) shows a prototypical implementation of *proxy*.

*Example 13-22. A trivial implementation of the proxy class*

```

template<typename charT, typename traits=char_traits<charT> >
class istreambuf_iterator<charT, traits>::proxy
{
    friend template<typename charT, typename traits>
        class istreambuf_iterator<charT,traits>;
    charT keep;
    basic_streambuf<charT,traits>* sbuf;
    proxy(charT c, basic_streambuf<charT,traits>* sbuf);
        : keep(c), sbuf(sbuf) {}
public:
    charT operator*( ) { return keep; }
};

```

In the following descriptions of the member functions of `istreambuf_iterator`, the data member *sbuf* is a pointer to the iterator's stream buffer. The *sbuf* member serves only to keep the function descriptions clear and simple; the class is not required to have such a member, nor is the class required to have a member with that name.

**istreambuf\_iterator** ( ) throw( )

Constructs the end-of-stream iterator.

**istreambuf\_iterator** (istream\_type& s) throw( ),  
**istreambuf\_iterator** (streambuf\_type\* sb) throw( ),  
**istreambuf\_iterator** (const proxy & p) throw( )

Constructs an `istreambuf_iterator` and initializes `sbuf` to `s.rdbuf( )`, `sb`, or `p.sbuf`. If `sb == 0`, an end-of-stream iterator is constructed.

`charT operator*` ( ) const

Returns `sbuf ->sgetc( )`.

`istreambuf_iterator<charT,traits>& operator++` ( )

Calls `sbuf ->sbumpc( )` and returns `*this`.

`proxy operator++` (int)

Returns `proxy ( sbuf ->sbumpc( ), sbuf)`.

`bool equal` (istreambuf\_iterator& b) const

Returns true if both iterators are end-of-stream iterators or if neither iterator is an end-of-stream iterator. The iterators do not have to use the same stream buffer.

## See Also

[istream\\_iterator class template](#), [ostreambuf\\_iterator class template](#), `basic_streambuf` in [<streambuf>](#)

# Name

iterator class template — Iterator base-class template

## Synopsis

```
template<typename Category, typename T, typename Difference =  
ptrdiff_t,  
        typename Pointer = T*, typename Reference = T&>  
struct iterator{  
    typedef T value_type;  
    typedef Difference difference_type;  
    typedef Pointer pointer;  
    typedef Reference reference;  
    typedef Category iterator_category;  
};
```

The `iterator` class template is a convenient base-class template to use when implementing your own iterator.

The following are the template parameters, which are all type parameters:

### Category

Must be one of the five iterator category tags: `bidirectional_iterator_tag`, `forward_iterator_tag`, `input_iterator_tag`, `output_iterator_tag`, or `random_access_iterator_tag`.

### T

The element type. It can be `void` for an output iterator because you cannot dereference an output iterator.

### Difference

An integral type that represents the distance between two iterators. It can be `void` for an output iterator because you cannot measure the distance between two output iterators. This parameter is optional; the default



is `ptrdiff_t`, which is suitable for typical pointer-like iterators.

### Pointer

The pointer-to-element type. This parameter is optional; the default is `T*`, which is correct for most iterators.

### Reference

The reference-to-element type. This parameter is optional; the default is `T&`, which is correct for most iterators.

## See Also

[iterator\\_traits class template](#), [reverse\\_iterator class template](#)

# Name

iterator\_traits class template — Iterator traits

## Synopsis

```
template<typename Iterator>
struct iterator_traits{
    typedef typename Iterator::difference_type difference_type;
    typedef typename Iterator::value_type value_type;
    typedef typename Iterator::pointer pointer;
    typedef typename Iterator::reference reference;
    typedef typename Iterator::iterator_category iterator_category;
};
```

The `iterator_traits` class template declares traits for an iterator. If you use the `iterator` class template as the base for your custom iterator, you don't need to specialize `iterator_traits`. If you are writing a custom container or algorithm, you should always use `iterator_traits` to obtain the traits of an iterator. If you use a plain pointer as an iterator, the standard library specializes `iterator_traits` for you. See the next subsection.

If you write your own specialization, the `iterator_category` type must be one of the five iterator tag classes. (See the `iterator` class template.) For an output iterator, `difference_type` and `value_type` are `void`.

When writing a generic algorithm or other function that uses iterators, you can use `iterator_traits` to specialize the behavior for certain kinds of iterators. See [Example 13-19](#) (under `distance`), which shows how iterator traits can be used to improve the performance of a function.

## See Also

[iterator class template](#)

## Name

`iterator_traits<T*>` template specialization — Iterator traits specialized for pointers

## Synopsis

```
template<typename T>
struct iterator_traits<T*>{
    typedef ptrdiff_t difference_type;
    typedef T value_type;
    typedef T* pointer;
    typedef T& reference;
    typedef random_access_iterator_tag iterator_category;
};
```

The `iterator_traits` class template is specialized for pointers. This specialization lets you use a pointer as a random access iterator.

## See Also

[iterator\\_traits class template](#)

## Name

`iterator_traits<const T*>` template specialization —  
Iterator traits specialized for pointers to const

## Synopsis

```
template<typename T>
struct iterator_traits<const T*>{
    typedef ptrdiff_t difference_type;
    typedef T value_type;
    typedef const T* pointer;
    typedef const T& reference;
    typedef random_access_iterator_tag iterator_category;
};
```

The `iterator_traits` class template is specialized for pointers to const. This specialization lets you use a pointer as a random access iterator.

## See Also

[iterator\\_traits class template](#)

# Name

ostream\_iterator class template — Output iterator to write items to an ostream

# Synopsis

```
template <typename T, typename charT = char,
          typename traits = char_traits<charT> >
class ostream_iterator :
    public iterator<output_iterator_tag, void, void, void, void>
{
public:
    typedef charT char_type;
    typedef traits traits_type;
    typedef basic_ostream<charT, traits> ostream_type;
    ostream_iterator(ostream_type& s);
    ostream_iterator(ostream_type& s, const charT* delimiter);
    ostream_iterator(const ostream_iterator<T, charT, traits>& x);
    ~ostream_iterator( );
    ostream_iterator<T, charT, traits>& operator=(const T& value);
    ostream_iterator<T, charT, traits>& operator*( );
    ostream_iterator<T, charT, traits>& operator++( );
    ostream_iterator<T, charT, traits>& operator++(int);
};
```

The `ostream_iterator` class template wraps an output iterator around an output stream (instance of `basic_ostream`), making the stream appear to be a sequence of items, each of type `T`. For example, suppose you have a vector of data. You can print the data, one number per line, by using an `ostream_iterator`:

```
std::vector<int> data;
. . . // Acquire data.
std::copy(data.begin( ), data.end( ),
          std::ostream_iterator(std::cout, "\n"));
```

The following are the member functions of `ostream_iterator`:

```
ostream_iterator (ostream_type& stream),
ostream_iterator (ostream_type& stream, const
```

charT\* delimiter), **ostream\_iterator** (const ostream\_iterator<T, charT, traits>& x)

Prepares to write items to stream. If delimiter is present, it will be written after each item. The copy constructor copies the reference to the stream and to the delimiter from x.

ostream\_iterator<T, charT, traits>& **operator=** (const T& value)

Writes value to the stream using operator<<. If the ostream\_iterator has a delimiter, it is written after value. The return value is \*this.

ostream\_iterator<T, charT, traits>& **operator\*** ( )

Returns \*this.

ostream\_iterator<T, charT, traits>& **operator++** ( ),  
ostream\_iterator<T, charT, traits>& **operator++** (int)

Returns \*this with no side effects.

## See Also

[istream\\_iterator class template](#), [ostreambuf\\_iterator class template](#), basic\_ostream in [<ostream>](#)

## Name

`ostreambuf_iterator` class template — Output iterator to write characters to a streambuf

## Synopsis

```
template<typename charT, typename traits=char_traits<charT> >
class ostreambuf_iterator :
    public iterator<output_iterator_tag,void,void,void,void>
{
public:
    typedef charT char_type;
    typedef traits traits_type;
    typedef basic_streambuf<charT,traits> streambuf_type;
    typedef basic_ostream<charT,traits> ostream_type;

    ostreambuf_iterator(ostream_type& s) throw( );
    ostreambuf_iterator(streambuf_type* s) throw( );
    ostreambuf_iterator& operator=(charT c);
    ostreambuf_iterator& operator*( );
    ostreambuf_iterator& operator++( );
    ostreambuf_iterator& operator++(int);
    bool failed( ) const throw( );
};
```

The `ostreambuf_iterator` class template wraps a `basic_streambuf` object as an output iterator to write characters to the stream buffer. [Example 13-21](#) (under `istreambuf_iterator`) shows how to use `streambuf` iterators to copy files.

In the following descriptions of the member functions of `ostreambuf_iterator`, the data member *sbuf* is a pointer to the iterator's stream buffer. The *sbuf* member serves only to keep the function descriptions clear and simple; the class is not required to have such a member, or a member with that name.

```
ostreambuf_iterator (ostream_type& s) throw( ),
ostreambuf_iterator (streambuf_type* sb) throw( )
```

Saves the stream buffer `s.rdbuf( )` or `sb` in `sbuf`.

`ostreambuf_iterator& operator= (charT c)`

Calls `sbuf ->sputc(c)` (only if `failed( )` returns false) and returns `*this`.

`ostreambuf_iterator& operator* ( )`

Returns `*this`.

`ostreambuf_iterator& operator++ ( )`,  
`ostreambuf_iterator& operator++ (int)`

Returns `*this`.

`bool failed ( ) const throw( )`

Returns true if `sbuf ->sputc( )` ever returned `traits::eof( )`. Otherwise, it returns false.

## See Also

[istreambuf\\_iterator class template](#), [ostream\\_iterator class template](#), `basic_streambuf` in [<streambuf>](#)



## Name

`output_iterator_tag` class — Tag for an output iterator

## Synopsis

```
struct output_iterator_tag {};
```

Use the `output_iterator_tag` class as the iterator category when declaring a new output iterator class. When writing a generic algorithm or similar function, you can use the iterator's category to write specialized implementations for different kinds of iterators. See [Example 13-19](#) (under the distance function template).

## See Also

[bidirectional\\_iterator\\_tag class](#), [forward\\_iterator\\_tag class](#), [input\\_iterator\\_tag class](#), [iterator class template](#), [random\\_access\\_iterator\\_tag class](#)

## Name

`random_access_iterator_tag` class — Tag for a random access iterator

## Synopsis

```
struct random_access_iterator_tag :  
    public bidirectional_iterator_tag {};
```

Use the `random_access_iterator_tag` class as the iterator category when declaring a new random access iterator class. When writing a generic algorithm or similar function, you can use the iterator's category to write specialized implementations for different kinds of iterators. See [Example 13-19](#) (under the `distance` function template).

## See Also

[bidirectional\\_iterator\\_tag\\_class](#), [forward\\_iterator\\_tag\\_class](#), [input\\_iterator\\_tag\\_class](#), [iterator class template](#), [output\\_iterator\\_tag\\_class](#)

# Name

reverse\_iterator class template — Iterator wrapper to reverse direction

# Synopsis

```
template <typename Iterator>
class reverse_iterator : public iterator<
    typename iterator_traits<Iterator>::iterator_category,
    typename iterator_traits<Iterator>::value_type,
    typename iterator_traits<Iterator>::difference_type,
    typename iterator_traits<Iterator>::pointer,
    typename iterator_traits<Iterator>::reference>
{
protected:
    Iterator current;
public:
    typedef Iterator iterator_type;
    typedef typename iterator_traits<Iterator>::difference_type
        difference_type;
    typedef typename iterator_traits<Iterator>::reference
reference;
    typedef typename iterator_traits<Iterator>::pointer pointer;
    reverse_iterator( );
    explicit reverse_iterator(Iterator x);
    template <typename U>
reverse_iterator(const reverse_iterator<U>& u);
    Iterator base( ) const; // Explicit
    reference operator*( ) const;
    pointer operator->( ) const;
    reverse_iterator& operator++( );
    reverse_iterator operator++(int);
    reverse_iterator& operator--( );
    reverse_iterator operator--(int);
    reverse_iterator operator+(difference_type n) const;
    reverse_iterator& operator+=(difference_type n);
    reverse_iterator operator-(difference_type n) const;
    reverse_iterator& operator-=(difference_type n);
    reference operator[](difference_type n) const;
};
```

The reverse\_iterator class template is an adapter for a bidirectional or random access iterator to iterate the

sequence in the opposite direction of the adapted iterator. In other words, if the adapted iterator advances from the first to the last item in a container, a reverse iterator starts with the last items and "advances" to the first item. In addition to the `reverse_iterator` template, there are also several function templates for the comparison and arithmetic (+ and -) operators.

The standard containers return `reverse_iterator` objects from the `rbegin()` and `rend()` functions. The following example shows a simple implementation of these functions:

```
reverse_iterator rbegin( ) { return reverse_iterator(end( )); }  
reverse_iterator rend( ) { return reverse_iterator(begin( )); }
```

Because an iterator can point to one past the last item in a container but cannot point to one item before the first, a reverse iterator conceptually points to one item *before* the position to which the adapted iterator points. In other words, given an iterator, `iter`, which points to 42 in [Example 13-23](#), if you construct a reverse iterator that adapts `iter`, the reverse iterator appears to point to 29. When you increment the reverse iterator, it decrements the adapted iterator. Thus, the "next" element in the sequence of the reverse iterator is 12. (The adapted iterator points to 29.) The adapted iterator is also called the *base iterator*. See [Chapter 10](#) for a discussion of some of the interesting ramifications of using `reverse_iterator`.

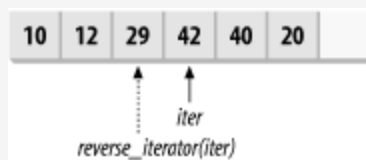


Figure 13-23. How a reverse iterator works

The following member functions of `reverse_iterator` refer to *adapted* as the data member that stores the `base()` iterator. The *adapted* member is actually named `current`, and is a protected data member of `reverse_iterator`:

## **reverse\_iterator** ( )

Initializes the *adapted* data member with its default constructor.

## explicit **reverse\_iterator** (Iterator iter)

Initializes the *adapted* data member to *iter*.

## template <typename U> **reverse\_iterator** (const reverse\_iterator<U>& ri)

Initializes the *adapted* data member to *ri.base*( ).

## Iterator **base** ( ) const

Returns the adapted iterator, *adapted*. Note that the adapted iterator points to one position *after* the reverse iterator's logical position.

## reference **operator** \*( ) const

Returns a reference to the item that the reverse iterator logically points to, which is one item *before* the item to which *base*( ) points. Thus, you can think of the dereference function working as follows:

```
reference operator*( ) const {
    iterator_type tmp = base( );
    --tmp;
    return *tmp;
}
```

## pointer **operator** ->( ) const

Returns a pointer to the item that the reverse iterator logically points to, that is, *&(operator\*( ))*.

## reverse\_iterator& **operator** ++( )

Decrements *adapted* and returns *\*this*.

## reverse\_iterator **operator**++ (int)

Saves a copy of (*\*this*), decrements *adapted*, and returns the saved copy of *\*this*.

`reverse_iterator& operator -- ( )`

Increments *adapted* and returns *\*this*.

`reverse_iterator operator-- (int)`

Saves a copy of the current item (`operator*( )`), increments *adapted*, and returns the saved item.

`reverse_iterator operator+ (difference_type n)`  
`const`

Returns `reverse_iterator(base( ) - n)`.

`reverse_iterator& operator+= (difference_type n)`

Subtracts *n* from *adapted* and returns *\*this*.

`reverse_iterator operator- (difference_type n)`  
`const`

Returns `reverse_iterator(base( ) + n)`.

`reverse_iterator& operator-= (difference_type n)`

Adds *n* to *adapted* and returns *\*this*.

`reference operator[] (difference_type n) const`

Returns `base( )[-n-1]`.

The following are several nonmember functions that compare reverse iterators and perform basic arithmetic such as finding the distance between two reverse iterators and advancing a reverse iterator by an integer:

```
template <typename Iterator>, bool operator==  
(const reverse_iterator<Iterator>& x, const  
reverse_iterator<Iterator>& y)
```

Returns true when the base iterators are equal, that is, `x.base( ) == y.base( )`.

```
template <typename Iterator>, bool operator!=  
(const reverse_iterator<Iterator>& x, const
```

```
reverse_iterator<Iterator>& y)
```

Returns true when x and y have different base iterators, that is, `x.base( ) != y.base( )`.

```
template <typename Iterator>, bool operator< (const reverse_iterator<Iterator>& x, const reverse_iterator<Iterator>& y)
```

Returns true when x is closer than y to the beginning of the sequence. Because x and y are reverse iterators, the function returns `y.base( ) < x.base( )`.

```
template <typename Iterator>, bool operator> (const reverse_iterator<Iterator>& x, const reverse_iterator<Iterator>& y)
```

Returns true when x is farther than y from the beginning of the sequence. Because x and y are reverse iterators, the function returns `y.base( ) > x.base( )`.

```
template <typename Iterator>, bool operator>= (const reverse_iterator<Iterator>& x, const reverse_iterator<Iterator>& y)
```

Returns true when x is farther than y from the beginning of the sequence or x equals y. Because x and y are reverse iterators, the function returns `y.base( ) >= x.base( )`.

```
template <typename Iterator>, bool operator<= (const reverse_iterator<Iterator>& x, const reverse_iterator<Iterator>& y)
```

Returns true when x is closer than y to the beginning of the sequence or x equals y. Because x and y are reverse iterators, the function returns `y.base( ) <= x.base( )`.

```
template <typename Iterator>, typename reverse_iterator<Iterator>::difference_type,
```

**operator-** (const reverse\_iterator<Iterator>& x,  
const reverse\_iterator<Iterator>& y)

Returns the distance between two reverse iterators, that is, `y.base( ) - x.base( )`.

```
template <typename Iter>, reverse_iterator<Iter>,  
operator+ (typename  
reverse_iterator<Iter>::difference_type n,  
const reverse_iterator<Iter>& ri)
```

Advances a reverse iterator `ri` by `n`. This function is a counterpart to the `operator+` member function, allowing you to write `ri + n` and `n + ri`, which yield the same result.

## See Also

[iterator class template](#), [Chapter 10](#)



## **<limits>**

The `<limits>` header declares the `numeric_limits` class template and related types and specializations that define the limits and characteristics of the fundamental arithmetic types (such as the largest possible `int`). It is the C++ equivalent of the C headers `<float>` and `<climits>` (and the `wchar_t` limits in `<wchar>`).

The `<limits>` header has a number of advantages over the `<float>` and `<climits>` declarations. In particular, by using a template, you can write your own template-based classes that depend on the characteristics of a template parameter; this is not possible with the macro-based C headers.

## Name

`float_denorm_style` type — Represents existence of denormalized, floating-point values

## Synopsis

```
enum float_denorm_style {  
    denorm_indeterminate = -1;  
    denorm_absent = 0;  
    denorm_present = 1;  
};
```

The `float_denorm_style` type is an enumerated type that represents whether denormalized floating-point values are supported.

## Name

`float_round_style` type — Represents the floating-point rounding style

## Synopsis

```
enum float_round_style {  
    round_indeterminate = -1,  
    round_toward_zero = 0,  
    round_to_nearest = 1,  
    round_toward_infinity = 2,  
    round_toward_neg_infinity = 3  
};
```

The `float_round_style` type is an enumerated type that represents how floating-point numbers are rounded.

# Name

`numeric_limits` class template — Represents the limits and characteristics of an arithmetic type

# Synopsis

```
template<typename T>
class numeric_limits{
public:
    static const bool is_specialized = false;
    static T min( ) throw( );
    static T max( ) throw( );
    static const int digits = 0;
    static const int digits10 = 0;
    static const bool is_signed = false;
    static const bool is_integer = false;
    static const bool is_exact = false;
    static const int radix = 0;
    static T epsilon( ) throw( );
    static T round_error( ) throw( );
    static const int min_exponent = 0;
    static const int min_exponent10 = 0;
    static const int max_exponent = 0;
    static const int max_exponent10 = 0;
    static const bool has_infinity = false;
    static const bool has_quiet_NaN = false;
    static const bool has_signaling_NaN = false;
    static const float_denorm_style has_denorm = denorm_absent;
    static const bool has_denorm_loss = false;
    static T infinity( ) throw( );
    static T quiet_NaN( ) throw( );
    static T signaling_NaN( ) throw( );
    static T denorm_min( ) throw( );
    static const bool is_iec559 = false;
    static const bool is_bounded = false;
    static const bool is_modulo = false;
    static const bool traps = false;
    static const bool tinyness_before = false;
    static const float_round_style round_style = round_toward_zero;
};
```

The `numeric_limits` class template represents the limits and characteristics of an arithmetic type. The data

members that are shown as `static const` are constants that you can use in other integral constant expressions.

The default is for all members to be `0` or `false`. The header has specializations for all fundamental types, and only for the fundamental types. Every specialization in the standard library defines every member, even if the member does not pertain to the type (e.g., floating-point characteristics of an integral type). Meaningless members are defined as `0` or `false`.

You can specialize `numeric_limits` for your own types. For example, suppose you write a class, `bigint`, to represent integers of arbitrary size. You can define your specialization to show that the type is unbounded, signed, integral, etc. You should follow the convention of the standard library, namely, by defining all members, even if they do not apply to your type. Be sure to define `is_specialized` as `true`.

Use `numeric_limits` to query the properties or traits of a numeric type. For example, suppose you are writing a data analysis program. Among the data are points you want to ignore, but you need to keep their places in the data array. You decide to insert a special marker value. Ideally, the marker value (such as infinity) cannot possibly appear in the actual data stream. If the floating-point type that you are using does not support infinity, you can use the maximum finite value. [Example 13-23](#) lists the `no_data` function, which returns the value used for the no-data marker.

## Example

*Example 13-23. Using infinity or maximum finite value to mean "no data"*

```
// Define a template that will differentiate types that have a
// specialized
// numeric_limits and an explicit value for infinity.
```

```

template<typename T, bool is_specialized, bool has_infinity>
struct max_or_infinity
{};

// Specialize the template to obtain the value of infinity.
template<typename T>
struct max_or_infinity<T, true, true>
{
    static T value( )
        { return std::numeric_limits<T>::infinity( ); }
};

// Specialize the template if infinity is not supported.
template<typename T>
struct max_or_infinity<T, true, false>
{
    static T value( ) { return std::numeric_limits<T>::max( ); }
};

// Note that a type without a numeric_limits specialization does
// not have a
// max_or_infinity specialization, so the no_data function would
// result in
// compile-time errors when applied to such a type.
//
// The no_data function returns a value that can be used to mark
// points that do
// not have valid data.
template<typename T>
Tno_data( )
{
    return max_or_infinity<T,
        std::numeric_limits<T>::is_specialized,
        std::numeric_limits<T>::has_infinity>::value( );
}

```

The C++ standard mandates that all integers are binary and use two's complement, ones' complement, or signed magnitude representation. The representation of floating-point numbers is not specified. The `numeric_limits` template assumes that a number is represented as a sign, a significand (sometimes called the *mantissa*), a base, and an exponent:

$$x = \text{sign} \times \text{significand} \times \text{base}^{\text{exponent}}$$

In everyday arithmetic, we are used to working with a *base* of 10. (The base is also called the *radix*.) The most common bases for computer arithmetic, however, are 16 and 2. Many modern workstations use the IEC 60559 (IEEE 754) standard for floating-point arithmetic, which uses a base of 2.

The *significand* is a string of digits in the given base. There is an implied radix point at the start of the significand, so the value of the significand is always less than one. (A *radix point* is the generalization of a decimal point for any radix.)

A finite floating-point value is *normalized* if the first digit of its significand is nonzero, or if the entire value is 0. The term *denormalized* means a finite value is not normalized.

The *precision* of a floating-point type is the maximum number of places in the significand. The *range* of a floating-point type depends primarily on the minimum and maximum values for the exponent.

The following are descriptions of the members of `numeric_limits`:

```
static T denorm_min ( ) throw( )
```

Returns the smallest positive, denormalized floating-point value. If `has_denorm` is `false`, it returns the smallest positive normalized value. For non-floating-point types, it returns 0.

```
static const int digits
```

The number of radix digits that can be represented. For integer types, it is the number of non-sign bits; for floating-point types, it is the number of places in the significand.

```
static const int digits10
```

The number of decimal digits that can be represented. If `is_bounded` is false, `digits10` is 0.

`static T epsilon ( ) throw( )`

Returns the difference between 1.0 and the smallest representable value greater than 1.0. For integral types, `epsilon` returns 0.

`static const float_denorm_style has_denorm`

Indicates the denormalized, floating-point style. It is `denorm_indeterminate` if the style cannot be determined at compile time. It is meaningful for all floating-point types.

`static const bool has_denorm_loss`

Indicates whether the loss of accuracy in a floating-point computation is a denormalization loss rather than an inexact result.

`static const bool has_infinity`

Indicates whether the floating-point type can represent positive infinity. In particular, `has_infinity` is true when `is_iec559` is true.

`static const bool has_quiet_NaN`

Indicates whether the floating-point type can represent a quiet (nonsignaling) NaN (not-a-number). In particular, this is true when `is_iec559` is true.

`static const bool has_signaling_NaN`

Indicates whether the floating point type can represent a signaling NaN. In particular, this is true when `is_iec559` is true.

`static T infinity ( ) throw( )`

Returns the value of positive infinity if `has_infinity` is true.



static const bool **is\_bounded**

Indicates whether the type represents a finite set of values. This is true for all fundamental types.

static const bool **is\_exact**

Indicates whether the type represents values exactly. It is true for all integral types and false for the fundamental floating-point types.

static const bool **is\_iec559**

Indicates whether the type follows the IEC 60559 (IEEE 754) standard for floating-point arithmetic. It is meaningful only for floating-point types. Among the requirements of the IEC 60559 standard are support for positive and negative infinity, and for values that are NaN.

static const bool **is\_integer**

true for all integral types.

static const bool **is\_modulo**

Indicates whether the type uses modulo arithmetic. This is always true for unsigned integral types and often true for signed integral types. It is false for typical floating-point types.

static const bool **is\_signed**

Indicates whether the type is signed, that is, supports positive and negative values.

static const bool **is\_specialized**

Indicates whether `numeric_limits` is specialized for the type. It is false by default, so you can detect whether `numeric_limits<>` has been specialized for a particular type, and therefore determine whether a false or 0 value is meaningful.

static T **max** ( ) throw( )

Returns the maximum finite value when `is_bounded` is true.

static const int **max\_exponent**

The largest allowable exponent for a finite floating-point number.

static const int **max\_exponent10**

The largest allowable decimal exponent for a finite floating-point number.

static T **min** ( ) throw( )

Returns the minimum finite value. It is meaningful when `is_bounded` is true or when `is_bounded` and `is_signed` are both false.

static const int **min\_exponent**

The smallest allowable exponent for a floating-point number such that `radix` raised to `min_exponent - 1` is representable as a normalized floating-point number.

static const int **min\_exponent10**

The smallest negative decimal exponent such that 10 raised to `min_exponent10` is representable as a normalized floating-point number.

static T **quiet\_NaN** ( ) throw( )

Returns a quiet NaN value if `has_quiet_NaN` is true.

static const int **radix**

The base used in the representation of a numeric value. For floating-point numbers, it is the base of the exponent.

static T **round\_error** ( ) throw( )

Returns the maximum rounding error.

static const float\_round\_style **round\_style**

Indicates the rounding style used by the floating-point type. (See the `float_round_style` type for a definition of the possible return values.) For integral types, the return value is always `round_toward_zero`.

static T **signaling\_NaN** ( ) throw( )

Returns a signaling NaN value if `has_signaling_NaN` is true.

static const bool **tinyness\_before**

Indicates whether a floating-point type tests for denormalized values before rounding.

static const bool **traps**

Indicates whether arithmetic errors *trap*, that is, result in signals or exceptions. It is false if errors are quietly ignored.

The `numeric_limits` template is specialized for all the fundamental numeric types and for no other types in the C++ standard. In each case, `is_specialized` is true, and other members are set as appropriate. The C++ standard (by way of the C standard) defines the minimum requirements for an implementation. The requirements for integral types are given in `<climits>`, and for floating-point types in `<cfloat>`.

The following are the standard specializations:

```
numeric_limits<bool>  
numeric_limits<char>  
numeric_limits<double>  
numeric_limits<float>  
numeric_limits<int>
```

```
numeric_limits<long>  
numeric_limits<long double>  
numeric_limits<short>  
numeric_limits<signed char>  
numeric_limits<unsigned char>  
numeric_limits<unsigned int>  
numeric_limits<unsigned long>  
numeric_limits<unsigned short>  
numeric_limits<wchar_t>
```

## **<list>**

The `<list>` header is one of the standard container template headers. It declares the `list` class template and a few global function templates that operate on `list` objects.

A `list` is a sequence container that has constant performance when adding to or removing from any point in the container. It supports bidirectional iterators, but not random access. Although the standard does not mandate any particular implementation, the obvious choice is to use a doubly-linked list to implement the `list` class template.

See [Chapter 10](#) for information about containers.

# Name

list class template — List container

## Synopsis

```
template <typename T, typename Alloc = allocator<T> >
class list{
public:
    // Types
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef T value_type;
    typedef Alloc allocator_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;

    // Construct/copy/destroy
    explicit list(const Alloc& = Alloc( ));
    explicit list(size_type n, const T& value = T( ), const Alloc&
= Alloc( ));
    template <class InputIterator>
list(InputIterator first, InputIterator last, const Alloc& =
Alloc( ));
    list(const list<T,Alloc>& x);
    ~list( );
    list<T,Alloc>& operator=(const list<T,Alloc>& x);
    template <class InputIterator>
    void assign(InputIterator first, InputIterator last);
    void assign(size_type n, const T& t);
    allocator_type get_allocator( ) const;
    // Iterators
    iterator begin( );
    const_iterator begin( ) const;
    iterator end( );
    const_iterator end( ) const;
    reverse_iterator rbegin( );
    const_reverse_iterator rbegin( ) const;
```

```

reverse_iterator rend( );
const_reverse_iterator rend( ) const;
// Capacity
bool empty( ) const;
size_type size( ) const;
size_type max_size( ) const;
void resize(size_type sz, T c = T( ));
// Element access
reference front( );
const_reference front( ) const;
reference back( );
const_reference back( ) const;
// Modifiers
void push_front(const T& x);
void pop_front( );
void push_back(const T& x);
void pop_back( );
iterator insert(iterator position, const T& x);
void insert(iterator position, size_type n, const T& x);
template <class InputIterator>
void insert(iterator position, InputIterator first,
InputIterator last);
iterator erase(iterator position);
iterator erase(iterator position, iterator last);
void swap(list<T,Alloc>&);
void clear( );
// List operations
void splice(iterator position, list<T,Alloc>& x);
void splice(iterator position, list<T,Alloc>& x, iterator i);
void splice(iterator position, list<T,Alloc>& x, iterator
first,
            iterator last);
void remove(const T& value);
template <class Predicate>
void remove_if(Predicate pred);
void unique( );
template <class BinaryPredicate>
void unique(BinaryPredicate binary_pred);
void merge(list<T,Alloc>& x);
template <class Compare>
void merge(list<T,Alloc>& x, Compare comp);
void sort( );
template <class Compare> void sort(Compare comp);
void reverse( );
};

```

The `list` class template is one of the standard container types, like `deque` and `vector`. A `list` stores a sequence of items such that inserting or erasing an item at any position requires constant time. The `list` template supports all the usual operations for a sequence container plus some functions that are unique to `list`.

When an item is erased from the list (by calling `pop_back`, `erase`, `remove`, etc.), all iterators that point to that item become invalid. All pointers and references to the item become invalid. No other iterators, pointers, or references are invalidated when inserting or erasing any items.



The `size` function can have constant or linear complexity. The standard encourages library vendors to implement the `list` class template so that `size` has constant complexity, but it permits worse performance (namely, linear in the size of the list). If `size` does not have constant complexity, you should expect all versions of `splice` to have constant complexity in all cases. (The last constraint is not mandated by the standard, but by common sense.)

The following are the member functions of `list`:

```
explicit list (const Alloc& = Alloc( ))
```

Initializes an empty list that uses the given allocator.

```
explicit list (size_type n, const T& value = T( ),  
const Alloc& = Alloc( ))
```

Initializes a list that contains `n` copies of `value`.

```
template < typename InputIterator>, list  
(InputIterator first, InputIterator last, const  
Alloc& = Alloc( ))
```



Initializes the list with a copy of the items in the range [first, last), unless InputIterator is an integral type, in which case the list is constructed as though the arguments were cast:

```
list(static_cast<size_type>(first), static_cast<value_type>
      (last),
      alloc);
```

**list** (const list<T,Alloc>& x)

Constructs a copy of the contents and allocator of the list x.

list<T,Alloc>& **operator=** (const list<T,Alloc>& x)

Replaces the list's contents with copies of the contents of x.

template <typename InputIterator>, void **assign**  
(InputIterator first, InputIterator last)

Replaces the list's contents with the items in the range [first, last), unless InputIterator is an integral type, in which case the arguments are interpreted as though they were cast:

```
assign(static_cast<size_type>(first), static_cast<value_type>
      (last));
```

void **assign** (size\_type n, const T& value)

Replaces the list's contents with n copies of value.

reference **back** ( ), const\_reference **back** ( ) const

Returns the last item in the list. The behavior is undefined if the list is empty.

iterator **begin** ( ), const\_iterator **begin** ( ) const

Returns an iterator that points to the first item in the list.

void **clear** ( )

Erases all the items in the list, invalidating all iterators that point to the list.

bool **empty** ( ) const

Returns true if the list is empty. Note that `empty( )` has constant complexity even if `size( )` does not.

iterator **end** ( ), const\_iterator **end** ( ) const

Returns an iterator that points one past the last item in the list.

iterator **erase** (iterator position), iterator **erase** (iterator first, iterator last)

Erases the item at position or all the items in the range [first, last).

reference **front** ( ), const\_reference **front** ( ) const

Returns the first item in the list. The behavior is undefined if the list is empty.

allocator\_type **get\_allocator** ( ) const

Returns the list's allocator.

iterator **insert** (iterator position, const T& x),  
void **insert** (iterator position, size\_type n, const T& x),  
template <typename InputIterator>, void **insert** (iterator position, InputIterator first, InputIterator last)

Inserts one or more items before position. The performance is linear in the number of items inserted, and the T copy constructor is invoked once for each item inserted in the list. The first form inserts the item x; the second form inserts n copies of x; the third form copies the items in the range [first, last), unless InputIterator is an integral type, in which case the arguments are interpreted as though they were cast:

```
insert(position, static_cast<size_type>(first),
        static_cast<value_type>(last));
```

If an exception is thrown, such as `bad_alloc` when there is insufficient memory for a new element, the list is unchanged.

```
size_type max_size ( ) const
```

Returns the size of the largest possible list.

```
void merge (list<T,Alloc>& x), template <class  
Compare>, void merge (list<T,Alloc>& x, Compare  
comp)
```

Merges another sorted list, `x`, into the current list, which must also be sorted. Items are erased from `x`, so after `merge` returns, `x` is empty. Items are compared using the `<` operator or `comp`. The same function used to sort the items must be used to compare items. The merge is stable, so the relative order of items is unchanged; if the same item is already in the list and in `x`, the item from `x` is added after the item already in the list.

The performance of the merge is linear: exactly `size( ) + x.size( ) - 1` comparisons are performed.

```
void pop_back ( )
```

Erases the last item from the list. The behavior is undefined if the list is empty.

```
void pop_front ( )
```

Erases the first item from the list. The behavior is undefined if the list is empty.

```
void push_back (const T& x)
```

Inserts `x` at the end of the list.

```
void push_front (const T& x)
```

Inserts `x` at the beginning of the list.

reverse\_iterator **rbegin** ( ), const\_reverse\_iterator **rbegin** ( ) const

Returns a reverse iterator that points to the last item in the list.

void **remove** (const T& value)

Erases all occurrences of value from the list. The performance is linear; exactly size( ) comparisons are performed.

template <typename Predicate>, void **remove\_if** (Predicate pred)

Erases all items for which pred(item) returns true. The performance is linear: pred is called exactly size( ) times.

reverse\_iterator **rend** ( ), const\_reverse\_iterator **rend** ( ) const

Returns a reverse iterator that points to one position before the first item in the list.

void **resize** (size\_type sz, T c = T( ))

Changes the size of the list to n. If n > size( ), one or more copies of c are added to the end of the list to reach the desired size. If the new size is smaller than the current size, elements are erased from the end to reach the new size.

void **reverse** ( )

Reverses the order of the entire list. The performance is linear.



size\_type **size** ( ) const

Returns the number of elements in the list. The complexity of size( ) can be constant or linear,

depending on the implementation.

```
void sort ( ), template <typename Compare>, void sort  
(Compare comp)
```

Sorts the items in the list, comparing items with the < operator or by calling comp. The sort is stable, so the relative positions of the items do not change. The performance is  $N \log N$ , in which  $N$  is `size( )`.

### **Not e**

You must call the `sort` member function to sort a list. The generic sort algorithm requires a random access iterator, but `list` provides only a bidirectional iterator.

```
void splice (iterator position, list<T,Alloc>& x),  
void splice (iterator position, list<T,Alloc>& x,  
iterator i), void splice (iterator position,  
list<T,Alloc>& x, iterator first, iterator last)
```

Moves one or more items from `x`, inserting the items just before `position`. The first form moves every item from `x` to the list. The second form moves the item at position `i`. The third form moves all items in the range `[first, last)`; `position` must not be in that range. The third form requires no more than linear time when `&x != this`; all other cases work in constant time. If `size( )` has linear complexity, you should expect `splice( )` to have constant complexity in all cases.

```
void swap (list<T,Alloc>& x)
```

Swaps all the items in this list with all the items in `x`. The performance should be constant.

```
void unique ( ), template <typename  
BinaryPredicate>, void unique (BinaryPredicate  
pred)
```

Erases adjacent duplicate items from the list. Items are compared with the `==` operator or by calling `pred`. When adjacent equal items are found in the list, the first one is retained, and the second and subsequent items are erased. The performance is linear: `size( ) - 1` comparisons are performed (unless the list is empty).

## See Also

deque in [<deque>](#), vector in [<vector>](#)

## Name

`operator==` function template — Compares lists for equality

## Synopsis

```
template <typename T, typename A>  
bool operator==(const list<T,A>& x, const list<T,A>& y);
```

The `==` operator returns true if `x` and `y` have the same size and their elements are equal, that is, `x.size( ) == y.size( ) && equals(x.begin( ), x.end( ), y.begin( ))`.

## See Also

`equals` in [<algorithm>](#)

## Name

`operator!=` function template — Compares lists for inequality

## Synopsis

```
template <typename T, typename A>  
bool operator!=(const list<T,A>& x, const list<T,A>& y);
```

The `!=` operator is equivalent to `!(x == y)`.



## Name

operator< function template — Compares lists for less-than

## Synopsis

```
template <typename T, typename A>  
bool operator<(const list<T,A>& x, const list<T,A>& y);
```

The < operator determines whether x is less than y using the same algorithm as

```
lexicographical_compare(x.begin( ), x.end( ),  
y.begin( ), y.end( )).
```

## See Also

lexicographical\_compare in [<algorithm>](#)

## Name

`operator<=` function template — Compares lists for less-than-or-equal

## Synopsis

```
template <typename T, typename A>  
bool operator<=(const list<T,A>& x, const list<T,A>& y);
```

The `<=` operator is equivalent to `!(y < x)`.

## Name

`operator>` function template — Compares lists for greater-than

## Synopsis

```
template <typename T, typename A>  
bool operator>(const list<T,A>& x, const list<T,A>& y);
```

The `>` operator is equivalent to `(y < x)`.

## Name

`operator>=` function template — Compares lists for greater-than-or-equal

## Synopsis

```
template <typename T, typename A>  
bool operator>=(const list<T,A>& x, const list<T,A>& y);
```

The `>=` operator is equivalent to `! (x < y)`.

## Name

swap function template — Swaps the contents of two lists

## Synopsis

```
template<typename T, typename A>  
void swap(list<T, A>& x, list<T, A>& y)
```

The swap function template specialization is equivalent to calling `x.swap(y)`.

## See Also

swap in [<algorithm>](#)

## <locale>

The <locale> header declares class and function templates for internationalization and localization. It supports conversion between narrow and wide character sets, character classification and collation, formatting and parsing numbers, currency, dates and times, and retrieving messages. For example, every I/O stream has a locale, which it uses to parse formatted input or to format output.

A *locale* is an embodiment of a set of cultural conventions, including information about the native character set, how dates are formatted, which symbol to use for currency, and so on. Each set of related attributes is called a *facet*, which are grouped into categories.

The categories are fixed and defined by the standard (see [Table 13-20](#), under `locale::category`, for a complete list), and each category has several predefined facets. For example, one of the facets in the time category is `time_get<charT, InputIter>`, which specifies rules for parsing a time string. You can define additional facets; see the description of the `locale::facet` class in this section for details.

Many of the facets come in two flavors: plain and named. The plain versions implement default behavior, and the named versions implement the behavior for a named locale. See the `locale` class later in this section for a discussion of locale names.

When a program starts, the global locale is initialized to the "C" locale, and the standard I/O streams use this locale for character conversions and formatting. A program can change the locale at any time; see the `locale` class in this section for details.

The C++ `<locale>` header provides more functionality than the C `<locale>`, [`<ctype>`](#), and `<wchar>` headers, especially the ability to extend a locale with your own facets. On the other hand, facets and the locale class template are more complicated than the C functions. For simple character classification in a single locale, you are probably better off using the C functions. If a program must work with multiple locales simultaneously, use the C++ locale template.

# Name

codecvt class template — Facet for mapping one character set to another

## Synopsis

```
template <typename internT,typename externT,typename stateT>
class codecvt : public locale::facet, public codecvt_base
{
public:
    typedef internT intern_type;
    typedef externT extern_type;
    typedef stateT state_type;
    explicit codecvt(size_t refs = 0);
    result out(stateT& state, const internT* from, const internT*
from_end,
               const internT*& from_next, externT* to, externT*
to_limit,
               externT*& to_next) const;
    result unshift(stateT& state, externT* to, externT* to_limit,
externT*& to_next) const;
    result in(stateT& state, const externT* from, const externT*
from_end,
               const externT*& from_next, internT* to, internT*
to_limit,
               internT*& to_next) const;
    int encoding( ) const throw( );
    bool always_noconv( ) const throw( );
    int length(stateT&, const externT* from, const externT* end,
size_t max)
        const;
    int max_length( ) const throw( );
    static locale::id id;
protected:
    virtual ~codecvt( );
    virtual result do_out(stateT& state, const internT* from,
                           const internT* from_end, const internT*&
from_next,
                           externT* to, externT* to_limit, externT*&
to_next)
        const;
    virtual result do_in(stateT& state, const externT* from,
                           const externT* from_end, const externT*&
from_next,
```



```

        internT* to, internT* to_limit, internT*&
to_next)
    const;
    virtual result do_unshift(stateT& state, externT* to, externT*
to_limit,
                                externT*& to_next) const;
    virtual int do_encoding( ) const throw( );
    virtual bool do_always_noconv( ) const throw( );
    virtual int do_length(stateT&, const externT* from, const
externT* end,
                                size_t max) const;
    virtual int do_max_length( ) const throw( );
};

```

The `codecvt` template converts characters from one character encoding to another. It is most often used to convert multibyte characters to and from wide characters.

The following template specializations are required by the standard:

```

codecvt<wchar_t, char, mbstate_t>

```

Converts multibyte narrow characters to wide characters (in) and wide to multibyte (out)

```

codecvt<char, char, mbstate_t>

```

A no-op, "converting" characters to themselves

As with other facets, the public members of `codecvt` call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `in` and `out`, which in turn call `do_in` and `do_out`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual function description, there is a corresponding description for a public, nonvirtual function, such as:

```

bool always_noconv ( ) const throw( )

```

Returns `do_always_noconv( )`

The following are the virtual, protected members of `codecvt`:

```
virtual bool do_always_noconv ( ) const throw( )
```

Returns true if the `codecvt` object does not actually perform any conversions, that is, in and out are no-ops. For example, for the specialization `codecvt<char, char, mbstate_t>`, `do_always_noconv` always returns true.

```
virtual int do_encoding ( ) const throw( )
```

Returns the number of `externT` characters needed to represent a single `internT` character. If this number is not a fixed constant, the return value is 0. The return value is -1 if `externT` character sequences are not state-dependent.

```
virtual result do_in (stateT& state, const externT* from, const externT* from_end, const externT*& from_next, internT* to, internT* to_limit, internT*& to_next) const
```

Converts `externT` characters to `internT` characters. The characters in the range `[from, from_end)` are converted and stored in the array starting at `to`. The number of characters converted is the minimum of `from_end - from` and `to_limit - to`.

The `from_next` parameter is set to point to the value in `[from, from_end)` where the conversion stopped, and `to_next` points to the value in `[to, to_limit)` where the conversion stopped. If no conversion was performed, `from_next` is the same as `from`, and `to_next` is equal to `to`.

The return value is a `result`, as described in [Table 13-19](#) (under the `codecvt_base` class).

virtual int **do\_length** (stateT& state, const externT\* from, const externT\* from\_end, size\_t max) const

Returns the number of externT characters in the range [from, from\_end) that are used to convert to internT characters. At most, max internT characters are converted.

virtual int **do\_max\_length** ( ) const throw( )

Returns the maximum number of externT characters needed to represent a single internT character, that is, the maximum value that do\_length can return when max is 1.

virtual result **do\_out** (stateT& state, const internT\* from, const internT\* from\_end, const internT\*& from\_next, externT\* to, externT\* to\_limit, externT\*& to\_next) const

Converts internT characters to externT characters. The characters in the range [from, from\_end) are converted and stored in the array starting at to. The number of characters converted is the minimum of from\_end - from and to\_limit - to.

The from\_next parameter is set to point to the value in [from, from\_end) where the conversion stopped, and to\_next points to the value in [to, to\_limit) where the conversion stopped. If no conversion was performed, from\_next is the same as from, and to\_next is equal to to.

The return value is a result, as described in [Table 13-19](#) (under codecvt\_base class).

virtual result **do\_unshift** (stateT& state, externT\* to, externT\* to\_limit, externT\*& to\_next) const

Ends a shift state by storing characters in the array starting at to such that the characters undo the state

shift given by state. Up to `to_limit` - `to` characters are written, and `to_next` is set to point to one past the last character written into `to`.

The return value is a result, as described in [Table 13-19](#) (under `codecvt_base` class).

## See Also

[`codecvt\_base` class](#), [`codecvt\_byname` class template](#),  
[`locale::facet` class](#)

# Name

codecvt\_base class — Base class for the codecvt template

# Synopsis

```
class codecvt_base {  
public:  
    enum result { ok, partial, error, noconv };  
};
```

The codecvt\_base class is the base class for the codecvt and codecvt\_byname class templates. It declares the result type, which is the type returned by the do\_in and do\_out conversion functions. [Table 13-19](#) lists the literals of the result enumerated type.

*Table 13-19. codecvt\_base::result literals*

L i t e r a l	D e s c r i p t i o n
error	Error in conversion (e.g., invalid state or multibyte character sequence)
noconv	No conversion (or unshift terminated) needed
ok	Conversion finished successfully
partial	Not all source characters converted, or unshift sequence is incomplete

# See Also

[codecvt class template](#), [codecvt\\_byname class template](#)

## Name

`codecvt_byname` class template — Facet for mapping one character set to another

## Synopsis

```
template<typename internT, typename externT, typename stateT>
class codecvt_byname :
    public codecvt<internT, externT, stateT>
{
public:
    explicit codecvt_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in codecvt
};
```

The `codecvt_byname` class template converts characters from one character encoding to another using the rules of a named locale. The

`codecvt_byname<char, char, mbstate_t>` and `codecvt_byname<wchar_t, char, mbstate_t>` instantiations are standard.

## See Also

[codecvt class template](#), [locale::facet class](#)

# Name

collate class template — Facet for comparing strings in collation order

# Synopsis

```
template <typename charT>
class collate : public locale::facet
{
public:
    typedef charT char_type;
    typedef basic_string<charT> string_type;
    explicit collate(size_t refs = 0);
    int compare(const charT* low1, const charT* high1, const charT*
low2,
                const charT* high2) const;
    string_type transform(const charT* low, const charT* high)
const;
    long hash(const charT* low, const charT* high) const;
    static locale::id id;
protected:
    virtual ~collate( );
    virtual int do_compare(const charT* low1, const charT* high1,
                const charT* low2, const charT* high2)
const;
    virtual string_type do_transform (const charT* low, const
charT* high) const;
    virtual long do_hash (const charT* low, const charT* high)
const;
};
```

The `collate` class template is a facet used to compare strings. In some locales, the collation order of characters is not the same as the numerical order of their encodings, and some characters might be logically equivalent even if they have different encodings.

You can use a `locale` object as a comparator for algorithms that need a comparison function; the locale's `operator( )` function uses the `collate` facet to perform the comparison.

The standard mandates the `collate<char>` and `collate<wchar_t>` instantiations, which perform lexicographical (element-wise, numerical) comparison. See `lexicographical_compare` in `<algorithm>` earlier in this chapter.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `compare`, which calls `do_compare`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual function description, there is a corresponding description for a public, nonvirtual function, such as:

```
int compare (const charT* low1, const charT* high1,  
             const charT* low2, const charT* high2) const
```

Returns `do_compare(low1, high1, low2, high2)`

The following are the virtual, protected members of `collate`:

```
virtual int do_compare (const charT* low1, const  
                        charT* high1, const charT* low2, const charT* high2)  
const
```

Compares the character sequences `[low1, high1)` with the character sequence `[low2, high2)`. The return value is one of the following:

- -1 if sequence 1 is less than sequence 2
- 0 if the sequences are equal
- 1 if sequence 1 is greater than sequence 2

```
virtual long do_hash (const charT* low, const charT*  
                      high) const
```



Returns a hash value for the character sequence [low, high). If do\_compare returns 0 for two character sequences, do\_hash returns the same value for the two sequences. The reverse is not necessarily the case.

virtual string\_type **do\_transform** (const charT\* low, const charT\* high) const

Transforms the character sequence [low, high) into a string that can be compared (as a simple lexicographical comparison) with another transformed string to obtain the same result as calling do\_compare on the original character sequences. The do\_transform function is useful if a program needs to compare the same character sequence many times.

## See Also

[collate\\_byname class template](#), [locale class](#), [locale::facet class](#)

## Name

`collate_byname` class template — Facet for comparing strings in collation order

## Synopsis

```
template <typename charT>
class collate_byname : public collate<charT>
{
public:
    typedef basic_string<charT> string_type;
    explicit collate_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in collate
};
```

Compares strings using a named locale's collation order. The `collate_byname<char>` and `collate_byname<wchar_t>` instantiations are standard.

## See Also

[collate class template](#), [locale::facet class](#)

# Name

ctype class template — Facet for classifying characters

## Synopsis

```
class ctype : public locale::facet, public ctype_base
{
public:
    typedef charT char_type;
    explicit ctype(size_t refs = 0);
    bool is(mask m, charT c) const;
    const charT* is(const charT* low, const charT* high, mask* vec)
const;
    const charT* scan_is(mask m, const charT* low, const charT*
high) const;
    const charT* scan_not(mask m, const charT* low, const charT*
high) const;
    charT toupper(charT c) const;
    const charT* toupper(charT* low, const charT* high) const;
    charT tolower(charT c) const;
    const charT* tolower(charT* low, const charT* high) const;
    charT widen(char c) const;
    const char* widen(const char* low, const char* high, charT* to)
const;
    char narrow(charT c, char dfaul) const;
    const charT* narrow(const charT* low, const charT*, char
dfaul, char* to)
const;
    static locale::id id;
protected:
    virtual ~ctype( );
    virtual bool do_is(mask m, charT c) const;
    virtual const charT* do_is(const charT* low, const charT* high,
mask* vec)
const;
    virtual const charT* do_scan_is(mask m, const charT* low,
const charT* high) const;
    virtual const charT* do_scan_not(mask m, const charT* low,
const charT* high) const;
    virtual charT do_toupper(charT) const;
    virtual const charT* do_toupper(charT* low, const charT* high)
const;
    virtual charT do_tolower(charT) const;
    virtual const charT* do_tolower(charT* low, const charT* high)
```

```

const;
    virtual charT do_widen(char) const;
    virtual const char* do_widen(const char* low, const char* high,
charT* dest)
        const;
    virtual char do_narrow(charT, char dfault) const;
    virtual const charT* do_narrow(const charT* low, const charT*
high,
                                char dfault, char* dest) const;
};

```

The `ctype` class template is a facet for classifying characters.



The `ctype<char>` specialization is described in its own section later in this chapter. The standard also mandates the `ctype<wchar_t>` instantiation. Both instantiations depend on the implementation's native character set.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `narrow`, which calls `do_narrow`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual function description, there is a corresponding description for a public, nonvirtual function, such as:

```
bool is (mask m, charT c) const
```

Returns `do_is(m, c)`

The following are the virtual, protected members of `ctype`:

```
virtual bool do_is (mask m, charT c) const, virtual
const charT* do_is (const charT* low, const charT*
high, mask* dest) const
```

Classifies a single character `c` or a sequence of characters [`low`, `high`). The first form tests the

classification mask,  $M$ , of  $c$  and returns  $(M \& m) \neq 0$ . The second form determines the mask for each character in the range and stores the mask values in the `dest` array (which must be large enough to hold high - low masks), returning high. See [Table 13-19](#) (under the `ctype_base` class) for a description of the mask type.

```
virtual char do_narrow (charT c, char default)
const, virtual const charT* do_narrow (const charT*
low, const charT* high, char default, char* dest)
const
```

Converts a character  $c$  or a sequence of characters [`low`, `high`) to narrow characters of type `char`. The first form returns the narrow character, and the second form stores the characters in the array `dest` (which must be large enough to hold high - low characters), returning high. If a `charT` source character cannot be converted to a narrow character, the first function returns `default`, and the second function stores `default` in `dest` as the narrow version of that character.

```
virtual const charT* do_scan_is (mask m, const
charT* low, const charT* high) const
```

Searches the sequence of characters [`low`, `high`) for the first character that matches  $m$ , that is, for which `do_is(m, c)` is true. The return value is a pointer to the first matching character, or `high` if no characters match  $m$ .

```
virtual const charT* do_scan_not (mask m, const
charT* low, const charT* high) const
```

Searches the sequence of characters [`low`, `high`) for the first character that does not match  $m$ , that is, for which `do_is(m, c)` is false. The return value is a pointer to the first matching character, or `high` if every character matches  $m$ .

virtual charT **do\_tolower** (charT c) const, virtual  
const charT\* **do\_tolower** (charT\* low, const charT\*  
high) const

Converts a character *c* or a sequence of characters [*low*,  
*high*) to lowercase. The first form returns the lowercase  
version of *c*, or it returns *c* if *c* does not have a lowercase  
counterpart.

The second form modifies the character sequence: each  
character in [*low*, *high*) is replaced by its lowercase  
counterpart; if a character cannot be converted to  
lowercase, it is not touched. The function returns *high*.

virtual charT **do\_toupper** (charT c) const, virtual  
const charT\* **do\_toupper** (charT\* low, const charT\*  
high) const

Converts a character *c* or a sequence of characters [*low*,  
*high*) to uppercase. The first form returns the uppercase  
version of *c*, or it returns *c* if *c* does not have a  
uppercase counterpart.

The second form modifies the character sequence: each  
character in [*low*, *high*) is replaced by its uppercase  
counterpart; if a character cannot be converted to  
uppercase, it is not touched. The function returns *high*.

virtual charT **do\_widen** (char c) const, virtual const  
char\* **do\_widen** (const char\* low, const char\* high,  
charT\* dest) const

Converts a narrow character *c* or a sequence of narrow  
characters [*low*, *high*) to characters of type charT. The  
first form returns the new character, and the second  
form stores the characters in the array *dest* (which must  
be large enough to hold *high* - *low* characters), returning  
*high*.

## See Also

[ctype\\_base class](#), [ctype\\_byname class template](#),  
[locale::facet class](#)

# Name

`ctype<char>` class — Facet for classifying narrow characters

# Synopsis

```
template <>
class ctype<char> : public locale::facet, public ctype_base
{
    ...
public:
    explicit ctype(const mask* tab = 0, bool del = false, size_t
refs = 0);
    static const size_t table_size = . . . ;
    inline bool is(mask m, char c) const;
    inline const char* is(const char* low, const char* high, mask*
vec) const;
    inline const char* scan_is(mask m, const char* low, const char*
high) const;
    inline const char* scan_not(mask m, const char* low, const
char* high) const;
protected:
    virtual ~ctype( );
    inline const mask* table( ) const throw( );
    inline static const mask* classic_table( ) throw( );
};
```

The `ctype<>` class template is specialized for type `char` (but not `signed char` or `unsigned char`) so the member functions can be implemented as inline functions. The standard requires the implementation to have the protected member functions `table` and `classic_table`. Each of these functions returns an array of mask values indexed by characters cast to `unsigned char`. The number of elements in a table must be at least `table_size`, which is an implementation-defined constant value.

The following are the key member functions:

```
explicit ctype (const mask* tab = 0, bool del =
false, size_t refs = 0)
```



Initializes the `table( )` pointer with `tab`. If `tab` is a null pointer, `table( )` is set to `classic_table( )`. If `tab` is not null, and `del` is true, the `ctype` object owns the table, and when the `ctype` destructor is called, it will delete the table. The `refs` parameter is passed to the base class, as with any facet.

`virtual ~ctype ( )`

If the constructor's `del` flag was true, and `tab` was not a null pointer, performs `delete[] tab`.

```
inline bool is (mask m, charT c) const, inline  
const charT* is (const charT* low, const charT* high,  
mask* dest) const
```

Tests character classifications. The first form returns:

```
(table( )[static_cast<unsigned char>(c)] & m) != 0
```

The second form stores the following in `dest` for each element `c` of the range `[low, high)`:

```
table( )[static_cast<unsigned char>(c)]
```

Note that `is` does not call `do_is`, so `is` can be implemented as an inline function.

```
inline static const mask* classic_table ( ) throw(  
)
```

Returns a table that corresponds to the "C" locale.

```
inline const char* scan_is (mask m, const char*  
low, const char* high) const
```

Searches the sequence of characters `[low, high)` for the first character that matches `m`, that is, for which `is(m, c)` is true. The return value is a pointer to the first matching character, or `high` if no characters match `m`.

```
inline const char* scan_not (mask m, const char*  
low, const char* high) const
```

Searches the sequence of characters [low, high) for the first character that does not match m, that is, for which is(m, c) is false. The return value is a pointer to the first matching character, or high if every character matches m.

```
inline const mask* table ( ) throw( )
```

Returns the value that was passed to the constructor as the tab parameter, or, if tab was null, classic\_table( ) is returned.

## See Also

[ctype class template](#), [locale::facet class](#), [<cctype>](#), [<cwctype>](#)

# Name

ctype\_base class — Base class for ctype facet

# Synopsis

```
class ctype_base{
public:
    enum mask {
        space, print, cntrl, upper, lower, alpha, digit, punct,
        xdigit,
        alnum=alpha|digit, graph=alnum|punct
    };
};
```

The ctype\_base class is the base class for the ctype and ctype\_byname class templates. It declares the mask enumerated type, which is used for classifying characters. [Table 13-20](#) describes the mask literals and their definitions for the classic "C" locale.

*Table 13-20. mask literals for classifying characters*

Literal	Description	"C" locale
alpha	Alphabetic (a letter)	lower OR upper
alnum	Alphanumeric (letter or digit)	alpha OR digit
cntrl	Control (nonprintable)	Not print
digit	'0'-'9'	All locales
graph	Character that occupies graphical space	print but not space
lower	Lowercase letter	'a'-'z'

Literal	Description	"C" locale
print	Printable character (alphanumeric, punctuation, space, etc.)	Depends on character set; in ASCII: '\x20'-'\x7e')
space	Whitespace	' ', '\f', '\n', '\r', '\t', '\v'
upper	Uppercase letter	'A' - 'Z'
xdigit	Hexadecimal digit ('0'-'9', 'a'-'f', 'A'-'F')	All locales

## See Also

[ctype class template](#), [ctype\\_byname class template](#)

## Name

`ctype_byname` class template — Facet for classifying characters

## Synopsis

```
template <typename charT>
class ctype_byname : public ctype<charT>
{
public:
    typedef ctype<charT>::mask mask;
    explicit ctype_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in ctype
};
```

The `ctype_byname` class template is a facet for classifying characters; it uses a named locale. The `ctype_byname<char>` and `ctype_byname<wchar_t>` instantiations are standard.

## See Also

[ctype class template](#), [ctype\\_byname<char> class](#)

## Name

`ctype_byname<char>` class — Facet for classifying narrow characters

## Synopsis

```
template <>
class ctype_byname<char> : public ctype<char>
{
public:
    explicit ctype_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in ctype<char>
};
```

The `ctype_byname<char>` class specializes the `ctype_byname` template for type `char`. (No specialization exists for signed `char` and unsigned `char`.) It derives from `ctype<char>`, so it inherits its table-driven implementation.

## See Also

[ctype<char> class](#), [ctype\\_byname class template](#)

## Name

`has_facet` function template — Test for existence of a facet in a locale

## Synopsis

```
template <typename Facet>
bool has_facet(const locale& loc) throw( );
```

The `has_facet` function determines whether the locale `loc` supports the facet `Facet`. It returns `true` if the facet is supported or `false` if it is not. Call `has_facet` to determine whether a locale supports a user-defined facet. (Every locale must support the standard facets that are described in this section.) [Example 13-24](#) shows how `has_facet` is used.

## Example

*Example 13-24. Testing whether a locale supports a facet*

```
// The units facet is defined under the locale::facet class
// (later in this
// section).

using std::locale;
if (std::has_facet<units>(locale( ))) {
    // Get a reference to the units facet of the locale.
    const units& u = std::use_facet<units>(locale( ));
    // Construct a value of 42 cm.
    units::value_t len = u.make(42, units::cm);
    // Print the length (42 cm) in the locale's preferred units.
    u.length_put(std::cout, len);
}
```

## See Also

[locale class](#), [use\\_facet function template](#)

## Name

isalnum function template — Determines whether a character is alphanumeric in a locale

## Synopsis

```
template <typename charT>  
bool isalnum(charT c, const locale& loc);
```

The `isalnum` function determines whether the character `c` is an alphanumeric character in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::alnum, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isalpha function template](#), [isdigit function template](#)



## Name

isalpha function template — Determines whether a character is a letter in a locale

## Synopsis

```
template <typename charT>  
bool isalpha(charT c, const locale& loc);
```

The `isalpha` function determines whether the character `c` is a letter in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::alpha, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isalnum function template](#), [islower function template](#), [isupper function template](#)

## Name

isctrl function template — Determines whether a character is a control character in a locale

## Synopsis

```
template <typename charT>  
bool isctrl(charT c, const locale& loc);
```

The `isctrl` function determines whether the character `c` is a control character in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::ctrl, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isprint function template](#)

## Name

isdigit function template — Determines whether a character is a digit in a locale

## Synopsis

```
template <typename charT>  
bool isdigit(charT c, const locale& loc);
```

The `isdigit` function determines whether the character `c` is a digit in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::digit, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isalnum function template](#), [isxdigit function template](#)

## Name

isgraph function template — Determines whether a character is graphical in a locale

## Synopsis

```
template <typename charT>  
bool isgraph(charT c, const locale& loc);
```

The `isgraph` function determines whether the character `c` is graphical (alphanumeric or punctuation) in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::graph, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isalnum function template](#), [isprint function template](#), [ispunct function template](#)

## Name

`islower` function template — Determines whether a character is lowercase in a locale

## Synopsis

```
template <typename charT>
bool islower(charT c, const locale& loc);
```

The `islower` function determines whether the character `c` is a lowercase letter in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::lower, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isalpha function template](#), [isupper function template](#)

## Name

isprint function template — Determines whether a character is printable in a locale

## Synopsis

```
template <typename charT>  
bool isprint(charT c, const locale& loc);
```

The `isprint` function determines whether the character `c` is printable in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::print, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [iscntrl function template](#), [isgraph function template](#), [isspace function template](#)

## Name

ispunct function template — Determines whether a character is punctuation in a locale

## Synopsis

```
template <typename charT>  
bool ispunct(charT c, const locale& loc);
```

The `ispunct` function determines whether the character `c` is punctuation in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::punct, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isgraph function template](#)

## Name

isspace function template — Determines whether a character is whitespace in a locale

## Synopsis

```
template <typename charT>  
bool isspace(charT c, const locale& loc);
```

The `isspace` function determines whether the character `c` is whitespace in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::space, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isgraph function template](#), [isprint function template](#)



## Name

isupper function template — Determines whether a character is uppercase in a locale

## Synopsis

```
template <typename charT>  
bool isupper(charT c, const locale& loc);
```

The `isupper` function determines whether the character `c` is an uppercase letter in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::upper, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isalpha function template](#), [islower function template](#)

## Name

isxdigit function template — Determines whether a character is a hexadecimal digit in a locale

## Synopsis

```
template <typename charT>  
bool isxdigit(charT c, const locale& loc);
```

The `isxdigit` function determines whether the character `c` is a hexadecimal digit in the locale `loc`. It returns the following:

```
use_facet<ctype<charT> >(loc).is(ctype_base::xdigit, c)
```

## See Also

[ctype\\_base class](#), [ctype class template](#), [isdigit function template](#)

# Name

locale class — Represents a locale as a set of facets

## Synopsis

```
class locale{
public:
    class facet;
    class id;
    typedef int category;
    static const category
        none, collate, ctype, monetary, numeric, time, messages,
        all = collate|ctype|monetary|numeric|time|messages;
    // Construct/copy/destroy
    locale( ) throw( );
    locale(const locale& other) throw( );
    explicit locale(const char* std_name);
    locale(const locale& other, const char* std_name,category);
    template <typename Facet>
    locale(const locale& other, Facet* f);
    locale(const locale& other, const locale& one, category);
    ~locale( ) throw( );

    const locale& operator=(const locale& other) throw( );
    template <typename Facet>
    locale combine(const locale& other) const;

    basic_string<char> name( ) const;
    bool operator==(const locale& other) const;
    bool operator!=(const locale& other) const;
    template <typename charT, typename Traits, typename Alloc>
    bool operator( )(const basic_string<charT,Traits,Alloc>& s1,
                    const basic_string<charT,Traits,Alloc>& s2)
const;

    static locale global(const locale&);
    static const locale& classic( );
};
```

The locale class template represents the information for a locale. This information is stored as a set of facets. Several facets are defined by the C++ standard, and user-defined facets can be added to any locale. The `has_facet` function

template tests whether a locale supports a particular facet. The `use_facet` function template retrieves a facet of a locale.

References to a facet are safe until all locale objects that use that facet have been destroyed. New locales can be created from existing locales, with modifications to a particular facet.

Some locales have names. A locale can be constructed for a standard name, or a named locale can be copied or combined with other named locales to produce a new named locale.

### **Not e**

When interacting with the user, either through the standard I/O streams or through a graphical user interface, you should use the native locale, that is, `locale( "" )`. When performing I/O—especially to external files where the data must be portable to other programs, systems, or environments—always use the "C" locale (`locale::classic( )`).

[Example 13-25](#) shows locales that control input and output formats.

## **Example**

*Example 13-25. Using locales for input and output*

```
// Open a file and read floating-point numbers from it, computing
the mean.
// Return the mean or 0 if the file contains no data. The data is
in the classic
// format, that is, the same format used by C++.
double mean(const char* filename)
{
    std::ifstream in(filename);
    // Force the datafile to be interpreted in the classic locale,
    so the same
```

```

// datafile can be used everywhere.
in.imbue(std::locale::classic( ));
double sum = 0;
unsigned long count = 0;
std::istream_iterator<double> iter(in), end;
for ( ; iter != end; ++iter) {
    ++count;
    sum += *iter;
}
return count == 0 ? 0.0 : sum / count;
}

int main( )
{
    // Print results in the user's native locale.
    std::cout.imbue(std::locale(""));
    std::cout << mean("data.txt") << '\n';
}

```

The following are the member functions of `locale`:

**locale** ( ) throw( )

Initializes the locale with a copy of the current global locale. The initial global locale is `locale::classic( )`.

**locale** (const locale& other) throw( )

Copies the other locale.



**explicit locale** (const char\* std\_name)

Initializes the locale using a standard name. The names "C" and "" (empty string) are always defined, in which "C" is the locale returned by the `classic( )` function, and "" identifies the implementation-defined native locale.

An implementation can define additional names. Many C++ implementations use ISO language codes and country codes to identify a locale. For example, the ISO 639 language code for English is "en", and the ISO 3166 country code for the United States is "US", so "en\_US" could identify the locale for U.S. English.

**locale** (const locale& other, const char\* std\_name, category mask)

Copies the locale from other, except for those categories identified by mask, which are copied from the standard locale identified by std\_name. The new locale has a name only if other has a name.

template <typename Facet>, **locale** (const locale& other, Facet\* f)

Copies the locale from other except for the facet Facet, which is obtained from f if f is not null.

**locale** (const locale& other, const locale& one, category mask)

Copies the locale from other except for those categories identified by mask, which are copied from one. The new locale has a name only if other and one have names.

template <typename Facet>, locale **combine** (const locale& other) const

Returns a new locale that is a copy of \*this, except for Facet, which is copied from other. If other does not support the facet—that is, has\_facet<Facet>(other) is false—runtime\_error is thrown. The new locale has no name.



**basic\_string<char> name** ( ) const

Returns the locale's name or "\*" if the locale has no name. The exact contents of the name string are implementation-defined, but you can use the string to construct a new locale that is equal to \*this—that is, \*this == locale(name( ).c\_str( )).

const locale& **operator=** (const locale& other)  
throw( )

Copies other and returns \*this.

bool **operator==** (const locale& other) const

Returns true if the two locales are the same object, one locale object is a copy of the other, or the two locales are named and have the same name. Otherwise, the function returns false.

bool **operator!=** (const locale& other) const

Returns !(\*this == other).

template <typename charT, typename Tr, typename A>, bool **operator**( ) (const basic\_string<charT, Tr, A>& s1, const basic\_string<charT, Tr, A>& s2) const

Compares two strings using the collate<charT> facet and returns true if s1 < s2. You can use the locale object as a comparator predicate to compare strings. See <string> for more information.



static locale **global** (const locale& loc)

Sets the global locale to loc and returns the previous global locale. If the new locale has a name, the C locale is set by calling setlocale(LC\_ALL, loc.name( ).c\_str( )); if the locale does not have a name, the effect on the C locale is implementation-defined.

static const locale& **classic** ( )

Returns a locale that implements the "C" locale.

## See Also

[has\\_facet function template](#), [use\\_facet function template](#),  
setlocale in [<locale>](#), [<string>](#)

# Name

locale::category type — Bitmask of facet categories

## Synopsis

```
typedef int category;  
static const category  
    none, collate, ctype, monetary, numeric, time, messages,  
    all = collate|ctype|monetary|numeric|time|messages;
```

The `category` type is an `int` and represents a bitmask of category identifiers, as listed in [Table 13-21](#). Each category represents a set of one or more related facets. When combining locales, you can copy all the facets in one or more categories. Category identifiers can be combined using bitwise operators.

*Table 13-21. Standard categories and their facets*

Literal	Facets
collate	collate<char> collate<wchar_t>
ctype	ctype<char> ctype<wchar_t> codecvt<char, char, mbstate_t> codecvt<wchar_t, char, mbstate_t>
messages	messages<char> messages<wchar_t>



Literal	Facets
monetary	money_get<char> money_get<wchar_t> money_put<char> money_put<wchar_t> moneypunct<char> moneypunct<wchar_t> moneypunct<char,true> moneypunct<wchar_t,true>
numeric	num_get<char> num_get<wchar_t> num_put<char> num_put<wchar_t> numpunct<char> numpunct<wchar_t>
time	time_get<char> time_get<wchar_t> time_put<char> time_put<wchar_t>

## See Also

[locale class](#)

# Name

locale::facet class — Base class for locale facets

## Synopsis

```
class locale::facet{
protected:
    explicit facet(size_t refs = 0);
    virtual ~facet( );
private:
    facet(const facet&);           // Not defined
    void operator=(const facet&); // Not defined
};
```

The facet class is the base class for all facets. A derived class must also declare a public, static data member of type `locale::id` whose name is `id`. Even a derived class must declare its own `id` member because it must have an identifier that is distinct from that of the base-class facet. Any other members for a custom facet are entirely up to the programmer; the derived class does not need to provide a copy or default constructor or an assignment operator.

The `locale` class assigns a value to `id` when the facet object is added to a locale. You never need to examine or alter the `id` member; it is managed entirely by `locale`.

The explicit constructor for `facet` takes a single argument, `ref`. If `ref == 0`, the facet object is not deleted until the last locale that uses the facet is destroyed. If `ref == 1`, the facet object is never destroyed. The standard facet classes (`ctype`, etc.) also take a `ref` parameter and pass it directly to the inherited facet constructor. Custom facets can do whatever the programmer wants, such as relying on the default value of `0` to manage the facet's lifetime automatically.

For example, suppose you want to define a facet that captures the locale-specific preferences for units of

measure, such as length and weight. A program can store and manipulate values in a common base unit and convert to the preferred unit for output. [Example 13-26](#) shows a units facet that allows you to do these things.

## Example

*Example 13-26. A simple facet for working with units of measure*

```
class units : public std::locale::facet
{
public:
    enum length_units {
        length_base=1,
        mm=10, cm=10*mm, m=10*cm, km=1000*m,
        in=254, ft=12*in, yd=3*ft, mi=5280*ft };

    typedef double value_t;

    // All facets must have a static ID member.
    static std::locale::id id;

    // Constructor initializes length_units_ according to local
    preferences.
    units( );

    // Read a length and its units, and return the length in base
    units.
    value_t length_get(std::istream& stream) const;

    // Convert value to the preferred units, and print the
    converted value followed
    // by the unit name.
    void length_put(std::ostream& stream, value_t value) const;

    // Make a base unit value from a value in src_units.
    value_t make(value_t src_value, length_units src_units)
        const;
    // Convert base units to dst_unit.
    value_t convert(value_t src_value, length_units dst_units)
        const;
    // Return the name of a unit.
    const char* unit_name(length_units units) const;
    // Return the preferred unit for length.
```

```

    length_units get_length_unit( ) const;
private:
    length_units length_units_;
};

int main( )
{
    // Add the units facet to the global locale:
    // 1. Construct a new locale that is a copy of the global
locale, with the new
    //    units facet added to it.
    // 2. Set the new locale as the global locale.
    std::locale loc(std::locale(std::locale( )), new units));
    std::locale::global(loc);

    // Now anyone can get the units facet from the global locale.
    const units& u = std::use_facet<units>(std::locale( ));
    units::value_t size = u.make(42, units::cm);
    u.length_put(std::cout, size);
}

```

## See Also

[locale class](#), [locale::id class](#)

# Name

locale::id class — Facet identification

# Synopsis

```
class locale::id{
public:
    id( );
private:
    void operator=(const id&); // Not defined
    id(const id&);             // Not defined
};
```

The `id` class identifies a facet. It is used only to declare a public, static member of type `locale::id` in every facet class. See `locale::facet` for more information.

# See Also

[locale::facet class](#)

# Name

messages class template — Facet for retrieving strings from a message catalog

# Synopsis

```
template <typename charT>
class messages : public locale::facet, public messages_base
{
public:
    typedef charT char_type;
    typedef basic_string<charT> string_type;
    explicit messages(size_t refs = 0);
    catalog open(const basic_string<char>& fn, const locale&
const;
    string_type get(catalog c, int set, int msgid, const
string_type& dfault)
        const;
    void close(catalog c) const;
    static locale::id id;
protected:
    virtual ~messages( );
    virtual catalog do_open(const basic_string<char>&, const
locale&) const;
    virtual string_type do_get(catalog, int set, int msgid,
                                const string_type& dfault) const;
    virtual void do_close(catalog) const;
};
```

The `messages` class template is a facet for a message catalog. A message catalog is a database of textual messages that can be translated into different languages. The `messages<char>` and `messages<wchar_t>` instantiations are standard.



How a message catalog is found is implementation-defined. For example, a catalog name could be the name of an external file, or it could be the name of a special resource section in the program's executable file. The mapping of

message identifiers to a particular message is also implementation-defined.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `get`, which calls `do_get`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual-function description, there is a corresponding description for a public, nonvirtual function, such as:

```
void close (catalog cat) const
```

Calls `do_close(cat)`.

The following are the virtual, protected members of messages:

```
virtual void do_close (catalog cat) const
```

Closes the message catalog `cat`.

```
virtual string_type do_get (catalog cat, int set, int msgid, const string_type& default) const
```

Gets a message that is identified by `set`, `msgid`, and `default` from catalog `cat`. If the message cannot be found, `default` is returned.

```
virtual catalog do_open (const basic_string<char>& name, const locale& loc) const
```

Opens a message catalog `name`. If the catalog cannot be opened, a negative value is returned. Otherwise, the catalog value can be passed to `get` to retrieve messages. Call `close` to close the catalog.

## See Also

[messages\\_base class](#), [messages\\_byname class template](#)



## Name

messages\_base class — Base class for message facets

## Synopsis

```
class messages_base {  
public:  
    typedef int catalog;  
};
```

The message\_base class is the base class for messages and message\_byname. It declares the catalog type, which stores a handle for an open message catalog.

## See Also

[messages class template](#), [messages\\_byname class template](#)

## Name

messages\_byname class template — Facet for retrieving strings from a message catalog

## Synopsis

```
template <typename charT>
class messages_byname : public messages<charT>
{
public:
    typedef messages_base::catalog catalog;
    typedef basic_string<charT> string_type;
    explicit messages_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in messages
};
```

The messages\_byname class template is a facet for a message catalog; it uses a named locale. The messages\_byname<char> and messages\_byname<wchar\_t> instantiations are standard.

## See Also

[messages class template](#), [messages\\_base class](#)

## Name

money\_base class — Base class for moneypunct facet

## Synopsis

```
class money_base {  
public:  
    enum part { none, space, symbol, sign, value };  
    struct pattern { char field[4]; };  
};
```

The money\_base class is a base class for the moneypunct and moneypunct\_byname class templates. It declares the part and pattern types. A pattern actually stores four part values, but they are stored as four char values for space efficiency. See moneypunct for an explanation of how part and pattern are used.

## See Also

[moneypunct class template](#), [moneypunct\\_byname class template](#)

# Name

money\_get class template — Facet for input of monetary values

# Synopsis

```
template <typename charT,  
          typename InputIterator = istreambuf_iterator<charT> >  
class money_get : public locale::facet  
{  
public:  
    typedef charT char_type;  
    typedef InputIterator iter_type;  
    typedef basic_string<charT> string_type;  
    explicit money_get(size_t refs = 0);  
    iter_type get(iter_type s, iter_type end, bool intl, ios_base&  
f,  
                 ios_base::iostate& err, long double& units)  
const;  
    iter_type get(iter_type s, iter_type end, bool intl, ios_base&  
f,  
                 ios_base::iostate& err, string_type& digits)  
const;  
    static locale::id id;  
protected:  
    virtual ~money_get( );  
    virtual iter_type do_get(iter_type begin, iter_type end, bool  
intl,  
                             ios_base& stream, ios_base::iostate&  
err,  
                             long double& units) const;| virtual  
iter_type  
                             do_get(iter_type begin, iter_type end, bool  
intl,  
                             ios_base& stream, ios_base::iostate&  
err,  
                             string_type& digits) const;  
};
```

The `money_get` class template is a facet for parsing monetary values from an input stream. The `money_get<char>` and `money_get<wchar_t>` instantiations

are standard. [Example 13-27](#) shows a simple use of `money_get` and `money_put`.

## Example

*Example 13-27. Reading and writing monetary values*

```
#include <iostream>
#include <locale>
#include <ostream>

int main( )
{
    std::ios_base::iostate err = std::ios_base::goodbit;
    long double value;
    std::cout << "What is your hourly wage? ";
    std::use_facet<std::money_get<char> >(std::locale( )).get(
        std::cin, std::istreambuf_iterator<char>( ),
        false, std::cin, err, value);
    if (err)
        std::cerr << "Invalid input\n";
    else {
        std::cout << value << '\n';
        std::cout << "You make ";
        std::use_facet<std::money_put<char> >(std::locale( )).put(
            std::cout, false, std::cout, std::cout.fill( ),
            value * 40);
        std::cout << " in a 40-hour work week.\n";
    }
}
```

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public function `get`, which calls `do_get`. The description below is for the virtual functions because they do the real work. Imagine that for each virtual-function description, there is a corresponding description for a public, nonvirtual function, such as:

```
iter_type get (iter_type begin, iter_type end, bool
intl, ios_base& stream, ios_base::iostate& err,
long double& units) const
```

Calls `do_get(begin, end, intl, stream, err, units)`

The following are the virtual, protected members of `money_get`:

```
virtual iter_type do_get (iter_type begin,  
iter_type end, bool intl, ios_base& stream,  
ios_base::iostate& err, long_double& units) const,  
virtual iter_type do_get (iter_type begin,  
iter_type end, bool intl, ios_base& stream,  
ios_base::iostate& err, string_type& digits) const
```

Reads characters in the range `[begin, end)` and interprets them as a monetary value. If `intl` is true, the value is parsed using international format; otherwise, local format is used. That is, the `intl` value is used as the `Intl` template parameter to `moneypunct<char_type, Intl>`. If a valid monetary value is read from the input stream, the integral value is stored in `units` or is formatted as a string in `digits`. (For example, the input "\$1,234.56" yields the units 123456 or the digits "123456".) The digit string starts with an optional minus sign ('-') followed by digits ('0'-'9'), in which each character `c` is produced by calling `c_type<char_type>.widen(c)`.

If a valid sequence is not found, `err` is modified to include `stream.failbit`. If the end of the input is reached without forming a valid monetary value, `stream.eofbit` is also set.

If the `showbase` flag is set (`stream.flags( ) & stream.showbase` is not 0), the currency symbol is required; otherwise, it is optional. Thousands grouping, if the local format supports it, is optional.

The sign of the result is dictated by `positive_sign( )` and `negative_sign( )` from the `moneypunct` facet.

The return value is an iterator that points to one past the last character of the monetary value.

## **See Also**

[money\\_put class template](#), [moneypunct class template](#),  
[num\\_get class template](#)

# Name

money\_put class template — Facet for output of monetary values

# Synopsis

```
template <typename charT,  
    typename OutputIterator = ostreambuf_iterator<charT> >  
class money_put : public locale::facet  
{  
public:  
    typedef charT char_type;  
    typedef OutputIterator iter_type;  
    typedef basic_string<charT> string_type;  
    explicit money_put(size_t refs = 0);  
    iter_type put(iter_type s, bool intl, ios_base& f, char_type  
fill,  
                long double units) const;  
    iter_type put(iter_type s, bool intl, ios_base& f, char_type  
fill,  
                const string_type& digits) const;  
    static locale::id id;  
protected:  
    virtual ~money_put( );  
    virtual iter_type do_put(iter_type, bool, ios_base&, char_type  
fill,  
                            long double units) const;  
    virtual iter_type do_put(iter_type, bool, ios_base&, char_type  
fill,  
                            const string_type& digits) const;  
};
```

The `money_put` class template is a facet for formatting and printing monetary values. See [Example 13-27](#) (under `money_get`), which shows how to use the `money_put` facet. The `money_put<char>` and `money_put<wchar_t>` instantiations are standard.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `put`, which calls `do_put`. The descriptions below are for the



virtual functions because they do the real work. Imagine that for each virtual function description, there is a corresponding description for a public, nonvirtual function, such as:

```
iter_type put (iter_type iter, bool intl, ios_base&
stream, char_type fill, long double units) const
```

Returns `do_put(iter, intl, stream, fill, units)`

The following are the virtual, protected members of `money_put`:

```
virtual iter_type do_put (iter_type iter, bool intl,
ios_base& stream, char_type fill, long double units)
const, virtual iter_type do_put (iter_type iter, bool
intl, ios_base& stream, char_type fill, const
string_type& digits) const
```

Formats a monetary value and writes the formatted characters to `iter`. The value to format is either an integer, `units`, or a string of digit characters in `digits`. If the first character of `digits` is `widen(' -')`, the remaining digits are interpreted as a negative number.

The formatting pattern and punctuation characters are obtained from the `money_punct` facet. For positive values, `pos_format( )` is used; for negative values, `neg_format( )` is used. The pattern dictates the output format. (See `money_punct` later in this section for information on patterns.) The currency symbol is printed only if the `showbase` flag is set (that is, `stream.flags( ) & stream.showbase` is nonzero). Thousands separators and a decimal point are inserted at the appropriate places in the formatted output.

If necessary, `fill` characters are inserted until the formatted width is `stream.width( )`. The stream's

adjustfield flag dictates how fill characters are inserted. That is, `stream.flags( ) & stream.adjustfield` is tested, and if it is equal to:

`ios_base::internal`

Fill characters are inserted where the pattern is none or space.

`ios_base::left`

Fill characters are appended to the end of the formatted field.

None of the above

Fill characters are inserted at the start of the formatted field.

Finally, `stream.width(0)` is called to reset the field width to 0. The return value is an iterator that points to one past the last output character.

## See Also

[money\\_get class template](#), [moneypunct class template](#), [num\\_put class template](#)

# Name

money\_punct class template — Facet for punctuation of monetary values

# Synopsis

```
template <typename charT, bool International = false>
class money_punct : public locale::facet, public money_base
{
public:
    typedef charT char_type;
    typedef basic_string<charT> string_type;
    explicit money_punct(size_t refs = 0);
    charT decimal_point( ) const;
    charT thousands_sep( ) const;
    string grouping( ) const;
    string_type curr_symbol( ) const;
    string_type positive_sign( ) const;
    string_type negative_sign( ) const;
    int frac_digits( ) const;
    pattern pos_format( ) const;
    pattern neg_format( ) const;
    static locale::id id;
    static const bool intl = International;
protected:
    virtual ~money_punct( );
    virtual charT do_decimal_point( ) const;
    virtual charT do_thousands_sep( ) const;
    virtual string do_grouping( ) const;
    virtual string_type do_curr_symbol( ) const;
    virtual string_type do_positive_sign( ) const;
    virtual string_type do_negative_sign( ) const;
    virtual int do_frac_digits( ) const;
    virtual pattern do_pos_format( ) const;
    virtual pattern do_neg_format( ) const;
};
```

The money\_punct class template is a facet that describes the punctuation characters used to format a monetary value.

The money\_punct<char, false>, money\_punct<wchar\_t, false>, money\_punct<char, true>,

and `moneypunct<wchar_t, true>` instantiations are standard.

Specify `true` for the `International` template parameter to obtain an international format, or `false` to obtain a local format. In an international format, the currency symbol is always four characters, usually three characters followed by a space.

The `money_get` and `money_put` facets use a pattern to parse or format a monetary value. The pattern specifies the order in which parts of a monetary value must appear. Each pattern has four fields, in which each field has type `part` (cast to `char`). The `symbol`, `sign`, and `value` parts must appear exactly once, and the remaining field must be `space` or `none`. The value `none` cannot be first (`field[0]`); `space` cannot be first or last (`field[3]`).

Where `sign` appears in the pattern, the first character of the sign string (`positive_sign( )` or `negative_sign( )`) is output, and the remaining characters of the sign string appear at the end of the formatted output. Thus, if `negative_sign( )` returns `"( )"`, the value `-12.34` might be formatted as `"(12.34)"`.

As with other facets, the public members call `virtual`, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `grouping`, which calls `do_grouping`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual-function description, there is a corresponding description for a public, nonvirtual function, such as:

```
string_type curr_symbol ( ) const
```

```
Returns do_curr_symbol( )
```

The following are the virtual, protected members of `money_punct`:

`virtual string_type do_curr_symbol ( ) const`

Returns the currency symbol, such as "\$" (which is used by some U.S. locales) or "USD " (which is the international currency symbol for U.S. dollars). In the "C" locale, the currency symbol is "" or "L".

`virtual charT do_decimal_point ( ) const`

Returns the character used before the fractional part when `do_frac_digits` is greater than 0. For example, in the U.S., this is typically '.', and in Europe, it is typically ','. In the "C" locale, the decimal point is '.' or "L'.".

`virtual int do_frac_digits ( ) const`

Returns the number of digits to print after the decimal point. This value can be 0. In the "C" locale, the number of digits is `std::numeric_limits<char>.max( )`.

`virtual string do_grouping ( ) const`

Returns a string that specifies the positions of thousands separators. The string is interpreted as a vector of integers, in which each value is a number of digits, starting from the right. Thus, the string "\3" means every three digits form a group. In the "C" locale, the grouping is "" or "L".

`virtual pattern do_neg_format ( ) const`

Returns the pattern used to format negative values. In the "C" locale, the negative format is { symbol, sign, none, value }.

`virtual string_type do_negative_sign ( ) const`

Returns the string (which may be empty) used to identify negative values. The position of the sign is dictated by

the `do_neg_format` pattern. In the "C" locale, the negative sign is "-" or L"-".

virtual pattern **do\_pos\_format** ( ) const

Returns the pattern used to format positive values. In the "C" locale, the positive format is { symbol, sign, none, value }.

virtual string\_type **do\_positive\_sign** ( ) const

Returns the string (which may be empty) used to identify positive values. The position of the sign is dictated by the `do_pos_format` pattern. In the "C" locale, the positive sign is "" or L"".

virtual charT **do\_thousands\_sep** ( ) const

Returns the character used to separate groups of digits, in which the groups are specified by `do_grouping`. In the U.S., the separator is typically ',', and in Europe, it is often '.'. In the "C" locale, the thousands separator is '\0' or L'\0'.

## See Also

[money\\_base class](#), [money\\_get class template](#), [money\\_put class template](#), [moneypunct\\_byname class template](#), [numpunct class template](#)

## Name

money\_punct\_byname class template — Facet for punctuation of monetary values

## Synopsis

```
template <typename charT, bool Intl = false>
class money_punct_byname : public money_punct<charT, Intl>
{
public:
    typedef money_base::pattern pattern;
    typedef basic_string<charT> string_type;
    explicit money_punct_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in money_punct
};
```

The money\_punct\_byname class template provides formatting characters and information for monetary values using the rules of a named locale. The money\_punct\_byname<char, International> and money\_punct\_byname<wchar\_t, International> instantiations are standard.

## See Also

[money\\_punct class template](#)

# Name

num\_get class template — Facet for input of numbers

## Synopsis

```
template <typename charT,  
          typename InputIterator = istreambuf_iterator<charT> >  
class num_get : public locale::facet  
{  
public:  
    typedef charT char_type;  
    typedef InputIterator iter_type;  
    explicit num_get(size_t refs = 0);  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, bool& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, long& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, unsigned short& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, unsigned int& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, unsigned long& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, float& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, double& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, long double& v) const;  
    iter_type get(iter_type in, iter_type end, ios_base&,  
                 ios_base::iostate& err, void*& v) const;  
    static locale::id id;  
protected:  
    virtual ~num_get( );  
    virtual iter_type do_get(iter_type, iter_type, ios_base&,  
                            ios_base::iostate& err, bool& v)  
const;  
    virtual iter_type do_get(iter_type, iter_type, ios_base&,  
                            ios_base::iostate& err, long& v)  
const;  
    virtual iter_type do_get(iter_type, iter_type, ios_base&,  
                            ios_base::iostate& err, unsigned  
short& v) const;  
    virtual iter_type do_get(iter_type, iter_type, ios_base&,
```



```

        ios_base::iostate& err, unsigned int&
v) const;
    virtual iter_type do_get(iter_type, iter_type, ios_base&,
        ios_base::iostate& err, unsigned long&
v) const;
    virtual iter_type do_get(iter_type, iter_type, ios_base&,
        ios_base::iostate& err, float& v)
const;
    virtual iter_type do_get(iter_type, iter_type, ios_base&,
        ios_base::iostate& err, double& v)
const;
    virtual iter_type do_get(iter_type, iter_type, ios_base&,
        ios_base::iostate& err, long double&
v) const;
    virtual iter_type do_get(iter_type, iter_type, ios_base&,
        ios_base::iostate& err, void*& v)
const;
};

```

The `num_get` class template is a facet for parsing and reading numeric values from an input stream. The `istream` extraction operators (`>>`) use `num_get`. The `num_get<char>` and `num_get<wchar_t>` instantiations are standard.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `get`, which calls `do_get`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual function description, there is a corresponding description for a public, nonvirtual function, such as:

```

iter_type get (iter_type begin, iter_type end,
ios_base& stream, ios_base::iostate& err, bool& v)
const

```

Returns `do_get(begin, end, stream, err, v)`

The following are the virtual, protected members of `num_get`:

virtual iter\_type **do\_get** (iter\_type begin, iter\_type end, ios\_base& stream, ios\_base::iostate& err, bool& v) const

Reads a bool value, which can be represented as a number or as a character string. The function first tests the boolalpha flag, that is, stream.flags( ) & stream.boolalpha. If the flag is 0, a numeric value is read; if the flag is 1, a string is read from [begin, end).

If boolalpha is false, the input is interpreted as a long int. If the numeric value is 1, v is assigned true; if the value is 0, v is assigned false; otherwise, failbit is set in err, and v is not modified.

If boolalpha is true, characters are read from begin until one of the following happens:

- The input matches truename( ) from the numpunct facet:

```
use_facet<numpunct<char_type> >(stream.getloc(
)).truename( )
```

v is assigned true, and err is assigned goodbit. A match is determined by the shortest input sequence that uniquely matches truename( ) or falsename( ).

- The input matches falsename( ): v is assigned false, and err is assigned goodbit.
- begin == end, in which case eofbit is set in err.
- The input does not match truename( ) or falsename( ); failbit is set in err.

virtual iter\_type **do\_get** (iter\_type begin, iter\_type end, ios\_base& stream, ios\_base::iostate& err, type & v) const

Reads a single value. The `do_get` function is overloaded for most of the fundamental types. The behavior of each function is essentially the same (except for the `bool` version described earlier) and depends on `stream.flags( )`, the `ctype` facet, and the `numunct` facet. Both facets are obtained for the locale `stream.getloc( )`.

First, input characters are collected from the range `[begin, end)` or until the input character is not part of a valid number according to the flags and `numunct` facet. A locale-dependent decimal point is replaced with the character `'.'`. Thousands separators are read but not checked for valid positions until after the entire number has been read. The set of valid characters depends on the type of `v` and the flags, in particular the `basefield` flags. If `stream.flags( ) & basefield` is `hex`, hexadecimal characters are read; if it is `oct`, only octal characters are read (`'0'-'7'`). If the `basefield` is `0`, the prefix determines the radix: `0x` or `0X` for hexadecimal, `0` for octal, and anything else for decimal. Floating-point numbers can use fixed or exponential notation, regardless of the flags.



If `v` is of type `void*`, the format is implementation-defined in the same manner as the `%p` format for `scanf` (in `<cstdio>`).

After all the valid characters have been read, they are interpreted as a numeric value. If the string is invalid, or if the thousands groupings are incorrect, `failbit` is set in `err` and `v` is not changed. If the string is valid, its numeric value is stored in `v` and `err` is set to `goodbit`. If the entire input stream is read (up to `end`), `eofbit` is set in `err`.

## See Also

[money\\_get class template](#), [num\\_put class template](#),  
[numput class template](#), `basic_istream` in [<istream>](#)

# Name

num\_put class template — Facet for output of numbers

## Synopsis

```
template <typename charT,  
          typename OutputIterator = ostreambuf_iterator<charT> >  
class num_put : public locale::facet  
{  
public:  
    typedef charT char_type;  
    typedef OutputIterator iter_type;  
    explicit num_put(size_t refs = 0);  
    iter_type put(iter_type s, ios_base& f, char_type fill, bool v)  
const;  
    iter_type put(iter_type s, ios_base& f, char_type fill, long v)  
const;  
    iter_type put(iter_type s, ios_base& f, char_type fill,  
unsigned long v) const;  
    iter_type put(iter_type s, ios_base& f, char_type fill, double  
v) const;  
    iter_type put(iter_type s, ios_base& f, char_type fill, long  
double v) const;  
    iter_type put(iter_type s, ios_base& f, char_type fill, const  
void* v) const;  
    static locale::id id;  
protected:  
    virtual ~num_put( );  
    virtual iter_type do_put(iter_type, ios_base&, char_type fill,  
bool v) const;  
    virtual iter_type do_put(iter_type, ios_base&, char_type fill,  
long v) const;  
    virtual iter_type do_put(iter_type, ios_base&, char_type fill,  
unsigned long)  
const;  
    virtual iter_type do_put(iter_type, ios_base&, char_type fill,  
double v) const;  
    virtual iter_type do_put(iter_type, ios_base&, char_type fill,  
long double v)  
const;  
    virtual iter_type do_put(iter_type, ios_base&, char_type fill,  
const void* v)  
const;  
};
```

The `num_put` class template is a facet for formatting and outputting a numeric value. The ostream output operators (`>>`) use `num_put`. The `num_put<char>` and `num_put<wchar_t>` instantiations are standard.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `put`, which calls `do_put`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual function description, there is a corresponding description for a public, nonvirtual function, such as:

```
iter_type put (iter_type out, ios_base& stream,  
char_type fill, bool v) const
```

Returns `do_put(out, stream, fill, v)`

The following are the virtual, protected members of `num_put`:

```
virtual iter_type do_put (iter_type out, ios_base&  
stream, char_type fill, bool v) const
```

Writes a `bool` value to `out`. If the `boolalpha` flag is clear, that is, `stream.flags( ) & stream.boolalpha` is `0`, the integer value of `v` is written as a number. If `boolalpha` is set, `v` is written as a word: if `v` is true, `truename( )` is written; if `v` is false, `falsename( )` is written using the `num_punct` facet. For example:

```
const num_punct<charT>& n = use_facet<num_punct<charT> >;  
string_type s = v ? n.truename( ) : n.falsename( );  
// Write characters of s to out.
```

```
virtual iter_type do_put (iter_type out, ios_base&  
stream, char_type fill, type v) const
```

Formats `v` as a string and writes the string contents to out using the flags of `stream` to control the formatting and the imbued locale of `stream` to obtain the `ctype` and `numpunct` facets for punctuation rules and characters. The format also depends on *type*:

Integral types (long, unsigned long)

The format depends on the `basefield` flags (`stream.flags( ) & basefield`). If `oct`, the number is formatted as octal; if `hex`, the number is formatted as hexadecimal (using 'a'-'f' for the digits 10-16, or 'A'-'F' if the uppercase flag is set); or else the number is decimal. If the `showbase` flag is set, a prefix is used: `0` for octal, or `0x` or `0X` for hexadecimal (depending on the uppercase flag).

Floating-point types (double, long double)

The format depends on the `floatfield` flags (`stream.flags( ) & floatfield`). If `fixed`, the format is fixed-point: an integer part, a decimal point, and a fractional part. If `scientific`, the format is exponential.

If the `floatfield` flags do not indicate `fixed` or `scientific`, the general format is used: exponential if the exponent is -4 or less or greater than the precision (number of places after the decimal point), or `fixed` otherwise. Trailing zeros are dropped, as is the decimal point if would be the last character in the string.

If the uppercase flag is set, the exponent is introduced by 'E', or else by 'e'. If the `showpoint` flag is set, the decimal point is always present.

 Pointer (`void*`)

The output format for a pointer is implementation-defined.

If the number is negative, it is prefaced with a minus sign ('-'). If the number is positive, no sign is output unless the `showpos` flag is set, in which case a positive sign ('+') appears at the start of the string.

If a decimal point character is needed, it is obtained from the `numpunct` facet's `decimal_point( )` function. Integers have thousands separators inserted according to the `grouping( )` function. See `numpunct` later in this section for more information.

If necessary, fill characters are inserted until the formatted width is `stream.width( )`. The stream's `adjustfield` flag dictates how fill characters are inserted. That is, `stream.flags( ) & stream.adjustfield` is tested, and if it is equal to:

`ios_base::internal`

Fill characters are inserted after a sign (if present) or, if there is no sign, after a leading `0x` or `0X`, or else at the start of the field.

`ios_base::left`

Fill characters are appended to the end of the formatted field.

Any other value

Fill characters are inserted at the start of the formatted field.

Finally, `stream.width(0)` is called to reset the field width to `0`. The return value is an iterator that points to one past the last output character.

## See Also



[money\\_put class template](#), [num\\_get class template](#),  
[numput class template](#), basic\_ostream in [<ostream>](#)

# Name

numpunct class template — Facet for punctuation of numbers

## Synopsis

```
template <typename charT>
class numpunct : public locale::facet
{
public:
    typedef charT char_type;
    typedef basic_string<charT> string_type;
    explicit numpunct(size_t refs = 0);
    char_type decimal_point( ) const;
    char_type thousands_sep( ) const;
    string grouping( ) const;
    string_type truename( ) const;
    string_type falsename( ) const;
    static locale::id id;
protected:
    virtual ~numpunct( );
    virtual char_type do_decimal_point( ) const;
    virtual char_type do_thousands_sep( ) const;
    virtual string do_grouping( ) const;
    virtual string_type do_truename( ) const;
    virtual string_type do_falsename( ) const;
};
```

The `numpunct` class template is a facet for numeric formatting and punctuation. The `num_get` and `num_put` facets use `numpunct`. The `numpunct<char>` and `numpunct<wchar_t>` instantiations are standard.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `grouping`, which calls `do_grouping`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual-function description, there is a corresponding description for a public, nonvirtual function, such as:

`char_type decimal_point ( ) const`

Returns `do_decimal_point( )`

The following are the virtual, protected members of `num_punct`:

`virtual char_type do_decimal_point ( ) const`

Returns the decimal point character, which is typically `'.'` in U.S. locales and `','` in European locales. In the "C" locale, the decimal point is `'.'` or `L'.'`

`virtual string_type do_falsename ( ) const`

Returns the textual representation for the value `false`. In the standard instantiations (`num_punct<char>` and `num_punct<wchar_t>`), the value is `"false"` or `L"false"`.

`virtual string do_grouping ( ) const`

Returns a string that specifies the positions of thousands separators. The string is interpreted as a vector of integers, in which each value is a number of digits, starting from the right. Thus, the string `"\3"` means every three digits form a group. In the "C" locale, the grouping is `""` or `L""`.

`virtual char_type do_thousands_sep ( ) const`

Returns the character used to separate digit groups. (See `do_grouping` earlier in this section.) In U.S. locales, this is typically `','`, and in European locales, it is typically `'.'`. In the "C" locale, the thousands separator is `'\0'` or `L'\0'`.

`virtual string_type do_truename ( ) const`

Returns the textual representation for the value `true`. In the standard instantiations (`num_punct<char>` and `num_punct<wchar_t>`), the value is `"true"` or `L"true"`.

## See Also

[moneypunct class template](#), [num\\_get class template](#),  
[num\\_put class template](#)

## Name

`numpunct_byname` class template — Facet for punctuation of numbers

## Synopsis

```
template <typename charT>
class numpunct_byname : public numpunct<charT>
{
// This class is specialized for char and wchar_t.
public:
    typedef charT char_type;
    typedef basic_string<charT> string_type;
    explicit numpunct_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in numpunct
};
```

The `numpunct_byname` class template is a facet for numeric formatting and punctuation; it uses the rules of a named locale. The `numpunct_byname<char>` and `numpunct_byname<wchar_t>` instantiations are standard.

## See Also

[numpunct class template](#)

## Name

`time_base` class — Base class for time facets

## Synopsis

```
class time_base {  
public:  
    enum dateorder { no_order, dmy, mdy, ymd, ydm };  
};
```

The `time_base` class is a base class for the `time_get` class template. It declares the `dateorder` type. See `time_get` for more information.

## See Also

[time\\_get class template](#)

# Name

time\_get class template — Facet for input of dates and times

## Synopsis

```
template <typename charT, typename InputIterator =
istreambuf_iterator<charT> >
class time_get : public locale::facet, public time_base
{
public:
    typedef charT char_type;
    typedef InputIterator iter_type;
    explicit time_get(size_t refs = 0);
    dateorder date_order( ) const;
    iter_type get_time(iter_type s, iter_type end, ios_base& f,
        ios_base::iostate& err, tm* t) const;
    iter_type get_date(iter_type s, iter_type end, ios_base& f,
        ios_base::iostate& err, tm* t) const;
    iter_type get_weekday(iter_type s, iter_type end,
        ios_base& f, ios_base::iostate& err, tm*
t) const;
    iter_type get_monthname(iter_type s, iter_type end,
        ios_base& f, ios_base::iostate& err, tm*
t) const;
    iter_type get_year(iter_type s, iter_type end, ios_base& f,
        ios_base::iostate& err, tm* t) const;
    static locale::id id;
protected:
    virtual ~time_get( );
    virtual dateorder do_date_order( ) const;
    virtual iter_type do_get_time(iter_type s, iter_type end,
ios_base&,
        ios_base::iostate& err, tm* t)
const;
    virtual iter_type do_get_date(iter_type s, iter_type end,
ios_base&,
        ios_base::iostate& err, tm* t)
const;
    virtual iter_type do_get_weekday(iter_type s, iter_type end,
ios_base&,
        ios_base::iostate& err, tm* t)
const;
    virtual iter_type do_get_monthname(iter_type s, iter_type end,
```

```

ios_base&,
                                                    ios_base::iostate& err, tm*
t) const;
    virtual iter_type do_get_year(iter_type s, iter_type end,
ios_base&,
                                                    ios_base::iostate& err, tm* t)
const;
};

```

The `time_get` class template is a facet for parsing and reading dates and times from an input stream. The components of the date and time value are stored in a `tm` structure. (See `<ctime>` for more information about `tm`.) The `time_get<char>` and `time_get<wchar_t>` instantiations are standard.

Most of the `time_get` functions take an `err` parameter in much the same way other facets and their functions do. Unlike other facets, however, the `time_get` functions do not set `err` to `goodbit` upon success. Instead, they only set `failbit` for an error. They do not set `eofbit` if the end of the input sequence is reached.

Most of the `time_get` functions take a `t` parameter, which is a pointer to a `tm` object, which is filled in with the relevant parts of the date and time. If a function fails, the state of the `tm` object is undefined.

As with other facets, the public members call virtual, protected members with the same name prefaced by `do_`. Thus, to use the facet, call the public functions, such as `get_date`, which calls `do_get_date`. The descriptions below are for the virtual functions because they do the real work. Imagine that for each virtual-function description, there is a corresponding description for a public, nonvirtual function, such as:

```

dateorder date_order ( ) const
    Returns do_date_order( )


```



The following are the virtual, protected members of `time_get`:

```
virtual dateorder do_date_order ( ) const
```

Returns the order in which the day, month, and year appear in a locale-specific date. If the formatted date includes additional elements, the return value is `no_order`. See the `time_base` class for the declaration of the `dateorder` type.

```
 virtual iter_type do_get_time (iter_type begin,  
iter_type end, ios_base& stream, ios_base::iostate&  
err, tm* t) const
```

Reads characters from `[begin, end)` and interprets them as a time, according to the format of `time_put<>::put`, using the 'X' format. The time elements are stored in `*t`. If the input is invalid, the state of `t`'s members is undefined, and `err` is set to `failbit`. The return value is an iterator that points to one past where the input stopped.

```
virtual iter_type do_get_date (iter_type begin,  
iter_type end, ios_base& stream, ios_base::iostate&  
err, tm* t) const
```

Reads characters from `[begin, end)` and interprets them as a date, according to the format of `time_put<>::put`, using the 'x' format. The date elements are stored in `*t`. If the input is invalid, the state of `t`'s members is undefined, and `err` is set to `failbit`. The return value is an iterator that points to one past where the input stopped.

```
virtual iter_type do_get_weekday (iter_type begin,  
iter_type end, ios_base& stream,  
ios_base::iostate& err, tm* t) const
```

Reads characters from [begin, end) until it reads the name of a day of the week, either abbreviated or spelled out. The appropriate date elements are stored in \*t. If the input is invalid, the state of t's members is undefined, and err is set to failbit. The return value is an iterator that points to one past where the input stopped.

```
virtual iter_type do_get_monthname (iter_type  
begin, iter_type end, ios_base& stream,  
ios_base::iostate& err, tm* t) const
```

Reads characters from [begin, end) until it reads the name of a month, either abbreviated or spelled out. The appropriate date elements are stored in \*t. If the input is invalid, the state of t's members is undefined, and err is set to failbit. The return value is an iterator that points to one past where the input stopped.

```
virtual iter_type do_get_year (iter_type begin,  
iter_type end, ios_base& stream,  
ios_base::iostate& err, tm* t) const
```

Reads characters from [begin, end) until it reads a year. It is up to the implementation to determine whether two-digit years are accepted, and if so, which century to apply to the abbreviated year. The t->tm\_year member is set appropriately. If the input is invalid, the state of t's members is undefined, and err is set to failbit. The return value is an iterator that points to one past where the input stopped.

## See Also

[time\\_base class](#), [time\\_get\\_byname class template](#), [time\\_put class template](#), tm in [<ctime>](#)

## Name

`time_get_byname` class template — Facet for input of dates and times

## Synopsis

```
template <typename charT, typename InputIterator =
istreambuf_iterator<charT> >
class time_get_byname : public time_get<charT, InputIterator>
{
public:
    typedef time_base::dateorder dateorder;
    typedef InputIterator iter_type;
    explicit time_get_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in time_get
};
```

The `time_get_byname` class template is a facet for reading dates and times from an input stream using a named locale. The `time_get_byname<char>` and `time_get_byname<wchar_t>` instantiations are standard.

## See Also

[time\\_get class template](#)

# Name

time\_put class template — Facet for output of dates and times

# Synopsis

```
template <typename charT, typename OutputIterator =
ostreambuf_iterator<charT> >
class time_put : public locale::facet
{
public:
    typedef charT char_type;
    typedef OutputIterator iter_type;
    explicit time_put(size_t refs = 0);
    iter_type put(iter_type s, ios_base& f, char_type fill, const
tm* tmb,
                const charT* pattern, const charT* pat_end)
const;
    iter_type put(iter_type s, ios_base& f, char_type fill, const
tm* tmb,
                char format, char modifier = 0) const;
    static locale::id id;
protected:
    virtual ~time_put( );
    virtual iter_type do_put(iter_type s, ios_base&, char_type,
const tm* t,
                            char format, char modifier) const;
};
```

The `time_put` class template is a facet for formatting and writing dates and times. The `time_put<char>` and `time_put<wchar_t>` instantiations are standard.

Note that `time_put` is unlike other facets. The public `put` function does not always directly call `do_put`. Here are the complete descriptions of `put` and `do_put`:

```
iter_type put (iter_type out, ios_base& stream,
char_type fill, const tm* t, const charT*
pattern, const charT* pat_end) const
```

Reads the pattern in [pattern, pat\_end) and writes formatted date and time information to out. The pattern contains ordinary characters (which are written directly to out) interspersed with format specifiers. A format specifier starts with '%' and is followed by an optional modifier character, which is followed in turn by a format specifier character. The put function checks format characters by first calling narrow from the ctype<charT> facet, then checking the narrowed character.

For each format specifier, put calls do\_put(out, stream, fill, t, format, modifier), in which format is the format specifier and modifier is the modifier character or 0 if no modifier is present.



The use of modifier characters is implementation-defined. The standard does not define any modifiers. See the do\_put member function for more information.

```
iter_type put (iter_type out, ios_base& stream,  
char_type fill, const tm* t, char format, char  
modifier = 0) const
```

Returns do\_put(out, stream, fill, t, format, modifier).

```
virtual iter_type do_put (iter_type out, ios_base&  
stream, char_type fill, const tm* t, char format,  
char modifier) const
```

Formats a single date or time element and writes the formatted characters to out. The format character specifies what to output (as shown in [Table 13-22](#)). The date and time information is obtained from t.

The do\_put function, unlike some of the other output facets, does not use stream's flags or field width. The

fill parameter is used by implementation-defined formatting.

*Table 13-22. Format specifiers for do\_put*

Specified	Description
a	Abbreviated weekday name
A	Full weekday name
b	Abbreviated month name
B	Full month name
C	Complete date and time
D	Day of the month (01-31)
H	Hour (00-23); 24-hour clock
I	Hour (01-12); 12-hour clock
j	Day of the year (001-366)
m	Month (01-12)
M	Minutes (00-59)
P	A.M./P.M. designation for use with a 12-hour clock
S	Second (00-61); up to two leap seconds
U	Week number (00-53); week 1 starts with the first Sunday
w	Weekday (0-6); Sunday is day 0
W	Week number (00-53); week 1 starts with first Monday

Specified	Description
x	Date
X	Time
y	Year in century (00-99)
Y	Year
Z	Time zone name or abbreviation, or empty string if time zone is unknown
%	Literal %

## Notes



The use of `modifier` is implementation-defined. The C++ standard recommends the use of POSIX modifiers, which are 'E' and 'O'. These modifiers request the use of an alternative format if the locale has one. The 'E' modifier applies to certain format specifiers to request an alternative representation for dates and times. The 'O' modifier applies to certain format specifiers to request the use of alternative numeric symbols. If a locale cannot honor the modified request, it uses the unmodified format specifier.

## See Also

[time\\_get class template](#), [time\\_put\\_byname class template](#), [<ctime>](#)

## Name

`time_put_byname` class template — Facet for output of dates and times

## Synopsis

```
template <typename charT, typename OutputIterator =
ostreambuf_iterator<charT> >
class time_put_byname : public time_put<charT,OutputIterator>
{
public:
    typedef charT char_type;
    typedef OutputIterator iter_type;
    explicit time_put_byname(const char*, size_t refs = 0);
protected:
    // . . . Same virtual functions as in time_put
};
```

The `time_put` class template is a facet for formatting and writing dates and times using a named locale. The `time_put_byname<char>` and `time_put_byname<wchar_t>` instantiations are standard.

## See Also

[time\\_put class template](#)



## Name

tolower function template — Converts a character to lowercase in a locale

## Synopsis

```
template <typename charT>  
charT tolower(charT c, const locale& loc);
```

The `tolower` function converts the character `c` to lowercase using the locale `loc`:

```
use_facet<ctype<charT> >(loc).tolower(c)
```

## See Also

[ctype class template](#), [islower function template](#), [toupper function template](#)

## Name

toupper function template — Converts a character to uppercase in a locale

## Synopsis

```
template <typename charT>  
charT toupper(charT c, const locale& loc);
```

The `toupper` function converts the character `c` to uppercase using the locale `loc`:

```
use_facet<ctype<charT> >(loc).toupper(c)
```

## See Also

[ctype class template](#), [isupper function template](#), [tolower function template](#)

## Name

`use_facet` function template — Retrieves a facet for a locale

## Synopsis

```
template <typename Facet>  
const Facet& use_facet(const locale& loc)
```

The `use_facet` function template obtains a facet from locale `loc`. See [Example 13-24](#) and [Example 13-27](#), earlier in this section.

## See Also

[has\\_facet function template](#), [locale::facet class](#)

## <map>

The <map> header is one of the standard container template headers. It declares the map and multimap class templates and a few global function templates that operate on map and multimap objects.

A *map* is a container that stores pairs of keys and values. Looking up keys, inserting keys, and deleting keys can all be performed in logarithmic or better time. Maps support bidirectional iterators (no random access). In other languages and libraries, maps are also called dictionaries and associative arrays.

See [Chapter 10](#) for information about containers. See the [<utility>](#) section later in this chapter for information about the pair class template.

# Name

map class template — Associative map container with unique keys

## Synopsis

```
template <typename Key, typename T, typename Compare = less<Key>,
          typename Alloc = allocator<pair<const Key, T> > >
class map {
public:
    typedef Key key_type;
    typedef T mapped_type;
    typedef pair<const Key, T> value_type;
    typedef Compare key_compare;
    typedef Alloc allocator_type;
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;
    class value_compare :
        public binary_function<value_type,value_type,bool>
    {
        friend class map;
    protected:
        Compare comp;
        value_compare(Compare c) : comp(c) {}
    public:
        bool operator( )(const value_type& x, const value_type& y)
const
            { return comp(x.first, y.first); }
    };
    explicit map(const Compare& comp = Compare( ), const Alloc& =
Alloc( ));
    template <class InputIterator>
map(InputIterator first, InputIterator last,
        const Compare& comp = Compare( ), const Alloc& = Alloc(
));
```

```

map(const map<Key,T,Compare,Alloc>& x);
~map( );
map<Key,T,Compare,Alloc>& operator=(const
map<Key,T,Compare,Alloc>& x);

allocator_type get_allocator( ) const;
// Iterators
iterator begin( );
const_iterator begin( ) const;
iterator end( );
const_iterator end( ) const;
reverse_iterator rbegin( );
const_reverse_iterator rbegin( ) const;
reverse_iterator rend( );
const_reverse_iterator rend( ) const;
// Capacity
bool empty( ) const;
size_type size( ) const;
size_type max_size( ) const;
// Element access
T& operator[(const key_type& x)];
// Modifiers
pair<iterator, bool> insert(const value_type& x);
iterator insert(iterator hintpos, const value_type& x);
template <class InputIterator>
void insert(InputIterator first, InputIterator last);
void erase(iterator position);
size_type erase(const key_type& x);
void erase(iterator first, iterator last);
void swap(map<Key,T,Compare,Alloc>&);
void clear( );
// Observers
key_compare key_comp( ) const;
value_compare value_comp( ) const;
// Map operations
iterator find(const key_type& x);
const_iterator find(const key_type& x) const;
size_type count(const key_type& x) const;
iterator lower_bound(const key_type& x);
const_iterator lower_bound(const key_type& x) const;
iterator upper_bound(const key_type& x);
const_iterator upper_bound(const key_type& x) const;
pair<iterator,iterator> equal_range(const key_type& x);
pair<const_iterator,const_iterator>
equal_range(const key_type& x) const;
};

```

The map class template represents a map container. A map stores pairs of unique keys and associated objects, in which the key type is specified by the Key template parameter, and the associated type is the T template parameter. The values stored in the map are of type `pair<const Key, T>` (which has the convenience typedef `value_type`).

A map's iterators are bidirectional. They return `value_type` references; use the first member to access the key or second to access the associated object.

Note that keys are `const` in the map. You must not change the key while it is stored in a map. More precisely, you must not change the key in a way that alters its relative order with the other keys in the map. If you need to modify a key, you can erase the key from the map, modify the key, and insert the new key with its original associated value, as shown in [Example 13-28](#).

## Examples

*Example 13-28. One way to modify a key in a map*

```
template <typename Key, typename T, typename C, typename A>
void change_key(std::map<Key, T, C, A>& m,
               const Key& oldkey, const Key& newkey)
{
    using std::map;
    typedef typename map<Key, T, C, A>::iterator map_iterator;
    map_iterator i = m.find(oldkey);
    if (i != m.end( )) {
        // Save a copy of i->second because erase invalidates i.
        T tmp = i->second;
        m.erase(i);
        m[newkey] = tmp;
    }
    // Exercise: What if newkey is already in m?
}
```

Within a map, keys are ordered in ascending order, according to the Compare template parameter (which can be a function pointer or functor that compares two objects

of type `Key` and returns `true` if the first argument should come before the second). Keys must be unique, but note that uniqueness is determined only by calling `Compare`, not by using the `==` operator. That is, two objects, `a` and `b`, are different (and therefore can both be present in a single map object) if `Compare(a, b)` is `true` or `Compare(b, a)` is `true`. See `multimap` later in this section for a map container that can store non-unique keys.

Inserting into a map does not invalidate any iterators for that map or references to pairs in the map. Erasing an element invalidates only iterators and references that refer to that element.

Inserting into a map and searching for an element in a map usually take logarithmic time. Erasing a single element, given an iterator, takes amortized constant time.

The subscript operator (`[]`) lets you use a map as an associative array, for which the array indices are keys. If a key is not already present in the map, it is added. Using operator `[]` allows for compact, easy-to-read code, as you can see in [Example 13-29](#), which shows how to use map to count word frequencies in the standard input.

*Example 13-29. Using a map to count word frequencies*

```
#include <cstdlib>
#include <iostream>
#include <istream>
#include <map>
#include <ostream>
#include <string>

typedef std::map<std::string, std::size_t> freqmap;

// Print a single word and its count.
void print(const freqmap::value_type info)
{
    std::cout << info.first << '\t' << info.second << '\n';
}
```



```

int main( )
{
    freqmap fm;
    std::string word;
    // Count words. If a word is not in the map, add it. When a new
word is added,
    // its count is initially 0. Each time, including the first,
increment the
    // count.
    while (std::cin >> word)
        ++fm[word];
    // Print the frequencies of each word, in order.
    std::for_each(fm.begin( ), fm.end( ), print);
}

```

The following are the member functions of map:

explicit **map** (const Compare& comp = Compare( ),  
const Alloc& = Alloc( ))

Creates an empty map.

template <class InputIterator>, **map** (InputIterator  
first, InputIterator last, const Compare& comp =  
Compare( ), const Alloc& = Alloc( ))

Creates an empty map and then copies all pairs in the  
range [first, last) into the new map.

**map** (const map<Key,T,Compare,Alloc>& x)

Creates a new map and copies the allocator and all the  
pairs from x to the new map.

iterator **begin** ( ), const\_iterator **begin** ( ) const

Returns an iterator that points to the first item in the  
map.

void **clear** ( )

Erases every item in the map.

size\_type **count** (const key\_type& x) const

Returns the number of pairs whose keys are equivalent to x. This value is always 0 or 1.

bool **empty** ( ) const

Returns `size( ) == 0`.

iterator **end** ( ), const\_iterator **end** ( ) const

Returns an iterator that points to one past the last item in the map.

pair<iterator,iterator> **equal\_range** (const key\_type& x), pair<const\_iterator,const\_iterator> **equal\_range** (const key\_type& x) const

Returns the lower bound and upper bound as a pair:

`std::make_pair(lower_bound(x), upper_bound(x))`

void **erase** (iterator position), size\_type **erase** (const key\_type& x), void **erase** (iterator first, iterator last)

Erases one or more pairs from the map. The first version erases the pair at position in constant time (amortized over many calls). The second version erases the pair equivalent to x, if it is present, returning a count of the number of pairs erased, that is, 0 or 1. It runs in logarithmic time. The third version erases all elements in the range [first, last) in a time proportional to `log size( ) + (last - first)`.

iterator **find** (const key\_type& x), const\_iterator **find** (const key\_type& x) const

Searches for a pair whose key is equivalent to x and returns an iterator that points to that pair or `end( )` if it is not found. It runs in logarithmic time.

allocator\_type **get\_allocator** ( ) const

Returns the map's allocator.

```
pair<iterator, bool> insert (const value_type& x),  
iterator insert (iterator hintpos, const value_type&  
x), template <class InputIterator> void insert  
(InputIterator first, InputIterator last)
```

Inserts one or more pairs into the map, but only if an equivalent key is not already present in the map. If the key is already present, the insert attempt is ignored. The first version attempts to insert the pair *x* in logarithmic time.

The second version inserts the pair *x* using *hintpos* as a position hint. If *x* is inserted immediately after *hintpos*, the performance is constant (amortized over many insertions); at any other position, the performance is logarithmic. Use this form when inserting many items that are already in the desired order.

The third version copies all the pairs in the range [*first*, *last*), which must be pointing to a different map object. If the items are already in the desired order, the performance is linear; otherwise, it is  $N \log (\text{size}(\ ) + N)$ , in which *N* is *last* - *first*.

```
key_compare key_comp ( ) const
```

Returns the comparator function pointer or object, which compares keys. The *key\_compare* type is the same as the *Compare* template parameter. See also the *value\_comp* member.

```
iterator lower_bound (const key_type& x),  
const_iterator lower_bound (const key_type& x)  
const
```

Returns an iterator that points to the first pair in the map that does not come before *x*. That is, if *x* is in the map, the iterator points to its position; otherwise, the iterator points to the first position where *x* should be inserted. Performance is logarithmic.

`size_type` **max\_size** ( ) const

Returns the largest number of pairs that can be in the map.

reverse\_iterator **rbegin** ( ), const\_reverse\_iterator **rbegin** ( ) const

Returns a reverse iterator that points to the last element of the map.

reverse\_iterator **rend** ( ), const\_reverse\_iterator **rend** ( ) const

Returns a reverse iterator that points to one position before the first element of the map.

`size_type` **size** ( ) const

Returns the number of pairs in the map.

void **swap** (map<Key,T,Compare,Alloc>&)

Swaps the contents of the map with the contents of x.

iterator **upper\_bound** (const key\_type& x),  
const\_iterator **upper\_bound** (const key\_type& x)  
const

Returns an iterator that points to the first pair in the map that comes after x. Performance is logarithmic.

value\_compare **value\_comp** ( ) const

Returns a `value_compare` object, which can be used to compare pairs. The `value_compare` object takes two `value_type` arguments and compares their keys, returning true if the first should come before the second in the map.

map<Key,T,Compare,Alloc>& **operator=** (const  
map<Key,T,Compare,Alloc>& x)

Erases all the elements of the map and replaces them with copies of the elements of `x`.

**T& operator[]** (const key\_type& x)

Returns a reference to the object associated with the key `x`. If `x` is not in the map, it is added with a default associated object, and a reference to that new object is returned. That is, `operator[]` returns:

```
(*((insert(std::make_pair(x, T( )))).first)).second
```

Note that there is no `const` version of this operator.

## See Also

[multimap class template](#), set in [<set>](#)

# Name

multimap class template — Associative map container with duplicate keys

## Synopsis

```
template <class Key, class T, class Compare = less<Key>,
          class Alloc = allocator<pair<const Key, T> > >
class multimap {
public:
    typedef Key key_type;
    typedef T mapped_type;
    typedef pair<const Key,T> value_type;
    typedef Compare key_compare;
    typedef Alloc allocator_type;
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;
    class value_compare :
        public binary_function<value_type,value_type,bool>
    {
        friend class multimap;
    protected:
        Compare comp;
        value_compare(Compare c) : comp(c) {}
    public:
        bool operator( )(const value_type& x, const value_type& y)
            const { return comp(x.first, y.first); }
    };

    explicit multimap(const Compare& comp = Compare( ),
                     const Alloc& = Alloc( ));
    template <class InputIterator>
multimap(InputIterator first, InputIterator last,
          const Compare& comp = Compare( ), const Alloc& = Alloc( ));
multimap(const multimap<Key,T,Compare,Alloc>& x);
```

```

~multimap( );
multimap<Key,T,Compare,Alloc>&
  operator=(const multimap<Key,T,Compare,Alloc>& x);
allocator_type get_allocator( ) const;
// Iterators
iterator begin( );
const_iterator begin( ) const;
iterator end( );
const_iterator end( ) const;
reverse_iterator rbegin( );
const_reverse_iterator rbegin( ) const;
reverse_iterator rend( );
const_reverse_iterator rend( ) const;
// Capacity
bool empty( ) const;
size_type size( ) const;
size_type max_size( ) const;
// Modifiers
iterator insert(const value_type& x);
iterator insert(iterator hintpos, const value_type& x);
template <class InputIterator>
void insert(InputIterator first, InputIterator last);
void erase(iterator position);
size_type erase(const key_type& x);
void erase(iterator first, iterator last);
void swap(multimap<Key,T,Compare,Alloc>&);
void clear( );
// Observers
key_compare key_comp( ) const;
value_compare value_comp( ) const;
// Map operations
iterator find(const key_type& x);
const_iterator find(const key_type& x) const;
size_type count(const key_type& x) const;
iterator lower_bound(const key_type& x);
const_iterator lower_bound(const key_type& x) const;
iterator upper_bound(const key_type& x);
const_iterator upper_bound(const key_type& x) const;
pair<iterator,iterator> equal_range(const key_type& x);
pair<const_iterator,const_iterator>
equal_range(const key_type& x) const;
};

```

The `multimap` class template represents a map container that can store duplicate keys. A map stores pairs of keys and associated objects, in which the key type is specified by the `Key` template parameter, and the associated type is the

T template parameter. The values stored in the map are of type `pair<const Key, T>` (which has the convenience typedef `value_type`).

A map's iterators are bidirectional. They return `value_type` references; use the `first` member to access the key or `second` to access the associated object.

Note that keys are `const` in the map. You must not change the key while it is stored in a map. More precisely, you must not change the key in a way that alters its relative order with the other keys in the map. See [Example 13-28](#) (earlier in this section), which shows how to change a key by erasing and reinserting an object.

Within a map, keys are ordered in ascending order, according to the `Compare` template parameter (which can be a function pointer or functor that compares two objects of type `Key` and returns `true` if the first argument should come before the second).

Inserting into a map does not invalidate any iterators for that map. Erasing an element invalidates only iterators that refer to that element.

Inserting a single item into a map and searching for an element in a map usually take logarithmic time. Erasing a single element, given an iterator, takes amortized constant time.

The following are the member functions of `multimap`. Note that `multimap` does not have a subscript operator.

```
explicit multimap (const Compare& comp = Compare(
), const Alloc& = Alloc( ))
```

Creates an empty map.

```
template <class InputIterator>, multimap
(InputIterator first, InputIterator last, const
```



Compare& comp = Compare( ), const Alloc& = Alloc( ))

Creates an empty map and then copies all pairs in the range [first, last) into the new map.

**multimap** (const multimap<Key,T,Compare,Alloc>& x)

Creates a new map and copies all the pairs from x to the new map.

iterator **begin** ( ), const\_iterator **begin** ( ) const

Returns an iterator that points to the first item in the map.

void **clear** ( )

Erases every item in the map in linear time.

size\_type **count** (const key\_type& x) const

Returns the number of pairs whose keys are equivalent to x. Complexity is proportional to  $\log(\text{size}())$  + the return value.

bool **empty** ( ) const

Returns  $\text{size}() == 0$ .

iterator **end** ( ), const\_iterator **end** ( ) const

Returns an iterator that points to one past the last item in the map.

pair<iterator,iterator> **equal\_range** (const key\_type& x), pair<const\_iterator,const\_iterator> **equal\_range** (const key\_type& x) const

Returns the lower bound and upper bound as a pair:

`std::make_pair(lower_bound(x), upper_bound(x))`

void **erase** (iterator position), size\_type **erase** (const key\_type& x), void **erase** (iterator first, iterator last)

Erases one or more pairs from the map. The first version erases the pair at position in constant time (amortized over many calls). The second version erases all the pairs equivalent to x if any are present, returning a count of the number of pairs erased. The third version erases all elements in the range [first, last). The last two forms take time proportional to  $\log(\text{size}()) +$  the number of elements erased.

```
iterator find (const key_type& x), const_iterator  
find (const key_type& x) const
```

Searches for a pair whose key is equivalent to x and returns an iterator that points to that pair or end() if it is not found. If x occurs more than once, the iterator might point to any of the equivalent pairs.

```
allocator_type get_allocator () const
```

Returns the map's allocator.

```
pair<iterator, bool> insert (const value_type& x),  
iterator insert (iterator hintpos, const value_type&  
x), template <class InputIterator> void insert  
(InputIterator first, InputIterator last)
```

Inserts one or more pairs into the map. The first version inserts the pair x in logarithmic time.

The second version inserts the pair x using hintpos as a position hint. If x is inserted immediately after hintpos, the performance is constant (amortized over many insertions); at any other position, the performance is logarithmic. Use this form when inserting many items that are already in the desired order.

The third version copies all the pairs in the range [first, last), which must be pointing to a different multimap object. If the items are already in the desired order, the

performance is linear; otherwise, it is  $N \log (\text{size}() + N)$ , in which  $N$  is last - first.

key\_compare **key\_comp** ( ) const

Returns the comparator function pointer or object, which compares keys. The key\_compare type is the same as the Compare template parameter. See also the value\_comp member.

iterator **lower\_bound** (const key\_type& x),  
const\_iterator **lower\_bound** (const key\_type& x)  
const

Returns an iterator that points to the first pair in the map that does not come before x. That is, if x is in the map, the iterator points to the position of its first occurrence; otherwise, the iterator points to the first position where x should be inserted. Performance is logarithmic.

size\_type **max\_size** ( ) const

Returns the largest number of pairs that can be in the map.

reverse\_iterator **rbegin** ( ), const\_reverse\_iterator  
**rbegin** ( ) const

Returns a reverse iterator that points to the last element of the map.

reverse\_iterator **rend** ( ), const\_reverse\_iterator  
**rend** ( ) const

Returns a reverse iterator that points to one position before the first element of the map.

size\_type **size** ( ) const

Returns the number of pairs in the map.

void **swap** (multimap<Key,T,Compare,Alloc>&)

Swaps the contents of the map with the contents of x.

```
iterator upper_bound (const key_type& x),  
const_iterator upper_bound (const key_type& x)  
const
```

Returns an iterator that points to the first pair in the map that comes after all occurrences of x. Performance is logarithmic.

```
value_compare value_comp ( ) const
```

Returns a value\_compare object, which can be used to compare pairs. The value\_compare object takes two value\_type arguments and compares their keys, returning true if the first should come before the second in the map.

```
multimap<Key,T,Compare,Alloc>& operator= (const  
multimap<Key,T,Compare,Alloc>& x)
```

Erases all the elements of the map and replaces them with copies of the elements of x.

## See Also

[map class template](#), multiset in [<set>](#)

## Name

`operator==` function template — Compares maps for equality

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
bool operator==(const map<Key,T,Comp,Alloc>& x,
                const map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
bool operator==(const multimap<Key,T,Comp,Alloc>& x,
                const multimap<Key,T,Comp,Alloc>& y);
```

The `==` operator returns true if `x` and `y` have the same size and their elements are equal, that is, `x.size( ) == y.size( ) && equals(x.begin( ), x.end( ), y.begin( ))`.

## See Also

`equals` in [<algorithm>](#)

## Name

`operator!=` function template — Compares maps for inequality

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
bool operator!=(const map<Key,T,Comp,Alloc>& x,
                const map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
bool operator!=(const multimap<Key,T,Comp,Alloc>& x,
                const multimap<Key,T,Comp,Alloc>& y);
```

The `!=` operator is equivalent to `!(x == y)`.

## Name

operator< function template — Compares maps for less-than

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
bool operator<(const map<Key,T,Comp,Alloc>& x,
               const map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
bool operator<(const multimap<Key,T,Comp,Alloc>& x,
               const multimap<Key,T,Comp,Alloc>& y);
```

The < operator determines whether x is less than y, using the same algorithm as

lexicographical\_compare(x.begin( ), x.end( ), y.begin( ), y.end( )).

## See Also

lexicographical\_compare in [<algorithm>](#)

## Name

`operator<=` function template — Compares maps for less-than-or-equal

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
bool operator<=(const map<Key,T,Comp,Alloc>& x,
                const map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
bool operator<=(const multimap<Key,T,Comp,Alloc>& x,
                const multimap<Key,T,Comp,Alloc>& y);
```

The `<=` operator is equivalent to `!(y < x)`.



## Name

`operator>` function template — Compares maps for greater-than

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
bool operator>(const map<Key,T,Comp,Alloc>& x,
               const map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
bool operator>(const multimap<Key,T,Comp,Alloc>& x,
               const multimap<Key,T,Comp,Alloc>& y);
```

The `>` operator is equivalent to `(y < x)`.

## Name

`operator>=` function template — Compares maps for greater-than-or-equal

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
bool operator>=(const map<Key,T,Comp,Alloc>& x,
                const map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
bool operator>=(const multimap<Key,T,Comp,Alloc>& x,
                const multimap<Key,T,Comp,Alloc>& y);
```

The `>=` operator is equivalent to `!(x < y)`.

## Name

swap function template — Swaps the contents of two maps

## Synopsis

```
template <class Key, class T, class Comp, class Alloc>
    void swap(map<Key,T,Comp,Alloc>& x, map<Key,T,Comp,Alloc>& y);
template <class Key, class T, class Comp, class Alloc>
    void swap(multimap<Key,T,Comp,Alloc>& x,
multimap<Key,T,Comp,Alloc>& y);
```

The swap function template specialization is equivalent to calling `x.swap(y)`.

## See Also

swap in [<algorithm>](#)

## **<memory>**

The `<memory>` header declares function and class templates for allocating and using memory, such as the `auto_ptr<>` smart pointer, algorithms for working with uninitialized memory, and a standard allocator for use with the standard containers.

The `auto_ptr` class template provides a simple ownership model for working with pointers. It can be extremely useful for writing exception-safe code. On the other hand, copying an `auto_ptr<>` object does not produce an exact copy (ownership of the pointer is transferred), so `auto_ptr<>` objects cannot be stored in standard containers.

Several functions work with uninitialized memory, which can be helpful when implementing a container. For example, an implementation of `vector` must allocate an uninitialized array of objects and initialize elements of the array as they are needed. The `uninitialized_ . . .` functions can come in handy for that purpose.

The `allocator` class template manages memory allocation and deallocation and the construction and destruction of objects in the memory it manages. It is the default allocator for all the standard containers.

# Name

allocator class template — Encapsulates memory allocation and deallocation

# Synopsis

```
template <class T>
class allocator {
public:
    typedef size_t size_type;
    typedef ptrdiff_t difference_type;
    typedef T* pointer;
    typedef const T* const_pointer;
    typedef T& reference;
    typedef const T& const_reference;
    typedef T value_type;
    template <class U> struct rebind {
        typedef allocator<U> other;
    };
    allocator( ) throw( );
    allocator(const allocator&) throw( );
    template <class U> allocator(const allocator<U>&) throw( );
    ~allocator( ) throw( );
    pointer address(reference x) const;
    const_pointer address(const_reference x) const;
    pointer allocate(size_type, allocator<void>::const_pointer hint
= 0);
    void deallocate(pointer p, size_type n);
    size_type max_size( ) const throw( );
    void construct(pointer p, const T& val);
    void destroy(pointer p);
};
```

The `allocator` class template encapsulates basic allocation and deallocation functions. The standard containers rely on allocators for memory management and use `allocator` as the default allocator.

Most programmers do not need to use `allocator`, which offers few advantages over plain `new` and `delete`. However, if you want to write your own container, or provide a

custom allocator for the standard containers, you should take the time to understand allocator.

Perhaps the easiest way to understand allocator is to take a look at a trivial implementation in [Example 13-30](#). Note that a library might have a more complicated implementation to handle multithreading, improve performance, etc. Some libraries offer allocators for special purposes, such as allocating memory that can be shared among multiple processes. This particular implementation is just a sample.

## Example

### *Example 13-30. Sample allocator implementation*

```
template<typename T>
class myallocator {
public:
    typedef std::size_t size_type;
    typedef std::ptrdiff_t difference_type;
    typedef T* pointer;
    typedef const T* const_pointer;
    typedef T& reference;
    typedef const T& const_reference;
    typedef T value_type;
    template <class U> struct rebind {
        typedef myallocator<U> other;
    };
    myallocator( ) throw( ) {}
    myallocator(const myallocator&) throw( ) {}
    template <class U>
    myallocator(const myallocator<U>&) throw( ) {}
    ~myallocator( ) throw( ) {}
    pointer address(reference x) const {return &x;}
    const_pointer address(const_reference x) const {return &x;}
    pointer allocate(size_type n, void* hint = 0) {
        return static_cast<T*> (::operator new (n * sizeof(T)) );
    }
    void deallocate(pointer p, size_type n) {
        ::operator delete(static_cast<void*>(p));
    }
    size_type max_size( ) const throw( ) {
        return std::numeric_limits<size_type>::max( ) / sizeof(T);
    }
};
```

```

    }
    void construct(pointer p, const T& val) {
        new(static_cast<void*>(p)) T(val);
    }
    void destroy(pointer p) {
        p->~T( );
    }
};

template<typename T>
bool operator==(const myallocator<T>&, const myallocator<T>&)
{
    return true;
}

template<typename T>
bool operator!=(const myallocator<T>&, const myallocator<T>&)
{
    return false;
}

template<>
class myallocator<void> {
public:
    typedef void* pointer;
    typedef const void* const_pointer;
    typedef void value_type;
    template <class U> struct rebind {
        typedef myallocator<U> other;
    };
};
};

```

The following are the members of allocator:

**allocator** ( ) throw( ), **allocator** (const allocator&) throw( ), **template<class U> allocator** (const allocator<U>&) throw( )

Constructs a new allocator object, possibly copying an existing allocator. Remember that all instances must be equivalent, so an allocator typically does not have any data members to initialize.

**pointer address** (reference x) const, **const\_pointer address** (const\_reference x) const

Returns the address of `x`, that is, `&x`.

```
pointer allocate (size_type n,  
allocator<void>::const_pointer hint = 0)
```

Calls the global new operator to allocate enough memory to hold `n` objects of type `T`. The `hint` argument must be `0` or a pointer obtained from a prior call to `allocate` that has not been passed to `deallocate`. The return value is a pointer to the newly allocated memory. If the memory cannot be allocated, `bad_alloc` is thrown.

An implementation might use `hint` to improve performance.

```
 typedef const T* const_pointer
```

A type for a pointer to `const`. In a custom allocator, the type should be equivalent to `const T*`.

```
typedef const T& const_reference
```

A type for a `const` lvalue. In a custom allocator, the type should be equivalent to `const T&`.

```
void construct (pointer p, const T& val)
```

Calls the global new operator to construct an instance of `T` with value `val` using the memory that `p` points to. That is, it calls `new(static_cast<void*>(p)) T(val)`.

```
void deallocate (pointer p, size_type n)
```

Calls the global delete operator to free the memory that `p` points to. The `n` argument is the number of items of type `T`—the same value passed to `allocate`.

```
void destroy (pointer p)
```

Calls the destructor for the object at address `p`. That is, it calls `reinterpret_cast<T*>(p)->~T( )`.

```
typedef ptrdiff_t difference_type
```



A type that represents the difference of any two pointers that the allocator returns from `allocate( )`.

```
size_type max_size ( ) const throw( )
```

Returns the maximum size that can be passed to `allocate`.

```
typedef T* pointer
```

A pointer type. In a custom allocator, the type should be equivalent to `T*`.

```
template <class U> struct rebind
```

Binds the allocator object to a different value type. The `rebind` class has a single typedef, `other`, which is an instance of the same allocator template, but with `U` as the template parameter. The `rebind` template is necessary for standard containers that allocate helper objects, such as link nodes, rather than allocating values directly. If you are not implementing a standard container, you probably don't need to understand `rebind`.

```
typedef T& reference
```

An lvalue type. In a custom allocator, the type should be equivalent to `T&`.

```
typedef size_t size_type
```

A type that can represent the size of the largest allocation request.

```
typedef T value_type
```

The type of allocated values, which is typically `T`.

## See Also

[allocator<void> class](#), [<new>](#), new operator, delete operator

## Name

`allocator<void>` class — Specializes allocator for void pointers

## Synopsis

```
template <> class allocator<void> {
public:
    typedef void* pointer;
    typedef const void* const_pointer;
    typedef void value_type;
    template <class U> struct rebind {
        typedef allocator<U> other;
    };
};
```

The `allocator<void>` specialization is necessary to represent pointers to `void` without permitting the allocation of objects of type `void`.

## See Also

[allocator class template](#)

# Name

auto\_ptr class template — Smart pointer to manage ownership of pointers

## Synopsis

```
template <class T> struct auto_ptr_ref {};  
template<class T>  
class auto_ptr {  
public:  
    typedef T element_type;  
    explicit auto_ptr(T* p = 0) throw( );  
    auto_ptr(auto_ptr&) throw( );  
    template<class U> auto_ptr(auto_ptr<U>&) throw( );  
    auto_ptr(auto_ptr_ref<T>) throw( );  
    ~auto_ptr( ) throw( );  
    auto_ptr& operator=(auto_ptr&) throw( );  
    template<class U>  
    auto_ptr& operator=(auto_ptr<U>&) throw( );  
    auto_ptr& operator=(auto_ptr_ref<T> r) throw( );  
  
    T& operator*( ) const throw( );  
    T* operator->( ) const throw( );  
    T* get( ) const throw( );  
    T* release( ) throw( );  
    void reset(T* p = 0) throw( );  
  
    template<class U> operator  
        auto_ptr_ref<U>( ) throw( );  
    template<class U> operator  
        auto_ptr<U>( ) throw( );  
};
```

The `auto_ptr` class template implements a smart pointer to manage ownership of pointers. Proper use of `auto_ptr` ensures that a pointer has exactly one owner (which prevents accidental double deletes), and the owner automatically frees the memory when the owner goes out of scope (which prevents memory leaks). Assignment of `auto_ptr` values transfers ownership from the source to the target of the assignment.

The `auto_ptr_ref` type holds a reference to an `auto_ptr`. Implicit conversions between `auto_ptr` and `auto_ptr_ref` facilitate the return of `auto_ptr` objects from functions. Usually, you can ignore the `auto_ptr_ref` type and let the implicit conversions handle the details for you. All you need to do is use `auto_ptr` as a return type. The details of `auto_ptr_ref` are implementation-defined.

Some of the typical uses for `auto_ptr` are:

### Data members

Data members of pointer type that point to dynamically allocated objects are prime candidates for `auto_ptr`, which ensures that the memory is properly freed when the owning object is destroyed. Be sure to implement a copy constructor and assignment operator for any class that uses `auto_ptr<>` for its data members.

### Local variables

If a function must dynamically allocate a temporary object, store the pointer in an `auto_ptr` variable. When the function returns (normally or as the result of an exception), the object is destroyed automatically. This can drastically reduce the need for `try-catch` statements.

### Transferring ownership

In a complex program, objects are frequently allocated in one part of the program and freed in another. It can be difficult to keep track of when it is safe or proper to free an object. Using `auto_ptr`, you can safely ensure that each object has exactly one owner, and ownership is properly passed via assignment and function calls. When the object is no longer needed, it is freed automatically.

Because you cannot simply copy or assign an `auto_ptr`, you cannot use `auto_ptr` objects in a standard container. Another limitation is that `auto_ptr` cannot hold a pointer to an array. Allocating and freeing a single object (e.g., `new int`) is different from allocating and freeing an array of objects, (e.g., `new int[42]`), and `auto_ptr` is designed to work only with single objects.

### **Not e**

A useful guideline is that a program should avoid bare pointers (e.g., `int* x`). Bare pointers are error-prone: they are subject to memory leaks, double-freeing, and dangling references. Instead, use some form of ownership, such as `auto_ptr<>`, or one of the Boost smart pointers.

The Boost project has additional smart-pointer class templates that permit copying, arrays, and shared ownership. See [Appendix B](#) for more information about Boost.

[Example 13-31](#) shows some uses of `auto_ptr`.

## **Example**

*Example 13-31. Sample uses of `auto_ptr`*

```
class brush { . . . };
class pen { . . . };

// A function can return an auto_ptr<> object.
std::auto_ptr<brush> default_brush( )
{
    return std::auto_ptr<brush>(new brush);
}

class DisplayContext {
    // Display or graphics context for drawing on a window.
public:
    DisplayContext( )
```

```

    : brush_(default_brush( )), pen_(new pen) {...}
    . . .
private:
    // Make sure caller never tries to copy or assign
    // DisplayContext, but uses only objects that are
    // managed by auto_ptr<>.
    DisplayContext(const DisplayContext& dc);
    DisplayContext& operator=(const DisplayContext& dc);
    // Automatically manage lifetime of the pen and brush.
    // When the DisplayContext is freed, so are the pen
    // and brush instances.
    std::auto_ptr<brush> brush_;
    std::auto_ptr<pen> pen_;
};

void repaint( )
{
    // Allocate a new display context. Use auto_ptr to ensure
    // that it will be freed automatically.
    std::auto_ptr<DisplayContext> dc(new DisplayContext( ));
    // Draw stuff on the display context.
    dc->draw( . . . );
    // No need to call release; the display context is
    // automatically released when repaint( ) returns.
}

int main( )
{
    std::auto_ptr<DisplayContext> dc1(new DisplayContext);
    std::auto_ptr<DisplayContext> dc2(dc1);

    dc1 = dc2;
    repaint( );
}

```

The following are the members of `auto_ptr`:

`explicit auto_ptr (T* p = 0) throw( )`

Initializes the `auto_ptr` object to own the pointer `p`.

`auto_ptr (auto_ptr& x) throw( ), template<class U>`

`auto_ptr (auto_ptr<U>& x) throw( )`

Initializes the `auto_ptr` object with the pointer returned from `x.release( )`. In the second version, the type `U*` must be implicitly convertible to `T*`. Note that `x` is not

const. It is not possible to copy a const auto\_ptr because to do so would break the ownership rules.

**auto\_ptr** (auto\_ptr\_ref<T> r) throw( )

Initializes the auto\_ptr object with the pointer obtained from calling release on r's auto\_ptr.

**~auto\_ptr** ( ) throw( )

Deletes the owned pointer (e.g., delete get( )).

T\* **get** ( ) const throw( )

Returns the owned pointer.

T\* **release** ( ) throw( )

Returns get( ) and resets the owned pointer to 0.

void **reset** (T\* p = 0) throw( )

Deletes the owned pointer (if it is not equal to p) and saves p as the new owned pointer.

template<class U> **operator auto\_ptr\_ref<U>**( )  
throw( )

Returns a temporary auto\_ptr\_ref object that owns the pointer. The pointer must be convertible to U\*. Ownership is released and transferred to the new auto\_ptr\_ref object.

template<class U> **operator auto\_ptr<U>**( ) throw( )

Returns a new auto\_ptr object. The owned pointer is converted to type U\*, and ownership is transferred to the new auto\_ptr object.

auto\_ptr& **operator=** (auto\_ptr& x) throw( ),

template<class U>, auto\_ptr& **operator=**  
(auto\_ptr<U>& x) throw( ), auto\_ptr& **operator=**  
(auto\_ptr\_ref<T> r) throw( )

Transfers ownership of the pointer that is owned by `x` or by the `auto_ptr` object held by `r` to `*this`. That is, it calls `reset(x.release( ))`.

**T& operator\*** ( ) const throw( )

Returns `*get( )`. If the owned pointer is a null pointer, the behavior is undefined.

**T\* operator ->**( ) const throw( )

Returns `get( )`.

## See Also

`new operator`



## Name

`get_temporary_buffer` function template — Allocates temporary memory buffer

## Synopsis

```
template <class T>  
pair<T*, ptrdiff_t> get_temporary_buffer(ptrdiff_t n);
```

The `get_temporary_buffer` function template allocates memory for temporary use. The request is for up to `n` adjacent objects of type `T`. The return value is a pair of the pointer to the newly allocated memory and the actual size of the memory allocated (in units of `sizeof(T)`). If the memory cannot be allocated, the return value is a pair of `0`s. The allocated memory must be freed by calling `return_temporary_buffer`. The temporary buffer is uninitialized.

This function has limited usefulness. You must test the return value to see how much memory was allocated and ensure that the memory is properly freed if an exception is thrown. It is usually simpler to call `new` and save the pointer in an `auto_ptr<>`.

## See Also

[auto\\_ptr class template](#), [return\\_temporary\\_buffer function template](#), `pair` in [<utility>](#)

## Name

`operator==` function template — Compares allocators for equality

## Synopsis

```
template <class T1, class T2>  
bool operator==(const allocator<T1>&, const allocator<T2>&)  
    throw( );
```

The `operator==` function template always returns true. In other words, any object of type `allocator` is considered to be the same as every other `allocator`.

## See Also

[allocator class template](#)

## Name

`operator!=` function template — Compares allocators for inequality

## Synopsis

```
template <class T1, class T2>  
bool operator!=(const allocator<T1>&, const allocator<T2>&)  
    throw( );
```

The `operator!=` function template always returns false. In other words, any object of type `allocator` is considered to be the same as every other `allocator`.

## See Also

[allocator class template](#)

## Name

`raw_storage_iterator` class template — Iterator for uninitialized memory

## Synopsis

```
template <class OutputIterator, class T>
class raw_storage_iterator :
    public iterator<output_iterator_tag,void,void,void,void> {
public:
    explicit raw_storage_iterator(OutputIterator x);
    raw_storage_iterator<OutputIterator,T>& operator*( );
    raw_storage_iterator<OutputIterator,T>&
        operator=(const T& element);
    raw_storage_iterator<OutputIterator,T>& operator++( );
    raw_storage_iterator<OutputIterator,T> operator++(int);
};
```

The `raw_storage_iterator` class template implements an output iterator that writes to uninitialized memory. It adapts another output iterator that must have `operator&` return a pointer to T. The adapted iterator is typically used as a pointer to uninitialized memory.

Use the `raw_storage_iterator` as you would any other output iterator.

## See Also

[uninitialized\\_copy function template](#), [uninitialized\\_fill function template](#), [uninitialized\\_fill\\_n function template](#), [<iterator>](#)

## Name

`return_temporary_buffer` function template — Frees temporary memory buffer

## Synopsis

```
template <class T>  
void return_temporary_buffer(T* p);
```

The `return_temporary_buffer` function reclaims the memory that was previously allocated by `get_temporary_buffer`.

## See Also

[get\\_temporary\\_buffer function template](#)

## Name

`uninitialized_copy` function template — Copies into uninitialized memory

## Synopsis

```
template <class InputIter, class FwdIter>
FwdIter uninitialized_copy(InputIter first, InputIter last,
FwdIter result);
```

The `uninitialized_copy` function template is like the `copy` algorithm, except the result iterator is assumed to point to uninitialized memory. The range `[first, last)` is copied to `result` using placement `new`.

## See Also

[raw\\_storage\\_iterator class template](#), [uninitialized\\_fill function template](#), `copy` in [<algorithm>](#), `new` operator

## Name

`uninitialized_fill` function template — Fills uninitialized memory

## Synopsis

```
template <class FwdIter, class T>  
void uninitialized_fill(FwdIter first, FwdIter last, const T& x);
```

The `uninitialized_fill` function template is like the `fill` algorithm, except that it fills uninitialized memory. Every item in the range `[first, last)` is constructed as a copy of `x` using placement `new`.

## See Also

`raw_storage_iterator` class template, [uninitialized\\_copy function template](#), [uninitialized\\_fill\\_n function template](#), `fill` in [<algorithm>](#), `new` keyword

## Name

`uninitialized_fill_n` function template — Fills uninitialized memory

## Synopsis

```
template <class FwdIter, class Size, class T>  
void uninitialized_fill_n(FwdIter first, Size n, const T& x);
```

The `uninitialized_fill_n` function template is like the `fill_n` algorithm, except it fills uninitialized memory. Starting with `first`, `n` copies of `x` are constructed using placement `new`.

## See Also

[raw\\_storage\\_iterator class template](#), [uninitialized\\_fill function template](#), `fill_n` in [<algorithm>](#), `new` keyword



## <new>

The <new> header declares types and functions related to dynamic memory management. (See [Chapter 3](#) for more information about the new and delete expressions, including placement new and delete, and the operator new and operator delete functions.) Most programs do not need to use <new>. The header is typically used by libraries and programs that implement their own operator new and operator delete functions or otherwise provide custom management of dynamic memory.

If a source file uses the standard new and delete expressions, it does not need to `#include <new>`. You can also use the pointer placement new without including this header. In order to use the nothrow placement new, or `catch bad_alloc`, you must include this header.

Most programs do not call the operators directly, but instead use new and delete expressions, and the compiler generates calls using the appropriate operators. Library implementors sometimes make direct calls to the operators, especially to allocate uninitialized memory. See <memory> earlier in this chapter for examples.

Some specialized applications might implement the global operator new and operator delete functions or provide additional overloaded operators for specialized circumstances, such as allocating memory that is shared across process boundaries. If you write your own operator new, you should obey the following guidelines:

- Implement operator new and operator new[ ].
- Implement operator delete and operator delete[ ]. Even if your operator new is a placement new function, you should have a corresponding placement delete

function (which is called if a new expression throws an exception).

- Return a pointer that meets the strictest alignment requirements of any type. (Note that `malloc` in `<cstdlib>` and the standard operator `new` function return aligned pointers.)
- Handle out-of-memory situations by throwing `bad_alloc` (or a class that derives from `bad_alloc`) or returning a null pointer. If you return a null pointer, your operator `new` function must have an empty exception specification.
- If operator `new` and operator `delete` are member functions, include the `static` keyword as a reminder to the human reader. The compiler always treats these functions as static, even if you omit the keyword.

[Example 13-32](#) shows a trivial implementation of the global operator `new` and operator `delete` functions.

*Example 13-32. Implementing operator new and operator delete with malloc and free*

```
#include <cstdlib>
#include <new>

void* operator new(std::size_t size) throw(std::bad_alloc)
{
    void* ptr = std::malloc(size);
    if (ptr == 0)
        throw std::bad_alloc( );
    return ptr;
}

void* operator new(std::size_t size, const std::nothrow_t&
throw( ))
{
    return std::malloc(size);
}

void* operator new[](std::size_t size) throw(std::bad_alloc)
{
```

```
    return operator new(size);
}

void* operator new[](std::size_t size, const std::nothrow_t&)
throw( )
{
    return operator new(size, std::nothrow);
}

void operator delete(void* ptr) throw( )
{
    std::free(ptr);
}

void operator delete(void* ptr, const std::nothrow_t&)
throw( )
{
    std::free(ptr);
}

void operator delete[](void* ptr) throw( )
{
    operator delete(ptr);
}

void operator delete[](void* ptr, const std::nothrow_t&)
throw( )
{
    operator delete(ptr);
}
}
```

## Name

`bad_alloc` class — Exception class for failed memory allocation

## Synopsis

```
class bad_alloc : public exception {
public:
    bad_alloc( ) throw( );
    bad_alloc(const bad_alloc&) throw( );
    bad_alloc& operator=(const bad_alloc&) throw( );
    virtual ~bad_alloc( ) throw( );
    virtual const char* what( ) const throw( );
};
```



The `bad_alloc` class is an exception class that is thrown when `operator new` is unable to fulfill a request to allocate memory. As with any of the standard exception classes, `what( )` returns an implementation-defined character string.

## See Also

[operator new](#), `set_new_handler` exception in [<exception>](#)

## Name

nothrow object — Requests null pointer return when out of memory

## Synopsis

```
struct nothrow_t {};  
extern const nothrow_t nothrow;
```

The `nothrow` object is used in placement `new` expressions to request that the `new` operator return a null pointer instead of throwing `bad_alloc` if the memory allocation request cannot be fulfilled.

The `nothrow_t` type does not do anything; it is used only in overloaded placement `new` and `delete` operators.

Note that `nothrow` is also accepted by overloaded operator `delete` for symmetry with `new`. The `nothrow` version of operator `delete` behaves just like the ordinary operator `delete`. Like any placement `delete` function, it is called only if the placement `new` expression throws an exception.

## See Also

[operator delete](#), [operator new](#)

## Name

operator delete — Global operator delete

## Synopsis

```
void operator delete(void* ptr) throw( );  
void operator delete[](void* ptr) throw( );  
void operator delete(void* ptr, const std::nothrow_t&) throw( );  
void operator delete[](void* ptr, const std::nothrow_t&) throw( );  
void operator delete(void* ptr, void*) throw( );  
void operator delete[](void* ptr, void*) throw( );
```

The global operator delete function is called from a delete expression to free memory. The memory, which ptr points to, must have been returned by a corresponding call to operator new or be a null pointer. You must not call operator delete more than once for the same pointer. If ptr is null, operator delete returns without doing anything.

The first two versions of operator delete free the memory that ptr points to, which must have been allocated by calling the plain form of operator new. These forms of operator delete are called from a delete expression. The first is called for a scalar delete, and the second is called for an array delete[ ].

The remaining forms are called only when the corresponding placement new expression throws an exception during construction. The nothrow functions free the memory that ptr points to. The last two forms do nothing. See the new expression in [Chapter 3](#) to learn how and when placement operator delete is called.

Unlike other identifiers in the standard library, operator delete is global and is not in the std namespace. Also, unlike with other functions in the standard library, you can

provide your own implementation of operator delete, which replaces the standard implementation. You cannot, however, replace the last two versions of delete with your own implementation.

## **See Also**

[operator\\_new](#)

# Name

operator new — Global operator new

## Synopsis

```
void* operator new(std::size_t size) throw(std::bad_alloc);
void* operator new(std::size_t size, const std::nothrow_t&
throw( ));
void* operator new[](std::size_t size) throw(std::bad_alloc);
void* operator new[](std::size_t size, const std::nothrow_t&
throw( ));
void* operator new(std::size_t size, void* ptr) throw( );
void* operator new[](std::size_t size, void* ptr) throw( );
```

The global operator new function allocates memory and returns a pointer to the newly allocated memory. The memory must later be released by a corresponding delete expression or an explicit call to operator delete.

The first version of new allocates at least size bytes of memory, suitably aligned to store any type, and returns a pointer to the memory. If the request cannot be fulfilled, it throws bad\_alloc.

The second version is like the first, but it returns a null pointer instead of throwing bad\_alloc if sufficient memory cannot be allocated.

The third version is like the first, but it allocates memory for storing an array of objects. It might allocate more than size bytes to permit the library to store additional bookkeeping information. You must use the array form of delete[] to free this memory.

The fourth version is like the third, but it returns a null pointer instead of throwing bad\_alloc if sufficient memory cannot be allocated.

To allocate memory, the operator new functions first try to allocate size bytes. If they cannot, they call the handler



function set by the most recent call to `set_new_handler`. Then they try again to allocate `size` bytes. This loop repeats until the request is fulfilled or the handler function fails to return. The `nothrow` versions of the function return `0` if the most recent call to `set_new_handler` was a null pointer or if the new handler function throws `bad_alloc`.

The final two versions do nothing except return `ptr`. These forms permit placement new expressions to specify a memory location where an object will be constructed.

Unlike other identifiers in the standard library, `operator new` is global and is not in the `std` namespace. Also, unlike with other functions in the standard library, you can provide your own implementation of `operator new`, which replaces the standard implementation. You cannot, however, replace the last two versions of `new` with your own implementation.

## See Also

[nothrow object](#), [operator delete](#), [set\\_new\\_handler function](#)

## Name

`set_new_handler` function — Sets handler for obtaining memory

## Synopsis

```
typedef void (*new_handler)( );  
new_handler set_new_handler(new_handler new_p) throw( );
```

The `set_new_handler` function stores a function pointer for a function that obtains additional memory from the operating system for use by the `new` operator. When the default operator `new` is unable to fulfill a request to allocate memory, it calls the handler that was set by the most recent call to `set_new_handler`. This handler must do one of the following:

- Obtain more memory from the host environment
- Throw `bad_alloc` (or a type that derives from `bad_alloc`)
- Call `abort( )` or `exit( )` to halt the program

The return value is the pointer to the previous handler, or `0` for the first call to `set_new_handler`.

## **<numeric>**

The `<numeric>` header declares several function templates for numerical algorithms. See `<algorithm>` earlier in this chapter for most of the standard algorithms. See `<cmath>` for math functions that operate on scalar quantities. The `<functional>` section contains the standard functors, which might be useful when calling the numeric algorithms.

Refer to [Chapter 10](#) for general information about algorithms and iterators. See [Appendix B](#) for information about other libraries that provide additional numeric functions.

## Name

accumulate function template — Computes a value from all items in a range

## Synopsis

```
template <typename InputIter, typename T>
T accumulate(InputIter first, InputIter last, T init);
template < typename InputIter, typename T, typename BinaryOp>
T accumulate(InputIter first, InputIter last, T init, BinaryOp
binary_op);
```

The `accumulate` function template sums all the values in the range `[first, last)` added with `init` and returns the result. The result and intermediate sum have the same type as `init`. The second version calls `binary_op` instead of using the addition (+) operator.

## Technical Notes

The result is computed as follows: for each `i` in the range `[first, last)`, `tmp = binary_op(tmp, *i)`, in which `tmp` is initialized to `init`. The final value of `tmp` is returned.

The `binary_op` function or functor must not have any side effects.

Complexity is linear: `binary_op` is called exactly `last - first` times.

## Name

`adjacent_difference` function template — Computes differences of adjacent elements in a range

## Synopsis

```
template <typename InIter, typename OutIter>
OutIter adjacent_difference(InIter first, InIter last, OutIter
result);
template <typename InIter, typename OutIter, typename BinOp>
OutIter adjacent_difference(InIter first, InIter last, OutIter
result,
                             BinOp binary_op);
```

The `adjacent_difference` function computes the differences of adjacent elements in the range `[first, last)` and assigns those differences to the output range starting at `result`. The second version calls `binary_op` instead of using the subtraction `(-)` operator.

## Technical Notes

For each `i` in `[first + 1, last)` and `j` in `[result, result + (last - first))`, assign `*j = *i - tmp`, in which `tmp` is initially `*first`; it becomes `*i` after each assignment to `*j`.

The return value is the `result` iterator pointing to one past the last element written.

The `binary_op` function or functor must not have any side effects. The `result` iterator can be the same as `first`.

Complexity is linear: `binary_op` is called exactly `last - first - 1` times.

## Name

`inner_product` function template — Computes inner product of two ranges

## Synopsis

```
template <typename InIter1, typename InIter2, typename T>
T inner_product(InIter1 first1, InIter1 last1, InIter2 first2, T
init);
template <typename InIter1, typename InIter2, typename T,
          typename BinaryOp1, typename BinaryOp2>
T inner_product(InIter1 first1, InIter1 last1, InIter2 first2, T
init,
                BinaryOp1 binary_op1, BinaryOp2 binary_op2);
```

The `inner_product` function template computes an inner product of two ranges. It accumulates the products of corresponding items in `[first1, last1)` and `[first2, last2)`, in which  $last2 = first2 + (last1 - first1)$ . The second version calls `binary_op1` as the accumulator operator (instead of addition) and `binary_op2` as the multiplication operator.

## Technical Notes

The result is computed as follows: for each `i` in the range `[first1, last1)`, and for each `j` in `[first2, last2)`, in which  $last2 = first2 + (last1 - first1)$ , assign `tmp = binary_op1(tmp, binary_op2(*i, *j))`, in which `tmp` is initialized to `init`. The final value of `tmp` is returned.

The `binary_op1` and `binary_op2` functions or functors must not have side effects.

Complexity is linear: `binary_op1` and `binary_op2` are called exactly `last - first` times.

## See Also

[accumulate function template](#)

## Name

`partial_sum` function template — Compute sums of subranges in a range

## Synopsis

```
template <typename InIter, typename OutIter>
OutIter partial_sum(InIter first, InIter last, OutIter result);
template <typename InIter, typename OutIter, typename BinOp>
OutIter partial_sum(InIter first, InIter last, OutIter result,
BinOp binary_op);
```

The `partial_sum` function template assigns partial sums to the range that starts at `result`. The partial sums are computed by accumulating successively larger subranges of `[first, last)`. Thus, the first result item is `*first`, the second is `*first + *(first + 1)`, and so on. The second version calls `binary_op` instead of using the addition operator (+).

## Technical Notes

For each `i` in `[first, last)`, assign `*(result + k) = sum(first, i)`, in which `k = i - first`, and `sum(a, b)` computes the sum in the manner of `accumulate(a + 1, b, *a, binary_op)`.

The return value is the `result` iterator, pointing to one past the last item written.

The `binary_op` function or functor must not have any side effects. The `result` iterator can be the same as `first`.

Complexity is linear: `binary_op` is called exactly `(last - first) - 1` times.



## **<ostream>**

The `<ostream>` header declares the output stream class template, specializations, and manipulators.

See `<fstream>` for derived classes that write to files and `<sstream>` for derived classes that write to strings. See `<ios>` for the base-class declarations. See `<string>` for information about the `char_traits` template. Refer to [Chapter 9](#) for general information about I/O.

# Name

`basic_ostream` class template — Base class for output streams

# Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_ostream : virtual public basic_ios<charT,traits>
{
public:
    // Types (inherited from basic_ios)
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    explicit basic_ostream(basic_streambuf<charT,traits>* sb);
    virtual ~basic_ostream( );

    class sentry;

    // Formatted output
    basic_ostream<charT,traits>&
operator<<(basic_ostream<charT,traits>&
    (*pf)(basic_ostream<charT,traits>&));
    basic_ostream<charT,traits>&
operator<<(basic_ios<charT,traits>&
    (*pf)(basic_ios<charT,traits>&));
    basic_ostream<charT,traits>& operator<<(ios_base&(*pf)
    (ios_base&));
    basic_ostream<charT,traits>& operator<<(bool n);
    basic_ostream<charT,traits>& operator<<(short n);
    basic_ostream<charT,traits>& operator<<(unsigned short n);
    basic_ostream<charT,traits>& operator<<(int n);
    basic_ostream<charT,traits>& operator<<(unsigned int n);
    basic_ostream<charT,traits>& operator<<(long n);
    basic_ostream<charT,traits>& operator<<(unsigned long n);
    basic_ostream<charT,traits>& operator<<(float f);
    basic_ostream<charT,traits>& operator<<(double f);
    basic_ostream<charT,traits>& operator<<(long double f);
    basic_ostream<charT,traits>& operator<<(const void* p);
    basic_ostream<charT,traits>& operator<<
    (basic_streambuf<char_type,traits>* sb);
```

```

// Unformatted output
basic_ostream<charT,traits>& put(char_type c);
basic_ostream<charT,traits>& write(const char_type* s,
streamsize n);
basic_ostream<charT,traits>& flush( );

pos_type tellp( );
basic_ostream<charT,traits>& seekp(pos_type);
basic_ostream<charT,traits>& seekp(off_type,
ios_base::seekdir);
};

```

The `basic_ostream` class template is the base for all output streams. It declares members for writing to streams and for managing streams. The act of writing to a stream is also known as *inserting* to the stream.

All writes go through a stream buffer, which provides low-level access to the stream data. (See the `sputc` function for `basic_streambuf` in the `<streambuf>` header.) If the stream buffer object throws an exception, the stream sets `badbit`.

Before performing an output operation (e.g., `operator<<`, `put`, or `write`), a stream constructs a sentry object. If the sentry evaluates to `true`, the write operation continues. If the write throws an exception, `badbit` is set. The sentry object is destroyed before the output function returns. See `basic_ostream::sentry` later in this section for more information.

When an output operation throws an exception, the stream sets `badbit`. If `badbit` is set in the `exceptions( )` mask, the stream does not throw `ios_base::failure`, but instead rethrows the original exception.

The following are the `basic_ostream` member functions:

```

explicit basic_ostream
(basic_streambuf<char_type,traits>* sb)

```

Constructs a `basic_ostream` object and then initializes it by calling `init(sb)`.

`virtual ~basic_ostream ( )`

Destroys the `basic_ostream` object without calling any functions of the stream buffer. Derived classes that might have buffered, unflushed data must take appropriate action to ensure that the buffer is flushed before the stream object is destroyed.

`basic_ostream<charT,traits>& flush ( )`

Flushes the output buffer. If `rdbuf( )` is not null, `flush` calls `rdbuf( )->pubsync( )`. If `pubsync` returns `-1`, `flush` sets `badbit`. The return value is `*this`.

`basic_ostream<charT,traits>& put (char_type c)`

Writes a single character `c`. If the write fails, `put` sets `badbit`. The return value is `*this`.

`basic_ostream<charT,traits>& seekp (pos_type pos),`  
`basic_ostream<charT,traits>& seekp (off_type off,`  
`ios_base::seekdir dir)`

Tries to seek to a new position in the stream. The first form specifies the position explicitly; the second form specifies the new position as an offset from a known position (start-of-file, current position, or end-of-file). If `fail( )` is false, `seekp` calls `rdbuf( )->pubseekoff(pos)` or `rdbuf( )->pubseekoff(off, dir)`. If `fail( )` is true, `seekp` does nothing. The return value is `*this`.

`pos_type tellp ( )`

Returns the current position in the stream. If `fail( )` is true, the return value is `pos_type(-1)`; otherwise, the return value is `rdbuf( )->pubseekoff(0,ios_base::cur,ios_base::out)`.

```
basic_ostream<charT,traits>& write (const
char_type* s, streamsize n)
```

Writes n characters from s. If the output fails after any character, write sets badbit and stops writing. The return value is \*this.

```
basic_ostream<charT,traits>& operator<<
(basic_ostream<charT,traits>& (*pf)
(basic_ostream<charT,traits>&))
```

Calls pf(\*this) and returns \*this. See endl later in this section for an example of a manipulator that uses this operator.

```
basic_ostream<charT,traits>& operator<<
(basic_ios<charT,traits>& (*pf)
(basic_ios<charT,traits>&)),
basic_ostream<charT,traits>& operator<< (ios_base&
(*pf)(ios_base&))
```

Calls pf(\*this) and returns \*this. See the dec function in <ios> for an example of a manipulator that uses this operator.

```
basic_ostream<charT,traits>& operator<< (bool n),
basic_ostream<charT,traits>& operator<< (short n),
basic_ostream<charT,traits>& operator<< (unsigned
short n), basic_ostream<charT,traits>& operator<<
(int n), basic_ostream<charT,traits>& operator<<
(unsigned int n), basic_ostream<charT,traits>&
operator<< (long n), basic_ostream<charT,traits>&
operator<< (unsigned long n),
basic_ostream<charT,traits>& operator<< (float f),
basic_ostream<charT,traits>& operator<< (double f),
basic_ostream<charT,traits>& operator<< (long
double f), basic_ostream<charT,traits>& operator<<
(const void* p)
```

Formats a value and writes the formatted characters to the output stream. These functions start by creating a sentry object; they then use the `num_put` facet of the stream's imbued locale as shown in [Example 13-33](#). If the formatting fails, `failbit` is set. If an exception is thrown, `badbit` is set. See the `<locale>` header for information about `num_put` and locales.

*Example 13-33. Using the `num_put` facet to format output*

```
typedef
    std::num_put<char_type,
    std::ostreambuf_iterator<char_type, traits_type> >
    numput;
std::ostreambuf_iterator<char_type, traits_type> iter =
    std::use_facet<numput>(getloc( )).(*this,*this,fill(
),val);
if (iter.failed( ))
    setstate(ios_base::badbit);
```

```
basic_ostream<charT,traits>& operator<<
(basic_streambuf<char_type,traits>* sb)
```

Writes characters from the stream buffer `sb`. If `sb` is null, `badbit` is set. Otherwise, characters are read from `sb` and written to `*this` until one of the following happens:

- The end-of-file is reached on `sb`
- Writing fails (`badbit` is set)
- An exception is thrown when reading from `sb` (`failbit` is set)

If no characters are written, `failbit` is set.

## See Also

[ostream class](#), [wostream class](#), `ostream` in [<ostream>](#)

## Name

`basic_ostream::sentry` class — Sentry class for output streams

## Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_ostream<charT, traits>::sentry {
public:
    explicit sentry(basic_ostream<charT, traits>& os);
    ~sentry( );
    operator bool( ) const;
private:
    sentry(const sentry&); // Not defined
    sentry& operator=(const sentry&); // Not defined
};
```

A `basic_ostream` object constructs a temporary sentry object prior to each output operation. The sentry object is destroyed when the output operation finishes and the function returns. The sentry manages tied streams and unit buffering.

The stream passes itself to the sentry's constructor. If `stream.good( )` is true, the sentry first flushes any tied stream. That is, if `stream.tie( )` is not null, the sentry calls `stream.tie( )->flush( )`.

If sentry preparation fails, `badbit` is set.

The sentry destructor flushes the buffer if the `unitbuf` flag is on and the output function did not throw an exception:

```
if ((os.flags( ) & ios_base::unitbuf) && !uncaught_exception(
))
    os.flush( );
```

## See Also

[basic\\_ostream class template](#), `basic_ios` in [<ios>](#)

## Name

`endl` function template — Manipulator to write an end-of-line character

## Synopsis

```
template <class charT, class traits>
basic_ostream<charT,traits>& endl(basic_ostream<charT,traits>&
os);
```

The `endl` function template is a manipulator that writes a newline to `os` and then calls `os.flush( )`:

```
std::cout << "Hello, world." << std::endl;
```

If you do not need to flush the output stream, do not use `endl`; write a plain `'\n'` character instead. If you feel you need to flush the output stream after finishing a line of output, consider using unit buffering for the stream, or if you need to flush the output prior to reading an input stream, you can tie the streams instead of using `endl`. See the `basic_ios` class template (in `<ios>`) for information about tied streams and the `ios_base` class (also in `<ios>`) for information about unit buffering.

## See Also

`basic_ios` in [<ios>](#), `ios_base::fmtflags` in [<ios>](#)



## Name

ends function template — Manipulator to write an end-of-string character

## Synopsis

```
template <class charT, class traits>
basic_ostream<charT,traits>& ends(basic_ostream<charT,traits>&
os);
```

The ends function template is a manipulator that writes a null character (defined by `charT( )`) to `os` to mark the end of a string. Typically, ends is used only when writing to a character array stream, that is, `ostrstream`:

```
std::ostrstream out1;
out1 << "Hi" << std::ends; // out1.str( ) has length 2.
```

## See Also

[<strstream>](#)

## Name

flush function template — Manipulator to flush output buffer

## Synopsis

```
template <class charT, class traits>  
basic_ostream<charT,traits>& flush(basic_ostream<charT,traits>&  
os);
```

The flush function template is a manipulator that calls `os.flush`:

```
std::cout << "This is important!" << std::flush;
```

## See Also

[basic\\_ostream class template](#)

# Name

operator<< function template — Character output operator

## Synopsis

```
template<class charT, class traits>
basic_ostream<charT,traits>&
    operator<<(basic_ostream<charT,traits>& out, charT c);
template<class charT, class traits>
basic_ostream<charT,traits>&
    operator<<(basic_ostream<charT,traits>& out, char c);
template<class traits>
basic_ostream<char,traits>&
    operator<<(basic_ostream<char,traits>& out, char c);
template<class traits>
basic_ostream<char,traits>&
    operator<<(basic_ostream<char,traits>& out, signed char c);
template<class traits>
basic_ostream<char,traits>&
    operator<<(basic_ostream<char,traits>& out,
                unsigned char c);
template<class charT, class traits>
basic_ostream<charT,traits>&
    operator<<(basic_ostream<charT,traits>& out,
                const charT* s);
template<class charT, class traits>
basic_ostream<charT,traits>&
    operator<<(basic_ostream<charT,traits>& out,
                const char* s);
template<class traits>
basic_ostream<char,traits>&
    operator<<(basic_ostream<char,traits>& out, const char* s);
template<class traits>
basic_ostream<char,traits>&
    operator<<(basic_ostream<char,traits>& out,
                const signed char* s);
template<class traits>
basic_ostream<char,traits>&
    operator<<(basic_ostream<char,traits>& out,
                const unsigned char* s);
```

The << operator writes a single character *c*, or a character string *s*, to the output stream *out*. As with other formatted output functions, a sentry object is created, and the

character or string is written with appropriate padding. Each character is converted to the stream's character type by calling `widen`. Finally, `width(0)` is called.

## **See Also**

[basic\\_ostream class template](#), [basic\\_ostream::sentry class](#)

## Name

ostream class — Output stream

## Synopsis

```
typedef basic_ostream<char> ostream;
```

The ostream class specializes basic\_ostream for the char type.

## See Also

[basic\\_ostream class template](#), [wostream class](#), ostream in [<istream>](#)

## Name

wostream class — Wide output stream

## Synopsis

```
typedef basic_ostream<wchar_t> wostream;
```

The wostream class specializes basic\_ostream for the wchar\_t type.

## See Also

[basic\\_ostream class template](#), [ostream class](#), wiostream in [<istream>](#)

## <queue>

The <queue> header declares the `queue` and `priority_queue` container adapters. These class templates are not containers in their own rights, but they adapt containers to present the behavior of a queue or priority queue.

A *queue* is a sequence of items that supports insertion at one end and removal from the other end. Because the first item inserted into a queue is the first item removed, a queue is sometimes called a FIFO (first-in, first-out) container.

Instead of preserving FIFO order, a *priority queue* maintains heap order, which ensures that the largest item is always first. In strict C++ terms, the first item in a priority queue is not less than any other item in the queue. This is called the "largest" item, but you can also think of it as the most important item or the one with the highest priority.

See [Chapter 10](#) for information about containers.

## Name

`operator==` function template — Compares queues for equality

## Synopsis

```
template <typename T, typename Container>  
bool operator==(const queue<T, Container>& x, const queue<T,  
Container>& y);
```

The `==` operator compares two queues for equality by comparing the adapted containers (e.g., the return value is `x.c == y.c`).



## Name

`operator!=` function template — Compares queues for inequality

## Synopsis

```
template <typename T, typename Container>  
bool operator!=(const queue<T, Container>& x, const queue<T,  
Container>& y);
```

The `!=` operator compares two queues for inequality by comparing the adapted containers (e.g., the return value is `x.c != y.c`).

## Name

`operator<` function template — Compares queues for less-than

## Synopsis

```
template <typename T, typename Container>  
bool operator<(const queue<T, Container>& x, const queue<T,  
Container>& y);
```

The `<` operator compares two queues by comparing the adapted containers (e.g., the return value is `x.c < y.c`).

## Name

`operator<=` function template — Compares queues for less-than-or-equal

## Synopsis

```
template <typename T, typename Container>  
bool operator<=(const queue<T, Container>& x, const queue<T,  
Container>& y);
```

The `<=` operator compares two queues by comparing the adapted containers (e.g., the return value is `x.c <= y.c`).

## Name

`operator>` function template — Compares queues for greater-than

## Synopsis

```
template <typename T, typename Container>  
bool operator>(const queue<T, Container>& x, const queue<T,  
Container>& y);
```

The `>` operator compares two queues by comparing the adapted containers (e.g., the return value is `x.c >= y.c`).

## Name

`operator>=` function template — Compares queues for greater-than-or-equal

## Synopsis

```
template <typename T, typename Container>  
bool operator>=(const queue<T, Container>& x, const queue<T,  
Container>& y);
```

The `>=` operator compares two queues by comparing the adapted containers (e.g., the return value is `x.c >= y.c`).

# Name

priority\_queue class template — Priority queue container adapter

# Synopsis

```
template <typename T, typename Container = vector<T>,
          typename Compare = less<typename Container::value_type>
>
class priority_queue {
public:
    typedef typename Container::value_type value_type;
    typedef typename Container::size_type size_type;
    typedef Container container_type;

    explicit priority_queue(const Compare& x = Compare( ),
                           const Container& = Container( ));
    template <class InputIterator>
priority_queue(InputIterator first, InputIterator last,
                const Compare& x = Compare( ),
                const Container& = Container( ));
    bool empty( ) const { return c.empty( ); }
    size_type size( ) const { return c.size( ); }
    const value_type& top( ) const { return c.front( ); }
    void push(const value_type& x);
    void pop( );
protected:
    Container c;
    Compare comp;
};
```

The `priority_queue` class template is an adapter for any sequence container that supports random access, such as `deque` and `vector`. (The default is `vector`.) The priority queue keeps its elements in heap order, so it requires a comparator (the `Compare` template parameter).

Because `priority_queue` is not itself a standard container, it cannot be used with the standard algorithms. (In particular, note the lack of `begin` and `end` member

functions.) Thus, the `priority_queue` adapter is useful only for simple needs.

Unlike `queue`, `priority_queue` has no comparison operators.

Most of the members of `priority_queue` are straightforward mappings from a simple queue protocol to the underlying container protocol. The members are:

```
explicit priority_queue (const Compare& cmp =  
Compare( ), const Container& cont = Container( ))
```

Copies `cont` to the data member `c`, copies `cmp` to `comp`, and then calls `make_heap(c.begin( ), c.end( ), comp)` to initialize the priority queue.

```
template <class InputIter>, priority_queue  
(InputIter first, InputIter last, const Compare&  
cmp = Compare( ), const Container& cont =  
Container( ))
```

Copies `cont` to the data member `c`, copies `cmp` to `comp`, and then adds the elements `[first, last)` to the container by calling `c.insert(c.end( ), first, last)`. Finally, this method initializes the priority queue by calling `make_heap(c.begin( ), c.end( ), comp)`.

```
bool empty ( ) const
```

Returns `true` if the priority queue is empty.

```
void pop ( )
```

Erases the largest (last) item from the priority queue by calling `pop_heap` and then erasing the last element in the container.

```
void push (const value_type& x)
```

Inserts `x` in the container and then calls `push_heap` to restore priority queue order.

size\_type **size** ( ) const

Returns the number of items in the priority queue.

const value\_type& **top** ( ) const

Returns the largest (last) item in the priority queue.

## See Also

make\_heap, pop\_heap, and push\_heap in [<algorithm>](#),  
list in [<list>](#), vector in [<vector>](#)



# Name

queue class template — Queue container adapter

## Synopsis

```
template <class T, class Container = deque<T> >
class queue {
public:
    typedef typename Container::value_type value_type;
    typedef typename Container::size_type size_type;
    typedef Container container_type;

    explicit queue(const Container& = Container( ));
    bool empty( ) const { return c.empty( ); }
    size_type size( ) const { return c.size( ); }
    value_type& front( ) { return c.front( ); }
    const value_type& front( ) const { return c.front( ); }
    value_type& back( ) { return c.back( ); }
    const value_type& back( ) const { return c.back( ); }
    void push(const value_type& x) { c.push_back(x); }
    void pop( ) { c.pop_front( ); }
protected:
    Container c;
};
```

The queue class template is an adapter for any sequence container that supports the `front( )`, `back( )`, `push_back( )`, and `pop_front( )` members. See the `list` and `deque` class templates for the standard containers that are suitable. (The default is `deque`.)

Because `queue` is not itself a standard container, it cannot be used with the standard algorithms. (In particular, note the lack of `begin` and `end` member functions.) Thus, the queue adapter is useful only for simple needs.

Most of the members of `queue` are straightforward mappings from a simple queue protocol to the underlying container protocol. The members are:

explicit **queue** (const Container& cont = Container( ))

Takes an existing container cont and copies its contents into the queue. With no argument, the constructor creates a new, empty container for the queue.

value\_type& **back** ( ), const value\_type& **back** ( )  
const

Returns the last item in the queue, that is, the item that was added most recently to the queue.

bool **empty** ( ) const

Returns true if the queue is empty.

value\_type& **front** ( ), const value\_type& **front** ( )  
const

Returns the first item in the queue.

void **pop** ( )

Erases the first item from the queue.

void **push** (const value\_type& x)

Inserts x at the end of the queue.

size\_type **size** ( ) const

Returns the number of items in the queue.

## See Also

deque in [<deque>](#), list in [<list>](#), stack in [<stack>](#)

## **<set>**

The `<set>` header is one of the standard container template headers. It declares the `set` and `multiset` class templates and a few global function templates that operate on `set` and `multiset` objects.

A `set` is a container that stores keys. Looking up keys, inserting keys, and deleting keys can all be performed in logarithmic or better time. Sets support bidirectional iterators (not random access).

See [Chapter 10](#) for information about containers.

# Name

multiset class template — Set container with duplicate keys

## Synopsis

```
template <typename Key, typename Compare = less<Key>,
          typename Alloc = allocator<Key> >
class multiset {
public:
    typedef Key key_type;
    typedef Key value_type;
    typedef Compare key_compare;
    typedef Compare value_compare;
    typedef Alloc allocator_type;
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;

    explicit multiset(const Compare& comp = Compare( ), const
Alloc& = Alloc( ));
    template <class InputIterator>
multiset(InputIterator first, InputIterator last,
          const Compare& comp = Compare( ), const Alloc& =
Alloc( ));
    multiset(const multiset<Key,Compare,Alloc>& x);
    ~multiset( );
    multiset<Key,Compare,Alloc>& operator=(const
multiset<Key,Compare,Alloc>& x);
    allocator_type get_allocator( ) const;
    // Iterators
    iterator begin( );
    const_iterator begin( ) const;
    iterator end( );
    const_iterator end( ) const;
    reverse_iterator rbegin( );
    const_reverse_iterator rbegin( ) const;
```

```

reverse_iterator rend( );
const_reverse_iterator rend( ) const;

bool empty( ) const;
size_type size( ) const;
size_type max_size( ) const;

iterator insert(const value_type& x);
iterator insert(iterator hintpos, const value_type& x);
template <class InputIterator>
void insert(InputIterator first, InputIterator last);
void erase(iterator position);
size_type erase(const key_type& x);
void erase(iterator first, iterator last);
void swap(multiset<Key, Compare, Alloc>&);
void clear( );

key_compare key_comp( ) const;
value_compare value_comp( ) const;

iterator find(const key_type& x) const;
size_type count(const key_type& x) const;
iterator lower_bound(const key_type& x) const;
iterator upper_bound(const key_type& x) const;
pair<iterator, iterator> equal_range(const key_type& x) const;
};

```

The multiset class template is a standard container that contains an ordered set of keys of type T. The keys can be duplicated, that is, the multiset can contain more than one instance of a particular key.

A multiset's iterators are bidirectional. Note that keys are const in the set. You must not change the key while it is stored in a set. More precisely, you must not change the key in a way that alters its relative order with the other keys in the set. If you need to modify a key, erase the key from the set, modify the key, and insert the new key, as shown in [Example 13-34](#) (in the set class template).

Within a multiset, keys are ordered in ascending order, according to the Compare template parameter (which can be a function pointer or functor that compares two objects of type Key and returns true if the first argument should

come before the second). Keys do not need to be unique. When searching for keys, they are compared using the function or functor specified by the Compare template parameter. Two objects, a and b, are different if Compare(a, b) is true or Compare(b, a) is true. See set later in this section for a set container that stores unique keys.

Inserting into a multiset does not invalidate any iterators for that set or references to items in the set. Erasing an element invalidates only iterators and references that refer to that element.

Inserting into a set and searching for an element in a set usually take logarithmic time. Erasing a single element, given an iterator, takes constant time, amortized over many erasures.

The following are the member functions of multiset:

```
explicit multiset (const Compare& comp = Compare(
), const Alloc& = Alloc( ))
```

Constructs an empty multiset.

```
template <class InputIterator>, multiset
(InputIterator first, InputIterator last, const
Compare& comp = Compare( ), const Alloc& = Alloc( ))
```

Constructs an empty multiset and then copies all items in the range [first, last) into the new set. The complexity is linear if the keys in [first, last) are already sorted. If they are not sorted, the complexity is  $N \log N$ , in which  $N$  is last - first.

```
multiset (const multiset<Key, Compare, Alloc>& x)
```

Constructs a new multiset and copies the allocator and all the items from x to the new set.

```
iterator begin ( ), const_iterator begin ( ) const
```

Returns an iterator that points to the first element of the set.

`void clear ( )`

Erases every item in the set.

`size_type count (const key_type& x) const`

Returns the number of keys that are equivalent to `x`. The complexity is  $\log(\text{size}(\ )) + r$ , in which `r` is the return value.

`bool empty ( ) const`

Returns `size( ) == 0`.

`iterator end ( ), const_iterator end ( ) const`

Returns an iterator that points to one past the last element of the set.

`pair<iterator, iterator> equal_range (const key_type& x) const`

Returns the lower bound and upper bound as a pair:

```
std::make_pair(lower_bound(x), upper_bound(x))
```

The complexity is  $\log(\text{size}(\ ))$ .

`void erase (iterator position), size_type erase (const key_type& x), void erase (iterator first, iterator last)`

Erases one or more elements from the set. The first version erases the item at `position` in constant time (amortized over many calls). The second version erases the items equivalent to `x`, if any are present, returning a count of the number of items erased. The third version erases all elements in the range `[first, last)`. The last two forms run in time proportional to  $\log(\text{size}(\ )) + r$ , in which `r` is the number of items erased.

`iterator find (const key_type& x) const`

Searches for a key that is equivalent to `x` and returns an iterator that points to that key or `end( )` if it is not found. If `x` occurs more than once, the iterator might point to any of its occurrences in the multiset. The complexity is  $\log(\text{size}( ))$ .

`allocator_type` **get\_allocator** ( ) const

Returns the set's allocator.

`pair<iterator, bool>` **insert** (const `value_type&` `x`),  
`iterator` **insert** (`iterator` `hintpos`, const `value_type&` `x`),  
`template <class InputIterator>`, `void` **insert**  
(`InputIterator` `first`, `InputIterator` `last`)

Inserts one or more items into the set. The first version inserts `x` in logarithmic time.

The second version inserts `x` using `hintpos` as a position hint. If `x` is inserted immediately after `hintpos`, the performance is constant (amortized over many insertions); at any other position, the performance is logarithmic.

The third version copies all the items in the range `[first, last)`, which must be pointing to a different multiset object. If the items are already in the desired order, the performance is linear; otherwise, it is  $N \log(\text{size}( ) + N)$ , in which  $N$  is `last - first`.

`key_compare` **key\_comp** ( ) const

Returns the comparator function pointer or object. The `key_compare` type is the same as the `Compare` template parameter.

`iterator` **lower\_bound** (const `key_type&` `x`) const

Returns an iterator that points to the first item in the set that does not come before `x`. That is, if `x` is in the set, the iterator points to the position of its first occurrence;



otherwise, the iterator points to the first position where  $x$  should be inserted. The complexity is  $\log(\text{size}(\ ))$ .

`size_type` **max\_size** ( ) const

Returns the largest number of items that can be in the set.

`reverse_iterator` **rbegin** ( ), `const_reverse_iterator` **rbegin** ( ) const

Returns a reverse iterator that points to the last element of the set.

`reverse_iterator` **rend** ( ), `const_reverse_iterator` **rend** ( ) const

Returns a reverse iterator that points to one position before the first element of the set.

`size_type` **size** ( ) const

Returns the number of items in the set.

`void` **swap** (multiset<Key,Compare,Alloc>&)

Swaps the contents of the set with the contents of  $x$ .

`iterator` **upper\_bound** (const key\_type&  $x$ ) const

Returns an iterator that points to the first item in the set that comes after all occurrences of  $x$ . The complexity is  $\log(\text{size}(\ ))$ .

`value_compare` **value\_comp** ( ) const

Returns the comparator function pointer or functor object. The `value_compare` type is the same as the `Compare` template parameter.

`multiset<Key,Compare,Alloc>&` **operator=** (const `multiset<Key,Compare,Alloc>&`  $x$ )

Erases all the elements of the set and replaces them with copies of the elements of  $x$ .

## See Also

[set class template](#), multimap in [<map>](#)

## Name

`operator==` function template — Compares sets for equality

## Synopsis

```
template <typename Key, typename T, typename C, typename A>
bool operator==(const set<Key,T,C,A>& x, const set<Key,T,C,A>&
y);
template <typename Key, typename T, typename C, typename A>
bool operator==(const multiset<Key,T,C,A>& x, const
multiset<Key,T,C,A>& y);
```

The `==` operator returns true if `x` and `y` have the same size and their elements are equal, that is, `x.size( ) == y.size( ) && equals(x.begin( ), x.end( ), y.begin( ))`.

## See Also

`equals` in [<algorithm>](#)

## Name

`operator!=` function template — Compares sets for inequality

## Synopsis

```
template <typename Key, typename T, typename C, typename A>
bool operator!=(const set<Key,T,C,A>& x, const set<Key,T,C,A>&
y);
template <typename Key, typename T, typename C, typename A>
bool operator!=(const multiset<Key,T,C,A>& x, const
multiset<Key,T,C,A>& y);
```

The `!=` operator returns `! (x == y)`.

## Name

operator< function template — Compares sets for less-than

## Synopsis

```
template <typename Key, typename T, typename C, typename A>  
bool operator<(const set<Key,T,C,A>& x, const set<Key,T,C,A>& y);  
template <typename Key, typename T, typename C, typename A>  
bool operator<(const multiset<Key,T,C,A>& x, const  
multiset<Key,T,C,A>& y);
```

The < operator determines whether x is less than y using the same algorithm as

lexicographical\_compare(x.begin( ), x.end( ), y.begin( ), y.end( )).

## See Also

lexicographical\_compare in [<algorithm>](#)

## Name

`operator<=` function template — Compares sets for less-than-or-equal

## Synopsis

```
template <typename Key, typename T, typename C, typename A>
bool operator<=(const set<Key,T,C,A>& x, const set<Key,T,C,A>&
y);
template <typename Key, typename T, typename C, typename A>
bool operator<=(const multiset<Key,T,C,A>& x, const
multiset<Key,T,C,A>& y);
```

The `<=` operator returns `! (y < x)`.

## Name

`operator>` function template — Compares sets for greater-than

## Synopsis

```
template <typename Key, typename T, typename C, typename A>
bool operator>(const set<Key,T,C,A>& x, const set<Key,T,C,A>& y);
template <typename Key, typename T, typename C, typename A>
bool operator>(const multiset<Key,T,C,A>& x, const
multiset<Key,T,C,A>& y);
```

The `>` operator returns `(y < x)`.

## Name

`operator>=` function template — Compares sets for greater-than-or-equal

## Synopsis

```
template <typename Key, typename T, typename C, typename A>
bool operator>=(const set<Key,T,C,A>& x, const set<Key,T,C,A>&
y);
template <typename Key, typename T, typename C, typename A>
bool operator>=(const multiset<Key,T,C,A>& x, const
multiset<Key,T,C,A>& y);
```

The `>=` operator returns `! (x < y)`.



# Name

set class template — Set container with unique keys

## Synopsis

```
template <typename Key, typename Compare = less<Key>,
          typename Alloc = allocator<Key> >
class set {
public:
    typedef Key key_type;
    typedef Key value_type;
    typedef Compare key_compare;
    typedef Compare value_compare;
    typedef Alloc allocator_type;
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;

    explicit set(const Compare& comp = Compare( ),
                const Alloc& = Alloc( ));
    template <class InputIterator>
set(InputIterator first, InputIterator last,
     const Compare& comp = Compare( ), const Alloc& = Alloc(
));
    set(const set<Key,Compare,Alloc>& x);
    ~set( );
    set<Key,Compare,Alloc>& operator=(const set<Key,Compare,Alloc>&
x);

    allocator_type get_allocator( ) const;

    iterator begin( );
    const_iterator begin( ) const;
    iterator end( );
    const_iterator end( ) const;
    reverse_iterator rbegin( );
```

```

const_reverse_iterator rbegin( ) const;
reverse_iterator rend( );
const_reverse_iterator rend( ) const;

bool empty( ) const;
size_type size( ) const;
size_type max_size( ) const;

pair<iterator,bool> insert(const value_type& x);
iterator insert(iterator hintpos, const value_type& x);
template <class InputIterator>
void insert(InputIterator first, InputIterator last);
void erase(iterator position);
size_type erase(const key_type& x);
void erase(iterator first, iterator last);
void swap(set<Key,Compare,Alloc>&);
void clear( );
// Observers
key_compare key_comp( ) const;
value_compare value_comp( ) const;
// Set operations
iterator find(const key_type& x) const;
size_type count(const key_type& x) const;
iterator lower_bound(const key_type& x) const;
iterator upper_bound(const key_type& x) const;
pair<iterator,iterator> equal_range(const key_type& x) const;
};

```

The set class template is a standard container that contains an ordered set of unique keys of type T.

A set's iterators are bidirectional. Note that keys are const in the set. You must not change the key while it is stored in a set. More precisely, you must not change the key in a way that alters its relative order with the other keys in the set. If you need to modify a key, erase the key from the set, modify the key, and insert the new key, as shown in [Example 13-34](#).

## Example

*Example 13-34. One way to modify a key in a set*

```

template <typename T, typename C, typename A>
void change_key(std::set<T, C, A>& s,

```

```

const T& oldkey, const T& newkey)
{
using std::set;
typedef typename set<T, C, A>::iterator set_iterator;
set_iterator i = s.find(oldkey);
if (i != s.end( )) {
    m.erase(i);
    m.insert(newkey);
}
// Exercise: What if newkey is already in s?
}

```

Within a set, keys are ordered in ascending order, according to the Compare template parameter (which can be a function pointer or functor that compares two objects of type Key and returns true if the first argument should come before the second). Keys must be unique, but note that uniqueness is determined only by calling Compare, not by using the == operator. That is, two objects, a and b, are different (and therefore can both be present in a single set object) if Compare(a, b) is true or Compare(b, a) is true. See multiset earlier in this section for a set container that can store non-unique keys.

Inserting into a set does not invalidate any iterators for that set or any references to items in the set. Erasing an element invalidates only iterators or references that refer to that element.

Inserting into a set and searching for an element in a set usually take logarithmic time. Erasing a single element, given an iterator, takes constant time, amortized over many erasures.

The following are the member functions of set:

```

explicit set (const Compare& comp = Compare( ),
const Alloc& = Alloc( ))

```

Constructs an empty set.

```
template <class InputIterator>, set (InputIterator  
first, InputIterator last, const Compare& comp =  
Compare( ), const Alloc& = Alloc( ))
```

Constructs an empty set and then copies all items in the range [first, last) into the new set. The complexity is linear if the keys in [first, last) are already sorted. If they are not sorted, the complexity is  $N \log N$ , in which  $N$  is last - first.

```
set (const set<Key, Compare, Alloc>& x)
```

Constructs a new set and copies the allocator and all the items from x to the new set.

```
iterator begin ( ), const_iterator begin ( ) const
```

Returns an iterator that points to the first element of the set.

```
void clear ( )
```

Erases every item in the set.

```
size_type count (const key_type& x) const
```

Returns the number of keys that are equivalent to x. This value is always 0 or 1. The complexity is  $\log(\text{size}())$ .

```
bool empty ( ) const
```

Returns  $\text{size}() == 0$ .

```
iterator end ( ), const_iterator end ( ) const
```

Returns an iterator that points to one past the last element of the set.

```
pair<iterator, iterator> equal_range (const  
key_type& x) const
```

Returns the lower bound and upper bound as a pair:

```
std::make_pair(lower_bound(x), upper_bound(x))
```

The complexity is  $\log(\text{size}())$ .

```
void erase (iterator position), size_type erase  
(const key_type& x), void erase (iterator first,  
iterator last)
```

Erases one or more elements from the set. The first version erases the item at position in constant time (amortized over many calls). The second version erases the item equivalent to x, if it is present, returning a count of the number of items erased, that is, 0 or 1. It runs in logarithmic time. The third version erases all elements in the range [first, last) in a time proportional to  $\log(\text{size}()) + r$ , in which r is last - first.

```
iterator find (const key_type& x) const
```

Searches for a key that is equivalent to x and returns an iterator that points to that key or end() if it is not found. The complexity is  $\log(\text{size}())$ .

```
allocator_type get_allocator () const
```

Returns the set's allocator.

```
pair<iterator, bool> insert (const value_type& x),  
iterator insert (iterator hintpos, const value_type&  
x), template <class InputIterator>, void insert  
(InputIterator first, InputIterator last)
```

Inserts one or more items into the set, but only if an equivalent key is not already present in the set. If the key is already present, the insert attempt is ignored. The first version attempts to insert x in logarithmic time.

The second version inserts x using hintpos as a position hint. If x is inserted immediately after hintpos, the performance is constant (amortized over many insertions); at any other position, the performance is logarithmic.

The third version copies all the items in the range [first, last), which must be pointing to a different set object. If the items are already in the desired order, the performance is linear; otherwise, it is  $N \log(\text{size}() + N)$ , in which  $N$  is last - first.

key\_compare **key\_comp** ( ) const

Returns the comparator function pointer or object. The key\_compare type is the same as the Compare template parameter.

iterator **lower\_bound** (const key\_type& x) const

Returns an iterator that points to the first item in the set that does not come before x. That is, if x is in the set, the iterator points to its position; otherwise, the iterator points to the first position where x should be inserted. The complexity is  $\log(\text{size}())$ .

size\_type **max\_size** ( ) const

Returns the largest number of items that can be in the set.

reverse\_iterator **rbegin** ( ), const\_reverse\_iterator **rbegin** ( ) const

Returns a reverse iterator that points to the last element of the set.

reverse\_iterator **rend** ( ), const\_reverse\_iterator **rend** ( ) const

Returns a reverse iterator that points to one position before the first element of the set.

size\_type **size** ( ) const

Returns the number of items in the set.

void **swap** (set<Key, Compare, Alloc>&)

Swaps the contents of the set with the contents of x.

iterator **upper\_bound** (const key\_type& x) const

Returns an iterator that points to the first item in the set that comes after x. The complexity is  $\log(\text{size}())$ .

value\_compare **value\_comp** ( ) const

Returns the comparator function pointer or functor object. The value\_compare type is the same as the Compare template parameter.

set<Key, Compare, Alloc>& **operator=** (const set<Key, Compare, Alloc>& x)

Erases all the elements of the set and replaces them with copies of the elements of x.

## See Also

[multiset class template](#), multimap in [<map>](#)

## Name

swap function template — Swaps the contents of two sets

## Synopsis

```
template <typename Key, typename T, typename C, typename A>  
void swap(set<Key,T,C,A>& x,  
          set<Key,T,C,A>& y);  
template <typename Key, typename T, typename C, typename A>  
void swap(multiset<Key,T,C,A>& x,  
          multiset<Key,T,C,A>& y);
```

The swap function template specialization is equivalent to calling `x.swap(y)`.

## See Also

swap in [<algorithm>](#)



## <sstream>

The <sstream> header declares classes, templates, and other types for reading from and writing to strings in the same manner as reading from and writing to files.

See [Chapter 9](#) for a general discussion of I/O, [Chapter 1](#) for more information about character sets, and the <iostream> section in this chapter for information about the base-class templates required by the stringstream class templates. Refer to [Chapter 8](#) for information about traits in general and to the <string> section in this chapter for detailed information about the char\_traits template. Refer to the <streambuf> section in this chapter for information about the basic\_streambuf template. See also <strstream> for classes that are similar to the string streams, except they work with arrays of narrow characters.

To read from a string, use istream; for writing, use ostream; for reading and writing, use stringstream. For wide character I/O, use wistream, wostream, or wstringstream. [Example 13-35](#) shows tostring, a simple use of ostream to convert a value to a string. (Think of tostring as the inverse of strtol and friends.)

*Example 13-35. Converting a value to a string*

```
template<typename T>
std::string tostring(const T& x)
{
    std::ostream out;
    out << x;
    return out.str( );
}
```

[Example 13-36](#) shows a use of istream to interpret HTML colors. In HTML, a color can be a name, such as white, or a hexadecimal digit string that begins with #. The

digit string is interpreted as a triplet of red, green, and blue color elements, each expressed as two hexadecimal digits. For the sake of simplicity, the example omits error handling and assumes that the order of the color elements matches the order needed by the program. The known color names are stored in a map.

### *Example 13-36. Interpreting an HTML color string*

```
typedef std::map<std::string, unsigned long> colormap;
colormap colors;

unsigned long get_color(const std::string& text)
{
    unsigned long rgb;
    colormap::iterator i = colors.find(text);
    if (i != colors.end( ))
        return i->second;
    else if (text.length( ) == 0)
        return 0;
    else {
        std::istringstream in(text);
        if (in.peek( ) == '#')
            in.ignore( );
        in >> std::noskipws >> std::hex >> rgb;
        if (in)
            return rgb;
        else
            return 0;
    }
}

void initcolors(colormap& colors)
{
    . . .
    colors["black"] = 0x000000;
    colors["blue"] = 0x0000FF;
    colors["green"] = 0x00FF00;
    colors["red"] = 0xFF0000;
    colors["white"] = 0xFFFFFFFF;
}
```

# Name

`basic_istream` class template — Base class for input string streams

# Synopsis

```
template <class charT, class traits = char_traits<charT>,
          class Alloc = allocator<charT> >
class basic_istream: public basic_istream<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    explicit basic_istream(ios_base::openmode which =
ios_base::in);
    explicit basic_istream(const
basic_string<charT,traits,Alloc>& str,
                          ios_base::openmode which =
ios_base::in);

    basic_stringbuf<charT,traits,Alloc>* rdbuf( ) const;
    basic_string<charT,traits,Alloc> str( ) const;
    void str(const basic_string<charT,traits,Alloc>& s);
};
```

The `basic_istream` class template is the base class for input string streams. Typically, you would construct an `istream` with a string argument and then read from the string stream just as you would from any other input stream.

The following are the methods of `basic_istream`:

```
explicit basic_istream (ios_base::openmode
which = ios_base::in)
```

Initializes an empty input string stream by constructing an internal `basic_stringbuf` object, passing which |

`ios_base::in` to that object's constructor, and passing the address of the string buffer to the base-class constructor for `basic_istream`.

```
explicit basic_istream (const  
basic_string<charT,traits,Alloc>& str,  
ios_base::openmode which = ios_base::in)
```

Initializes a string stream with `str` as the initial string contents by constructing an internal `basic_stringbuf` object, passing `str` and `which | ios_base::in` to that object's constructor, and passing the address of the string buffer to the base-class constructor for `basic_istream`.

```
basic_stringbuf<charT,traits,Alloc>* rdbuf ( )  
const
```

Returns a pointer to the internal `basic_stringbuf` object.

```
basic_string<charT,traits,Alloc> str ( ) const
```

Returns the buffer contents as a string, that is, `rdbuf( )->str( )`.

```
void str (const basic_string<charT,traits,Alloc>&  
s)
```

Calls `rdbuf( )->str(s)` to set the buffer contents.

## See Also

[basic\\_stringstream class template](#), [istream class](#), [wistream class](#), `basic_ifstream` in [<fstream>](#), `basic_istream` in [<istream>](#)

# Name

`basic_ostringstream` class template — Base class for output string streams

# Synopsis

```
template <class charT, class traits = char_traits<charT>,
          class Alloc = allocator<charT> >
class basic_ostringstream: public basic_ostream<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    explicit basic_ostringstream(ios_base::openmode which =
ios_base::out);
    explicit basic_ostringstream(const
basic_string<charT,traits,Alloc>& str,
                                ios_base::openmode which =
ios_base::out);

    basic_stringbuf<charT,traits,Alloc>* rdbuf( ) const;
    basic_string<charT,traits,Alloc> str( ) const;
    void str(const basic_string<charT,traits,Alloc>& s);
};
```

The `basic_ostringstream` class template is the base class for output string streams. Typically, you would construct an `ostringstream` with no string and let the string stream allocate the string as you write to the stream. You would then call the `str` member function to read the resulting string.

The following are the methods of `basic_ostringstream`:

```
explicit basic_ostringstream (ios_base::openmode
which = ios_base::out)
```

Initializes an empty output string stream by constructing an internal `basic_stringbuf` object, passing which | `ios_base::out` to that object's constructor, and passing the address of the string buffer to the base-class constructor for `basic_ostream`.

```
explicit basic_ostream (const  
basic_string<charT,traits,Alloc>& str,  
ios_base::openmode which = ios_base::out)
```

Initializes a string stream with `str` as the initial string contents by constructing an internal `basic_stringbuf` object, passing `str` and which | `ios_base::out` to that object's constructor, and passing the address of the string buffer to the base-class constructor for `basic_ostream`.

```
basic_stringbuf<charT,traits,Alloc>* rdbuf ( )  
const
```

Returns a pointer to the internal `basic_stringbuf` object.

```
basic_string<charT,traits,Alloc> str ( ) const
```

Returns the buffer contents as a string, that is, `rdbuf( )->str( )`.

```
void str (const basic_string<charT,traits,Alloc>&  
s)
```

Calls `rdbuf( )->str(s)` to set the buffer contents.

## See Also

[basic\\_stringstream class template](#), [ostream class](#), [wostringstream class](#), `basic_ofstream` in [<fstream>](#), `basic_ostream` in [<ostream>](#)

# Name

`basic_stringbuf` class template — Base class for string buffers

# Synopsis

```
template <class charT, class traits = char_traits<charT>,
          class Alloc = allocator<charT> >
class basic_stringbuf : public basic_streambuf<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;
    explicit basic_stringbuf(ios_base::openmode mode = ios_base::in
|
|                               ios_base::out);
    explicit basic_stringbuf(const
basic_string<charT,traits,Alloc>& str,
|
|                               ios_base::openmode mode = ios_base::in
|
|                               ios_base::out);
    basic_string<charT,traits,Alloc> str( ) const;
    void str(const basic_string<charT,traits,Alloc>& s);
protected:
    virtual int_type underflow( );
    virtual int_type pbackfail(int_type c = traits::eof( ));
    virtual int_type overflow (int_type c = traits::eof( ));
    virtual basic_streambuf<charT,traits>* setbuf(charT*,
streamsize);
    virtual pos_type seekoff(off_type off, ios_base::seekdir way,
ios_base::in |
|                               ios_base::out);
    virtual pos_type seekpos(pos_type sp,
ios_base::in |
|                               ios_base::out);
};
```

The `basic_stringbuf` class template implements a stream buffer for string-based streams. A string buffer maintains a

single character buffer with separate positions for reading and writing. That is, the buffer has `begin`, `next`, and `end` pointers for reading and separate `begin`, `next`, and `end` pointers for writing. The `begin` pointer points to the start of a buffer, and the `end` pointer points to one past the end of the buffer. The `next` pointer points to the position where the next character will be read or written. Refer to `basic_streambuf` in `<streambuf>` for details about buffer positions.

In the following descriptions of the member functions of `basic_stringbuf`, `mode` refers to a private copy of the `mode` parameter that is passed to the constructors. The implementation is not required to have such a data member, but the descriptions below are clearer with the assumption that it exists.

```
explicit basic_stringbuf (ios_base::openmode mode  
= ios_base::in | ios_base::out)
```

Initializes the buffer with an empty string and remembers the mode.

```
explicit basic_stringbuf (const  
basic_string<charT,traits,Alloc>& str,  
ios_base::openmode mode = ios_base::in |  
ios_base::out)
```

Initializes the buffer with a copy of `str` and remembers the mode. If `mode & ios_base::in` is nonzero, the input position is initialized to read from the start of the buffer. If `mode & ios_base::out` is nonzero, the output position is initialized to overwrite the buffer.

```
virtual int_type overflow (int_type c =  
traits::eof( ))
```

Attempts to append `c` to the end of the buffer as follows:



- If `c` is an end-of-file character (`c` is `traits::eof( )`), nothing happens, and a non-end-of-file character is returned to indicate success.
- If `c` is not end-of-file, and a write position is available, `c` is appended to the buffer by calling `sputc(c)`.
- If a write position is not available, and the mode allows writing (`mode & ios_base::out` is nonzero), the buffer is extended by one character and `c` is appended. If the mode allows reading (`mode & ios_base::in` is nonzero), the read end pointer `egptr( )` is set to point to one position past the end of the buffer.

The return value is `traits::not_eof(c)` for success or `traits::eof( )` for failure.

virtual int\_type **backfail** (int\_type c = `traits::eof( )`)

Attempts to push `c` back onto the buffer for reading as follows:

- If `c` is an end-of-file character (`c` is `traits::eof( )`), and a putback position is available, `gptr( )` is set to `gptr( ) - 1`.
- If `c` is not an end-of-file character, and a putback position is available, and `gptr( )[-1]` is equal to `c`, `gptr( )` is set to `gptr( ) - 1`.
- If `c` is not an end-of-file character, and a putback position is available, and the mode allows writing (`mode & ios_base::out` is nonzero), `gptr( )` is set to `gptr( ) - 1`, and `*gptr( )` is assigned `c`.

The return value is `traits::not_eof(c)` for success or `traits::eof( )` for failure.

virtual pos\_type **seekoff** (off\_type off,  
ios\_base::seekdir way, ios\_base::openmode which =  
ios\_base::in|ios\_base::out)

Sets the stream position. The input position, output position, or both can be set, depending on (which & (ios\_base::in | ios\_base::out)). The following are the possible results of this expression:

ios\_base::in

Sets the input position

ios\_base::out

Sets the output position

ios\_base::in | ios\_base::out, and way is either ios\_base::beg or ios\_base::end

Sets input and output positions

Otherwise


The function fails and returns pos\_type(-1)

The new position is determined as an offset off, which is added to the position at the start of the stream, the current position, or at the end of the stream, depending on way (ios\_base::beg, ios\_base::cur, or ios\_base::end). If the desired position is negative or past the end of the buffer, the function fails and returns pos\_type(-1). If the function succeeds, it returns the new position.

virtual pos\_type **seekpos** (pos\_type sp,  
ios\_base::openmode which =  
ios\_base::in|ios\_base::out)

Sets the stream position to sp. The input position is set if which & ios\_base::in is nonzero. The output position is set if which & ios\_base::out is nonzero. If sp is not a valid position, or if neither the input nor the output

position is set, seekpos fails and pos\_type(-1) is returned. The return value is sp for success. If sp was not returned from a prior call to a positioning function (that is, seekoff, seekpos, tellg, or tellp), the results are undefined.

 virtual basic\_streambuf<charT,traits>\* **setbuf**  
(charT\*, streamsize)

Calling setbuf(0, 0) has no effect other than to return this. The result of any other call to setbuf is implementation-defined.

basic\_string<charT,traits,Alloc> **str** ( ) const

Returns the contents of the buffer as a string. If the mode allows output (mode & ios\_base::out is nonzero), the buffer contents are taken from the output positions; otherwise, the buffer contents are copied from the input positions.

void **str** (const basic\_string<charT,traits,Alloc>& s)

Deallocates the current buffer if one exists and replaces it with a copy of s. If mode & ios\_base::in is nonzero, the input positions are set to read from the start of the buffer. If mode & ios\_base::out is nonzero, the output positions are set to overwrite the buffer.

virtual int\_type **underflow** ( )

Returns \*gptr( ) if more input is available, that is, if there is a read position. Otherwise, the function returns traits::eof( ).

## See Also

[stringbuf class](#), [wstringbuf class](#), basic\_filebuf in [<fstream>](#), basic\_streambuf in [<streambuf>](#)

# Name

`basic_stringstream` class template — Base class for input and output string streams

# Synopsis

```
template <class charT, class traits = char_traits<charT>,
          class Alloc = allocator<charT> >
class basic_stringstream: public basic_istream<charT,traits>
{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;

    explicit basic_stringstream(ios_base::openmode which =
                                ios_base::out|ios_base::in);
    explicit basic_stringstream(const
basic_string<charT,traits,Alloc>& str,
                                ios_base::openmode which =
                                ios_base::out|ios_base::in);

    basic_stringbuf<charT,traits,Alloc>* rdbuf( ) const;
    basic_string<charT,traits,Alloc> str( ) const;
    void str(const basic_string<charT,traits,Alloc>& str);
};
```

The `basic_stringstream` class template is the base class for string streams that permit input and output. You can start with an empty string and write to the stream, or start with a string and read from the stream. If you initialize the stream with a string and start writing, the output overwrites the string. You can switch between reading and writing at any time and set the read and write positions independently.

The following are the methods of `basic_stringstream`:

```
explicit basic_stringstream (ios_base::openmode  
which = ios_base::in | ios_base::out)
```

Initializes an empty string stream by constructing an internal `basic_stringbuf` object, passing `which | ios_base::in | ios_base::out` to that object's constructor and the address of the string buffer to the base-class constructor for `basic_istream`.

```
explicit basic_stringstream (const  
basic_string<charT,traits,Alloc>& str,  
ios_base::openmode which =  
ios_base::in|ios_base::out)
```

Initializes a string stream with `str` as the initial string contents by constructing an internal `basic_stringbuf` object, passing `str` and `which | ios_base::in | ios_base::out` to that object's constructor and the address of the string buffer to the base-class constructor for `basic_istream`.

```
basic_stringbuf<charT,traits,Alloc>* rdbuf ( )  
const
```

Returns a pointer to the internal `basic_stringbuf` object.

```
basic_string<charT,traits,Alloc> str ( ) const
```

Returns the buffer contents as a string, that is, `rdbuf( )->str( )`.

```
void str (const basic_string<charT,traits,Alloc>&  
s)
```

Calls `rdbuf( )->str(s)` to set the buffer contents.

## See Also

[basic\\_istream class template](#), [basic\\_ostringstream class template](#), [stringstream class](#), [wstringstream class](#),

basic\_fstream in [<fstream>](#), basic\_iostream in [<iostream>](#)

## Name

istringstream class — Input string stream

## Synopsis

```
typedef basic_istringstream<char> istringstream;
```

The istringstream class is a specialization of the basic\_istringstream template for char characters.

## See Also

[basic\\_istringstream class template](#), [wistringstream class](#), ifstream in [<fstream>](#), istream in [<istream>](#)

## Name

ostringstream class — Output string stream

## Synopsis

```
typedef basic_ostream<char> ostringstream;
```

The `ostringstream` class is a specialization of the `basic_ostream` template for `char` characters.

## See Also

[basic\\_ostream class template](#), [wostringstream class](#), `ofstream` in [<fstream>](#), `ostream` in [<ostream>](#)



## Name

stringbuf class — Narrow character string buffer

## Synopsis

```
typedef basic_stringbuf<char> stringbuf;
```

The stringbuf class is a specialization of the basic\_stringbuf template for char characters.

## See Also

[basic\\_stringbuf class template](#), [wstringbuf class](#), filebuf in [<fstream>](#), streambuf in [<streambuf>](#)

## Name

stringstream class — Input and output string stream

## Synopsis

```
typedef basic_stringstream<char> stringstream;
```

The stringstream class is a specialization of the basic\_stringstream template for char characters.

## See Also

[basic\\_stringstream class template](#), [wstringstream class](#),  
fstream in [<fstream>](#), istream in [<istream>](#)

## Name

wistringstream class — Wide input string stream

## Synopsis

```
typedef basic_istreamstream<wchar_t> wistringstream;
```

The wistringstream class is a specialization of the basic\_istreamstream template for wchar\_t characters.

## See Also

[basic\\_istreamstream class template](#), [istreamstream class](#), wifstream in [<fstream>](#), wistream in [<istream>](#)

## Name

wostringstream class — Wide output string stream

## Synopsis

```
typedef basic_ostringstream<wchar_t> wostringstream;
```

The `wostringstream` class is a specialization of the `basic_ostringstream` template for `wchar_t` characters.

## See Also

[basic\\_ostringstream class template](#), [ostringstream class](#),  
wofstream in [<fstream>](#), wostream in [<ostream>](#)

## Name

wstringbuf class — Wide character string buffer

## Synopsis

```
typedef basic_stringbuf<wchar_t> wstringbuf;
```

The `wstringbuf` class is a specialization of the `basic_stringbuf` template for `wchar_t` characters.

## See Also

[basic\\_stringbuf class template](#), [stringbuf class](#), `wfilebuf` in [<fstream>](#), `wstreambuf` in [<streambuf>](#)

## Name

wstringstream class — Wide input and output string stream

## Synopsis

```
typedef basic_stringstream<wchar_t> wstringstream;
```

The wstringstream class is a specialization of the basic\_stringstream template for wchar\_t characters.

## See Also

[basic\\_stringstream class template](#), [stringstream class](#),  
wfstream in [<fstream>](#), wiostream in [<iostream>](#)

## **<stack>**

The `<stack>` header declares the `stack` container adapter. This class template is not a container in its own right, but adapts other containers to present the behavior of a stack.

A stack is a sequence of items that supports insertion and removal at one end. Because the last item inserted into a stack is the first item removed, a stack is sometimes called a LIFO (last-in, first-out) container.

See [Chapter 10](#) for information about containers.

## Name

`operator==` function template — Compares stacks for equality

## Synopsis

```
template <typename T, typename Container>  
bool operator==(const stack<T, Container>& x, const stack<T,  
Container>& y);
```

The `==` operator compares two stacks for equality by comparing the adapted containers (e.g., the return value is `x.c == y.c`).



## Name

`operator!=` function template — Compares stacks for inequality

## Synopsis

```
template <typename T, typename Container>  
bool operator!=(const stack<T, Container>& x, const stack<T,  
Container>& y);
```

The `!=` operator compares two stacks for inequality by comparing the adapted containers (e.g., the return value is `x.c != y.c`).

## Name

`operator<` function template — Compares stacks for less-than

## Synopsis

```
template <typename T, typename Container>  
bool operator<(const stack<T, Container>& x, const stack<T,  
Container>& y);
```

The `<` operator compares two stacks by comparing the adapted containers (e.g., the return value is `x.c < y.c`).

## Name

`operator<=` function template — Compares stacks for less-than-or-equal

## Synopsis

```
template <typename T, typename Container>  
bool operator<=(const stack<T, Container>& x, const stack<T,  
Container>& y);
```

The `<=` operator compares two stacks by comparing the adapted containers (e.g., the return value is `x.c <= y.c`).

## Name

`operator>` function template — Compares stacks for greater-than

## Synopsis

```
template <typename T, typename Container>  
bool operator>(const stack<T, Container>& x, const stack<T,  
Container>& y);
```

The `>` operator compares two stacks by comparing the adapted containers (e.g., the return value is `x.c >= y.c`).

## Name

`operator>=` function template — Compares stacks for greater-than-or-equal

## Synopsis

```
template <typename T, typename Container>  
bool operator>=(const stack<T, Container>& x, const stack<T,  
Container>& y);
```

The `>=` operator compares two stacks by comparing the adapted containers (e.g., the return value is `x.c >= y.c`).

# Name

stack class template — Stack container adapter

## Synopsis

```
template <class T, class Container = deque<T> >
class stack {
public:
    typedef typename Container::value_type value_type;
    typedef typename Container::size_type size_type;
    typedef Container container_type;
protected:
    Container c;
public:
    explicit stack(const Container& = Container( ));
    bool empty( ) const { return c.empty( ); }
    size_type size( ) const { return c.size( ); }
    value_type& top( ) { return c.back( ); }
    const value_type& top( ) const { return c.back( ); }
    void push(const value_type& x) { c.push_back(x); }
    void pop( ) { c.pop_back( ); }
};
```

The `stack` class template is an adapter for any sequence container—such as `deque`, `list`, and `vector`—that supports the `back`, `push_back`, and `pop_back` members. (The default is `deque`.)

Because `stack` is not itself a standard container, it cannot be used with the standard algorithms. (In particular, note the lack of `begin` and `end` member functions.) Thus, the `stack` adapter is useful only for simple needs.

Most of the members of `stack` are straightforward mappings from a simple stack protocol to the underlying container protocol. The members are:

```
explicit stack (const Container& cont = Container(
))
```

Copies the elements from `cont` to the `c` data member

bool **empty** ( ) const

Returns true if the stack is empty

void **pop** ( )

Erases the item at the top of the stack

void **push** (const value\_type& x)

Adds x at the top of the stack

size\_type **size** ( ) const

Returns the number of items in the stack

value\_type& **top** ( ), const value\_type& **top** ( ) const

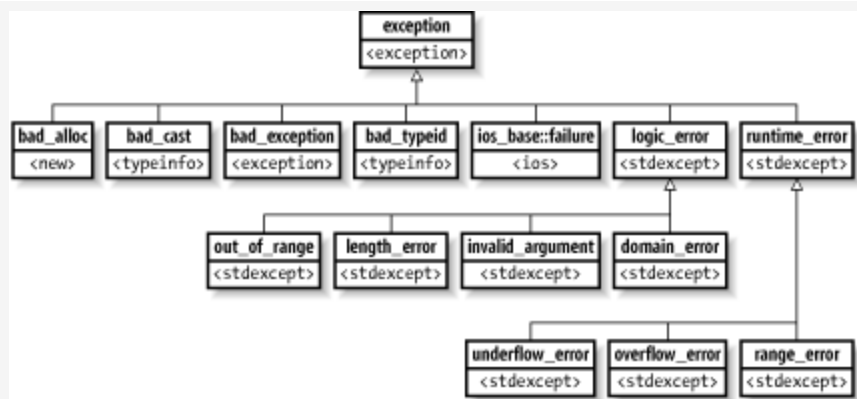
Returns the item at the top of the stack

## See Also

[<deque>](#), [<list>](#), [<queue>](#), [<vector>](#)

## <stdexcept>

The <stdexcept> header defines several standard exception classes. (Refer to <exception> for the base exception class.) [Figure 13-24](#) shows all the exception classes in the standard library, including a few that are declared in other headers. Note that the standard library has very few places that throw exceptions. The exceptions in <stdexcept> are available primarily for your use.



*Figure 13-24. All the standard exception classes*

See [Chapter 3](#) for more information about throw expressions and [Chapter 4](#) for information about handling exceptions with try statements.



## Name

`domain_error` class — Arithmetic domain error

## Synopsis

```
class domain_error : public logic_error {  
public:  
    explicit domain_error(const string& what_arg);  
};
```

The `domain_error` class is used to report domain errors, that is, arguments to functions that are outside the valid domain for input to the functions. For example, a function that converts a color from the Hue, Saturation, Lightness colorspace to the Red, Green, Blue colorspace might require a Saturation in the range [0.0, 1.0] and throw `domain_error` for any other value.

## See Also

[logic\\_error class](#)

## Name

`invalid_argument` class — Invalid function argument

## Synopsis

```
class invalid_argument : public logic_error {  
public:  
    explicit invalid_argument(const string& what_arg);  
};
```

The `invalid_argument` class is thrown to report invalid arguments to functions. Specific kinds of invalid arguments are covered by the other logic errors; use `invalid_argument` for any other situations. For example, constructing a `bitset` from a `string` throws `invalid_argument` if any character is other than '0' or '1'.

## See Also

[logic\\_error class](#)

## Name

`length_error` class — Exceed maximum size

## Synopsis

```
class length_error : public logic_error {  
public:  
    explicit length_error(const string& what_arg);  
};
```

The `length_error` class is used to attempt to set or change the size of an object that exceeds the maximum size. For example, the `string` class throws `length_error` if you attempt to create a string longer than `max_size( )` characters.

## See Also

[logic\\_error class](#)

## Name

`logic_error` class — Base class for logic errors

## Synopsis

```
class logic_error : public exception {  
public:  
    explicit logic_error(const string& what_arg);  
};
```

The `logic_error` class is a base class for logic-error exceptions. A *logic error* is a violation of the preconditions or other requirements for a function.

## See Also

[domain\\_error class](#), [invalid\\_argument class](#), [length\\_error class](#), [out\\_of\\_range class](#), [runtime\\_error class](#)

## Name

out\_of\_range class — Argument out of range

## Synopsis

```
class out_of_range : public logic_error {  
public:  
    explicit out_of_range(const string& what_arg);  
};
```

The `out_of_range` class is used when an index or similar value is out of its expected or allowed range. For example, the `at` member (of `deque`, `string`, and `vector`) throws `out_of_range` if the index is invalid.

## See Also

[logic\\_error class](#)

## Name

`overflow_error` class — Arithmetic overflow

## Synopsis

```
class overflow_error : public runtime_error {  
public:  
    explicit overflow_error(const string& what_arg);  
};
```

The `overflow_error` class can be used for arithmetic overflow. For example, `bitset::to_ulong` throws `overflow_error` if the arithmetic value of the `bitset` exceeds the maximum value of an unsigned long.



Note that overflow in most arithmetic expressions has undefined behavior; an implementation might throw `overflow_error`, but there is no guarantee that it will throw this or any other exception.

## See Also

[runtime\\_error class](#)

## Name

range\_error class — Arithmetic range error

## Synopsis

```
class range_error : public runtime_error {  
public:  
    explicit range_error(const string& what_arg);  
};
```

The range\_error class can be used when a function's results would fall outside its valid range. Note that the <cmath> functions do not throw any exceptions, but a third-party math library might throw range\_error for, say, computing a power that exceeds the limits of its return type.

## See Also

[runtime\\_error class](#)

## Name

`runtime_error` class — Base class for runtime errors

## Synopsis

```
class runtime_error : public exception {  
public:  
    explicit runtime_error(const string& what_arg);  
};
```

The `runtime_error` class is the base class for runtime errors, which are errors that cannot reasonably be detected by a static analysis of the code but can be revealed only at runtime.

## See Also

[overflow\\_error class](#), [range\\_error class](#), [underflow\\_error class](#)



## Name

`underflow_error` class — Arithmetic underflow

## Synopsis

```
class underflow_error : public runtime_error {  
public:  
    explicit underflow_error(const string& what_arg);  
};
```



The `underflow_error` class can be used for arithmetic underflow. Note that underflow in most arithmetic expressions has undefined behavior; an implementation might throw `underflow_error`, but there is no guarantee that it will throw this or any other exception.

## See Also

[runtime\\_error class](#)

## **<streambuf>**

The `<streambuf>` header declares the `basic_streambuf` class template and its two specializations: `streambuf` and `wstreambuf`. A stream buffer object manages low-level access to a sequence of characters. The characters might be stored in an external file or reside entirely in memory. See the `<fstream>`, `<sstream>`, and `<strstream>` headers for examples of different kinds of stream buffers that are derived from `basic_streambuf`.

Most programs do not use stream buffers directly, but use stream classes instead because they provide a higher-level interface. Each stream object has an associated stream buffer object.

See [Chapter 9](#) for general information on input and output using the stream classes.

# Name

basic\_streambuf class template — Stream buffer

## Synopsis

```
template <class charT, class traits = char_traits<charT> >
class basic_streambuf{
public:
    typedef charT char_type;
    typedef typename traits::int_type int_type;
    typedef typename traits::pos_type pos_type;
    typedef typename traits::off_type off_type;
    typedef traits traits_type;
    virtual ~basic_streambuf( );

    locale pubimbue(const locale &loc);
    locale getloc( ) const;
    basic_streambuf<char_type,traits>* pubsetbuf(char_type* s,
streamsize n);
    pos_type pubseekoff(off_type off, ios_base::seekdir way,
ios_base::openmode which = ios_base::in |
ios_base::out);
    pos_type pubseekoff(off_type off, ios_base::seekdir way,
ios_base::open_mode which=ios_base::in |
ios_base::out);
    pos_type pubseekpos(pos_type sp, ios_base::openmode which =
ios_base::in |
ios_base::out);
    pos_type pubseekpos(pos_type sp,ios_base::open_mode which);
    int pubsync( );
    // Input
    streamsize in_avail( );
    int_type snextc( );
    int_type sbumpc( );
    int_type stoss( );
    int_type sgetc( );
    streamsize sgetn(char_type* s, streamsize n);
    // Putback
    int_type sputbackc(char_type c);
    int_type sungetc( );
    // Output
    int_type sputc(char_type c);
    streamsize sputn(const char_type* s, streamsize n);
protected:
```

```

basic_streambuf( );
// Input
char_type* eback( ) const;
char_type* gptr( ) const;
char_type* egptr( ) const;
void gbump(int n);
void setg(char_type* gbeg, char_type* gnext, char_type* gend);
// Output
char_type* pbase( ) const;
char_type* pptr( ) const;
char_type* eptr( ) const;
void pbump(int n);
void setp(char_type* pbeg, char_type* pend);
// Locales
virtual void imbue(const locale &loc);
// Buffer management and positioning
virtual basic_streambuf<char_type,traits>*
    setbuf(char_type* s, streamsize n);
virtual pos_type seekoff(off_type off, ios_base::seekdir way,
                        ios_base::openmode which =
ios_base::in |
                        ios_base::out);
    virtual pos_type seekpos(pos_type sp, ios_base::openmode which
= ios_base::in |
                        ios_base::out);
virtual streamsize showmanyc( );
virtual int sync( );
virtual int_type underflow( );
virtual int_type uflow( );
virtual streamsize xsggetn(char_type* s, streamsize n);
// Putback
virtual int_type pbackfail(int_type c = traits::eof( ));
// Output
virtual int_type overflow(int_type c = traits::eof( ));
virtual streamsize xspn(const char_type* s, streamsize n);
};

```

The `basic_streambuf` class template manages an input buffer and an output buffer, in which each buffer is an array of characters. The base character type is a template parameter, `charT`. A buffer has the following three pointers to the array (the names below are not part of the standard but are used for informational purposes only):

`begin`

Points to the beginning of the array

next

Points to the next element in the array, that is, the next character to read or the position at which to write the next character

end

Points to one past the end of the array

The pointers in this list can all be null pointers, which makes the stream "buffer" unbuffered. If the pointers are not null, the following rules apply:

- If  $next < end$  for an output array, the stream buffer is said to have a *write position*. The next output character is written to  $* next$ .
- If  $begin < next$  for an input array, the stream buffer is said to have a *push back position*. When a character is pushed back, it is stored in  $next [- 1]$ .
- If  $next < end$  for an input array, the stream buffer is said to have a *read position*. The next character to read from the buffer is  $* next$ .

[Figure 13-25](#) depicts an input buffer, in which the characters "Hello, world." have been fetched from the input source. So far, the first six characters have been read from the input buffer, and the next character to read is the space between "Hello," and "world."



*Figure 13-25. Input stream buffer*

[Figure 13-26](#) depicts an output buffer, which has room for 12 characters. So far, the buffer contains the five

characters "Hello".

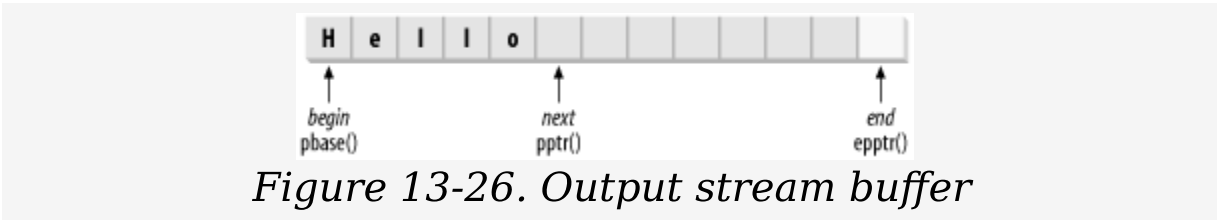


Figure 13-26. Output stream buffer

Several functions come in pairs: a public function that is the public interface and a protected, virtual function to do the actual work. For example, `pubsync` is a public function that simply calls `sync`, which is protected and virtual. The name of the public function is `pub` followed by the name of the protected function. The two exceptions are `sgetn` and `sputn`, which are public, with `xsgen` and `xspu` as their protected counterparts.

The `basic_streambuf` class template implements minimal functionality for managing the buffer pointers. Derived-class templates must override the virtual functions to provide concrete behavior. See `basic_filebuf` in `<fstream>`, `basic_stringbuf` in `<sstream>`, and `strstreambuf` in `<strstream>` for examples of derived classes and templates.

The following descriptions start with the expected public behavior. If you are writing a stream class that uses a stream buffer, you can rely on that behavior from the stream buffer object. If you are writing your own derived-class template, that is the behavior you must implement. Each description ends with the actual implementation in `basic_streambuf`.

When comparing characters, converting characters to integers or integers to characters, and obtaining the end-of-file marker, `basic_streambuf` uses the character traits specified by the `traits` template parameter. See `<string>` for information about the default character traits, `char_traits`.

The following are the public member functions of `basic_streambuf`:

locale **getloc** ( ) const

Returns the current locale, which can be set by calling `pubimbue`.

streamsize **in\_avail** ( )

Returns the number of characters available for input. The return value is `egptr( ) - gptr( )` if the stream buffer has a read position, or the result of calling `showmanyc( )` if the stream buffer does not have a read position.

locale **pubimbue** (const locale &loc)

Saves a copy of the locale, `loc`. Subsequent calls to `getloc` return a copy of the imbued locale. The `basic_streambuf` class does not use the locale, but a derived class can use the locale to interpret multibyte characters and for other purposes. When a stream buffer is created, it is initially imbued with the global locale.

pos\_type **pubseekoff** (off\_type off,  
ios\_base::seekdir way, ios\_base::openmode  
which=ios\_base::in|ios\_base::out), pos\_type  
**pubseekoff** (off\_type off, ios\_base::seekdir way,  
ios\_base::open\_mode  
which=ios\_base::in|ios\_base::out)

Returns `seekoff(off, way, which)`, which changes the stream position. The second form is deprecated. See `ios_base::openmode` in `<ios>` for details.

pos\_type **pubseekpos** (pos\_type sp,  
ios\_base::openmode  
which=ios\_base::in|ios\_base::out), pos\_type

**pubseekpos** (pos\_type sp, ios\_base::open\_mode which)

Returns seekpos(off, way, which), which changes the stream position. The second form is deprecated. See ios\_base::openmode in <ios> for details.

basic\_streambuf<char\_type,traits>\* **pubsetbuf**  
(char\_type\* s, streamsize n)

Returns setbuf(s, n), which typically sets the begin, next, and end array pointers.

int **pubsync** ( )

Returns sync( ), which synchronizes the internal arrays with the external I/O source, if any. The pubsync function returns -1 for a failure and any other value for success.

int\_type **snextc** ( )

Returns the next input character and advances the input pointers. Specifically, snextc calls sbumpc( ), and if sbumpc returns traits::eof( ), snextc returns traits::eof( ); otherwise, it returns sgetc( ).

int\_type **sbumpc** ( ), int\_type **stoss** ( )

Returns the next input character (if one is available) and increments the input array's next pointer. If there is no read position, sbumpc returns uflow( ); otherwise, it returns \*gptr( ).

The stoss function is deprecated. If it is implemented, it calls sbumpc( ).

int\_type **sgetc** ( )

Returns the next available input character without advancing the input pointer. If no read position is available, sgetc returns underflow( ); otherwise, it returns \*gptr( ).



streamsize **sgetn** (char\_type\* s, streamsize n)

Returns `xsggetn(s, n)`.

virtual streamsize **showmanyc** ( )

Returns an estimate of the number of characters immediately available for input. If `showmanyc` returns a positive value, at least that many characters can be read before underflow returns `traits::eof( )`. If `showmanyc` returns `-1`, calls to underflow or `uflow` will fail. If the return value is `0`, input might be available, but there is no guarantee. Read the function name as `s-how-many-c`.

int\_type **sputbackc** (char\_type c)

Tries to push back the character `c` so the next read will return `c`. If a push-back position is available (`gptr( ) > ebase( )`), and the most recently read character (`gptr( )[-1]`) is `c`, the input next pointer is decremented, and `*gptr( )` is returned. Otherwise, `pbackfail(c)` is returned.

int\_type **sputc** (char\_type c)

Writes the character `c` to the output array and returns `c` if an output position is available. Otherwise, `sputc` returns `overflow(c)`.

streamsize **sputn** (const char\_type\* s, streamsize n);

Calls `xsputn(s, n)` to write a string `s` of length `n`.

int\_type **sungetc** ( )

Pushes back the character that was read most recently from the input stream. If a push-back position is available, the function decrements the input next pointer and returns `*gptr( )`; otherwise, it returns `pbackfail( )`.

The following are the protected member functions of `basic_streambuf`. Those marked as `virtual` should probably be overridden in a derived class template.

**basic\_streambuf** ( )

Initializes all pointers to null pointers. The locale (as returned from `getloc( )`) is initialized to the current global locale.

`char_type*` **eback** ( ) const

Returns the begin pointer for the input array.

`char_type*` **egptr** ( ) const

Returns the end pointer for the input array.

`char_type*` **epptr** ( ) const

Returns the end pointer for the output array.

void **gbump** (int n)

Advances the next pointer for the input array by n characters.

`char_type*` **gptr** ( ) const

Returns the next pointer for the input array.

virtual void **imbue** (const locale &loc)

A derived class overrides the `imbue` function to take whatever action is needed when the stream buffer's locale is changed. For example, the stream buffer might cache facets or other information from the new locale.

The `pubimbue` function calls `imbue` *before* storing a new locale; thus, within the call to `imbue`, the `getloc` function returns the old locale, and the `loc` parameter is the new locale.

The default behavior for `imbue` is to do nothing.

```
virtual int_type overflow (int_type c =  
traits::eof( ))
```

The stream buffer calls `overflow` to flush the output buffer when it is full. The `overflow` function handles the sequence of characters stored in the buffer array, or an empty string if the `begin` and other pointers are null. If `c` is not `traits::eof( )`, `c` is appended to the output sequence. For example, a file stream buffer might write the array to the file. If an output error occurs, an exception is thrown or `traits::eof( )` is returned. For success, any value other than `traits::eof( )` is returned.

The default behavior is to do nothing and return `traits::eof( )`.

```
virtual int_type pbackfail (int_type c =  
traits::eof( ))
```

Handles a push-back failure. When the `sungetc` or `sputbackc` functions try to push back a character, they first attempt to adjust the current input next pointer. This can fail for one of three reasons:

- Input is unbuffered (`gptr( ) == 0`).
- The input buffer is empty (`gptr( ) == eback( )`).
- The push-back character, `c`, is not the same character that was most recently read (that is, `c != gptr( )[-1]`).

If the push-back attempt fails, `pbackfail` mediates with the external input source, if any, to try to complete the push-back operation. For example, a file stream buffer might adjust the external file position so the next character read from the input array will retrieve the pushed-back character.

If `c` is `traits::eof( )`, the input sequence is moved back one position to reread the character at that position. If `c` is not `traits::eof( )`, `pbackfail` tries to ensure that the next character read will be `c`; the derived class is free to modify the input source, change the stream position, or take another action to push back `c`.

The return value is `traits::eof( )` for failure or any other value for success.

The default behavior is to do nothing and return `traits::eof( )`.

`char_type* pbase ( ) const`

Returns the output buffer's begin pointer.

`void pbump (int n)`

Increments the output buffer's next pointer by `n` characters.

`char_type* pptr ( ) const`

Returns the output buffer's next pointer.

Virtual `pos_type seekoff (off_type off, ios_base::seekdir way, ios_base::openmode which=ios_base::in|ios_base::out)`

Changes the stream position. The default behavior is to do nothing except return `pos_type(-1)`.

virtual `pos_type seekpos (pos_type sp, ios_base::openmode which=ios_base::in|ios_base::out)`

Changes the stream position. The default behavior is to do nothing except return `pos_type(-1)`.

virtual `basic_streambuf<char_type,traits>* setbuf (char_type* s, streamsize n)`

Typically sets the internal array pointers. The default behavior is to do nothing and return this.

```
void setg (char_type* gbegin, char_type* gnext,  
char_type* gend)
```

Sets the input array pointers: `begin = gbegin`, `next = gnext`, and `end = gend`.

```
void setp (char_type* pbegin, char_type* pend)
```

Sets the output array pointers: `begin = next = pbegin` and `end = pend`.

```
virtual int sync ( )
```

Synchronizes the stream buffer with its external source, if any. Any characters in the output array are written, and all pointers are adjusted as needed. The return value is `-1` for failure and anything else for success.

The default behavior is to do nothing and return `0`.

```
virtual int_type uflow ( )
```

Fills the input array and fetches the first character from the array. It is called from `sputc` if the input is not buffered (`gptr( ) == 0`) or if the input buffer is empty (`gptr( ) >= egptr( )`). Input is obtained from the external source, and the input pointers are reset.

The key difference between `uflow` and `underflow` is that `uflow` advances the input next pointer after returning the first character in the buffer, but `underflow` does not.

The return value from `uflow` is `*gptr( )` if there is a read position or `traits::eof( )` if there is no more input or some other failure occurs.

The default behavior is to do nothing and return `traits::eof( )`.

```
virtual int_type underflow ( )
```

Fills the input array and returns the first character in the array. The `sgetc( )` function calls `underflow` when the input is not buffered (`gptr( ) == 0`) or when the input buffer is empty (`gptr( ) >= egptr( )`). Input is obtained from the external source, and the input pointers are reset.

The key difference between `uflow` and `underflow` is that `uflow` advances the input next pointer after returning the first character in the buffer, but `underflow` does not.

The return value is `*gptr( )` if there is a read position or `traits::eof( )` if there is no more input or some other failure occurs.

The default behavior is to do nothing and return `traits::eof( )`.

```
virtual streamsize xsggetn (char_type* s,  
streamsize n)
```

Reads up to `n` characters from the input stream into the character array that `s` points to. No null character is appended to `s`. It is equivalent to calling `sbumpc` up to `n` times or until it returns `traits::eof( )`. The number of characters actually read and stored in `s` is returned.

Derived classes can override this function to provide a more efficient implementation.

```
virtual streamsize xsputn (const char_type* s,  
streamsize n)
```

Writes up to `n` characters from `s` to the output stream. It is equivalent to calling `sputc` up to `n` times or until it returns `traits::eof( )`. The number of characters actually written is returned. Derived classes can override this function to provide a more efficient implementation.

## See Also

basic\_filebuf in [<fstream>](#), basic\_stringbuf in [<sstream>](#)

## Name

streambuf class — Stream buffer specialization

## Synopsis

```
typedef basic_streambuf<char> streambuf;
```

The streambuf class specializes basic\_streambuf for the char type.

## See Also

char\_traits<char> in [<string>](#)



## Name

wstreambuf class — Stream buffer, wide character specialization

## Synopsis

```
typedef basic_streambuf<wchar_t> wstreambuf;
```

The `wstreambuf` class specializes `basic_streambuf` for the `wchar_t` type.

## See Also

`char_traits<wchar_t>` in [<string>](#)

## <string>

The <string> header declares the class templates and functions that support the string and wstring types, which are specializations of the basic\_string class template. The string types are easier to use and safer than C-style character arrays. Another important class template is char\_traits, which describes a character type and is used throughout the standard library.

The complete declarations of the overloaded operators can be daunting to read. To help you, each function template declaration is followed by a comment that shows the equivalent declaration that uses the common typedefs for narrow characters (e.g., string instead of basic\_string<charT, traits, Allocator>).

[Example 13-37](#) shows a function that classifies a string as an identifier, integer, floating point, or other. The example demonstrates the use of the string class and several of its member functions.

### *Example 13-37. Classifying a string*

```
#include <iostream>
#include <string>

enum kind { empty, ident, integer, floatingpt, error };

kind classify(const std::string& s)
{
    using std::string;
    const string lower("abcdefghijklmnopqrstuvwxyz");
    const string upper("ABCDEFGHIJKLMNOPQRSTUVWXYZ");
    const string letters = lower + upper + '_';
    const string digits("0123456789");
    const string identchars = letters + digits;

    if (s.empty( ))
        return empty;
```

```

else if (letters.find_first_of(s[0]) != string::npos) {
    // Check for valid identifier.
    if (s.find_first_not_of(identchars, 1) == string::npos)
        return ident;
    else
        return error;
}

// Skip a leading sign, if present.
string::size_type pos;
if (s[0] == '+' or s[0] == '-')
    pos = 1;
else
    pos = 0;

// The number must start with a digit.
if (pos == s.length( ))
    return error;
if (not digits.find_first_of(s[pos]))
    return error;
// Find where the digit string ends.
pos = s.find_first_not_of(digits, pos);
if (pos == string::npos)
    // Only digits => integer
    return integer;

else if (s[pos] == '.') {
    // There is a decimal point.
    pos = s.find_first_not_of(digits, pos+1);
    if (pos == string::npos)
        // Integer part "." fractional part
        return floatingpt;
}

// Look for optional exponent.
if (s[pos] == 'e' or s[pos] == 'E') {
    if (pos == s.length( ) - 1)
        return error; // 'e' or 'E' is last char
    else if (s[pos+1] == '+' or s[pos+1] == '-')
        ++pos; // skip over sign;
    if (pos == s.length( ) - 1)
        return error; // Sign is last char.
    pos = s.find_first_not_of(digits, pos+1);
    if (pos == string::npos)
        return floatingpt;
}

```

```
    return error;  
}
```

# Name

basic\_string class template — Base class for string types

## Synopsis

```
template<class charT, class traits = char_traits<charT>,
         class Alloc = allocator<charT> >
class basic_string {
public:
    typedef traits traits_type;
    typedef typename traits::char_type value_type;
    typedef Alloc allocator_type;
    typedef typename Alloc::size_type size_type;
    typedef typename Alloc::difference_type difference_type;
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;
    static const size_type npos = -1;

    explicit basic_string(const Alloc& a = Alloc( ));
    basic_string(const basic_string& str);
    basic_string(const basic_string& str, size_type pos, size_type
n = npos,
                 const Alloc& a = Alloc( ));
    basic_string(const charT* s, size_type n, const Alloc& a =
Alloc( ));
    basic_string(const charT* s, const Alloc& a = Alloc( ));
    basic_string(size_type n, charT c, const Alloc& a=Alloc( ));
    template<class InputIterator>
basic_string(InputIterator begin, InputIterator end,
              const Alloc& a = Alloc( ));
    ~basic_string( );
    basic_string& operator=(const basic_string& str);
    basic_string& operator=(const charT* s);
    basic_string& operator=(charT c);

    iterator begin( );
    const_iterator begin( ) const;
```

```

iterator end( );
const_iterator end( ) const;
reverse_iterator rbegin( );
const_reverse_iterator rbegin( ) const;
reverse_iterator rend( );
const_reverse_iterator rend( ) const;
// Size and capacity
size_type size( ) const;
size_type length( ) const;
size_type max_size( ) const;
void resize(size_type n, charT c);
void resize(size_type n);
size_type capacity( ) const;
void reserve(size_type res_arg = 0);
void clear( );
bool empty( ) const;
// Element access
const_reference operator[](size_type pos) const;
reference operator[](size_type pos);
const_reference at(size_type n) const;
reference at(size_type n);
basic_string substr(size_type pos = 0, size_type n = npos)
const;
// Modifiers
basic_string& operator+=(const basic_string& str);
basic_string& operator+=(const charT* s);
basic_string& operator+=(charT c);
basic_string& append(const basic_string& str);
basic_string& append(const basic_string& str, size_type pos,
size_type n);
basic_string& append(const charT* s, size_type n);
basic_string& append(const charT* s);
basic_string& append(size_type n, charT c);
template<class InputIter>
basic_string& append(InputIter first, InputIter last);
void push_back(charT c);
basic_string& assign(const basic_string& str);
basic_string& assign(const basic_string& str, size_type pos,
size_type n);
basic_string& assign(const charT* s, size_type n);
basic_string& assign(const charT* s);
basic_string& assign(size_type n, charT c);
template<class InputIter>
basic_string& assign(InputIter first, InputIter last);
basic_string& insert(size_type pos1, const basic_string& str);
basic_string& insert(size_type pos1, const basic_string& str,
size_type pos2,
size_type n);

```

```

basic_string& insert(size_type pos, const charT* s, size_type
n);
basic_string& insert(size_type pos, const charT* s);
basic_string& insert(size_type pos, size_type n, charT c);
iterator insert(iterator p, charT c);
void insert(iterator p, size_type n, charT c);
template<class InputIter>
void insert(iterator p, InputIter first, InputIter last);
basic_string& erase(size_type pos = 0, size_type n = npos);
iterator erase(iterator position);
iterator erase(iterator first, iterator last);
basic_string& replace(size_type pos1, size_type n1, const
basic_string& str);
basic_string& replace(size_type pos1, size_type n1, const
basic_string& str,
size_type pos2, size_type n2);
basic_string& replace(size_type pos, size_type n1, const charT*
s,
size_type n2);
basic_string& replace(size_type pos, size_type n1, const charT*
s);
basic_string& replace(size_type pos, size_type n1, size_type
n2, charT c);
basic_string& replace(iterator i1, iterator i2, const
basic_string& str);
basic_string& replace(iterator i1, iterator i2, const charT* s,
size_type n);
basic_string& replace(iterator i1, iterator i2, const charT*
s);
basic_string& replace(iterator i1, iterator i2, size_type n,
charT c);
template<class InputIterator>
basic_string& replace(iterator i1, iterator i2, InputIterator
j1,
InputIterator j2);
size_type copy(charT* s, size_type n, size_type pos = 0) const;
void swap(basic_string& str);
// String operations
const charT* c_str( ) const;
const charT* data( ) const;
allocator_type get_allocator( ) const;
// Searching
size_type find(const basic_string& str, size_type pos = 0)
const;
size_type find(const charT* s, size_type pos, size_type n)
const;
size_type find(const charT* s, size_type pos = 0) const;
size_type find(charT c, size_type pos = 0) const;

```

```

    size_type rfind(const basic_string& str, size_type pos = npos)
const;
    size_type rfind(const charT* s, size_type pos, size_type n)
const;
    size_type rfind(const charT* s, size_type pos=npos) const;
    size_type rfind(charT c, size_type pos = npos) const;
    size_type find_first_of(const basic_string& str, size_type pos
= 0) const;
    size_type find_first_of(const charT* s, size_type pos,
size_type n) const;
    size_type find_first_of(const charT* s, size_type pos = 0)
const;
    size_type find_first_of(charT c, size_type pos = 0) const;
    size_type find_last_of(const basic_string& str, size_type pos =
npos) const;
    size_type find_last_of(const charT* s, size_type pos, size_type
n) const;
    size_type find_last_of(const charT* s, size_type pos = npos)
const;
    size_type find_last_of(charT c, size_type pos=npos) const;
    size_type find_first_not_of(const basic_string& str, size_type
pos = 0) const;
    size_type find_first_not_of(const charT* s, size_type pos,
size_type n) const;
    size_type find_first_not_of(const charT* s, size_type pos = 0)
const;
    size_type find_first_not_of(charT c, size_type pos = 0) const;
    size_type find_last_not_of(const basic_string& str, size_type
pos = npos)
const;
    size_type find_last_not_of(const charT* s, size_type pos,
size_type n) const;
    size_type find_last_not_of(const charT* s, size_type pos =
npos) const;
    size_type find_last_not_of(charT c, size_type pos = npos)
const;
    // Comparisons
    int compare(const basic_string& str) const;
    int compare(size_type pos1, size_type n1, const basic_string&
str) const;
    int compare(size_type pos1, size_type n1, const basic_string&
str,
                size_type pos2, size_type n2) const;
    int compare(const charT* s) const;
    int compare(size_type pos1, size_type n1, const charT* s)
const;
    int compare(size_type pos1, size_type n1, const charT* s,
size_type n2)

```



```
    const;  
};
```

The `basic_string` class template is the base for the `string` and `wstring` types. A `string` object holds a sequence, or string, of characters and provides a number of useful member functions for searching and modifying the string. You can also work with C-style, null-terminated character strings as arguments to `basic_string` members, including constructors. A `basic_string` object keeps track of an explicit length instead of using the C convention of null-terminated character arrays. The `string` and `wstring` types are therefore much easier to use and offer greater safety (see the `at` member function), while still offering ease-of-use with many functions that take C-style strings as arguments.

If you need a sequence of characters that you don't need to treat as a character string, you can use `vector<char>` or `vector<wchar_t>`, but in most cases you will probably find `string` or `wstring` to be more convenient. You can usually use a `string` or `wstring` as a container that supports random access iterators, so you can use strings with the standard algorithms.

Many of the member functions can throw exceptions. Specifying an index out of range often throws `out_of_range`. An attempt to construct a string or modify a string so its length exceeds `max_string( )` throws `length_error`. The `basic_string` class uses an allocator object for memory allocation, which can throw an exception (such as `bad_alloc`) almost any time the string is modified.

Iterators, pointers, and references to elements of a string become invalid in the following situations:

- The string is the target of the `swap` member function or an argument to the `swap` function template.
- The string is an argument to `operator>>` or `getline`.

- You call the `data` or `c_str` member function.
- You call any non-const member function except `operator[]`, `at`, `begin`, `end`, `rbegin`, or `rend`.
- You call the non-const version of `operator[]`, `at`, `begin`, `end`, `rbegin`, or `rend` after any of the above situations, except after calling a form of `insert` or `erase` that returns an iterator (so the returned iterator remains valid).

The following are the members of `basic_string`. Several small examples appear throughout this section, illustrating the use of some of the more complex member functions. Some of the functions are described in terms of temporary string objects or calls to other member functions. The actual implementation might be different, provided the behavior is the same.

`explicit basic_string (const Alloc& a = Alloc( ))`

Constructs an empty string.

`basic_string (const basic_string& str)`

Constructs a string that is a copy of `str`, with `Alloc( )` as the allocator.

`basic_string (const basic_string& str, size_type pos, size_type n = npos, const Alloc& a = Alloc( ))`

Copies a substring of `str`, starting at `pos`. If `pos` is out of range (that is, `pos > str.size( )`), `out_of_range` is thrown. The number of characters copied is `n` or the number of characters left in the string (`str.size( ) - pos`), whichever is smaller.

`basic_string (const charT* s, size_type n, const Alloc& a = Alloc( ))`

Copies the first `n` characters from `s`.

**basic\_string** (const charT\* s, const Alloc& a = Alloc( ))

Copies a null-terminated character array, s. More precisely, this constructor copies traits::length(s) characters from s.

**basic\_string** (size\_type n, charT c, const Alloc& a = Alloc( ))

Initializes the string with n copies of the character c.

template<class InputIterator>, **basic\_string**  
(InputIterator begin, InputIterator end, const Alloc& a = Alloc( ))

The constructor depends on the type of InputIterator:

- For any input iterator, the string is initialized by copying the contents of the range [begin, end).
- If InputIterator is an integral type, the string is initialized with static\_cast<size\_type>(begin) copies of the character static\_cast<value\_type>(end).

**basic\_string& append** (const basic\_string& str, size\_type pos, size\_type n)

Appends characters to the end of the string. If pos > str.size( ), out\_of\_range is thrown. Otherwise, up to n characters are copied from str, starting at position pos. The return value is \*this. See also operator+= later in this section.

**basic\_string& append** (const basic\_string& str)

Returns append(str, 0, npos).

**basic\_string& append** (const charT\* s, size\_type n),  
**basic\_string& append** (const charT\* s), **basic\_string& append** (size\_type n, charT c), template<class

InputIter>, basic\_string& **append** (InputIter first, InputIter last)

Constructs a temporary string *str*, passing the arguments to the constructor, and returns `append( str )`.

basic\_string& **assign** (const basic\_string& str, size\_type pos, size\_type n)

Erases the current contents of the string and replaces them with the substring of *str* that starts at `pos` and extends for up to `n` characters. The return value is `*this`. See also `operator=` later in this section.

basic\_string& **assign** (const basic\_string& str)

Returns `assign(str, 0, npos)`.

basic\_string& **assign** (const charT\* s, size\_type n),  
basic\_string& **assign** (const charT\* s),  
basic\_string& **assign** (size\_type n, charT c),  
template<class InputIter>, basic\_string& **assign**  
(InputIter first, InputIter last)

Constructs a temporary string *str*, passing the arguments to the constructor, and returns `assign( str )`.

const\_reference **at** (size\_type n) const, reference **at** (size\_type n)

Returns the character at position `n`. If `n >= size( )`, `out_of_range` is thrown. See also `operator[]` later in this section.

iterator **begin** ( ), const\_iterator **begin** ( ) const

Returns an iterator that points to the first character of the string.

const charT\* **c\_str** ( ) const

Returns a pointer to a null-terminated (C-style) character array that contains the same characters as the string followed by a terminating null character. The pointer becomes invalid after calling any non-const member function of the string. The typical use of `c_str` is to interface with C functions that require a null-terminated character string:

```
std::printf(fmtstr.c_str( ), value);
```

See also the data member function.

`size_type capacity ( ) const`

Returns the number of characters allocated for use by the string. The string grows as needed; `capacity` tells you how much you can put in the string before it must grow again.

`void clear ( )`

Erases all the characters in the string.

`int compare (const basic_string& str) const`

Returns `traits::compare(data( ), str.data( ), len )`, in which `len` is the smaller of `size( )` and `str.size( )`.

`int compare (size_type pos1, size_type n1, const basic_string& str) const`

Constructs a temporary string `tmp (*this, pos1, n1)`, and returns `tmp .compare(str)`.

`int compare (const charT* s) const`

Constructs a temporary string `tmp (s)`, and returns `this->compare( tmp )`.

`int compare (size_type pos1, size_type n1, const basic_string& str, size_type pos2, size_type n2) const, int compare (size_type pos1, size_type n1,`

`const charT* s) const, int compare (size_type pos1, size_type n1, const charT* s, size_type n2) const`

Constructs two temporary strings *tmp1* (*\*this*, pos1, n1) and *tmp2*: *tmp2* (str, pos2, n2), *tmp2* (s), or *tmp2* (s, n2). The function returns *tmp1* .compare( tmp2 ).

`size_type copy (charT* dst, size_type n, size_type pos = 0) const`

Copies up to n characters from the string, starting at position pos, to the character array dst. If pos > size( ), out\_of\_range is thrown. The number of characters copied, len, is the smaller of n and size( ) - pos. The return value is len.

`const charT* data ( ) const`

Returns a pointer to a character array that has the same character contents as the string. Note that the character array is not null-terminated. If size( ) == 0, data returns a valid, non-null pointer. Do not modify the contents of the data string. The pointer becomes invalid after calling any non-const member function of the string. See also the c\_str member function.

`bool empty ( ) const`

Returns true if the string is empty (size( ) == 0).

`iterator end ( ), const_iterator end ( ) const`

Returns an iterator that points to one position past the end of the string.

`basic_string& erase (size_type pos = 0, size_type n = npos)`

Erases characters from the string, starting at position pos and erasing n or size( ) - pos characters, whichever is smaller. If pos > size( ), out\_of\_range is thrown. The return value is \*this. For example:

```
std::string s("hello, world");
s.erase(9, 1) == "hello, wold"
s.erase(5)    == "hello"
```

iterator **erase** (iterator position)

Erases the character at position and returns an iterator that points to the next character (if there is one) or end( ).

iterator **erase** (iterator first, iterator last)

Erases characters in the range [first, last) and returns an iterator that points to the character that last pointed to (prior to the erasure) or end( ).

size\_type **find** (const basic\_string& str, size\_type pos = 0) const, size\_type **find** (const charT\* s, size\_type pos, size\_type n) const, size\_type **find** (const charT\* s, size\_type pos = 0) const, size\_type **find** (charT c, size\_type pos = 0) const

Returns the smallest index of a string or character, or npos if the search fails. The search starts at position pos. The string to search for is str or a temporary string *tmp* constructed as *tmp* (s, n), *tmp* (s), or *tmp* (1, c). In other words, find returns the smallest *i* such that  $i \geq \text{pos}$ ,  $i + \text{str.size}() \leq \text{size}()$ , and  $\text{at}(i + j) == \text{str.at}(j)$  for all *j* in  $[0, \text{str.size}())$ . For example:

```
string("hello").find('l')    == 2
string("hello").find("lo", 2) == 3
string("hello").find("low")  == string::npos
```

See also rfind later in this section.

size\_type **find\_first\_not\_of** (const basic\_string& str, size\_type pos = 0) const

Finds the first character at or after position pos that does not appear in str, or npos if every character appears in str. For example:

```
string("hello").find_first_not_of("aeiou")    == 0
string("hello").find_first_not_of("aeiou", 1) == 2
string("hello").find_first_not_of("aeiou", 6) == string::npos
```

size\_type **find\_first\_not\_of** (charT c, size\_type pos = 0) const, size\_type **find\_first\_not\_of** (const charT\* s, size\_type pos = 0) const, size\_type **find\_first\_not\_of** (const charT\* s, size\_type pos, size\_type n) const

Constructs a temporary string *tmp* and returns `find_first_not_of( tmp, pos)`, in which *tmp* is constructed as *tmp* (1, c), *tmp* (s), or *tmp* (s, n).

size\_type **find\_first\_of** (const basic\_string& str, size\_type pos = 0) const

Finds the first character at or after position *pos* that appears in *str*, or `npos` if no character appears in *str*. For example:

```
string("hello").find_first_of("aeiou")      = 1
string("hello").find_first_of("aeiou", 2)   = 4
string("hello").find_first_of("aeiou", 6)   = string::npos
```

size\_type **find\_first\_of** (charT c, size\_type pos = 0) const, size\_type **find\_first\_of** (const charT\* s, size\_type pos = 0) const, size\_type **find\_first\_of** (const charT\* s, size\_type pos, size\_type n) const

Constructs a temporary string *tmp* and returns `find_first_of( tmp, pos)`, in which *tmp* is constructed as *tmp* (1, c), *tmp* (s), or *tmp* (s, n).

size\_type **find\_last\_not\_of** (const basic\_string& str, size\_type pos = npos) const

Finds the last character at or before position *pos* that does not appear in *str*, or `npos` if every character appears in *str*. For example:

```
string("hello").find_last_not_of("aeiou")   == 3
string("hello").find_last_not_of("aeiou", 1) == 0
string("hello").find_last_not_of("aeiou", 0) == 0
```



size\_type **find\_last\_not\_of** (charT c, size\_type pos = npos) const, size\_type **find\_last\_not\_of** (const charT\* s, size\_type pos = npos) const, size\_type **find\_last\_not\_of** (const charT\* s, size\_type pos, size\_type n) const

Constructs a temporary string *tmp* and returns `find_last_not_of( tmp, pos)`, in which *tmp* is constructed as *tmp* (1, c), *tmp* (s), or *tmp* (s, n).

size\_type **find\_last\_of** (const basic\_string& str, size\_type pos = npos) const

Finds the last character at or before position *pos* that appears in *str*, or *npos* if no character appears in *str*. For example:

```
string("hello").find_last_of("aeiou")    == 4
string("hello").find_last_of("aeiou", 3) == 1
string("hello").find_last_of("aeiou", 0) == string::npos
```

size\_type **find\_last\_of** (charT c, size\_type pos = npos) const, size\_type **find\_last\_of** (const charT\* s, size\_type pos = npos) const, size\_type **find\_last\_of** (const charT\* s, size\_type pos, size\_type n) const

Constructs a temporary string *tmp* and returns `find_last_of( tmp, pos)`, in which *tmp* is constructed as *tmp* (1, c), *tmp* (s), or *tmp* (s, n).

allocator\_type **get\_allocator** ( ) const

Returns the string's allocator object.

basic\_string& **insert** (size\_type pos1, const basic\_string& str, size\_type pos2, size\_type n)

Inserts a substring of *str* into the string starting at position *pos1*. The substring to insert starts at *pos2* and extends for up to *n* characters. If *pos1* > `size( )` or *pos2* > *str.size( )*, `out_of_range` is thrown. The number of

characters inserted is the smaller of `n` and `str.size( ) - pos2`. The return value is `*this`. For example:

```
string s("hello");  
s.insert(5, " , world")           // s == "hello, world"  
s.insert(5, "out there", 3, 42) // s == "hello there, world"
```

`basic_string& insert (size_type pos, const basic_string& str)`, `basic_string& insert (size_type pos, const charT* s, size_type n)`, `basic_string& insert (size_type pos, const charT* s)`, `basic_string& insert (size_type pos, size_type n, charT c)`

Returns `insert(pos, str, 0, npos)`, in which the last three versions construct a temporary string *tmp* as *tmp* (*s*, *n*), *tmp* (*s*), or *tmp* (*n*, *c*), and then returns `insert(pos, tmp , 0, npos)`.

`iterator insert (iterator p, charT c)`, `void insert (iterator p, size_type n, charT c)`, `template<class InputIter>, void insert (iterator p, InputIter first, InputIter last)`

Inserts text before the character that *p* points to. The first version inserts the character *c* and returns an iterator that points to *c*, the second version inserts *n* copies of the character *c*, and the third version inserts the temporary string constructed from the arguments (*first*, *last*). If *InputIter* is an integral type, the temporary string contains `static_cast<size_type>(first)` copies of the character `static_cast<value_type>(last)`.

`size_type length ( ) const`

Returns `size( )`.

`size_type max_size ( ) const`

Returns the size of the largest possible string.

`const_reference` **operator[]** (size\_type pos) const,  
`reference` **operator[]** (size\_type pos)

Returns the character at position pos. If pos == size( ), the return value is charT( ), that is, a null character. The behavior is undefined if pos > size( ).

`basic_string&` **operator=** (const basic\_string& str)

If \*this and str are the same object, the assignment operator does nothing and returns \*this. If they are different objects, the operator replaces the current string contents with the contents of str and returns \*this.

`basic_string&` **operator=** (const charT\* s),  
`basic_string&` **operator=** (charT c)

Constructs a temporary string, *tmp*(s) or *tmp*(1, c), and assigns \*this = tmp. The return value is \*this.

`basic_string&` **operator+=** (const basic\_string& str),  
`basic_string&` **operator+=** (const charT\* s),  
`basic_string&` **operator+=** (charT c)

Calls append with the same arguments and returns \*this.

`void` **push\_back** (charT c)

Appends c to the end of the string. Its existence lets you use basic\_string with a back\_insert\_iterator.

`reverse_iterator` **rbegin** ( ), `const_reverse_iterator`  
**rbegin** ( ) const

Returns a reverse iterator that points to the last character of the string.

`reverse_iterator` **rend** ( ), `const_reverse_iterator`  
**rend** ( ) const

Returns a reverse iterator that points to one position before the first character of the string.

`basic_string& replace (size_type pos1, size_type n1, const basic_string& str, size_type pos2, size_type n2)`

Erases a substring and inserts another string in its place. The string to erase starts at `pos1` and extends for up to `n1` characters (the smaller of `n1` and `size( ) - pos1`). The string to insert is a substring of `str`, starting at `pos2` and extending for up to `n2` characters (the smaller of `n2` and `str.size( ) - pos2`). The replacement string is inserted at `pos1`. If `pos1 > size( )` or `pos2 > str.size( )`, `out_of_range` is thrown. The return value is `*this`.

`basic_string& replace (size_type pos, size_type n1, const basic_string& str), replace (size_type pos, size_type n1, const charT* str), replace (size_type pos, size_type n1, const charT* s, size_type n2), replace (size_type pos, size_type n1, size_type n2, charT c)`

Returns `replace(pos, n1, tmp, 0, npos)`, in which `tmp` is a temporary string constructed as `tmp (str)`, `tmp (s, n2)`, or `tmp (n2, c)`. For example:

```
std::string s("hello");
s.replace(1, 4, "appy")           s=="happy"
s.replace(5, 0, "your birthday !", 4, 10) s=="happy birthday"
s.replace(1, 1, 1, 'i')          s=="hippy birthday"
```

`basic_string& replace (iterator first, iterator last, const basic_string& str)`

Erases the text in the range `[first, last)` and inserts `str` at the position `first` pointed to. The return value is `*this`.

`basic_string& replace (iterator first, iterator last, const charT* s, size_type n), replace`

(iterator first, iterator last, const charT\* s),  
basic\_string& **replace** (iterator first, iterator last,  
size\_type n, charT c), template<class InputIterator>,  
basic\_string& **replace** (iterator first, iterator  
last, InputIterator i1, InputIterator i2)

Returns `replace(first, last, tmp)`, in which *tmp* is a temporary string constructed as *tmp*(s, n), *tmp*(s), *tmp*(n, c), or *tmp*(i1, i2).

void **reserve** (size\_type res\_arg = 0)

Ensures that the `capacity()` is at least as large as `res_arg`. Call `reserve` to avoid the need to reallocate the string data repeatedly when you know the string will grow by small increments to a large size. Note that `size()` does not change.

void **resize** (size\_type n, charT c), void **resize**  
(size\_type n)

Changes the size of the string to `n` characters. If `n <= size()`, the new string has the first `n` characters of the original string. If `n > size()`, the new string has `n - size()` copies of `c` appended to the end. The second version returns `resize(n, charT(c))`.

size\_type **rfind** (const basic\_string& str, size\_type  
pos = npos) const, size\_type **rfind** (const charT\* s,  
size\_type pos, size\_type n) const, size\_type **rfind**  
(const charT\* s, size\_type pos = npos) const,  
size\_type **rfind** (charT c, size\_type pos = npos) const

Returns the largest index at or before `pos` of a string or character or `npos` if the search fails. The string to search for is `str` or a temporary string *tmp* constructed as *tmp*(s, n), *tmp*(s), or *tmp*(1, c). In other words, `rfind` returns the largest *i* such that `i <= pos`, `i + str.size() <= size()`, and `at(i + j) == str.at(j)` for all *j* in `[0,`

`str.size( )`). (See also `find`, earlier in this section.) For example:

```
string("hello").rfind('l')    == 3
string("hello").rfind("lo", 2) == string::npos
string("hello").rfind("low")  == string::npos
```

`size_type` **size** ( ) const

Returns the number of characters (not bytes) in the string.

`basic_string` **substr** (size\_type pos = 0, size\_type n = npos) const

Returns a substring that starts at position `pos` and extends for up to `n` characters (the smaller of `n` and `size( ) - pos`). If `pos > size( )`, `out_of_range` is thrown.

void **swap** (basic\_string& str)

Exchanges string contents with `str` in constant time.

## See Also

[char\\_traits class template](#), [<sstream>](#), [<vector>](#)

## Name

`char_traits` class template — Base class for character traits

## Synopsis

```
template<typename charT> struct char_traits;
```

The `char_traits` template describes a character type and provides basic functions for comparing, converting, and copying character values and arrays. (See the `char_traits<char>` and `char_traits<wchar_t>` specializations later in this section for details.) If you create a custom character type, you should specialize `char_traits<>` or define your own traits class, which you can provide to `basic_string` and other templates as the traits template parameter. Your traits class should define the same members that `char_traits<char>` defines. See [Chapter 8](#) for an example.

## See Also

[char\\_traits<char> class](#), [char\\_traits<wchar\\_t> class](#)

# Name

`char_traits<char>` class — Character traits of `char` type

# Synopsis

```
template<> struct char_traits<char> {
    typedef char char_type;
    typedef int int_type;
    typedef streamoff off_type;
    typedef streampos pos_type;
    typedef mbstate_t state_type;

    static void assign(char_type& dst, const char_type& src);
    static char_type* assign(char_type* dst, size_t n, const
char_type& c);
    static bool eq(const char_type& c1, const char_type& c2);
    static bool lt(const char_type& c1, const char_type& c2);
    static size_t length(const char_type* str);
    static int compare(const char_type* s1, const char_type* s2,
size_t n);
    static const char_type* find(const char_type* str, size_t n,
                                const char_type& c);
    static char_type* copy(char_type* dst, char_type* src, size_t
n);
    static char_type* move(char_type* dst, char_type* src, size_t
n);
    static bool eq_int_type(const int_type& i1, const int_type&
i2);
    static int_type eof( );
    static int_type not_eof(const int_type& i);
    static char_type to_char_type(const int_type& i);
    static int_type to_int_type(const char_type& c);
};
```



The `char_traits<char>` class specializes `char_traits` for narrow characters. The `streamoff` type is implementation-defined. The `streampos` type is defined as `fpos<mbstate_t>` in `<iosfwd>`. The character traits are defined for the type `char` and have the same meaning in all



locales. The other types are self-explanatory. The following are the member functions:

```
static void assign (char_type& dst, const char_type& src)
```

Assigns `dst = src`.

```
static char_type* assign (char_type* dst, size_t n, const char_type& c)
```

Fills `dst` with `n` copies of `c`, that is, `dst[0]` through `dst[n-1] = c`.

```
static int compare (const char_type* s1, const char_type* s2, size_t n)
```

Compares the first `n` characters of the arrays `s1` and `s2`, returning an integer result:

- 0 if `eq(s1[ i ], s2[ i ])` is true for all `i` in `[0, n)`
- Negative if `eq(s1[ i ], s2[ i ])` is true for all `i` in `[0, k)`, and `lt(s1[ k ], s2[ k ])` is true for some `k` in `[0, n)`
- Positive otherwise

```
static char_type* copy (char_type* dst, char_type* src, size_t n)
```

Copies `n` characters from `src` to `dst`. The arrays `src` and `dst` must not overlap.

```
static int_type eof ( )
```

Returns the end-of-file marker, `EOF` (in `<cstdio>`). The end-of-file marker is different from all characters, that is, for all character values `c`, `eq_int_type(eof( ), to_int_type( c ))` is false.

```
static bool eq (const char_type& c1, const char_type& c2)
```

Returns `c1 == c2`.

```
static bool eq_int_type (const int_type& i1, const int_type& i2)
```

Returns true if `i1` is the same as `i2`. Specifically, for all character values `c1` and `c2`, `eq( c1, c2 )` has the same value as `eq_int_type(to_int_type( c1 ), to_int_type( c2 ))`. Also, `eof( )` is always equal to `eof( )` and not equal to `to_int_type( c )` for any character `c`. The value is unspecified for any other integer values.

```
static const char_type* find (const char_type* str, size_t n, const char_type& c)
```

Returns a pointer `p` to the first character in `str` such that `eq(* p, c)` is true. It returns a null pointer if there is no such character in the first `n` characters of `str`.

```
static size_t length (const char_type* str)
```

Returns the length of the null-terminated character string `str`, that is, it returns the smallest `i` such that `eq(str[ i ], charT( ))` is true.

```
static bool lt (const char_type& c1, const char_type& c2)
```

Returns `c1 < c2`.

```
static char_type* move (char_type* dst, char_type* src, size_t n)
```

Copies `n` characters from `src` to `dst`. The arrays `src` and `dst` are allowed to overlap.

```
static int_type not_eof (const int_type& i)
```

Returns a value that is guaranteed to be different from `eof( )`. If `i` is not `eof( )`, `i` is returned. Otherwise, some other value is returned.

```
static char_type to_char_type (const int_type& i)
```

Converts `i` to its equivalent character value (for which `eq_int_type(i, to_int_type(to_char_type(i)))` is true). If `i` is not equivalent to any character, the behavior is unspecified.

static int\_type **to\_int\_type** (const char\_type& c)

Converts `c` to its equivalent integer representation.

## See Also

[char\\_traits<wchar\\_t> class](#), `mbstate_t` in [<cwchar>](#), `fpos` in [<ios>](#), [<iosfwd>](#)

## Name

`char_traits<wchar_t>` class — Character traits of `wchar_t` type

## Synopsis

```
template<> struct char_traits<wchar_t> {
    typedef wchar_t char_type;
    typedef wint_t int_type;
    typedef streamoff off_type;
    typedef wstreampos pos_type;
    typedef mbstate_t state_type;

    static void assign(char_type& dst, const char_type& src);
    static char_type* assign(char_type* dst, size_t n, const
char_type& c);
    static bool eq(const char_type& c1, const char_type& c2);
    static bool lt(const char_type& c1, const char_type& c2);
    static size_t length(const char_type* str);
    static int compare(const char_type* s1, const char_type* s2,
size_t n);
    static const char_type* find(const char_type* str, size_t n,
                                const char_type& c);
    static char_type* copy(char_type* dst, char_type* src, size_t
n);
    static char_type* move(char_type* dst, char_type* src, size_t
n);
    static bool eq_int_type(const int_type& i1, const int_type&
i2);
    static int_type eof( );
    static int_type not_eof(const int_type& i);
    static char_type to_char_type(const int_type& i);
    static int_type to_int_type(const char_type& c);
};
```

The `char_traits<wchar_t>` class specializes `char_traits` for wide characters. The `wstreamoff` type is implementation-defined. The `wstreampos` type is defined as `fpos<mbstate_t>` in `<iosfwd>`. The other types are self-explanatory. The character traits are defined for the type `wchar_t` and have the same meaning in all locales.

See `char_traits<char>` earlier in this section for a description of the member functions. The `eof( )` function returns `WEOF` (in `<wchar>`).

## See Also

[char\\_traits<char> class](#), `mbstate_t` in [<wchar>](#), `fpos` in [<ios>](#), [<iosfwd>](#)

## Name

getline function template — Reads a line into a string

## Synopsis

```
template<class charT, class traits, class Allocator>
    basic_istream<charT,traits>&
getline(basic_istream<charT,traits>& in,
        basic_string<charT,traits,Allocator>& str,
            charT delim);
// istream& getline(istream& in, string& str, char delim);
template<class charT, class traits, class Allocator>
    basic_istream<charT,traits>&
getline(basic_istream<charT,traits>& in,
        basic_string<charT,traits,Allocator>&
            str);
// istream& getline(istream& in, string& str);
```

The `getline` function template reads a line of text from an input stream into the string `str`. It starts by creating a `basic_istream::sentry(in, true)` object. If the sentry object evaluates to true, `getline` erases `str` then reads characters from `in` and appends them to `str` until end-of-file is reached or `delim` is read. (The `delim` character is read from the stream but not appended to the string.) Reading also stops if `max_size( )` characters have been stored in the string, in which case `ios_base::failbit` is set. If no characters are read from the stream, `ios_base::failbit` is set. The return value is `in`.

The second form of `getline` uses a newline as the delimiter, that is, it returns `getline(in, str, in.widen('\n'))`.

## See Also

operator>> function template, basic\_istream in  
<istream>, basic\_istream::sentry in <istream>

# Name

operator+ function template — Concatenates two strings

## Synopsis

```
template<class charT, class traits, class Allocator>
    basic_string<charT,traits,Allocator> operator+(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// string& operator+(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    basic_string<charT,traits,Allocator> operator+(const charT* a,
        const basic_string<charT,traits,Allocator>& b);
// string& operator+(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    basic_string<charT,traits,Allocator> operator+(
        const basic_string<charT,traits,Allocator>& a, const charT*
    b);
// string& operator+(const string& a, const char* b);
template<class charT, class traits, class Allocator>
    basic_string<charT,traits,Allocator> operator+(
        const basic_string<charT,traits,Allocator>& a, charT b);
// string& operator+(const string& a, char b);
```

The + operator concatenates two strings and returns the result. It constructs a new string as a copy of a, then calls a.append(b) and returns the copy.

## See Also

[basic\\_string class template](#)



## Name

`operator==` function template — Compares strings for equality

## Synopsis

```
template<class charT, class traits, class Allocator>
    bool operator==(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// bool operator==(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator==(const charT* a, const
        basic_string<charT,traits,Allocator>& b);
// bool operator==(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator==(const basic_string<charT,traits,Allocator>& a,
        const charT* b);
// bool operator==(const string& a, const char* b);
```

The `==` operator compares two strings for equality or compares a string and a null-terminated character array. It returns `a.compare(b) == 0`, converting `a` or `b` from a character array to a string, as needed.

## See Also

[basic\\_string class template](#)

## Name

`operator!=` function template — Compares strings for inequality

## Synopsis

```
template<class charT, class traits, class Allocator>
    bool operator!=(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// bool operator!=(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator!=(const charT* a, const
        basic_string<charT,traits,Allocator>& b);
// bool operator!=(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator!=(const basic_string<charT,traits,Allocator>& a,
        const charT* b);
// bool operator!=(const string& a, const char* b);
```

The `!=` operator compares two strings for inequality or compares a string and a null-terminated character array. It returns `!(a == b)`.

## See Also

[basic\\_string class template](#)

## Name

operator< function template — Compares strings for less-than

## Synopsis

```
template<class charT, class traits, class Allocator>
    bool operator<(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// bool operator<(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator<(const charT* a, const
basic_string<charT,traits,Allocator>& b);
// bool operator<(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator<(const basic_string<charT,traits,Allocator>& a,
const charT* b);
// bool operator<(const string& a, const char* b);
```

The < operator compares two strings or compares a string and a null-terminated character array. It returns `a.compare(b) < 0`, converting a or b from a character array to a string, as needed.

## See Also

[basic\\_string class template](#)

## Name

`operator<=` function template — Compares strings for less-than-or-equal

## Synopsis

```
template<class charT, class traits, class Allocator>
    bool operator<=(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// bool operator<=(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator<=(const charT* a, const
        basic_string<charT,traits,Allocator>& b);
// bool operator<=(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator<=(const basic_string<charT,traits,Allocator>& a,
        const charT* b);
// bool operator<=(const string& a, const char* b);
```

The `<=` operator compares two strings or compares a string and a null-terminated character array. It returns `a.compare(b) <= 0`, converting `a` or `b` from a character array to a string, as needed.

## See Also

[basic\\_string class template](#),

## Name

`operator>` function template — Compares strings for greater-than

## Synopsis

```
template<class charT, class traits, class Allocator>
    bool operator>(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// bool operator>(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator>(const charT* a, const
basic_string<charT,traits,Allocator>& b);
// bool operator>(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator>(const basic_string<charT,traits,Allocator>& a,
const charT* b);
// bool operator>(const string& a, const char* b);
```

The `>` operator compares two strings or compares a string and a null-terminated character array. It returns `a.compare(b) > 0`, converting `a` or `b` from a character array to a string, as needed.

## See Also

[basic\\_string class template](#)

## Name

`operator>=` function template — Compares strings for greater-than-or-equal

## Synopsis

```
template<class charT, class traits, class Allocator>
    bool operator>=(
        const basic_string<charT,traits,Allocator>& a,
        const basic_string<charT,traits,Allocator>& b);
// bool operator>=(const string& a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator>=(const charT* a, const
        basic_string<charT,traits,Allocator>& b);
// bool operator>=(const char* a, const string& b);
template<class charT, class traits, class Allocator>
    bool operator>=(const basic_string<charT,traits,Allocator>& a,
        const charT* b);
// bool operator>=(const string& a, const char* b);
```

The `>=` operator compares two strings or compares a string and a null-terminated character array. It returns `a.compare(b) >= 0`, converting `a` or `b` from a character array to a string, as needed.

## See Also

[basic\\_string class template](#)

## Name

`operator<<` function template — Writes a string to an output stream

## Synopsis

```
template<class charT, class traits, class Allocator>
    basic_ostream<charT, traits>& operator<<(
        basic_ostream<charT, traits>& out,
        const basic_string<charT, traits, Allocator>& str);
// ostream& operator<<(ostream& out, const string& str);
```

The `<<` operator writes the string `str` to `out`. Like any formatted output function, it first creates a sentry object, and if the sentry evaluates to `true`, it writes the string contents by calling `out.rdbuf( )->sputn`. If `str.size( ) < out.width( )`, fill characters are added to achieve the desired width. If `sputn` fails, `ios_base::failbit` is set.

## See Also

`ios_base` in [<ios>](#), `basic_ostream` in [<ostream>](#),  
`basic_ostream::sentry` in [<ostream>](#)

## Name

`operator>>` function template — Reads a string from an input stream

## Synopsis

```
template<class charT, class traits, class Allocator>
    basic_istream<charT,traits>& operator>>(
        basic_istream<charT,traits>& in,
        basic_string<charT,traits,Allocator>& str);
// istream& operator>>(istream& in, string& str);
```

The `>>` operator reads a string from `in` and stores the string in `str`. Like any other formatted input operator, it first creates a sentry object `basic_istream::sentry(in)`, and if the sentry evaluates to `true`, it erases `str` and then reads characters from `in` and appends the characters to `str`. If `in.width( )` is greater than 0, no more than `in.width( )` characters are read from `in`; otherwise, up to `max_size( )` characters are read. Reading also stops at end-of-file or when reading a whitespace character (`isspace` is true for locale `in.getloc( )`). The whitespace character is left in the input stream. The return value is `in`.

## See Also

[getline function template](#), `basic_istream` in [<istream>](#),  
`basic_istream::sentry` in [<istream>](#)



## Name

string class — Narrow character string class

## Synopsis

```
typedef basic_string<char> string;
```

The string class specializes basic\_string for type char.

## See Also

[basic\\_string class template](#), [wstring class](#)

## Name

swap function template — Swaps two strings

## Synopsis

```
template<class charT, class traits, class Allocator>
    void swap(basic_string<charT,traits,Allocator>& a,
              basic_string<charT,traits,Allocator>& b);
// void swap(string& a, string& b);
```

The swap function template specialization is equivalent to calling `a.swap(b)`.

## See Also

swap in [<algorithm>](#)

## Name

wstring class — Wide character string class

## Synopsis

```
typedef basic_string<wchar_t> wstring;
```

The `wstring` class specializes `basic_string` for type `wchar_t`.

## See Also

[basic\\_string class template](#), [string class](#)

## **<strstream>**

The `<strstream>` header declares several classes for reading from character arrays and writing to character arrays in the same manner as reading from and writing to files.

This header and its classes are deprecated in the standard, meaning they might disappear from a future version of the standard. Instead, you are encouraged to use the `<sstream>` header and its class templates. Nonetheless, the `<strstream>` classes have their uses; when you are dealing exclusively with narrow characters, and are using character arrays instead of string objects, these classes sometimes offer better performance than their `<sstream>` counterparts.

See [Chapter 10](#) for a general discussion of I/O, and the `<istream>` and `<ostream>` sections in this chapter for information about the base classes from which the `strstream` classes derive. Refer to the `<streambuf>` section in this chapter for information about the `streambuf` class.

# Name

istream class — Input character array streams

# Synopsis

```
class istream: public istream
{
public:
    explicit istream(const char* str);
    explicit istream(char* str);
    istream(const char* str, streamsize n);
    istream(char* str, streamsize n);

    strstreambuf* rdbuf( ) const;
    char* str( );
};
```

The `istream` class represents an input string stream. To construct an `istream`, pass a character array (with an optional size). You can then read from the string stream just as you would from any other input stream.

The following are the methods of `istream`:

`explicit istream (const char* str), explicit istream (char* str)`

Initializes an input string stream by constructing an internal stream buffer as `strstreambuf(str, 0)` and passing the address of the stream buffer to the base-class constructor for `istream`.

`explicit istream (const char* str, streamsize n), explicit istream (char* str, streamsize n)`

Initializes an input string stream by constructing an internal stream buffer as `strstreambuf(str, n)` and passing the address of the stream buffer to the base-class constructor for `istream`.

`strstreambuf* rdbuf ( ) const`

Returns a pointer to the internal `strstreambuf` object.

`char* str ( )`

Returns the internal string, `rdbuf( )->str( )`.

## See Also

[ostream class](#), [strstream class](#), [strstreambuf class](#),  
istream in [<istream>](#), istringstream in [<sstream>](#)

# Name

ostream class — Output character array streams

## Synopsis

```
class ostream: public ostream
{
public:
    ostream( );
    ostream(char* str, int n, ios_base::openmode mode =
ios_base::out);

    ostreambuf* rdbuf( ) const;
    void freeze(bool flag = true);
    char* str( );
    int pcount( ) const;
};
```

The `ostream` class represents an output string stream. You can provide a character array, and the stream contents are written to that array. Another typical usage is to construct an `ostream` with no argument and let the string stream allocate the string as you write to the stream. Then call `str( )` to obtain the resulting character array. Once you call `str( )`, the stream is *frozen* and cannot be modified. The pointer returned from `str( )` remains valid until the `ostream` object is destroyed or until you *thaw* the stream to allow writing again.

The following are the methods of `ostream`:

### **ostream ( )**

Initializes an empty output string stream by constructing an internal `ostreambuf` object and passing the address of the string buffer to the base-class constructor for `ostream`.

**ostream** (char\* str, int n, ios\_base::openmode mode = ios\_base::out)

Initializes a string stream with `str` as the initial string contents by constructing an internal `strstreambuf` object and passing the address of the buffer to the base-class constructor for `ostream`. If the `ios_base::app` bit is set in mode, the buffer is constructed like this:

```
strstreambuf(str, n, str + std::strlen(str));
```

If the `ios_base::app` bit is clear in mode, the buffer is constructed like this:

```
strstreambuf(str, n, str);
```

```
void freeze (bool flag = true)
```

Freezes or thaws the buffer by calling `rdbuf( )->freeze(flag)`.

```
strstreambuf* rdbuf ( ) const
```

Returns a pointer to the internal `strstreambuf` object.

```
char* str ( )
```

Returns a pointer to the buffer's character array, that is, `rdbuf( )->str( )`.

```
int pcount ( ) const
```

Returns the number of bytes in the output buffer by calling `rdbuf( )->pcount( )`.

## See Also

[istrstream class](#), [strstream class](#), [strstreambuf class](#), `ostream` in [<ostream>](#), `ostream` in [<ostream>](#), `ostream` in [<ostream>](#)



# Name

stringstream class — Input and output character array streams

# Synopsis

```
class stringstream: public iostream
{
public:
    typedef char char_type;
    typedef typename char_traits<char>::int_type int_type;
    typedef typename char_traits<char>::pos_type pos_type;
    typedef typename char_traits<char>::off_type off_type;

    stringstream( );
    stringstream(char* s, int n,
                  ios_base::openmode mode =
ios_base::in|ios_base::out);
    virtual ~stringstream( );

    stringstreambuf* rdbuf( ) const;
    void freeze(bool freeze = true);
    int pcount( ) const;
    char* str( );
};
```

The `stringstream` class is a stream class that performs input and output to a character array. You can start with an empty string and write to the stream, or start with a string and read from the stream. You can switch between reading and writing at any time. If you use the default constructor and write to the stream, the stream buffer grows as needed. Then you can call `str( )` to obtain the resulting character array. Once you call `str( )`, the stream is frozen and cannot be modified. The pointer returned from `str( )` remains valid until the `stringstream` object is destroyed or until you thaw the stream to allow writing again.

The following are the methods of `stringstream`:

## **stringstream ( )**

Initializes an empty string stream by constructing an internal `stringstreambuf` object and passing the address of the string buffer to the base-class constructor for `iostream`.

```
basic_stringstream (char* str, int n,  
ios_base::openmode mode =  
ios_base::in|ios_base::out)
```

Initializes a string stream with `str` as the initial string contents by constructing an internal `stringstreambuf` object and passing the address of the buffer to the base-class constructor for `iostream`. If the `ios_base::app` bit is set in `mode`, the buffer is constructed like this:

```
stringstreambuf(str, n, str + std::strlen(str));
```

If the `ios_base::app` bit is clear in `mode`, the buffer is constructed like this:

```
stringstreambuf(str, n, str);
```

```
void freeze (bool flag = true)
```

Freezes or thaws the buffer by calling `rdbuf( )->freeze(flag)`.

```
stringstream* rdbuf ( ) const
```

Returns a pointer to the internal `stringstreambuf` object.

```
char* str ( )
```

Returns a pointer to the buffer's character array, that is, `rdbuf( )->str( )`.

```
int pcount ( ) const
```

Returns the number of bytes in the output buffer by calling `rdbuf( )->pcount( )`.

## **See Also**

istream class, ostream class, stringstream class,  
basic\_istream in <istream>, stringstream in  
<sstream>

# Name

strstreambuf class — I/O buffer for character array streams

## Synopsis

```
class strstreambuf : public basic_streambuf<char> {
public:
    explicit strstreambuf(streamsize a_size_arg = 0);
    strstreambuf(void* (*palloc_arg)(size_t), void (*pfree_arg)
(void*));
    strstreambuf(char* gnext_arg, streamsize n, char* pbeg_arg =
0);
    strstreambuf(const char* gnext_arg, streamsize n);
    strstreambuf(signed char* gnext_arg, streamsize n,
signed char* pbeg_arg = 0);
    strstreambuf(const signed char* gnext_arg, streamsize n);
    strstreambuf(unsigned char* gnext_arg, streamsize n,
unsigned char* pbeg_arg = 0);
    strstreambuf(const unsigned char* gnext_arg, streamsize n);
    virtual ~strstreambuf( );

    void freeze(bool freezefl = true);
    char* str( );
    int pcount( );

protected:
    virtual int_type overflow (int_type c = EOF);
    virtual int_type pbackfail(int_type c = EOF);
    virtual int_type underflow( );
    virtual pos_type seekoff(off_type off, ios_base::seekdir way,
ios_base::openmode which =
ios_base::in |
ios_base::out);
    virtual pos_type seekpos(pos_type sp,
ios_base::openmode which =
ios_base::in |
ios_base::out);
};
```

The `strstreambuf` class implements a stream buffer for character array streams. An internal buffer maintains a single character array with separate positions for reading and writing. That is, the buffer has `begin`, `next`, and `end`

pointers for reading and separate begin, next, and end pointers for writing. The begin pointer points to the start of a buffer, and the end pointer points to one past the end of the buffer. The next pointer points to the position where the next character is to be read or written. Refer to `basic_streambuf` in `<streambuf>` for details about buffer positions.

A `strstreambuf` object maintains a set of flags, an allocated buffer size, and two function pointers for an allocation and deallocation function. If the allocation function pointer is null, the `new[]` operator is used for allocating the character array; if the deallocation function pointer is null, the `delete[]` operator is used.

The flags are:

`allocated`

Indicates that the character array has been allocated, so the destructor should delete it

`constant`

Indicates that the character array is `const`, so it cannot be used for output

`dynamic`

Indicates that the character array has been dynamically allocated and can grow as needed to accommodate output

`frozen`

Indicates that the character array can no longer be modified, extended, or freed

The following are the public member functions of `strstreambuf`:

explicit **strstreambuf** (streamsize alloc\_size = 0)

Saves alloc\_size as the suggested size of the character array, and sets the dynamic flag. The allocation and deallocation functions are set to null pointers.

**strstreambuf** (void\* (\*palloc)(size\_t), void (\*pfree)(void\*))

Sets the dynamic flag and saves palloc and pfree as the allocation and deallocation functions.

**strstreambuf** (char\* gnext\_arg, streamsize n, char\* pbeg\_arg = 0), **strstreambuf** (signed char\* gnext\_arg, streamsize n, signed char\* pbeg\_arg = 0),  
**strstreambuf** (unsigned char\* gnext\_arg, streamsize n, unsigned char\* pbeg\_arg = 0)

Clears all flags and sets the allocation and deallocation functions to null pointers. If pbeg\_arg is null, the output pointers are null and the input buffer is set to n bytes starting at gnext\_arg by calling setg(gnext\_arg, gnext\_arg, gnext\_arg + N). If pbeg\_arg is not null, the input pointers are set by calling setg(gnext\_arg, gnext\_arg, pbeg\_arg), and the output pointers are set by calling setp(pbeg\_arg, pbeg\_arg + N). N is determined as follows:

n > 0

N is n.

n == 0

N is strlen(gnext\_arg).

n < 0

N is INT\_MAX.

**strstreambuf** (const char\* gnext\_arg, streamsize n),  
**strstreambuf** (const signed char\* gnext\_arg,

streamsize n), **strstreambuf** (const unsigned char\* gnext\_arg, streamsize n)

Initializes the buffer pointers in the same manner as constructing `strstreambuf(const_cast<char*>(gnext_arg), n)`. The only difference is that the constant flag is set.

virtual **~strstreambuf** ( )

The destructor frees the character array if the allocated flag is set and the frozen flag is clear.

void **freeze** (bool freezefl = true)

Freezes or thaws the character buffer. If the dynamic flag is set, the freeze function sets or clears the frozen flag to match the `freezefl` parameter. If the dynamic flag is clear, `freeze( )` does nothing.

char\* **str** ( )

Returns the internal character buffer by calling `freeze( )` and returning `gbase( )`.

int **pcount** ( )

Returns the number of output bytes in the buffer. If `pptr( )` is null, 0 is returned; otherwise, `pptr( ) - pbase( )` is returned.

The overridden virtual functions are:

virtual int\_type **overflow** (int\_type c = EOF)

Attempts to append `c` to the end of the character array as follows:

- If `c == EOF`, nothing happens and a non-end-of-file character is returned to indicate success.
- If `c != EOF`, and a write position is available, `c` is appended to the character array by calling `sputc(c)`.

- If a write position is not available, the dynamic flag is set, and the frozen flag is clear, then the character array is extended and `c` is appended to the array. The array is extended by allocating a new, larger character array; copying the old contents (if any); updating the read and write pointers. If the array is successfully extended, the allocated flag is set.
- Otherwise, the array cannot be extended, so the function fails.

The return value is `c` for success or `EOF` for failure. If `c` is `EOF`, a value other than `EOF` is returned for success.

```
virtual int_type putback (int_type c =
traits::eof( ))
```

Attempts to push back `c` onto the input array for reading as follows:

- If `c == EOF`, and a putback position is available, `gptr( )` is set to `gptr( ) - 1`.
- If `c != EOF`, a putback position is available, and `gptr( )[-1]` is equal to `c`, `gptr( )` is set to `gptr( ) - 1`.
- If `c != EOF`, the constant flag is clear, and a putback position is available, `gptr( )` is set to `gptr( ) - 1`, and `*gptr( )` is assigned `c`.
- Otherwise, the character cannot be put back, so the function fails.

The return value is `c` for success or `EOF` for failure. If `c` is `EOF`, a value other than `EOF` is returned for success.

```
virtual pos_type seekoff (off_type off,
ios_base::seekdir way, ios_base::openmode which =
ios_base::in|ios_base::out)
```

Sets the stream position. The input position, output position, or both can be set, depending on (`which &`



(`ios_base::in | ios_base::out`)). The following are the possible results of this expression:

`ios_base::in`

Sets the input position

`ios_base::out`

Sets the output position

`ios_base::in | ios_base::out`, and `way` is either `ios_base::beg` or `ios_base::end`

Sets input and output positions

Otherwise

The function fails and returns `pos_type(-1)`

The new position is determined by adding the offset `off` to a base position given by `way`, which must be one of the following:

`ios_base::beg`

The base position is the at start of the stream—that is, `off` is an absolute position.

`ios_base::cur`

The base position is the current stream position.


`ios_base::end`

The base position is at the end of the stream.

In all cases, a positive offset is toward the end of the stream, and a negative offset is toward the start of the stream. If the desired position is negative or past the end of the string, the function fails and returns `pos_type(-1)`. If the function succeeds, it returns the new position.

```
virtual pos_type seekpos (pos_type sp,  
ios_base::openmode which =  
ios_base::in|ios_base::out)
```

Sets the stream position to `sp`. The input position is set if `which & ios_base::in` is nonzero. The output position is set if `which & ios_base::out` is nonzero. If `sp` is not a valid position, or if neither the input nor the output position is set, `seekpos` fails and `pos_type(-1)` is returned. The return value is `sp` for success. If `sp` was not returned from a prior call to a positioning function (that is, `seekoff`, `seekpos`, `tellg`, or `tellp`), the results are undefined.

```
 virtual basic_streambuf<charT,traits>* setbu  
f(charT*, streamsize)
```

Calling `setbuf(0, 0)` has no effect. The result of any other call to `setbuf` is implementation-defined.

```
virtual int_type underflow ( )
```

Gets another character from the input range without moving the input pointer. If the stream has a read position, the function returns `*gptr( )`. If there is no read position, but there is a non-null write pointer past the end of the input range—that is, `pptr( ) > gend( )`—then the read end pointer (`gend( )`) is advanced at least one position but still less than or equal to `pptr( )`. The return value is `EOF` for failure or `*gnext( )` for success.

## See Also

`stringbuf` in [<sstream>](#), `basic_streambuf` in [<streambuf>](#)

## **<typeinfo>**

The <typeinfo> header declares the `type_info` class (for the `typeid` operator) and two exception classes related to type information and casting.

## Name

`bad_cast` class — Exception for `dynamic_cast<>`

## Synopsis

```
class bad_cast : public exception {  
public:  
    bad_cast( ) throw( );  
    bad_cast(const bad_cast&) throw( );  
    bad_cast& operator=(const bad_cast&) throw( );  
    virtual ~bad_cast( ) throw( );  
    virtual const char* what( ) const throw( );  
};
```

The `dynamic_cast<>` operator throws `bad_cast` when the cast of a reference fails. See `dynamic_cast` in [Chapter 3](#) for more information.

## See Also

`dynamic_cast` operator

## Name

`bad_typeid` class — Exception for null pointer in typeid expressions

## Synopsis

```
class bad_typeid : public exception {
public:
    bad_typeid( ) throw( );
    bad_typeid(const bad_typeid&) throw( );
    bad_typeid& operator=(const bad_typeid&) throw( );
    virtual ~bad_typeid( ) throw( );
    virtual const char* what( ) const throw( );
};
```

The typeid operator throws `bad_typeid` when it is applied to an expression of the form `*p`, in which `p` is a null pointer. See typeid in [Chapter 3](#) for more information.

## See Also

typeid operator

# Name

type\_info class — Type information

## Synopsis

```
class type_info {
public:
    virtual ~type_info( );
    bool operator==(const type_info& rhs) const;
    bool operator!=(const type_info& rhs) const;
    bool before(const type_info& rhs) const;
    const char* name( ) const;
private:
    type_info(const type_info& rhs);
    type_info& operator=(const type_info& rhs);
};
```



The typeid operator (described in [Chapter 3](#)) returns a static type\_info object. The type information includes the type's name and a collation order, both of which are implementation-defined. An implementation might derive classes from type\_info to provide additional information.

Note that the copy constructor and assignment operators are inaccessible, so you must store pointers if you want to use a standard container. [Example 13-38](#) shows how to store type\_info pointers in a set, where the order is determined by the before member function.

## Example

*Example 13-38. Storing type information*

```
#include <algorithm>
#include <functional>
#include <iostream>
#include <ostream>
#include <set>
#include <typeinfo>
```

```

typedef bool (*type_info_compare) (const std::type_info*, const
std::type_info*);

typedef std::set<const std::type_info*, type_info_compare>
typeset;

// Return true if *a comes before *b (comparison function to
store type_info
// pointers in an associative container).
bool type_info_less(const std::type_info* a, const
std::type_info* b)
{
    return a->before(*b);
}

// Print a type_info name on a line.
void print(const std::type_info* x)
{
    std::cout << x->name( ) << '\n';
}

void demo( )
{
    // Construct and initialize the set.
    typeset types(&type_info_less);

    types.insert(&typeid(int));
    types.insert(&typeid(char));
    types.insert(&typeid(std::type_info));
    types.insert(&typeid(std::bad_alloc));
    types.insert(&typeid(std::exception));
    . . .
    // Print the types in the set.
    std::for_each(types.begin( ), types.end( ), print);
}

```

The members of `type_info` are:

`bool before (const type_info& rhs) const`

Returns true if this `type_info` object comes before `rhs` in the implementation-defined order. The relative order of types can vary between programs, even for the same types.



`const char* name ( ) const`

Returns the type's name as a null-terminated string, which might be a multibyte string. The contents of the name string are implementation-defined.

`bool operator== (const type_info& rhs) const, bool operator!= (const type_info& rhs) const`

Compares `type_info` objects, which are equal when the types they describe are the same.

## See Also

`typeid operator`



## **<utility>**

The <utility> header declares the `pair<>` class template, which has many uses, especially by maps in the <map> header. It also defines the `rel_ops` namespace, which defines relational operators in terms of `==` and `<`.

# Name

make\_pair function template — Constructs a pair object

# Synopsis

```
template <typename T1, typename T2>
pair<T1,T2> make_pair(T1 a, T2 b);
```

Constructs a pair<T1,T2> object and initializes it with the values a and b. The advantage of using make\_pair over a simple pair<> constructor is that the compiler can deduce the types T1 and T2 from the values a and b. [Example 13-39](#) shows a typical use of make\_pair.

# Example

*Example 13-39. Making pairs of objects*

```
std::map<std::string, int> wordcounts;
wordcounts.insert(std::make_pair("hello", 1));

// Functor, suitable for passing to for_each to find minimum and
// maximum values
// in a range
template<typename T>
class minmax
{
public:
    minmax( ) : min_(std::numeric_limits<T>::max( )),
               max_(std::numeric_limits<T>::min( ))
    {}

    void operator( )(const T& x) {
        if (x < min_) min_ = x;
        if (max_ < x) max_ = x;
    }

    operator std::pair<T,T>( ) const {
        return std::make_pair(min_, max_);
    }
private:
    T min_;
    T max_;
};
```

```
int main( )
{
    std::vector<int> v;
    // Fill v with data.
    std::pair<int,int> mm =
        std::for_each(v.begin( ), v.end( ), minmax<int>( ));
    // Do something with mm.
}
```

## See Also

[pair class template](#)

## Name

`operator==` function template — Compares for equality

## Synopsis

```
template <typename T1, typename T2>  
bool operator==(const pair<T1,T2>& a, const pair<T1,T2>& b);
```

Returns true if `a` and `b` are equal, that is, `a.first == b.first && a.second == b.second`.

## Name

`operator!=` function template — Compares for inequality

## Synopsis

```
namespace rel_ops {  
    template<typename T>  
        bool operator!=(const T& a, const T& b);  
}  
template <typename T1, typename T2>  
    bool operator!=(const pair<T1,T2>& a, const pair<T1,T2>& b);
```

Returns true if `a` and `b` are not equal, that is, `!(a == b)`.

## Name

`operator<` function template — Compares for less-than

## Synopsis

```
template <typename T1, typename T2>  
bool operator<(const pair<T1,T2>& a, const pair<T1,T2>& b);
```

Returns true if a is less than b, assuming that the first member is more significant than second. That is, the return value is `a.first < b.first || (!(b.first < a.first) && a.second < b.second)`.

## Name

`operator<=` function template — Compares for less-than-or-equal

## Synopsis

```
namespace rel_ops {  
    template<typename T>  
        bool operator<=(const T& a, const T& b);  
}  
template <typename T1, typename T2>  
bool operator<=(const pair<T1,T2>& a, const pair<T1,T2>& b);
```

Returns true if a is less than or equal to b, that is, ! (b < a).

## Name

`operator>` function template — Compares for greater-than

## Synopsis

```
namespace rel_ops {  
    template<typename T>  
        bool operator>(const T& a, const T& b);  
}  
template <typename T1, typename T2>  
bool operator>(const pair<T1,T2>& a, const pair<T1,T2>& b);
```

Returns true if a is greater than b, that is,  $b < a$ .



## Name

`operator>=` function template — Compares for greater-than-or-equal

## Synopsis

```
namespace rel_ops {
    template<typename T>
        bool operator>=(const T& a, const T& b);
}
template <typename T1, typename T2>
    bool operator>=(const pair<T1,T2>& a, const pair<T1,T2>& b);
```

Returns true if a is greater than or equal to b, that is, ! (a < b).

# Name

pair class template — Represents a pair of related objects

## Synopsis

```
template <typename T1, typename T2>
struct pair {
    typedef T1 first_type;
    typedef T2 second_type;
    T1 first;
    T2 second;
    pair( );
    pair(const T1& x, const T2& y);
    template<typename U, typename V> pair(const pair<U, V> &p);
};
```

The `pair` class template represents a pair of related objects, in which the relationship is defined by the programmer. The most common use of pairs is by the `map` class template, which stores pairs of keys and associated objects.

The Boost project has a generalization of `pair`, called `tuple`. See [Appendix B](#) for information about Boost.

The `pair` constructors are straightforward:

**pair** ( )

Initializes `first` as `T1( )` and `second` as `T2( )`

**pair** (const T1& x, const T2& y)

Initializes `first` with `x` and `second` with `y`

template<typename U, typename V>, **pair** (const pair<U, V> &p)

Initializes `first` with `p.first` and `second` with `p.second`, performing implicit conversions as needed

## See Also

[make\\_pair function template](#)

# Name

rel\_ops namespace — Relational operators

## Synopsis

```
namespace std {  
    namespace rel_ops {  
        template<typename T> bool operator!=(const T&, const T&);  
        template<typename T> bool operator> (const T&, const T&);  
        template<typename T> bool operator<=(const T&, const T&);  
        template<typename T> bool operator>=(const T&, const T&);  
    }  
}
```

The `std::rel_ops` namespace declares four comparison operators. The four operators are implemented in terms of the `==` and `<` operators. The `rel_ops` namespace has limited utility. If you are using an unusual class, which has only `operator==` and `operator<`, you can add a `using namespace std::rel_ops` directive to a function that makes heavy use of comparison operators and this unusual class. Even better, though, is fixing the class declaration to provide all necessary comparison operators. If you are writing a class that represents an ordered value, you should provide all six operators and not force your users to rely on `rel_ops`. The Boost project has templates that you can derive from to fill in all the relational operators, based on equality and less-than. See [Appendix B](#) for information about Boost.

## See Also

[operator!= function template](#), [operator> function template](#), [operator<= function template](#), [operator>= function template](#)

## <valarray>

The <valarray> header declares types and functions for operating on arrays of numerical values. The intention is to provide types that could be optimized on certain hardware platforms for computationally-intensive programs. The consensus in the C++ user community seems to be that the standard failed to live up to the intentions. Several other numerical libraries, such as Blitz++ and MTL, provide high-performance matrix solutions. (See [Appendix B](#) for more information about Blitz++.) Most programs do not need <valarray>.

A `valarray` is a class template that represents a one-dimensional array of numerical values. The array can grow at runtime. All the arithmetic operators and mathematical functions are overloaded to work with two `valarray` arguments or with a `valarray` and a scalar. You can also work with parts of an array: slices, generalized slices, masks, and indirect arrays.

A *slice* is a set of elements of a `valarray`, with a starting index, a count, and a stride (an index interval). A *generalized slice* (`gslice`) lets the stride count and length vary, which can be used to implement multidimensional arrays. A *mask* is a `valarray` of flags, in which the flags indicate whether the corresponding item is part of the masked array. An *indirect array* is an array of indices. Each of these concepts is explained in this section.

The most important distinguishing feature of `valarrays` is that they do not allow aliasing, that is, an object cannot be an element of more than one `valarray`. This enables additional optimizations that are not possible on ordinary arrays.

Because `valarray` is optimized for performance, no error-checking is performed. Referring to an index out of range or operating on arrays of different size result in undefined behavior—the same as with ordinary arrays. Unlike ordinary arrays, a convenient `size( )` member function helps to ensure that you do not make mistakes.

See the `<cmath>` header for scalar mathematical functions and [<numeric>](#) for a few numeric algorithms. See `<complex>` for complex numbers.

Throughout this section, examples show `valarray` objects and subsets printed using operator `<<`, which is shown in [Example 13-40](#).

#### *Example 13-40. Printing a valarray or subset array*

```
// Print a valarray on one line, enclosed by curly braces. For
// example:
// "{ 1 2 3 }".
template<typename T>
void print_valarray(std::ostream& out, const std::valarray<T>& a)
{
    out << '{';
    for (size_t i = 0; i < a.size( ); ++i)
        out << ' ' << a[i];
    out << " }";
}

// Print a slice_array, gslice_array, etc. by converting to a
// valarray.
// Converting a valarray to a valarray is wasteful, but harmless
// for these simple
// examples.
template<template<typename T> class U, typename T>
std::ostream& operator<<(std::ostream& out, const U<T>& x)
{
    print_valarray(out, static_cast<std::valarray<T> >(x));
    return out;
}
```

## Name

abs function template — Computes absolute value

## Synopsis

```
template<typename T> valarray<T> abs(const valarray<T>& a);
```

The `abs` function computes the absolute value of each element of `a`.

## See Also

`abs` in [<cmath>](#), `abs` in [<cstdlib>](#)

## Name

acos function template — Computes inverse cosine

## Synopsis

```
template<typename T> valarray<T> acos(const valarray<T>& a);
```

The `acos` function computes the inverse cosine of each element of `a`.

## See Also

acos in [<cmath>](#)



## Name

asin function template — Computes inverse sine

## Synopsis

```
template<typename T> valarray<T> asin(const valarray<T>& a);
```

The `asin` function computes the inverse sine of each element of `a`.

## See Also

`asin` in [<cmath>](#)

## Name

atan function template — Computes inverse tangent

## Synopsis

```
template<typename T> valarray<T> atan(const valarray<T>& a);
```

The atan function computes the inverse tangent of each element of a.

## See Also

[atan2 function template](#), atan in [<cmath>](#)

## Name

atan2 function template — Computes inverse tangent of two arguments

## Synopsis

```
template<typename T>
valarray<T> atan2(const valarray<T>& b, const valarray<T>& a);
template<typename T> valarray<T> atan2(const valarray<T>& b,
const T& x);
template<typename T> valarray<T> atan2(const T& y, const
valarray<T>& a);
```

The `atan2` function computes the inverse tangent of  $y/x$ , in which  $y$  is a scalar or an element of `b`, and  $x$  is a scalar or an element of `a`.

## See Also

[atan function template](#), `atan2` in [<cmath>](#)

## Name

cos function template — Computes cosine

## Synopsis

```
template<typename T> valarray<T> cos(const valarray<T>& a);
```

The `cos` function computes the cosine of each element of `a`.

## See Also

cos in [<cmath>](#)

## Name

cosh function template — Computes hyperbolic cosine

## Synopsis

```
template<typename T> valarray<T> cosh(const valarray<T>& a);
```

The `cosh` function computes the hyperbolic cosine of each element of `a`.

## See Also

`cosh` in [<cmath>](#)

## Name

exp function template — Computes exponential

## Synopsis

```
template<typename T> valarray<T> exp(const valarray<T>& a);
```

The exp function computes the exponential  $e^x$  for each element  $x$  of  $a$ .

## See Also

exp in [<cmath>](#)

# Name

gslice class — Generalized slice

## Synopsis

```
class gslice {
public:
    gslice( );
    gslice(size_t start, const valarray<size_t>& size,
           const valarray<size_t>& stride);
    size_t start( ) const;
    valarray<size_t> size( ) const;
    valarray<size_t> stride( ) const;
};
```

The `gslice` class describes a generalized slice of a `valarray`. A generalized slice is a subset of the elements of a `valarray`, characterized by a starting index and a set of sizes and strides. The `size` and `stride` arrays must have the same size. Each size/stride pair denotes a set of elements at periodic indices. The number of elements in the generalized slice is equal to the product of the values in the `size` array. The elements are taken from a `valarray` at each index  $i$ :

$$i = \text{start} + \sum k_j \times \text{stride}_j$$

in which  $k_j$  takes all the values in the range  $[0, \text{size}[j])$ , and  $j$  is in the range  $[0, \text{stride.size}())$ . The highest value of  $j$  varies fastest. With a single element in `stride` and `size`, `gslice` is the same as plain `slice`. [Example 13-41](#) demonstrates `gslice` more clearly. Pay particular attention to the final `gslice`, where you can see how the indices advance, first with a stride of 3 ( $k_1$  ranges from 0 to 2), then with a stride of 2 ( $k_0$  ranges from 0 to 3)

## Example

### *Example 13-41. Generalized slicing of a valarray*

```
// Construct valarray objects from a few integers.
std::valarray<std::size_t> va(std::size_t a0)
{
    std::valarray<std::size_t> result(1);
    result[0] = a0;
    return result;
}

std::valarray<std::size_t> va(std::size_t a0, std::size_t a1)
{
    std::valarray<std::size_t> result(2);
    result[0] = a0;
    result[1] = a1;
    return result;
}

int main( )
{
    using namespace std;
    valarray<int> a(24);
    for (size_t i = 0; i < a.size( ); ++i)
        a[i] = i;
    cout << a << '\n';
    // Prints { 0 1 2 3 4 5 6 7 8 9 10 11 ... 20 21 22 23 }

    cout << a[slice(1, 4, 3)] << '\n';
    // Prints { 1 4 7 10 }
    cout << a[gslice(1, va(4), va(3))] << '\n';
    // Prints { 1 4 7 10 }

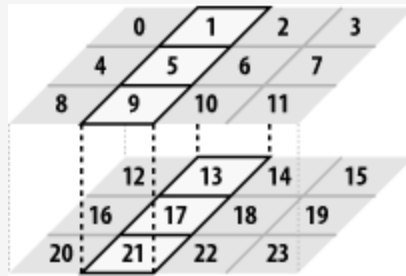
    const valarray<int> const_a(a);
    cout << const_a[gslice(2, va(4, 3), va(2, 3))] << '\n';
    // Prints { 2 5 8 4 7 10 6 9 12 8 11 14 }
}
```

Notice also that the final `gslice` requires a `const valarray`. This is because it contains degenerate slices, in which an element (e.g., 8) appears more than once in the result. The aliasing rules of a `valarray` prohibit multiple references to the same element, so if a `const valarray` were not used, the results would be undefined. By using a `const valarray`, the result is a copy of the sliced elements,



so the two occurrences of element 8 are separate objects, not aliases for the same object, and disaster is averted.

A generalized slice is most often used to represent a multidimensional array. For example, you can treat a `valarray` of 24 elements as a  $2 \times 3 \times 4$  matrix. To extract a plane of the matrix, you can use a `gslice`. [Figure 13-27](#) depicts the matrix and the plane. [Example 13-42](#) shows the code.



*Figure 13-27. A 3-D matrix stored in a `valarray`*

*Example 13-42. Using `gslice` for multidimensional arrays*

```
// See Example 13-41 for the va function.
int main( )
{
    using namespace std;
    valarray<int> a(24);
    for (size_t i = 0; i < a.size( ); ++i)
        a[i] = i;
    cout << a[gslice(1, va(2, 3), va(12, 4))] << '\n';
    // Prints: { 1 5 9 13 17 21 }
}
```

To create an  $n$ -dimensional submatrix of an  $m$ -dimensional matrix, the `size` and `stride` arrays must both have length  $n$ . The `size` array determines the dimensions of the result.

Use the subscript operator to take a generalized slice of a `valarray`. You can assign a `valarray` to a generalized slice, in which the righthand side of the assignment must have the same size as the size of the slice. You can also convert the slice to a `valarray`, which copies only those elements of the slice to the new `valarray`.

When you take a generalized slice of a `valarray`, the result is a `gslice_array` object, but the `gslice_array` type is mostly transparent to the programmer. See `gslice_array` later in this section for details.

## **See Also**

[gslice\\_array class template](#), [slice class](#)

# Name

gslice\_array class template — Helper class for generalized slices

## Synopsis

```
template <typename T>
class gslice_array {
public:
    typedef T value_type;
    void operator=(const valarray<T>&) const;
    void operator*=(const valarray<T>&) const;
    void operator/=(const valarray<T>&) const;
    void operator%=(const valarray<T>&) const;
    void operator+=(const valarray<T>&) const;
    void operator-= (const valarray<T>&) const;
    void operator^=(const valarray<T>&) const;
    void operator&=(const valarray<T>&) const;
    void operator|=(const valarray<T>&) const;
    void operator<<=(const valarray<T>&) const;
    void operator>>=(const valarray<T>&) const;
    void operator=(const T&);
    ~gslice_array( );
private:
    gslice_array( );
    gslice_array(const gslice_array&);
    gslice_array& operator=(const gslice_array&);
};
```

The `gslice_array` class template represents a subset of the elements of a `valarray`, called a generalized slice. To create a generalized slice, use `valarray`'s `operator[]` with an argument of type `gslice`.

For some operations, the `gslice_array` object is transparent. In particular, you can assign a `valarray` to a `gslice_array` object (provided they have the same size), or you can construct a new `valarray` from a `gslice_array`.

If you want to perform other operations, such as non-assignment arithmetic, you must explicitly convert the

gslice\_array to valarray, as demonstrated in [Example 13-43](#).

## Example

*Example 13-43. Using gslice\_array*

```
// See Example 13-41 for the va function.
int main( )
{
    using namespace std;
    const int data[] = { 1, 2, 3, 4, 5, 6, 7, 8 };
    valarray<int> a(data, sizeof(data)/sizeof(data[0]));
    cout << a << '\n';
    // Prints { 1 2 3 4 5 6 7 8 }

    cout << a[gslice(1, va(2, 2), va(4, 2))] << '\n'
          << a[gslice(0, va(2, 2), va(4, 2))] << '\n';
    // prints:
    // { 2 4 6 8 }
    // { 1 3 5 7 }

    // operator+ is not defined for gslice_array, so cast to
    valarray to perform
    // addition.
    cout <<
        static_cast<valarray<int> >(a[gslice(1, va(2,2), va(4,2))]) +
        static_cast<valarray<int> >(a[gslice(0, va(2,2), va(4,2))])
        << '\n';
    // Prints: { 3 7 11 15 }

    // Simple assignment does not require casting.
    a[gslice(0, va(2, 2), va(4, 2))] = 0;
    cout << a << '\n';
    // Prints: { 0 2 0 4 0 6 0 8 }

    // Computational assignment does not require casting.
    valarray<int> ten(10, 4);
    a[gslice(1, va(2, 2), va(4, 2))] *= ten;
    cout << a << '\n';
    // Prints: { 0 20 0 40 0 60 0 80 }
}
```

The members of gslice\_array are straightforward. When using any of the assignment operators, the valarray on the righthand side must be the same size as the gslice\_array

on the lefthand side. You can also assign a scalar to every element of the array. Note that the default constructor, copy constructor, and copy assignment operator are all private. The purpose of this is to restrict the use of `gslice_array` so it can be used only as a return value from `valarray`'s `operator[]`.

## See Also

[gslice class](#), [indirect\\_array\\_class template](#), [mask\\_array\\_class template](#), [slice\\_array\\_class template](#), [valarray\\_class template](#)

# Name

`indirect_array` class template — Helper class for indirect arrays

# Synopsis

```
template <typename T>
class indirect_array {
public:
    typedef T value_type;
    void operator=(const valarray<T>&) const;
    void operator*=(const valarray<T>&) const;
    void operator/=(const valarray<T>&) const;
    void operator%=(const valarray<T>&) const;
    void operator+=(const valarray<T>&) const;
    void operator-= (const valarray<T>&) const;
    void operator^=(const valarray<T>&) const;
    void operator&=(const valarray<T>&) const;
    void operator|=(const valarray<T>&) const;
    void operator<<=(const valarray<T>&) const;
    void operator>>=(const valarray<T>&) const;
    void operator=(const T&);
    ~indirect_array( );
private:
    indirect_array( );
    indirect_array(const indirect_array&);
    indirect_array& operator=(const indirect_array&);
};
```

The `indirect_array` class template represents a subset of the elements of a `valarray`. To create an indirect subset, use `valarray`'s `operator[]` with an argument of type `valarray<size_t>`. The elements of the argument are the desired indices in the subset.

For some operations, the `indirect_array` object is transparent. In particular, you can assign a `valarray` to an `indirect_array` object (provided they have the same size), or you can construct a new `valarray` from an `indirect_array`.

If you want to perform other operations, such as non-assignment arithmetic, you must explicitly convert the `indirect_array` to `valarray`, as demonstrated in [Example 13-44](#).

## Example

*Example 13-44. Using `indirect_array`*

```
int main( )
{
    using namespace std;
    const int data[] = { 1, 2, 3, 4, 5, 6, 7, 8 };
    valarray<int> a(data, sizeof(data)/sizeof(data[0]));
    cout << a << '\n';
    // Prints: { 1 2 3 4 5 6 7 8 }

    // Specify the indices into a.
    const size_t p[] = { 2, 3, 5, 7 };
    valarray<size_t> indices(p, sizeof(p)/sizeof(p[0]));
    cout << a[indices] << '\n';
    // Prints: { 3 4 6 8 }

    // Add 10 to the elements at the desired indices.
    valarray<int> ten(10, 4);
    a[indices] += ten;
    cout << a << '\n';
    // Prints: { 1 2 13 14 5 16 7 18 }

    // Must cast to perform ordinary arithmetic.
    cout << static_cast<valarray<int>> >(a[indices])
         * ten << '\n';
    // Prints: { 130 140 160 180 }
}
```

The members of `indirect_array` are straightforward. When using any of the assignment operators, the `valarray` on the righthand side must be the same size as the `indirect_array` on the lefthand side. You can also assign a scalar to every element of the array. Note that the default constructor, copy constructor, and copy assignment operator are all private. The purpose of this is to restrict

the use of `indirect_array` so it can be used only as a return value from `valarray`'s `operator[]`.

## See Also

[gslice\\_array class template](#), [mask\\_array class template](#),  
[slice\\_array class template](#), [valarray class template](#)



## Name

log function template — Computes natural logarithm

## Synopsis

```
template<typename T> valarray<T> log(const valarray<T>& a);
```

The `log` function computes the natural (base e) logarithm of each element of `a`.

## See Also

log in [<cmath>](#)

## Name

log10 function template — Computes common logarithm

## Synopsis

```
template<typename T> valarray<T> log10(const valarray<T>& a);
```

The `log10` function computes the common (base 10) logarithm of each element of `a`.

## See Also

log10 in [<cmath>](#)

# Name

mask\_array class template — Helper class for mask arrays

## Synopsis

```
template <typename T>
class mask_array {
public:
    typedef T value_type;
    void operator=(const valarray<T>&) const;
    void operator*=(const valarray<T>&) const;
    void operator/=(const valarray<T>&) const;
    void operator%=(const valarray<T>&) const;
    void operator+=(const valarray<T>&) const;
    void operator-=(const valarray<T>&) const;
    void operator^=(const valarray<T>&) const;
    void operator&=(const valarray<T>&) const;
    void operator|=(const valarray<T>&) const;
    void operator<<=(const valarray<T>&) const;
    void operator>>=(const valarray<T>&) const;
    void operator=(const T&);
    ~mask_array( );
private:
    mask_array( );
    mask_array(const mask_array&);
    mask_array& operator=(const mask_array&);
};
```

The `mask_array` class template represents a subset of the elements of a `valarray`. To create a mask subset, use `valarray`'s `operator[]` with an argument of type `valarray<bool>`. An element is included in the result set if the corresponding element in the argument is `true`.

For some operations, the `mask_array` object is transparent. In particular, you can assign a `valarray` to a `mask_array` object (provided they have the same size), or you can construct a new `valarray` from a `mask_array`.

If you want to perform other operations, such as non-assignment arithmetic, you must explicitly convert the

mask\_array to valarray, as demonstrated in [Example 13-45](#).

## Example

*Example 13-45. Using mask\_array*

```
// Simple average
template<typename T>
T avg(const std::valarray<T>& a)
{
    return a.sum( ) / a.size( );
}

int main( )
{
    using namespace std;
    const int data[] = { 1, -3, 10, 42, -12, 13, -7, 69 };
    valarray<int> a(data, sizeof(data)/sizeof(data[0]));
    cout << a << '\n';
    // Prints: { 1 -3 10 42 -12 13 -7 69 }

    // Print the values that are above average.
    cout << "avg=" << avg(a) << '\n';
    cout << a[a > avg(a)] << '\n';
    // Prints: { 42 69 }

    // Force all negative values to be 0. Notice how no cast is
    // needed for the
    // simple assignment.
    a[a < 0] = 0;
    cout << a << '\n';
    // Prints: { 1 0 10 42 0 13 0 69 }

    // Other operations, such as multiplication by a scalar, are
    // defined only for
    // valarray, so a cast is needed.
    cout << static_cast<valarray<int>>(a[a > 0]) * -1 << '\n';
    // Prints: { -1 -10 -42 -13 -69 }
}
```

The members of mask\_array are straightforward. When using any of the assignment operators, the valarray on the righthand side must be the same size as the mask\_array on the lefthand side. You can also assign a scalar to every

element of the array. Note that the default constructor, copy constructor, and copy assignment operator are all private. The purpose of this is to restrict the use of `mask_array` so it can be used only as a return value from `valarray`'s `operator[]`.

## See Also

[gslice\\_array class template](#), [indirect\\_array class template](#), [slice\\_array class template](#), [valarray class template](#)

# Name

operator\* function template — Performs multiplication

# Synopsis

```
template<typename T>
valarray<T> operator*(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator*(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator*(const T& x, const valarray<T>& b);
```

The \* operator performs elementwise multiplication. It multiplies each  $x * y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When multiplying two arrays, they must have the same size. The resulting array has the same size as the argument array(s).

# Name

operator/ function template — Performs division

# Synopsis

```
template<typename T>
valarray<T> operator/(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator/(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator/(const T& x, const valarray<T>& b);
```

The / operator performs elementwise division. It divides each  $x / y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When dividing two arrays, they must have the same size. The resulting array has the same size as the argument array(s).

# Name

operator+ function template — Performs addition

# Synopsis

```
template<typename T>
valarray<T> operator+(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator+(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator+(const T& x, const valarray<T>& b);
```

The + operator performs elementwise addition. It adds each  $x + y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When adding two arrays, they must have the same size. The resulting array has the same size as the argument array(s).



# Name

operator- function template — Performs subtraction

# Synopsis

```
template<typename T>
valarray<T> operator-(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator-(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator-(const T& x, const valarray<T>& b);
```

The `-` operator performs elementwise subtraction. It subtracts each  $x - y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When subtracting two arrays, they must have the same size. The resulting array has the same size as the argument array(s).

# Name

operator& function template — Performs bitwise and

# Synopsis

```
template<typename T>
valarray<T> operator&(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator&(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator&(const T& x, const valarray<T>& b);
```

The & operator performs *bitwise and* on each  $x$  &  $y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which operator & is defined.

# Name

operator| function template — Performs bitwise or

# Synopsis

```
template<typename T>
valarray<T> operator|(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator|(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator|(const T& x, const valarray<T>& b);
```

The | operator performs *bitwise inclusive or* on each  $x | y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which `operator |` is defined.

## Name

`operator^` function template — Performs bitwise exclusive or

## Synopsis

```
template<typename T>
valarray<T> operator^(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator^(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator^(const T& x, const valarray<T>& b);
```

The `^` operator performs *bitwise exclusive or* on each  $x \wedge y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which `operator ^` is defined.

## Name

operator>> function template — Performs right shift

## Synopsis

```
template<typename T>
valarray<T> operator>>(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator>>(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator>>(const T& x, const valarray<T>& b);
```

The >> operator performs right shift on each  $x \gg y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which `operator >>` is defined.

## Name

`operator<<` function template — Performs left shift

## Synopsis

```
template<typename T>
valarray<T> operator<<(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<T> operator<<(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> operator<<(const T& x, const valarray<T>& b);
```

The `<<` operator performs left shift on each `x << y`, in which `x` is a scalar or an element of `a`, and `y` is a scalar or an element of `b`. When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type `T` must be one for which `operator <<` is defined.

# Name

operator&& function template — Performs logical and

## Synopsis

```
template<typename T>  
valarray<bool> operator&&(const valarray<T>& a, const  
valarray<T>& b);  
template<typename T>  
valarray<bool> operator&&(const valarray<T>& a, const T& y);  
template<typename T>  
valarray<bool> operator&&(const T& x, const valarray<T>& b);
```

The && operator performs *logical and* on each  $x$  &&  $y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which operator && is defined. As with any other overloaded operator &&, short-cut evaluation is not supported.

# Name

operator|| function template — Performs logical or

## Synopsis

```
template<typename T>
valarray<bool> operator||(const valarray<T>& a, const
valarray<T>& b);
template<typename T>
valarray<bool> operator||(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator||(const T& x, const valarray<T>& b);
```

The || operator performs *logical or* on each  $x || y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which `operator ||` is defined and yields a `bool` result or a result that can be converted to `bool`. As with any other overloaded operator ||, short-cut evaluation is not supported.



# Name

`operator==` function template — Compares for equality

## Synopsis

```
template<typename T>
valarray<bool> operator==(const valarray<T>& a, const
valarray<T>& b);
template<typename T>
valarray<bool> operator==(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator==(const T& x, const valarray<T>& b);
```

The `==` operator compares each `x == y`, in which `x` is a scalar or an element of `a`, and `y` is a scalar or an element of `b`. When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type `T` must be one for which `operator ==` is defined and yields a `bool` result or a result that can be converted to `bool`.

## Name

operator!= function template — Compares for inequality

## Synopsis

```
template<typename T>
valarray<bool> operator!=(const valarray<T>& a, const
valarray<T>& b);
template<typename T>
valarray<bool> operator!=(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator!=(const T& x, const valarray<T>& b);
```

The != operator compares each  $x \neq y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which operator != is defined and yields a bool result or a result that can be converted to bool.

# Name

operator< function template — Compares for less-than

## Synopsis

```
template<typename T>
valarray<bool> operator<(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<bool> operator<(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator<(const T& x, const valarray<T>& b);
```

The < operator compares each  $x < y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which operator < is defined and yields a `bool` result or a result that can be converted to `bool`.

## Name

`operator<=` function template — Compares for less-than-or-equal

## Synopsis

```
template<typename T>
valarray<bool> operator<=(const valarray<T>& a, const
valarray<T>& b);
template<typename T>
valarray<bool> operator<=(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator<=(const T& x, const valarray<T>& b);
```

The `<=` operator compares each  $x \leq y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which `operator <=` is defined and yields a `bool` result or a result that can be converted to `bool`.

# Name

operator> function template — Compares for greater-than

## Synopsis

```
template<typename T>
valarray<bool> operator>(const valarray<T>& a, const valarray<T>&
b);
template<typename T>
valarray<bool> operator>(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator>(const T& x, const valarray<T>& b);
```

The > operator compares each  $x > y$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ . When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type  $T$  must be one for which operator > is defined and yields a `bool` result or a result that can be converted to `bool`.

## Name

`operator>=` function template — Compares for greater-than-or-equal

## Synopsis

```
template<typename T>
valarray<bool> operator>=(const valarray<T>& a, const
valarray<T>& b);
template<typename T>
valarray<bool> operator>=(const valarray<T>& a, const T& y);
template<typename T>
valarray<bool> operator>=(const T& x, const valarray<T>& b);
```

The `>=` operator compares each `x >= y`, in which `x` is a scalar or an element of `a`, and `y` is a scalar or an element of `b`. When operating on two arrays, they must have the same size. The resulting array has the same size as the argument array(s). The type `T` must be one for which `operator >=` is defined and yields a `bool` result or a result that can be converted to `bool`.

## Name

pow function template — Computes power

## Synopsis

```
template<typename T>
valarray<T> pow(const valarray<T>& a, const valarray<T>& b);
template<typename T>
valarray<T> pow(const valarray<T>& a, const T& y);
template<typename T>
valarray<T> pow(const T& x, const valarray<T>& b);
```

The pow function computes the power  $xy$ , in which  $x$  is a scalar or an element of  $a$ , and  $y$  is a scalar or an element of  $b$ .

## See Also

pow in [<cmath>](#)

## Name

sin function template — Computes sine

## Synopsis

```
template<typename T> valarray<T> sin(const valarray<T>& a);
```

The `sin` function computes the sine of the elements of `a`.

## See Also

sin in [<cmath>](#)



## Name

sinh function template — Computes hyperbolic sine

## Synopsis

```
template<typename T> valarray<T> sinh(const valarray<T>& a);
```

The `sinh` function computes the hyperbolic sine of the elements of `a`.

## See Also

`sinh` in [<cmath>](#)

# Name

slice class — Slice of an array

# Synopsis

```
class slice {  
public:  
    slice( );  
    slice(size_t, size_t, size_t);  
    size_t start( ) const;  
    size_t size( ) const;  
    size_t stride( ) const;  
};
```

The slice class describes a slice of a valarray. A slice is a subset of the elements of a valarray at periodic indices. The slice has a starting index, a size, and a stride, in which the stride is the index interval. [Figure 13-28](#) depicts slice(1,3,4) of a valarray.

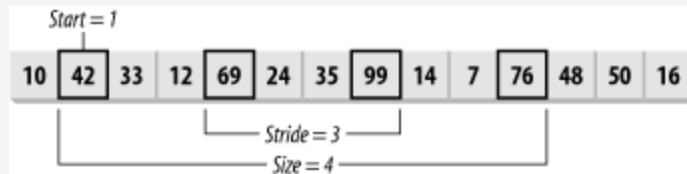


Figure 13-28. Slicing a valarray

Use the subscript operator to take a slice of a valarray. You can assign a valarray to a slice, in which the righthand side of the assignment must have the same size as the size of the slice. You can also convert the slice to a valarray, which copies only those elements of the slice to the new valarray.

When you take a slice of a valarray, the result is a slice\_array object, but the slice\_array type is mostly transparent to the programmer. See slice\_array later in this section for details.

You can use a slice to treat a valarray as a two-dimensional matrix. A slice can specify a row or column of the matrix. For an  $n \times m$  matrix, row  $r$  is `slice(r*m, m, 1)`, and column  $c$  is `slice(c, n, m)`, as you can see in [Example 13-46](#).

## Example

*Example 13-46. A simple 2-D matrix class*

```
template<typename T>
class matrix2D {
public:
    matrix2D(std::size_t rows, std::size_t columns) :
        rows_(rows), cols_(columns), data_(rows * columns) {}
    std::size_t rows( ) const { return rows_; }
    std::size_t cols( ) const { return cols_; }
    std::valarray<T> row(std::size_t r) const
        { return data_[std::slice(r*cols( ),cols( ), 1)]; }
    std::valarray<T> col(std::size_t c) const
        { return data_[std::slice(c, rows( ), cols( ))]; }
    std::slice_array<T> row(std::size_t r)
        { return data_[std::slice(r*cols( ),cols( ), 1)]; }
    std::slice_array<T> col(std::size_t c)
        { return data_[std::slice(c, rows( ), cols( ))]; }
    T& operator( )(std::size_t r, std::size_t c)
        { return data_[r*cols( )+c]; }
    T operator( )(std::size_t r, std::size_t c) const
        { return row(r)[c]; }
    matrix2D<T> transpose( ) {
        matrix2D<T> result(cols( ), rows( ));
        for (std::size_t i = 0; i < rows( ); ++i)
            result.col(i) = static_cast<std::valarray<T> >(row(i));
        return result;
    }
private:
    std::size_t rows_;
    std::size_t cols_;
    std::valarray<T> data_;
};
```

## See Also

[gslice class](#), [slice\\_array class template](#)

# Name

slice\_array class template — Helper class for slice

# Synopsis

```
template <typename T>
class slice_array {
public:
    typedef T value_type;
    void operator=(const valarray<T>&) const;
    void operator*=(const valarray<T>&) const;
    void operator/=(const valarray<T>&) const;
    void operator%=(const valarray<T>&) const;
    void operator+=(const valarray<T>&) const;
    void operator-=(const valarray<T>&) const;
    void operator^=(const valarray<T>&) const;
    void operator&=(const valarray<T>&) const;
    void operator|=(const valarray<T>&) const;
    void operator<<=(const valarray<T>&) const;
    void operator>>=(const valarray<T>&) const;
    void operator=(const T&);
    ~slice_array( );
private:
    slice_array( );
    slice_array(const slice_array&);
    slice_array& operator=(const slice_array&);
};
```

The `slice_array` class template represents a subset of the elements of a `valarray`, taken at periodic indices, called a slice. To create a slice, use `valarray`'s `operator[]` with an argument of type `slice`.

For some operations, the `slice_array` object is transparent. In particular, you can assign a `valarray` to a `slice_array` object (provided they have the same size), or you can construct a new `valarray` from a `slice_array`.

If you want to perform other operations, such as non-assignment arithmetic, you must explicitly convert the

`slice_array` to `valarray`, as demonstrated in [Example 13-47](#).

## Example

### *Example 13-47. Slicing a valarray*

```
int main( )
{
    using namespace std;
    const int data[] = { 1,2,3,4,5,6,7,8,9,10,11,12,13 };
    valarray<int> v(data, sizeof(data)/sizeof(data[0]));
    const int newdata[] = { 30, 70, 110 };
    valarray<int> rpl(newdata, 3);
    v[slice(2, 3, 4)] = rpl;
    cout << v << '\n';
    // Prints: { 1 2 30 4 5 6 70 8 9 10 110 12 13}
    v[slice(3, 4, 2)] = -1;
    cout << v << '\n';
    // Prints: { 1 2 30 -1 5 -1 70 -1 9 -1 110 12 13}
    valarray<int> mult(3, 2);
    v[slice(8, 2, 3)] *= mult;
    cout << v << '\n';
    // Prints: { 1 2 30 -1 5 -1 70 -1 27 -1 110 36 13}
    cout << static_cast<valarray<int> >(v[slice(1, 5, 2)])
        << '\n';
    // Prints: { 2 -1 -1 -1 -1}
    cout << static_cast<valarray<int> >(v[slice(4, 3, 2)]) +
        static_cast<valarray<int> >(v[slice(2, 3, 2)])
        << '\n';
    // Prints: { 35 75 97}
}
```

The members of `slice_array` are straightforward. When using any of the assignment operators, the `valarray` on the righthand side must be the same size as the `slice_array` on the lefthand side. You can also assign a scalar to every element of the array. Note that the default constructor, copy constructor, and copy assignment operator are all private. The purpose of this is to restrict the use of `slice_array` so it can be used only as a return value from `valarray`'s `operator[]`.

## See Also

[gslice\\_array\\_class template](#), [indirect\\_array class template](#), [mask\\_array class template](#), [slice class](#), [valarray class template](#)

## Name

sqrt function template — Computes square root

## Synopsis

```
template<typename T> valarray<T> sqrt(const valarray<T>& a);
```

The `sqrt` function computes the square root of the elements of `a`.

## See Also

sqrt in [<cmath>](#)

## Name

tan function template — Computes tangent

## Synopsis

```
template<typename T> valarray<T> tan(const valarray<T>& a);
```

The tan function computes the tangent of the elements of a.

## See Also

tan in [<cmath>](#)



## Name

tanh function template — Computes hyperbolic tangent

## Synopsis

```
template<typename T> valarray<T> tanh(const valarray<T>& a);
```

The `tanh` function computes the hyperbolic tangent of the elements of `a`.

## See Also

`tanh` in [<cmath>](#)

# Name

valarray class template — Array of values

## Synopsis

```
template<typename T>
class valarray {
public:
    typedef T value_type;

    valarray( );
    explicit valarray(size_t);
    valarray(const T&, size_t);
    valarray(const T*, size_t);
    valarray(const valarray&);
    valarray(const slice_array<T>&);
    valarray(const gslice_array<T>&);
    valarray(const mask_array<T>&);
    valarray(const indirect_array<T>&);
    ~valarray( );

    valarray<T>& operator=(const valarray<T>&);
    valarray<T>& operator=(const T&);
    valarray<T>& operator=(const slice_array<T>&);
    valarray<T>& operator=(const gslice_array<T>&);
    valarray<T>& operator=(const mask_array<T>&);
    valarray<T>& operator=(const indirect_array<T>&);

    T operator[](size_t) const;
    T& operator[](size_t);

    valarray<T> operator[](slice) const;
    slice_array<T> operator[](slice);
    valarray<T> operator[](const gslice&) const;
    gslice_array<T> operator[](const gslice&);
    valarray<T> operator[](const valarray<bool>&) const;
    mask_array<T> operator[](const valarray<bool>&);
    valarray<T> operator[](const valarray<size_t>&) const;
    indirect_array<T> operator[](const valarray<size_t>&);

    valarray<T> operator+( ) const;
    valarray<T> operator-( ) const;
    valarray<T> operator~( ) const;
    valarray<bool> operator!( ) const;
```

```

valarray<T>& operator*= (const T&);
valarray<T>& operator/= (const T&);
valarray<T>& operator%= (const T&);
valarray<T>& operator+= (const T&);
valarray<T>& operator-= (const T&);
valarray<T>& operator^= (const T&);
valarray<T>& operator&= (const T&);
valarray<T>& operator|= (const T&);
valarray<T>& operator<<= (const T&);
valarray<T>& operator>>= (const T&);
valarray<T>& operator*= (const valarray<T>&);
valarray<T>& operator/= (const valarray<T>&);
valarray<T>& operator%= (const valarray<T>&);
valarray<T>& operator+= (const valarray<T>&);
valarray<T>& operator-= (const valarray<T>&);
valarray<T>& operator^= (const valarray<T>&);
valarray<T>& operator|= (const valarray<T>&);
valarray<T>& operator&= (const valarray<T>&);
valarray<T>& operator<<= (const valarray<T>&);
valarray<T>& operator>>= (const valarray<T>&);

size_t size( ) const;
T sum( ) const;
T min( ) const;
T max( ) const;
valarray<T> shift (int) const;
valarray<T> cshift(int) const;
valarray<T> apply(T func(T)) const;
valarray<T> apply(T func(const T&)) const;
void resize(size_t sz, T c = T( ));
};

```

The `valarray` class template represents an array of numeric values, with restrictions that permit an implementation to optimize performance. In particular, an object cannot be an element of more than one array. A subset, such as a generalized slice or indirect array, cannot specify a single element more than once, or else the behavior is undefined.

You can instantiate `valarray` with any numerical type as its template parameter if you limit yourself to using only operations that are defined for that type. For example, you

cannot use operator<< on valarray<double> because you cannot use operator<< on scalars of type double.

Examples of using valarray can be found throughout this section.

The following are the members of valarray:

**valarray** ( )

Constructs an empty valarray.

explicit **valarray** (size\_t n)

Constructs a valarray of length n, in which all elements are initialized to T( ).

**valarray** (const T& x, size\_t n)

Construct a valarray that contains n copies of x.

**valarray** (const T\* x, size\_t n)

Constructs a valarray by copying n elements from the array x.

**valarray** (const valarray& a)

Constructs a valarray by copying a.

**valarray** (const slice\_array<T>& a), **valarray** (const gslice\_array<T>& a), **valarray** (const mask\_array<T>& a), **valarray** (const indirect\_array<T>& a)

Constructs a valarray by copying the elements referenced by a.

valarray<T> **apply** (T func(T)) const, valarray<T> **apply** (T func(const T&)) const

Returns a new array whose contents are the result of calling func for each element of the original array.

valarray<T> **cshift** (int n) const

Performs a circular shift (rotation) by  $n$  places. The return value is a new array that has the same size as the original, but the element at new index  $i$  comes from  $(i + n) \% \text{size}()$  in the original.

**T max** ( ) const

Returns the largest element of the array or an undefined value if the array is empty. Elements are compared using operator<.

**T min** ( ) const

Returns the smallest element of the array or an undefined value if the array is empty. Elements are compared using operator<.

**valarray<T>& operator=** (const valarray<T>& a)

Each element of this array is assigned the corresponding elements of  $a$ . If  $\text{size}() \neq a.\text{size}()$ , the behavior is undefined.

**valarray<T>& operator=** (const T& x)

Each element of this array is assigned the scalar value  $x$ .

**valarray<T>& operator=** (const slice\_array<T>& a),  
**valarray<T>& operator=** (const gslice\_array<T>& a),  
**valarray<T>& operator=** (const mask\_array<T>& a),  
**valarray<T>& operator=** (const indirect\_array<T>& a)

Each element of this array is assigned the corresponding element from the subset array  $a$ . The value being assigned to an element must not depend on any other value in this array, that is,  $a$  cannot be a subset of  $*\text{this}$ , or, if it is, the element assigned to index  $i$  must depend only on index  $i$  and not on values at any other index.

**T operator[]** (size\_t i) const, **T& operator[]** (size\_t i)

Returns the element at index *i*. The behavior is undefined for  $i \geq \text{size}()$ . The anti-aliasing rule means that for any two distinct `valarray` objects *a* and *b*, and for all valid indices *i* and *j*, you can safely assume that the following is true: `&a[ i ] != &b[ j ]`. All values in a single `valarray` object are stored contiguously, just as in an ordinary array. References become invalid after a call to `resize`.

```
valarray<T> operator[] (slice) const,  
slice_array<T> operator[] (slice), valarray<T>  
operator[] (const gslice&) const, gslice_array<T>  
operator[] (const gslice&), valarray<T> operator[]  
(const valarray<bool>&) const, mask_array<T>  
operator[] (const valarray<bool>&), valarray<T>  
operator[] (const valarray<size_t>&) const,  
indirect_array<T> operator[] (const  
valarray<size_t>&)
```

Returns a subset of this `valarray`. A subset of a `const valarray` is a new `valarray`. A subset of a non-`const valarray` is a helper object that maintains a reference to the original `valarray`. (See the descriptions of the helper classes earlier in this section for details.) Briefly, the four kinds of subsets are:

- A `slice` object specifies a simple slice, suitable for extracting a row or column of a 2-D array.
- A `gslice` is a generalized slice, which permits multidimensional matrices.
- A `mask_array` is created by specifying an array of `bool` as the argument. If an element is `true`, that element is part of the resulting subset.
- An `indirect_array` is created by specifying an array of `size_t` indices as the argument. Each element

specifies an index, and the element at that index is added to the resulting array.

```
valarray<T> operator+ ( ) const, valarray<T>  
operator- ( ) const, valarray<T> operator~ ( )  
const, valarray<bool> operator! ( ) const
```

Returns a new valarray in which each new element is the result of applying the unary operator to the corresponding element in the original array.

```
valarray<T>& operator*|= (const T& x), valarray<T>&  
operator/|= (const T& x), valarray<T>& operator%|=  
(const T& x), valarray<T>& operator+= (const T& x),  
valarray<T>& operator-|= (const T& x), valarray<T>&  
operator^|= (const T& x), valarray<T>& operator&|=  
(const T& x), valarray<T>& operator||= (const T& x),  
valarray<T>& operator<<|= (const T& x), valarray<T>&  
operator>>|= (const T& x)
```

Modifies this valarray by applying the assignment operator to each element of the array. The return value is \*this.

```
valarray<T>& operator*|= (const valarray<T>& a),  
valarray<T>& operator/|= (const valarray<T>& a),  
valarray<T>& operator%|= (const valarray<T>& a),  
valarray<T>& operator+= (const valarray<T>& a),  
valarray<T>& operator-|= (const valarray<T>& a),  
valarray<T>& operator^|= (const valarray<T>& a),  
valarray<T>& operator&|= (const valarray<T>& a),  
valarray<T>& operator<<|= (const valarray<T>& a),  
valarray<T>& operator>>|= (const valarray<T>& a)
```

Modifies this valarray by applying the assignment operator to each element of the array and to the corresponding element of a. The behavior is undefined if `size( ) != a.size( )`. The return value is \*this.

`void resize (size_t sz, T x = T( ))`

Changes the size of the array to `sz` and reinitializes every element of the array to `x`.

`valarray<T> shift (int n) const`

Performs a shift by `n` places. The return value is a new array with the same size as the original array, but the element at index `i` in the new array comes from index `i + n` in the original array. If `i + n < 0` or `≥ size( )`, the new element is set to `T( )`.

`size_t size ( ) const`

Returns the number of elements in the array.

`T sum ( ) const`

Returns the sum of all the elements in the array using operator`+=`. If the array is empty, the value is undefined.

## See Also

[gslice class](#), [gslice\\_array class template](#), [indirect\\_array class template](#), [mask\\_array class template](#), [slice class](#), [slice\\_array class template](#)



## **<vector>**

The `<vector>` header is one of the standard container template headers. It declares the `vector` class template and a few global function templates that operate on vector objects.

A vector is a sequence container that yields linear performance for inserting and erasing at any point in the container, except the end, for which performance is constant. A vector supports random access iterators. A vector is best thought of as a generalization of arrays.

See [Chapter 10](#) for information about containers in.

## Name

operator== function template — Compares for equality

## Synopsis

```
template <typename T, typename A>  
bool operator==(const vector<T,A>& x, const vector<T,A>& y);  
template <typename Alloc>  
bool operator==(const vector<bool,Alloc>& x, const  
vector<bool,Alloc>& y);
```

The == operator returns true if x and y have the same size and their elements are equal, that is, `x.size() == y.size() && equals(x.begin(), x.end(), y.begin( ))`.

## See Also

equals in [<algorithm>](#)

## Name

operator!= function template — Compares for inequality

## Synopsis

```
template <typename T, typename A>  
bool operator!=(const vector<T,A>& x, const vector<T,A>& y);  
template <typename Alloc>  
bool operator!=(const vector<bool,Alloc>& x, const  
vector<bool,Alloc>& y);
```

The != operator returns ! (x == y).

## Name

operator< function template — Compares for less-than

## Synopsis

```
template <typename T, typename A>
bool operator<(const vector<T,A>& x, const vector<T,A>& y);
template <typename Alloc>
bool operator<(const vector<bool,Alloc>& x, const
vector<bool,Alloc>& y);
```

The < operator determines whether x is less than y using the same algorithm as

lexicographical\_compare(x.begin( ), x.end( ), y.begin( ), y.end( )).

## See Also

lexicographical\_compare in [<algorithm>](#)

## Name

`operator<=` function template — Compares for less-than-or-equal

## Synopsis

```
template <typename T, typename A>  
bool operator<=(const vector<T,A>& x, const vector<T,A>& y);  
template <typename Alloc>  
bool operator<=(const vector<bool,Alloc>& x, const  
vector<bool,Alloc>& y);
```

The `<=` operator returns `!(y < x)`.

## Name

operator> function template — Compares for greater-than

## Synopsis

```
template <typename T, typename A>  
bool operator>(const vector<T,A>& x, const vector<T,A>& y);  
template <typename Alloc>  
bool operator>(const vector<bool,Alloc>& x, const  
vector<bool,Alloc>& y);
```

The > operator returns (y < x).

## Name

operator>= function template — Compares for greater-than-or-equal

## Synopsis

```
template <typename T, typename A>  
bool operator>=(const vector<T,A>& x, const vector<T,A>& y);  
template <typename Alloc>  
bool operator>=(const vector<bool,Alloc>& x, const  
vector<bool,Alloc>& y);
```

The >= operator returns ! (x < y).

## Name

swap function template — Swaps contents of two vectors

## Synopsis

```
template <typename T, typename Alloc>  
void swap(vector<T,Alloc>& x, vector<T,Alloc>& y);  
template <typename Alloc>  
void swap(vector<bool,Alloc>& x, vector<bool,Alloc>& y);
```

The swap function template specialization is equivalent to calling `x.swap(y)`.

## See Also

swap in [<algorithm>](#)



# Name

vector class template — Array-like container

## Synopsis

```
template <typename T, typename Alloc = allocator<T> >
class vector {
public:
    typedef typename Alloc::reference reference;
    typedef typename Alloc::const_reference const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef T value_type;
    typedef Alloc allocator_type;
    typedef typename Alloc::pointer pointer;
    typedef typename Alloc::const_pointer const_pointer;
    typedef std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;

    explicit vector(const Alloc& = Alloc( ));
    explicit vector(size_type n, const T& value = T( ), const
Alloc& = Alloc( ));
    template <class InpIt>
vector(InpIt first, InpIt last, const Alloc& = Alloc( ));
vector(const vector<T,Alloc>& x);
~vector( );
vector<T,Alloc>& operator=(const vector<T,Alloc>& x);
    template <class InputIterator>
void assign(InputIterator first, InputIterator last);
void assign(size_type n, const T& u);
    allocator_type get_allocator( ) const;

    iterator begin( );
    const_iterator begin( ) const;
    iterator end( );
    const_iterator end( ) const;
    reverse_iterator rbegin( );
    const_reverse_iterator rbegin( ) const;
    reverse_iterator rend( );
    const_reverse_iterator rend( ) const;
```

```

size_type size( ) const;
size_type max_size( ) const;
void resize(size_type sz, T c = T( ));
size_type capacity( ) const;
bool empty( ) const;
void reserve(size_type n);
// Element access
reference operator[](size_type n);
const_reference operator[](size_type n) const;
const_reference at(size_type n) const;
reference at(size_type n);
reference front( );
const_reference front( ) const;
reference back( );
const_reference back( ) const;
// Modifiers
void push_back(const T& x);
void pop_back( );
iterator insert(iterator position, const T& x);
void insert(iterator position, size_type n, const T& x);
template <class InpIt>
void insert(iterator position, InpIt first, InpIt last);
iterator erase(iterator position);
iterator erase(iterator first, iterator last);
void swap(vector<T,Alloc>&);
void clear( );
};

```

The vector class template is a standard sequence container that is like an array: adding or removing from the end of the vector takes constant time (amortized over many such operations), adding or removing from anywhere else takes linear time, and random access happens in constant time.

Elements of a vector are stored contiguously, just like an ordinary array. For most cases in which you need an array, you should use a vector instead because a vector offers greater safety (no need for dynamic memory and raw pointers, the `at` member function checks array bounds, etc.)

All iterators and references to a vector's elements become invalid when the vector's internal array is resized, which

can happen for an insertion when the size matches the capacity, or when you explicitly change the size (e.g., by calling `resize`). You can ensure that an insertion does not force a resize by calling `reserve` to set the capacity prior to inserting one or more items. Iterators and references also become invalid when they are past (at a higher index) the point where an item is inserted or erased.

If you need a vector of Boolean values, consider using `deque<bool>` instead of `vector<bool>`. See `vector<bool>` for an explanation.

The following are the members of `vector`:

```
explicit vector (const Alloc& = Alloc( ))
```

Constructs an empty vector.

```
explicit vector (size_type n, const T& value = T( ),  
const Alloc& = Alloc( ))
```

Constructs a vector of size `n`, in which each element is initialized to `value`.

```
template <class InpIt>, vector (InpIt first, InpIt  
last, const Alloc& = Alloc( ))
```

Constructs an empty vector and copies `[first, last)` into the new vector unless `InputIterator` is an integral type, in which case the vector is constructed as though the arguments were cast as follows:

```
vector(static_cast<size_type>(first), static_cast<value_type>  
(last),  
alloc);
```

```
vector (const vector<T,Alloc>& v)
```

Constructs a copy of `v`.

```
vector<T,Alloc>& operator= (const vector<T,Alloc>&  
v)
```

Erases all the elements of the vector, then copies the elements from `v` into the vector.

```
template <class InputIterator>, void assign  
(InputIterator first, InputIterator last)
```

Erases all the elements of the vector, then copies the elements from `[first, last)` into the vector, unless `InputIterator` is an integral type, in which case the arguments are interpreted as though they were cast as follows:

```
    assign(static_cast<size_type>(first), static_cast<value_type>  
          (last));
```

```
void assign (size_type n, const T& value)
```

Erases all the elements of the vector, then inserts `n` copies of `value`.

```
const_reference at (size_type n) const  
reference at (size_type n)
```

Returns the element at index `n`. If `n >= size()`, throws `out_of_range`.

```
reference back ( ), const_reference back ( ) const
```

Returns the last element of the vector. Behavior is undefined if the vector is empty.

```
iterator begin ( ), const_iterator begin ( ) const
```

Returns an iterator that points to the first element of the vector.

```
size_type capacity ( ) const
```

Returns the number of elements the vector can store before it resizes itself.

```
void clear ( )
```

Erases all elements of the vector.

iterator **end** ( ), const\_iterator **end** ( ) const

Returns an iterator that points to one past the last element of the vector.

bool **empty** ( ) const

Returns `size( ) == 0`.

iterator **erase** (iterator position)

Erases the element at position.

iterator **erase** (iterator first, iterator last)

Erases all the elements in the range `[first, last)`.

reference **front** ( ), const\_reference **front** ( ) const

Returns the first element of the vector. Behavior is undefined if the vector is empty.

locator\_type **get\_allocator** ( ) const

Returns the allocator object.

iterator **insert** (iterator position, const T& x)

Inserts `x` before position.

void **insert** (iterator pos, size\_type n, const T& x)

Inserts `n` copies of `x` at `pos`.

template <class InpIt>, void **insert** (iterator pos, InpIt first, InpIt last)

Inserts the elements in the range `[first, last)` starting at position `pos`, unless `InputIterator` is an integral type, in which case the arguments are interpreted as though they were cast as follows:

```
insert(pos, static_cast<size_type>(first),
       static_cast<value_type>(last));
```

If an exception is thrown, such as `bad_alloc` when there is insufficient memory for a new element, the vector is

unchanged, and all iterators and references remain valid. If the exception is thrown from an element's copy constructor or assignment operator, however, the behavior is unspecified.

`size_type` **max\_size** ( ) const

Returns the size of the largest possible vector.

reference **operator[]** (size\_type n), const\_reference **operator[]** (size\_type n) const

Returns the element at index n. If  $n \geq \text{size}()$ , the behavior is undefined.

void **pop\_back** ( )

Erases the last element of the vector. The behavior is undefined if the vector is empty.

void **push\_back** (const T& x)

Inserts x as the new last element of the vector.

reverse\_iterator **rbegin** ( ), const\_reverse\_iterator **rbegin** ( ) const

Returns a reverse iterator that points to the last element of the vector.

reverse\_iterator **rend** ( ), const\_reverse\_iterator **rend** ( ) const

Returns a reverse iterator that points to one position before the first element of the vector.

void **reserve** (size\_type n)

Ensures that `capacity()` is at least n. Call `reserve` to avoid the need to reallocate the vector repeatedly when you know the vector will grow by small increments to a large size, or when you want to ensure that iterators do not become invalid after inserting one or more items. Note that `size()` does not change.

`void resize (size_type sz, T c = T( ))`

Changes the size of this vector to `n`. If `n > size( )`, one or more copies of `c` are added to the end of the vector to reach the desired size. If the new size is smaller than the current size, elements are erased from the end to reach the new size.

`size_type size ( ) const`

Returns the number of elements in the vector.

`void swap (vector<T,Alloc>& that)`

Exchanges all the elements in this vector with all the elements in `that`.

## See Also

[vector<bool> class](#), deque in [<deque>](#), list in [<list>](#)

# Name

`vector<bool>` class — Specialized vector of `bool`

# Synopsis

```
template <typename Alloc>
class vector<bool, Alloc> {
public:
    typedef bool const_reference;
    typedef . . . iterator;
    typedef . . . const_iterator;
    typedef . . . size_type;
    typedef . . . difference_type;
    typedef bool value_type;
    typedef Alloc allocator_type;
    typedef . . . pointer;
    typedef . . . const_pointer typedef
std::reverse_iterator<iterator> reverse_iterator;
    typedef std::reverse_iterator<const_iterator>
const_reverse_iterator;

    class reference;
    static void swap(reference x, reference y);
    void flip( );
    . . . // Same as vector<> . . .
};
```

The `vector<bool>` specialization is an interesting beast. It is an attempt to demonstrate how to define a container that uses a *proxy* to represent the elements of the container. The `bool` elements are packed into integers, and the `vector<bool>::reference` type is a proxy that represents a single `bool` element by keeping track of the bit number within the integer and the integer's index in the vector.

However, by using a proxy, `vector<bool>` violates the constraints of a container, so it cannot be used in many situations that call for a standard container. In particular, the `pointer` type cannot point to an element of the container because C++ does not have a type that can point



to a single bit. Many algorithms require the pointer type, and so they cannot work with a `vector<bool>` object.

If you need to use a compact, fixed-size set of bits, use the `bitset` class template. If you need a standard container that contains `bool` elements, use `deque<bool>`.

In addition to the members of the `vector<>` template, `vector<bool>` also defines the following functions:

```
static void swap (reference x, reference y)
```

Swaps two bit values

```
void flip ( )
```

Flips all the bits in the vector

## See Also

[vector class template](#), [vector<bool>::reference class](#),  
`bitset` in [<bitset>](#), `deque` in [<deque>](#)

# Name

vector<bool>::reference class — Bit reference proxy

# Synopsis

```
class reference {
    friend class vector;
    reference( );
public:
    ~reference( );
    operator bool( ) const;
    reference& operator=(const bool x);
    reference& operator=(const reference& x);
    void flip( );
};
```

The reference class represents a single bit in a vector<bool>. The constructor is private, so only vector<bool> can create reference objects. The reference keeps track of the position of an individual bit in a vector<bool>, so you can get, set, or flip the bit. The following are the members of reference:

void **flip** ( )

Flips or toggles the bit, that is, performs the equivalent of \*this = ! \*this

operator **bool** ( ) const

Returns the bit value as a bool

reference& **operator=** (const bool x), reference& **operator=** (const reference& x)

Assigns x to \*this

# See Also

[vector<bool> class](#)

## Appendix A. Compiler Extensions

When writing portable code, you should stick to the standard, but sometimes you have to use compiler-specific extensions of the standard. This appendix lists some of the more interesting extensions in a few compilers. It is not a reference of compiler extensions, but more an illustration of the kinds of extensions that compiler writers choose to implement.

### Borland C++ Builder and Kylix

Borland has several extensions to C++ to support its Rapid Application Development products: C++ Builder (for Microsoft Windows) and Kylix (for Linux). This section presents highlights of the RAD extensions.

#### `__closure`

In C++ Builder, a closure is like a pointer to a member function that has been bound to a specific object. Given a closure, you can call it the way you would call an ordinary function. To declare a closure type or object, use `__closure` as a modifier for the name of a function pointer:

```
typedef int (* __closure MemFunc)(int);
MemFunc func;
struct demo {
    int sqr(int x) { return x * x; }
};
demo d;
func = d.sqr;
int n = func(10); // n = 100
```

#### `__declspec`

The `__declspec` keyword takes an attribute in parentheses and serves as a declaration specifier.

Depending on the attribute, it can be used to modify a function, object, or class. For example, `__declspec(noreturn)` is a function specifier that tells the compiler that the function does not return, which permits additional optimization and error-checking (for example, eliminating statements that follow a call to the `noreturn` function):

```
void __declspec(noreturn) abort( );
```

Other attributes include:

`thread`

A storage-class specifier that declares an object to be local to a thread; that is, each runtime thread has a separate copy of the object.

`dllexport`

A function specifier that tells the linker to export the function name from a dynamic-link library (DLL).

`uuid( string-literal )`

Modifies a class declaration. It associates a universally unique identifier (UUID) with the class, which is required for implementing COM objects in Windows. A class's UUID can be retrieved with the `__uuidof` operator.

`__int64`

The `__int64` type is a 64-bit integer type. In current releases of C++ Builder and Kylix, `long` is 32 bits. A 64-bit integer literal is written with a suffix of `i64` (e.g., `100000000000000i64`).

`__property`

A property is a class member that is used like a data member, but it can have the semantics of a member function. Properties are the foundation for the RAD

features of C++ Builder and Kylix. A property is associated with a reader and writer, which can be data member names or member function names:

```
class TControl {
private:
    int height_;
    void set_height(int h);
    . . .
__published:
    __property int height { read=height_, write=set_height };
};
TControl * ctl = new TControl;
ctl->height = 10;    // Calls ctl->set_height(10)
int h = ctl->height; // Gets ctl->height_
```

## \_\_published

The `__published` access specifier label yields the same accessibility as the `public` keyword, but it also directs the compiler to store additional runtime type information (RTTI) for the published declarations. The RAD features use the RTTI when the user designs an application

## \_\_thread

The `__thread` keyword is a synonym for `__declspec(thread)`.

## \_\_uuidof

The `__uuidof` operator takes an expression as an operand and returns the UUID of the expression's class. The class declares its UUID with `__declspec(uuid( . . . ))`. A class can implement the standard COM member function, `QueryInterface`, with `__uuidof`:

```
class demo {
    virtual HRESULT QueryInterface(const UUID& iid, void** obj)
    {
        if (iid == __uuidof(IUnknown)) {
            *obj = reinterpret_cast<IUnknown*>(this);
            static_cast<IUnknown*>(*obj)->AddRef( );
            return S_OK;
        }
    }
};
```

```
    return E_NOINTERFACE;  
  }  
};
```

# GNU Compiler Collection

The GNU C++ compiler has many extensions to the standard. The most widespread version, 2.95, is mature, stable, but very much outdated with regard to the C++ standard. The new 3.x version hews much closer to the standard, while retaining the familiar GNU extensions. This section presents highlights of only a few of the extensions.

## `__attribute__`

A function can be modified with attributes to tell the compiler about the function. For example, `__attribute__((noreturn))` tells the compiler that the function does not return, which enables the compiler to remove unneeded code, such as statements that follow calls to the function. Some attributes can apply to objects, labels, and types. Several different attributes are supported, such as:

### `always_inline`

Always expand this function inline, even if optimizations are disabled.

### `const`

The function has no side effects and does not depend on any global variables; therefore, the compiler can replace repeated calls to the function with a single call, saving the result.

### `deprecated`

The function, object, or type is deprecated. Using a deprecated entity results in a compiler warning.

### `dllexport`

On Windows, marks the function as exported from a DLL.

## pure

Slightly less strong than `const`, a pure function has no side effects but can depend on global variables. The compiler can optimize away repeated calls when it knows intervening code does not modify any global variables.

## case range

In a `switch` statement, a single case label can specify a range of values:

```
switch(c) {
case 'a' . . . 'z': case 'A' . . . 'Z':
    do_english_letter(c);
    break;
    // . . . Other cases, etc.
}
```

## long long

The `long long` type is an integral type that has at least 64 bits. A `long long` literal is written with a suffix of `LL` (e.g., `1000000000000000LL`).

## Minimum and maximum operators

The operators `<?` and `>?` return the minimum and maximum of their two operands. You can overload these operators the way you would any other operator:

```
template<typename T>
T max3(const T& a, const T& b, const T& c)
{
    return a >? b >? c;
}
```

## typeof

The `typeof` keyword takes an expression as an operand and yields the type of the expression. You can use `typeof`



wherever you can use a type identifier. Templates create a need for `typeof` because when writing the template, you have no way of knowing the return type of an operator or function. For example:

```
template<typename T, typename U>
typeof(T+U) incr(const T& t, const U& u)
{
    return t + u;
}
```

## Microsoft Visual C++

Microsoft's latest C++ compiler has a split personality. It can generate a conventional program following the standard (more or less), or it can produce .NET output using a modified language called Managed C++, which restricts some features and adds others. The following are some highlights of the Managed C++ extensions:

### `__box`

The `__box` operator takes a value object as an argument and returns a managed (`__gc`) object that wraps the value in a managed "box." You can also declare a pointer to a value type with the `__box` specifier so the pointer can store boxed values. The compiler treats boxed values specially and lets you access the members of the value transparently, as though the box were not present. Nonetheless, the box manages the lifetime of the value it contains.

### `__declspec`

The `__declspec` keyword takes a list of attributes in parentheses and serves as a declaration specifier. Depending on the attributes, it can be used to modify a function, object, or class. See [Section A.1](#) for more information.

### `__gc`

The key feature of Managed C++ is that objects are garbage-collected. This means the lifetime and memory are managed automatically by the runtime environment. As long as the object is in use, it remains alive. When the object is no longer used anywhere in the program, it is up for reclamation.

Declare a class using the `__gc` specifier to mark the class as managed. A managed class has numerous restrictions, such as:

- No more than one base class (which must also be managed)
- No unmanaged data members (except POD types)
- No friends
- No operator new or operator delete members
- No user-defined copy constructors
- No const or volatile qualified member functions

Managed objects are created with the global `new` operator. They are freed automatically when they are no longer needed. If the class has a destructor, you can invoke the `delete` operator, which calls the destructor but does not free the memory.

#### `__int64`

The `__int64` type is a 64-bit integer type. In the current releases of Visual C++, `long` is 32 bits. A 64-bit integer literal is written with a suffix of `i64` (e.g., `1000000000000000i64`).

#### `__pin`

Sometimes, a managed application must call an unmanaged library or system call. The `__pin` keyword locks a managed object at its address in memory and prevents the garbage collector from moving the object. The address of the object can safely be passed to the unmanaged function. When the `__pin` pointer goes out of scope, the managed object is no longer pinned and can be moved or reclaimed.

#### `__property`

A property is a pseudo-member that is used like a data member, but it can have the semantics of a member function. Properties are the foundation of the Visual Studio RAD features. A property is associated with a getter and setter, which have the forms `get_name ( )` and `set_name ( )` for the property named *name*:

```
__gc class Control {
private:
    int height_;
public:
    __property int get_height( ) const { return height_; }
    __property void set_height(int h);
};
Control * ctl = new Control;
ctl->height = 10;    // Calls ctl->set_height(10)
int h = ctl->height; // Calls ctl->get_height( )
```

\_\_value

A `__value` class is intended for small objects with short lifetimes, which are allocated on the runtime stack, not the managed heap (the way `__gc` objects are managed). A managed class can have `__value` data members, but not the other way around.

`#using`

The `#using` directive is similar to `#include`, except the included object is not a header or file but a .NET assembly, which contains all the information the compiler needs to use the classes, types, and functions that are defined in the assembly.

## Appendix B. Projects

This appendix lists three interesting C++ projects. Of course, many, many more projects exist. These were chosen because they best reflect the capabilities or future directions of C++. This book's web site (<http://www.tempest-sw.com/cpp/>) has links to these and other projects.

### Blitz++

The Blitz++ project brings high-performance numerical computing to C++. In some respects, it is what `valarray<>` should have been. Blitz++ has powerful array and matrix classes, operators, and functions, with strides, subarrays, and so on. The package is written to minimize the number of unnecessary temporary objects and take advantage of compile-time computation (via template metaprogramming) whenever possible.

One of the key optimizations is that arithmetic operators and mathematical functions involving Blitz++ arrays do not compute values immediately. Instead, they return expression objects. When an expression object is assigned to an array, the expression is computed, storing the results directly in the target of the assignment, without the need for large temporary arrays.

[Example B-1](#) shows a program that demonstrates a few of the features of Blitz++.

#### *Example B-1. Working with matrices*

```
#include <iostream>
#include <ostream>
#include "blitz/array.h"

// Blitz formats output in a manner that is best suited for
subsequent input by a
```

```

// program, not for reading by a human. The print( ) function
uses a format that
// is slightly better for human readers. (Further improvement is
left as an
// exercise.)
template<typename T>
void print(const blitz::Array<T, 3>& a)
{
    std::cout << a.extent(0) << " x " << a.extent(1) <<
        " x " << a.extent(2) << ":\n[";
    for (size_t i = a.lbound(0); i <= a.ubound(0); ++i)
    {
        // Blitz can print a 2-D array well enough, so extract each
2-D plane and
        // print it. Note that plane shares storage with a, so even
if a is large,
        // memory is not wasted by needlessly copying data.
        blitz::Array<T, 2>
            plane(a(i, blitz::Range::all( ), blitz::Range::all( )));
        std::cout << plane;
    }
    std::cout << "]\n";
}

int main( )
{
    // Math with TinyVector and TinyMatrix uses template
metaprogramming to produce
    // fast code, even for complicated operations, such as matrix
multiplication.
    blitz::TinyMatrix<double,2,4> a;
    blitz::TinyMatrix<double,4,3> b;
    blitz::TinyMatrix<double,2,3> c;
    a = 1, 2, 3, 4, // Set elements of a.
        0, 1, -2, -1;
    b = 3.14159; // Set all elements of b.
    c = blitz::product(a, b); // Matrix multiplication
    std::cout << a << b << c;

    // Arrays can have more than two dimensions and offer more
computing power.
    blitz::Array<double,3> d(2,3,4), e(2,3,4), f(2,3,4);
    // Set elements of d to values that depend on the indices.
    d = blitz::firstIndex( ) + blitz::secondIndex( ) +
        blitz::thirdIndex( );
    // Set a subarray of d to 42.
    d(blitz::Range::all( ),blitz::Range(1,2),blitz::Range(2,3))
    = 42;
}

```

```
print(d);

// Call sin for each element of d.
e = sin(d);
print(e);

// If an element of e is negative, set corresponding element of
f to 1; set
// other elements to -1.
f = blitz::where(e < 0, 1.0, -1.0);
print(f);

// Elementwise multiplication
e *= f;
print(e);
// Add all elements of f.
std::cout << blitz::sum(f) << '\n';
}
```

# Boost

The Boost project was started by members of the C++ Standard Committee as a way to explore future directions for the C++ library and to provide a high-quality library that extends what is available in the standard C++ library. Boost has grown to encompass many areas, such as dates and times, regular expressions, lambda expressions, mathematical functions, graphs (the mathematical kind, not the pictorial kind), type traits, multithreaded programming, and more.

Some of the Boost packages have already been proposed as extensions to the standard library. By the time you read this, the Standard Committee might already have accepted a formal Technical Specification for an extension to the standard library. See my web site at <http://www.tempest-sw.com/cpp/> for current information.

This section presents two of the packages in Boost that are part of the proposed library extension: tuples and smart pointers. A *tuple* is a generalization of the standard `pair` class template. Instead of being limited to 2 elements, a tuple can contain up to 10 elements (and the maximum can be extended if necessary). [Example B-2](#) shows one way to use `boost::tuple`.

## *Example B-2. Using a tuple*

```
#include <numeric>
#include <iostream>
#include <ostream>
#include <vector>
#include "boost/tuple/tuple.hpp"

// Store count, sum, and sum of squares. typedef
boost::tuple<std::size_t, double, double> Stats;

// Accumulate statistics.
Stats stats(Stats s, double x)
```



```

{
    ++s.get<0>( );
    s.get<1>( ) += x;
    s.get<2>( ) += x * x;
    return s;
}

int main( )
{
    std::vector<double> v;
    . . . fill v with data . . .

    Stats s = std::accumulate(v.begin( ), v.end( ),
                              boost::make_tuple(0U, 0.0, 0.0), stats);
    std::cout << "count = " << s.get<0>( ) << '\n';
    std::cout << "mean = " << s.get<1>( ) / s.get<0>( ) << '\n';
}

```

Boost has several smart pointer class templates. They solve a number of problems that the standard `auto_ptr<>` class template does not. For example, you cannot store an `auto_ptr<>` object in a standard container, but you can store a `boost::shared_ptr<>` object. Boost has several other smart pointer templates; for the sake of brevity, [Example B-3](#) shows only `shared_ptr<>`.

### *Example B-3. Using shared pointers*

```

#include <algorithm>
#include <iostream>
#include <iterator>
#include <ostream>
#include <string>
#include <vector>
#include "boost/smart_ptr.hpp"

// A company has employees. Each employee can be Exempt or
// NonExempt. Certain
// Exempt employees are Managers, which are distinguished by
// having a group of
// Employees. All the memory for said employees is managed
// automatically by Boost
// shared_ptr<> templates.

class Employee {
public:

```

```

Employee(const std::string& name) : name_(name) {}
virtual ~Employee( ) {}
const std::string name( ) const { return name_; }
virtual void print(std::ostream&);
private:
    const std::string name_;
}; typedef boost::shared_ptr<Employee> employee;

void Employee::print(std::ostream& out)
{
    out << name( ) << '\n';
}

class Exempt : public Employee {
public:
    Exempt(const std::string& name) : Employee(name) {}
};

class NonExempt : public Employee {
public:
    NonExempt(const std::string& name) : Employee(name) {}
};

class Manager : public Exempt {
public:
    Manager(const std::string& name) : Exempt(name) {}
    void add(Employee* e) { group_.push_back(employee(e)); }
    void add(employee e) { group_.push_back(e); }
    virtual void print(std::ostream&);
private:
    std::vector<employee> group_;
};
typedef boost::shared_ptr<Manager> manager;

void Manager::print(std::ostream& out)
{
    out << name( ) << " { ";
    std::copy(group_.begin( ), group_.end( ),
        std::ostream_iterator<employee>(out, ""));
    out << "}\n";
}

// Make it easier to print any kind of employee.
template<typename charT, typename traits>
std::basic_ostream<charT,traits>&
operator<<(std::basic_ostream<charT,traits>&
    out, employee e)
{

```

```
e->print(out);
return out;
}

int main( )
{
    manager ceo(new Manager("I. M. Portant"));
    manager muddle(new Manager("Muddled manager"));
    ceo->add(muddle);
    muddle->add(new Exempt("J. Doe"));
    muddle->add(new NonExempt("J. Dough"));
    ceo->print(std::cout);
}
```

## STLport

The STLport project is a free, open source implementation of the C++ standard library. Although every modern compiler comes with a more-or-less complete standard library, there remain differences, omissions, and errors in most vendor-supplied libraries. If portability across compilers and platforms is a major concern, you might want to use the same library implementation on all platforms.

You might also want to use STLport for its additional features, such as debug mode, which helps detect programming errors. You can also use the library safely in a multithreaded program; for example, reading from a standard container is thread-safe. Of course, writing to any shared data must be properly synchronized. See the Boost project for help with threads and synchronization.

Also included are extensions to the standard library: hashed containers, singly-linked lists, and *ropes* (strings that scale well for very large sizes).

See my web page at <http://www.tempest-sw.com/cpp/> for links and more information.

# Glossary

This brief glossary defines words and phrases that are used often in this book. Also included are terms that the C++ community often uses, even though this book does not use them. Frequent and judicious use of the terms in this glossary will convince all your friends and coworkers of your remarkable C++ skills.

## A-Z

### Algorithm

A generic function, usually one that operates on a sequence specified by a pair of iterators. See [Chapter 10](#).

### Argument

An expression that is used in a function call to initialize a function parameter ([Chapter 3](#)). Can also be a template argument ([Chapter 7](#)).

### cv -qualifier

A `const` or `volatile` qualifier, or both (in any order). See [Chapter 2](#).

### Deprecated

Obsolete. A language or library feature that is deprecated might be removed from the next version of the standard. You should avoid using deprecated features if you can.

### Explicit specialization

The C++ standard term for a template definition that defines a special case of an earlier template, supplying one or more template arguments that apply only to the

explicit specialization. In this book, plain *specialization* means the same thing. See [Chapter 7](#).

## Ill-formed

A source file or program that does not follow the rules of the C++ standard. This book uses informal terms, such as "illegal" or "invalid," to refer to programs that are ill-formed or that invoke undefined behavior. See also *Well-formed*.

## Implementation-defined behavior

Behavior that is well-defined but varies from one implementation to another. The vendor is required to document the precise behavior. For example, the size of an `int` is implementation-defined. It might be 16 bits, 32 bits, 48 bits, or some other size. The C++ standard can mandate limits on the implementation-defined behavior—for example, the minimum size for an `int` is 16 bits.

## Instance or instantiation

Can be a template instance or class instance. A template instance applies a set of template arguments to a template name, creating a specific function or class. Each template argument is matched with a template parameter; arguments can be explicit, or can be deduced for a function template or supplied by default arguments for a class template. See [Chapter 7](#).

A class instance is an object of class type. See *Object*, [Chapter 6](#).

## Iterator

An abstraction of a pointer that can point to an element of a container or other sequence. Most algorithms work with iterators. See [Chapter 10](#).

## Lvalue

An object reference. You can take the address of an lvalue, but you cannot take the address of an rvalue. The lefthand operand of the built-in assignment operator must be a non-const lvalue (hence the l in lvalue). See also *Rvalue*, [Chapter 3](#).

## NRVO

Named Return Value Optimization. A compiler is free to optimize away the call to a copy constructor when an object is returned from a function. Some compilers perform this optimization only when the return expression is the name of an object. See also *RVO*.

## Object

A region of storage with a type. Unlike some object-oriented languages, the term "object" is not restricted to instances of class type. An object of class type can be an lvalue or an rvalue. An object of fundamental or enumerated type is an lvalue. See [Chapter 2](#).

## ODR

One Definition Rule. A program can have any number of declarations of an entity, such as a function, template, or global variable. It must have at most one definition of these entities. See [Chapter 2](#) for details and exceptions to this rule.

## Opaque type

A type whose implementation is hidden. A pimpl is one way to implement an opaque type. See also *Pimpl idiom*.

## Parameter

Function parameter or template parameter. A function parameter is an object that is local to the function and is initialized with the value of a function argument in a function call. See [Chapter 5](#).

A template parameter can be a value, a type, or a template; it gets its "value" from a template argument in a template instantiation. See [Chapter 7](#).

## Parametric polymorphism

Generic programming, such that a function can have a single definition that applies regardless of the type of its parameters. See [Chapter 7](#).

## Pimpl idiom

A way of implementing an opaque type that uses a wrapper class that holds a pointer to a hidden implementation class. The name can mean pointer-to-implementation. See *Opaque type*, [Chapter 6](#).

## POD

Plain Old Data. A C-style type: a fundamental type, enumerated type, a pointer to a POD type, or an array of POD types. A class, structure, or union can be a POD type if its nonstatic data members all have POD types. See [Chapter 6](#).

## Polymorphism

Greek for "having many shapes." Usually means *type polymorphism*. See also *Parametric polymorphism*.

## RAII

Resource Allocation Is Initialization applies to a style of programming that ensures that resources are correctly allocated and deallocated. The `auto_ptr<>` template is an example of RAII.

## Rvalue

A value that does not necessarily have any storage or address. An rvalue of fundamental type can appear only on the right side of an assignment (hence the R in



Rvalue). An lvalue can be implicitly converted to an rvalue, but not the other way around. See [Chapter 3](#).

## RVO

Return Value Optimization. A compiler is free to optimize away the call to a copy constructor when an object is returned from a function. See also *NRVO*.

## SFINAE

Substitution Failure Is Not An Error. When the compiler looks for candidate functions for overload resolution, all function templates with the desired name are initially considered. If the compiler cannot generate a template instance to match the function call's arguments, that function template is not considered. That is, failure to substitute the arguments is not an error.

## Specialization

Defining a template for a specific set of template arguments. A specialization can be *total* (specifying all template arguments) or *partial* (specifying only some of the template arguments). Only class templates can have partial specializations. See [Chapter 7](#).

## Type polymorphism

An object can have a pointer or reference to a base class as its type. At runtime, it can take a pointer or reference to a derived class as its value. At runtime, calls to virtual functions bind to the functions in the derived class. The derived-class type is known as the object's *dynamic type*; the declared type of the object is the *static type*. See [Chapter 6](#).

## Undefined behavior

Anything can happen. Maybe an exception will be thrown; maybe a signal will be raised; maybe the

application will crash and burn. A particular implementation might have well-defined behavior, but different implementations are free to choose different behaviors. For example, dereferencing a null pointer produces undefined behavior. A desktop application might raise a segmentation fault signal. An embedded control for a toy robot might reset itself. An operating system kernel might panic and shut down, possibly corrupting filesystems and databases. Anything can happen.

### Well-formed

A program that obeys the syntactic rules, the diagnosable semantic rules, and the one-definition rule of the C++ standard. Note that the compiler is not required to diagnose all errors, so you can have a program that is well-formed but still incorrect. See also *Ill-formed*.

## Colophon

Our look is the result of reader comments, our own experimentation, and feedback from distribution channels. Distinctive covers complement our distinctive approach to technical topics, breathing personality and life into potentially dry subjects.

The animal on the cover of *C++ in a Nutshell* is an Eastern chipmunk, a striped ground squirrel found mostly in eastern North America. Eastern chipmunks have five dark and two light stripes on their backs, extending from head to rump, and two stripes on their long, bushy tails. They are distinguished from other ground squirrels by the white stripes above and below their eyes. The coloration of chipmunks throughout North America varies but is quite uniform within regions.

Chipmunks often make their homes in sparse forests or farms, where they can build the entrances to their lodges in stone walls, broken trees, or thick underbrush. A lodge consists of a maze of tunnels leading to a large, leaf-lined nest. Chipmunks spend most of the daylight hours outdoors but head for their lodges before nightfall. Although they are excellent climbers, chipmunks live primarily on the ground.

Chipmunks eat nuts, seeds, insects, and occasionally birds' eggs. Like all ground squirrels, they have large cheek pouches, sometimes extending as far back as their shoulders, in which they can store food. They collect and store nuts and seeds through the summer and fall. When the weather starts to get cool, all the chipmunks in a region suddenly disappear into their lodges, where they begin hibernation. On warm winter days one can often see chipmunk pawprints in the snow, as they will sometimes

wake up and leave their lodges for brief periods when the temperature rises.

Mating season for Eastern chipmunks is mid-March to early April. The gestation period is 31 days, after which a litter of three to six is born. Baby chipmunks leave the lodge after one month and are mature by July.

The chipmunk most likely got its name from the noise it makes, which sounds like a loud "cheep." You can occasionally see a chipmunk hanging upside down from a tree branch "cheeping" its call.

Matt Hutchinson was the production editor and copyeditor for *C++ in a Nutshell*. Sarah Sherman and Claire Cloutier provided quality control. Julie Hawks wrote the index. Derek Di Matteo and Mary Brady provided production assistance

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